Computer Labs: Processes 2º MIEIC

Pedro F. Souto (pfs@fe.up.pt)

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Minix 3 Notes: driver_receive() is not Polling

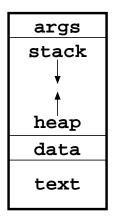
```
5: while(1) { /* You may want to use a different condition
 6:
        /* Get a request message. */
 7:
        if ( driver_receive(ANY, &msq, &ipc_status) != 0 ) {
 8:
            printf("driver_receive failed with: %d", r);
 9:
            continue;
10:
11:
        if (is_ipc_notify(ipc_status)) { /* received notificat
12:
            switch (_ENDPOINT_P(msq.m_source)) {
13:
            case HARDWARE: /* hardware interrupt notification
14:
                if (msq.NOTIFY_ARG & irg_set) { /* subscribed
15:
                    ... /* process it */
16:
17:
                break;
18:
            default:
19:
                break; /* no other notifications expected: do
20:
21:
        } else { /* received a standard message, not a notific
```

driver_receive() what if the process' "IPC queue" is empty?

(Sequential) Process

Abstracts a running program

int main(int argc, char *argv[], char* envp[])}



 0×0

args Command line args and environment variables de ambiente.

stack Activation frames/records corresponding to function calls

heap Dynamically allocated memory (e.g using malloc)

data Memory allocated statically (by the compiler) (e.g. the "Hello, World!" string)

text Machine instructions

Minix is a multiprogramming OS

```
$ ps ax | more
  PID TTY TIME CMD
 (-4) ? 0:46 idle
 (-3) ? 0:00 clock
(-2) ? 0:00 system
 (-1) ? 0:00 kernel
    5 ? 0:00 pm
7 ? 0:01 vfs
4 ? 0:00 rs
8 ? 0:00 memory
9 ? 0:00 log
   10 ? 0:00 tty
3 ? 0:00 ds
        ? 0:00 ds
   12 ? 0:00 vm
   13 ? 0:00 pfs
    6 ? 0:00 sched
    1 ? 0:00 init
   -- more (43 in all)
```

And so are Linux and all Windows OSs since XP (at least)

OS support multiple processes (multiprogramming) for reasons of **efficiency**



Multiprogramming and Efficiency

Problem Processes need to access to I/O devices (monitor, keyboard, mouse, disk, network ...)

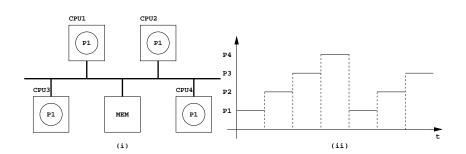
Parameter	Time
CPU cycle	1 ns (1 GHz)
Cache access	~ 2ns
Memory access	~ 10 ns
Disk access	~10 ms

Solution while a process waits for an I/O operation to complete, the OS can allocate the processor to another process:

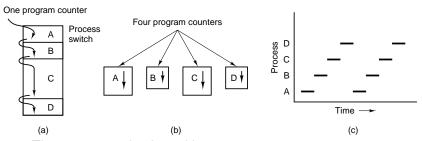
▶ Upon completion of the I/O operation, the I/O device can generate an interrupt

Multi-process Execution (1/2)

- In a multiprocessor/multicore system (i), each processor/core can execute a different process
- In a uniprocessor system (ii), the OS allocates the processor to the different processes (the processor is a resource shared by the different processes): pseudo-parallelism.



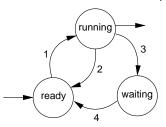
Multi-process Execution (2/2)



- The processor is shared by 4 processes;
- The OS creates the illusion that each process executes in its own CPU, i.e. that each process executes in its virtual CPU.
 - When the OS removes the CPU from a process, i.e. the OS preempts the CPU from a process, it must save the CPU state so that it can resume the process later, as if the CPU had not been taken away.
 - This is akin, but not exactly like, a function call, in which the CPU state of the caller is saved in the stack.

States of a Process

▶ In its lifetime, a process can be in 1 of 3 states:



- 1. The OS allocates a CPU to the process;
- The OS allocates the CPU to another process;
- The process blocks waiting for some event (usually I/O)
- An event the process was waiting for occurs

Running the CPU executes the process's instructions a executar as instruções do processo;

Waiting the process is waiting for an event (usually the end of an I/O operation)

Ready the process is waiting for the OS to allocate it a CPU, which is executing instructions of another process

Minix 3 Notes: driver_receive() is not Polling

driver_receive() is a blocking call.

If the process's "IPC queue" is empty:

- ► The OS will move the process to the WAIT state
- ► The process' state will be changed to READY, only when a message (or notification) is sent to the process

```
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               break;
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19:
               break; /* no other notifications expected: do
                                      20:
```

Further Reading

Sections 2, 2.1
 Andrew Tanenbaum, Modern Operating Systems, 2nd Ed.