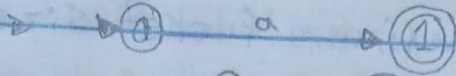
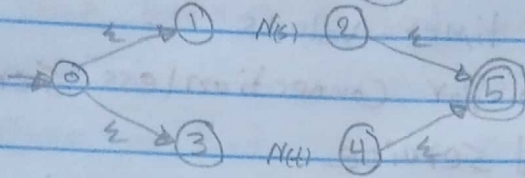
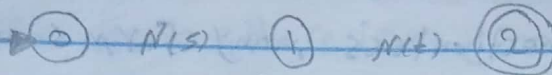
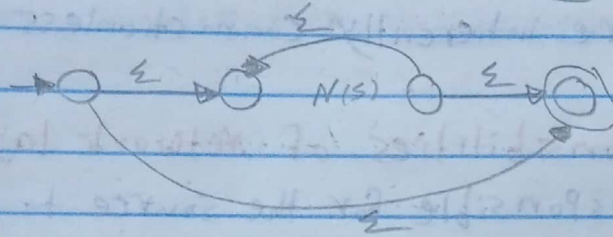


Convert NFA to DFA

RE  $\rightarrow$  NFA use basic $R = a$  $R = s|t$  $R = st$ 

نفاية أولوية ثابتة

 $R = s^*$ 

ch1

1 - What is the difference between a Compiler and an interpreter?

A Compiler is a program that can read a program in one language - the source language - and translate it into an equivalent program in another language.

Interpreter directly executes the operations specified in the source program on inputs supplied by the user.

2 - What are the advantages of (a) a compiler over an interpreter (b) an interpreter over a compiler?

(a) The machine language target program produced by a compiler is usually much faster than an interpreter at mapping inputs to outputs.

(b) An interpreter can usually give better error diagnostics than a compiler, because it executes the source program statement by statement.



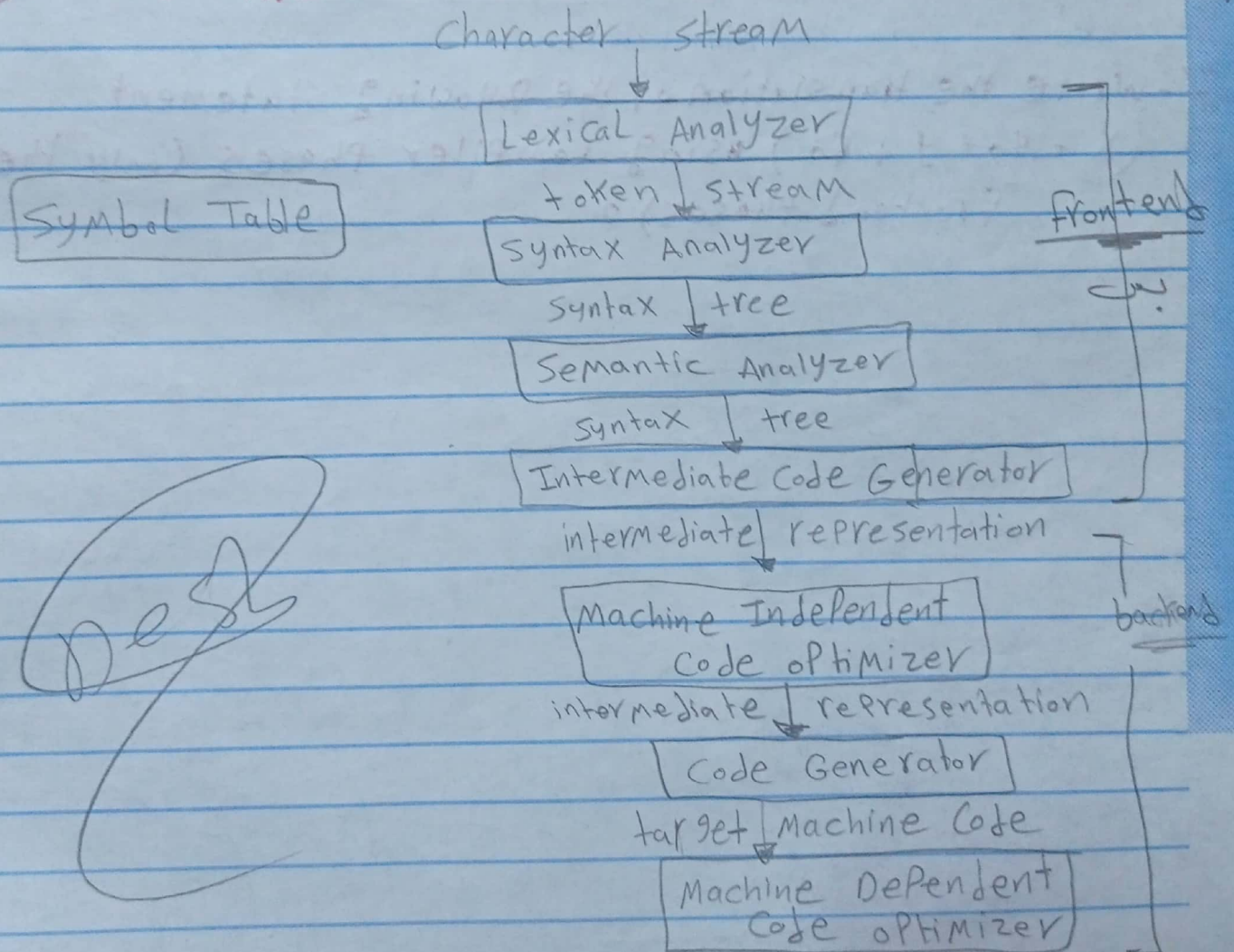
3- What advantages are there to a language processing system in which the compiler produces assembly language rather than machine language?

The compiler may produce an assembly language program as its output, because assembly language is easier to produce as output and is easier to debug.

4- Describe some of the tasks that an assembler needs.

It translates from the assembler language to machine code. This machine code is relocatable.

5- Describe the structure of compiler phases in brief.





6- Describe the language processing system in brief.

Source Program

PreProcessor

modified Source Program

Compiler

target assembly program

Assembler

relocatable machine code

Linker / Loader

relocatable object files

7- Write the translation of the following statement

$(X = 40 * Y + 50)$  using Compiler Phases (only the Frontend phases)?

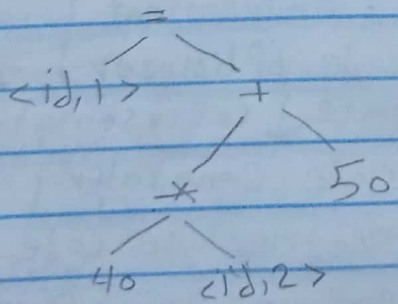
$X = 40 * Y + 50$

أي تعبير نقس الذاكرة  
total = 125 \* num + 180

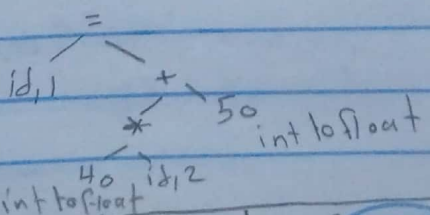
lexical analyzer

$\langle id, 1 \rangle \langle = \rangle \langle 40 \rangle \langle * \rangle \langle id, 2 \rangle \langle + \rangle \langle 50 \rangle$

Syntax Analyzer



Semantic analyzer



Intermediate code generator

$t_1 = \text{int to float}(50)$

$t_2 = \text{int to float}(40) * id, 2$

$t_3 = t_2 + t_1$

$id, 1 = t_3$



1. Write a simple intermediate code for the following Java segment?

```

int i; int j; float [100] a; float v; float x;
while (true) {
    do i = i + 1; while (a[i] < v);
    do j = j - 1; while (a[j] > v);
    if (i >= j) break;
    x = a[i]; a[i] = a[j]; a[j] = x;
}

```

```

1: i = i + 1
2: t1 = a[i]
3: if t1 < v goto 1
4: j = j - 1
5: t2 = a[j]
6: if t2 > v goto 4
7: if false i >= j goto 9
8: goto 14
9: x = a[i]
10: t3 = a[j]
11: a[i] = t3
12: a[j] = x
13: goto 1
14:

```

2.  $\text{expr} \rightarrow \text{expr} + \text{term} \mid \text{expr} - \text{term} \mid \text{term}$   
 $\text{term} \rightarrow \text{term} * \text{factor} \mid \text{term} / \text{factor} \mid \text{factor}$   
 $\text{factor} \rightarrow \text{digit} \mid (\text{expr})$   
 $\text{digit} \rightarrow 0 \sim 9$

3.  $S \rightarrow SS + \mid SS * \mid a$

(1) Show how the string  $aa + a^*$  can be generated by this

(2) Construct a Parse tree for this string

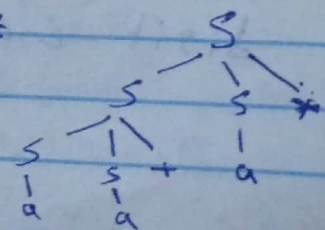
(3) What language does this grammar generate?

(1)  $S \Rightarrow SS * \Rightarrow SS + S * \Rightarrow aa + a *$

(2)

(3)  $L = \{ \text{Postfix expression consisting}$

of digits, plus and multiple signs  $\}$





(4)  $0S \rightarrow 0S1 \mid 01$

(2)  $S \rightarrow +SS \mid -SS \mid a$

(3)  $S \rightarrow S(S)S \mid \epsilon$

(4)  $S \rightarrow aSbS \mid bSas \mid \epsilon$

(5)  $S \rightarrow a \mid S+SS \mid SS \mid S* \mid (S)$

1.  $L = \{0^n 1^n \mid n \geq 1\}$

2.  $L = \{\text{Prefix expression consisting of Plus and minus signs}\}$

3.  $L = \{\text{matched brackets of arbitrary arrangement and nesting includes } \epsilon\}$

4.  $L = \{\text{String has the same amount of a \& b, include } \epsilon\}$

5.  $L = \{\text{Regular expressions used to describe regular language}\}$

4. Look book pg 2-14, 2-15 Infix to Postfix

5. Construct a syntax-directed translation scheme that translates arithmetic expressions from infix notation into prefix notation in which an operator appears before its operands -  $xy$  is the prefix notation for  $x.y$  Give annotated parse tree for the inputs  $9-5+2$  and  $9-5*2$

Productions  $\text{EXPR} \rightarrow \text{Grammar}$

$\text{EXPR} \rightarrow \text{EXPR} + \text{term}$   
 $\mid \text{EXPR} - \text{term}$   
 $\mid \text{term}$

$\text{term} \rightarrow \text{term} * \text{factor}$   
 $\mid \text{term} / \text{factor}$   
 $\mid \text{factor}$

$\text{factor} \rightarrow \text{digit} \mid (\text{EXPR})$

$\text{digit} \rightarrow 0 \mid \dots \mid 9$

translation schemes  
 $\text{EXPR} \rightarrow \{\text{Print}("+")\} \text{EXPR term}$   
 $\mid \{\text{Print}("-")\} \text{EXPR term}$   
 $\mid \text{term}$

$\text{term} \rightarrow \{\text{Print}("*")\} \text{EXPR} * \text{factor}$   
 $\mid \{\text{Print}("/")\} \text{EXPR} / \text{factor}$   
 $\mid \text{factor}$

$\text{factor} \rightarrow \text{digit} \{\text{Print}(\text{"digit"})\}$   
 $\mid (\text{EXPR})$

$\text{digit} \rightarrow 0 \mid \dots \mid 9$

Prefix  
 اشارة قبل

Infix  
 في الوسط

Postfix  
 بعد

Production حسب نوع  
 Prefix  
 Postfix  
 infix  
 اشارة قبل  
 اشارة بعد  
 اشارة في الوسط



```

graph TD
    EXPR1[EXPR] --- EXPR2[EXPR]
    EXPR1 --- PLUS1[+]
    EXPR1 --- TERM1[term]
    EXPR2 --- TERM2[term]
    EXPR2 --- TERM3[term]
    TERM2 --- FACTOR1[factor]
    FACTOR1 --- DIGIT1[digit]
    DIGIT1 --- 9[9]
    DIGIT1 --- PRINT9["{Print(9)}"]
    TERM3 --- FACTOR2[factor]
    FACTOR2 --- DIGIT2[digit]
    DIGIT2 --- 5[5]
    DIGIT2 --- PRINT5["{Print(5)}"]
    
```

**+ - 9 5 2**

## Predictive Parsing

```

Match(For);      Match('(');      optexpr();      Match(';');
optexpr();      Match(';');      optexpr();      Match(')');
stmt()

```

if (lookAhead == t)

```
else report("syntax error"); }
```

```
if (lookAhead == expr) match(expr); }
```

Case expr: match (expr); match(';'); break;

Case for :  $\text{match}(\text{for})$  كحل القاعة آخره break;

```
default : report('syntax error'); }
```



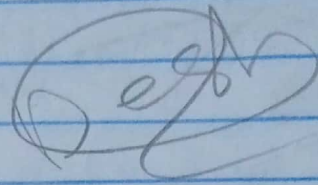
7- left Recursion  $expr \rightarrow expr + term$   
 $A \rightarrow A \alpha | B$

8- Construct recursive-descent Parsers, starting with the following grammars

(a)  $S \rightarrow +SS | -SS | a$

(b)  $S \rightarrow S(S)S | \epsilon$

(c)  $S \rightarrow 0S1 | 01$



void Match (Terminal t) {

if (lookahead == t) { lookahead = nextTerminal(); }

else { report("syntax error"); }

(a) void S() { switch (lookahead) {

case "+" Match("+"); S(); S(); break;

case "-" Match("-"); S(); S(); break;

case "a" Match("a"); break;

default report("syntax error"); }

(b) void S() {

if (lookahead == "(") {

Match("("); S(); Match(")"); S(); }

(c) void S() { switch (lookahead) {

case "0" Match("0"); S(); Match("1"); break;

case "1" // match(epsilon); break;

default report("syntax error"); }

9- in book Pg 2.23 & 2.24 Ref 2.5.2  $A \rightarrow XR$

الكود يقى للحل مشكلة الوب 2.5.2

$R \rightarrow \epsilon | R | PR | \epsilon$

- void expr() { term(); rest(); }

- void term() { if (lookahead is a digit) {

t = lookahead; Match(lookahead); Print(t); }

else { report("syntax error"); }

Debn

Mosta Abdelhakem



\* Fig 2.23 `void rest() { if (lookahead == "+") {  
    match("+"); term(); Print('+'); rest(); }  
    else if (lookahead == "-") {  
        match("-"); term(); Print('-'); rest(); }  
    else { } }`

Def `-void rest() { while (true) {  
    if (lookahead == '+') {  
        match('+'); term(); Print('+'); continue; }  
    else if (lookahead == '-') {  
        match('-'); term(); Print('-'); continue; }  
    break; } }`

10-Write a PseudoCode of a Program that takes input from File and remove the new line and multiple spaces?

1 Fig 2.20 or Skip White Space

`for ( ; ; peek = next input character ) {  
    if (peek is a blank or a tab) do nothing ;  
    else if (peek is a new line) line = line + 1 ;  
    else break ; }`

2 Grouping digits into integers

Fig 2.30 `if (peek holds a digit) { int v = 0  
    do { v = v * 10 + integer value of digit peek ;  
        peek = next input character ;  
    } while (peek holds a digit) ;  
    return token(num, v) }`

3 Recognizing Key Words and Identifiers

Fig 2.31 `if (peek holds a letter) {  
    Collect letter or digits into a buffer b ;  
    S = String formed from the characters in b ;  
    W = token returned by Words.get(s) ;  
    if (W is not Null) return W ;  
    else Enter the Key Value Pair (S, <id, s>) into Words ;  
    return token(id, S) ;  
}`



الكود السقي مطاى عمار فى سعاد حسب السؤال المطلوب نظامه كالتالى

Page: ce... number... Word - Comend - relation - number... float

**Public Token Scan()**

```

//space
for (; ; Peek = (char) system.in.read()) {
    if (Peek == ' ' || Peek == '\t') continue;
    else if (Peek == '\n') line = line + 1;
    else break;
}

//number
if (character.isDigit(Peek)) { int v = 0;
    do { v = v * 10 + character.digit(Peek, 10);
        Peek = (char) System.in.read();
    } while (character.isDigit(Peek));
    return new Num(v);
}

//Word
if (character.isLetter(Peek)) {
    String Buffer b = new StringBuffer();
    do { b.append(Peek);
        Peek = (char) System.in.read();
    } while (character.isLetter(Peek));
    String s = b.toString();
    Word w = (Word) words.get(s);
    if (w != null) return w;
    else w = new Word(Tag.ID, s);
    words.put(s, w);
    return w;
}

//Comment Single line Comment
if (Peek == '/') {
    Peek = (char) stream.read();
    if (Peek == '/') {
        for (; ; Peek = (char) stream.read()) {
            if (Peek == '\n') break;
        }
    }
} else throw new SyntaxException();

```



\*Comment / x x /

if (Peek == '/') {

Peek = (char) stream.read();

if (Peek == 'x') {

char PrevPeek = ' ';

for (; PrevPeek = Peek, Peek = (char) stream.read(); {

if (PrevPeek == 'x' &amp;&amp; Peek == '/') {

break; } }

} else Throw New SyntaxException();

// Relation sign {&lt;, &lt;=, ==, !=, &gt;, &gt;=}

if ("&lt;=!&gt;".indexOf(Peek) &gt; -1) {

StringBuffer b = new StringBuffer();

b.append(Peek);

Peek = (char) stream.read();

if (Peek == '=') b.append(Peek); }

return new Rel(b.toString()); }

// number no type Sensitive float

if (character.isDigit(Peek) || Peek == '.') {

Boolean isDotExist = false;

StringBuffer b = new StringBuffer();

do { if (Peek == '.') { isDotExist = true; }

b.append(Peek);

Peek = (char) stream.read();

} while (isDotExist == true ? character.isDigit(Peek):

character.isDigit(Peek) || Peek == '.');

return new Num(new Float(b.toString())); }

Token t = new Token(Peek);

Peek = ' '; // Peek = blank;

return t; }

آی اسامی عزیز کد احسن  
 حس الطوب کونین سواد  
 emad married 7 soad Code

Mostafa Abdelharem

01121811782

2022

Deshay

monsy

2022