

Heaven's Light is Our Guide
Computer Science & Engineering
Rajshahi University of Engineering & Technology

Lab Manual

Module- 01

Course Title: Sessional based on CSE 1201

Course No. : CSE 1202

Experiment No. 1**Name of the Experiment:** Complexity of Algorithms**Duration:** 1 cycle**Background Study:** Chapter 1-2 (Theory and Problems of Data Structures Written by Seymour Lipschutz)**Definition of Big-O Notation:** $f(n)$ is $O(g(n))$ if and only if there exist two constants c and n_0 such that $|f(n)| < c|g(n)|$ for all $n > n_0$. $f(n)$ will normally represent the computing time of some algorithm. When we say that the computing time of an algorithm is $O(g(n))$ we mean that its execution takes no more than a constant times $g(n)$. n is a parameter which characterizes the inputs and/or outputs. For example n might be the number of inputs or the number of outputs or their sum or the magnitude of one of them.**Example**

Let $f(n) = 3n+2$
 Question: Find the Big-O notation of $f(n)$.
 We have, $f(n)=3n+2$
 We can write, $f(n) \leq 3n + n$ if $n \geq 2$
 $f(n) \leq 4n$,
 From definition, $n_0 = 2$ and $c = 4$, $g(n) = n$
 $f(n) \leq cg(n)$
 $f(n)$ is $O(g(n))$ or $f(n)$ is $O(n)$

Problem I: Find the Complexity of a Loop.**Algorithm1.1:**

1. Repeat for $K = 1$ to n by 1
2. Write: K
[End of Step 1 loop]
3. Exit

Complexity: the computing time of the loop, $f(n) = n$, So the complexity of the above algorithm is $O(n)$. Here $f(n) \leq cg(n)$, $c = 2$, $n_0 = 0$, $g(n) = n$.**Complexity Table:**

n	f(n) [from Program, Count Statement]	cg(n) [Theoretical]
10		
20		

Graph: Draw a Graph.**Problem II:** Find the Complexity of the following Program.**Algorithm1.2:**

1. Repeat for $K = 1$ to n by 1
2. Repeat for $L = 1$ to n by 1
3. Write: L
[End of Step 2 loop]
4. Write: K
[End of Step 1 loop]
5. Exit

Complexity: the computing time of the loop, $f(n) = n^2$, So the complexity of the above algorithm is $O(n^2)$. (Let $c = 2$)**Complexity Table:**

n	f(n) [from Program, Count Statement]	cg(n) [Theoretical]
10		
20		

Graph: Draw a Graph.

Problem III: Find the Complexity of the elementary Sort algorithm.

Algorithm1.3: (Given a nonempty array A with n numerical values. This algorithm sorts the values)

1. Repeat for i = 2 to n by 1
2. Repeat for k = i to 1 by -1
3. If $A[k] < A[k-1]$ then:
 - Swap (A[k], A[k-1])
 - [End of If Structure]
 - [End of Step 2 loop]
 - [End of Step 1 loop]
4. Exit

Complexity: The complexity of the above algorithm is $O(n^2)$. (Let $c = 2$)

Complexity Table:

n	f(n) [from Program, Count Statement]	cg(n) [Theoretical]
10		
20		

Graph: Draw a Graph.

Problem IV: Find the largest element in Array.

Algorithm1.4: (Given a nonempty array A with n numerical values. This algorithm finds the location LOC and the value MAX of the largest element of A)

1. Set $K:=1$, $LOC:=1$ and $MAX:=A[1]$
2. Repeat steps 3 and 4 while $K \leq n$
3. IF $MAX < A[K]$ then:
 - Set $LOC:=K$ and $MAX:=A[K]$.
 - [End of If structure]
4. $K:=K+1$.
 - [End of step 2 loop]
5. Write: LOC, MAX.
6. Exit

Complexity: Not Required

Problem V: Linear Search.

Algorithm1.5: (Given a nonempty array A with n numerical values and a specific x of information is given. This algorithm finds the location LOC of x in the array A or Sets $LOC=-1$)

1. Set $K:=1$, $LOC:=-1$
2. Repeat steps 3 and 4 while $LOC = -1$ and $K \leq n$
3. IF $x = A[K]$ then:
 - Set $LOC:=K$.
 - [End of If structure]
4. $K:=K+1$.
 - [End of step 2 loop]
5. If $LOC = -1$ then: Write: x is not in the array A.
Else: Write: LOC is the location of x
[End of If structure]
6. Exit

Complexity: The worst case complexity is $C(n) = n$ and the average case complexity is $C(n) = (n+1)/2$.

MORE PROBLEMS

1. Programming Problems of Chapter 1 and 2 of "Data Structures" by Seymour Lipschutz.

LAB REPORT: You have to submit all assigned problems in next lab.