

The Art of Computer Programming

Explicit Solutions

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October 10, 2022

§ 1.2.4

42. re: Lower-bound on internal path length of extended binary trees (p. 400)

(a) Prove that $\sum_{k=1}^n a_k = na_n - \sum_{k=1}^{n-1} k(a_{k+1} - a_k)$, if $n > 0$.

Solution:

$$\begin{aligned}\sum_{k=1}^n a_k &= na_n - \sum_{k=1}^{n-1} k(a_{k+1} - a_k) \\&= na_n + \sum_{k=1}^{n-1} ka_k - \sum_{k=1}^{n-1} ka_{k+1} \\&= na_n + (a_1 + 2a_2 + \cdots + (n-1)a_{n-1}) \\&\quad - (a_2 + 2a_3 + \cdots + (n-2)a_{n-1} + (n-1)a_n) \\&= na_n - (n-1)a_n + (a_1 + 2a_2 + \cdots + (n-1)a_{n-1}) \\&\quad - (a_2 + 2a_3 + \cdots + (n-2)a_{n-1}) \\&= a_n + \sum_{k=1}^{n-1} ka_k - \sum_{k=1}^{n-2} ka_{k+1} \\&= a_n + \sum_{k=0}^{n-2} (k+1)a_{k+1} - \sum_{k=1}^{n-2} ka_{k+1} \\&= a_n + a_1 + \sum_{k=1}^{n-2} (k+1)a_{k+1} - \sum_{k=1}^{n-2} ka_{k+1} \\&= a_n + a_1 + \sum_{k=1}^{n-2} (k+1)a_{k+1} - ka_{k+1} \\&= a_n + a_1 + \sum_{k=1}^{n-2} a_{k+1} \\&= a_n + a_1 + \sum_{k=2}^{n-1} a_k \\&= \sum_{k=1}^n a_k\end{aligned}$$

- (b) The preceding formula is useful for evaluating certain sums involving the floor function. Prove that, if b is an integer ≥ 2 ,

$$\sum_{k=1}^n \lfloor \log_b k \rfloor = (n+1) \lfloor \log_b n \rfloor - (b^{\lfloor \log_b n \rfloor + 1} - b)/(b-1).$$

Aside:

$$(1) \quad f = \lfloor \log_b n \rfloor; b \geq 2$$

$$\begin{aligned} S_{n-1} &= \sum_{k=1}^{f-1} b^k \\ &= b^{\lfloor \log_b 1 \rfloor} - b^{\lfloor \log_b n \rfloor} + \sum_{k=1}^{f-1} b^{k+1} \\ &= 1 - b^{\lfloor \log_b n \rfloor} + b \sum_{k=1}^{f-1} b^k \\ &= 1 - b^{\lfloor \log_b n \rfloor} + b S_{n-1} \end{aligned}$$

$$\begin{aligned} (1-b)S_{n-1} &= 1 - b^{\lfloor \log_b n \rfloor} \\ \Rightarrow S_{n-1} &= \frac{b^{\lfloor \log_b n \rfloor} - 1}{b-1} \end{aligned}$$

$$\begin{aligned} (2) \quad \sum_1^{f-1} &= \sum_j^{f-1} + \sum_1^{j-1} \\ \Rightarrow \sum_j^{f-1} &= \sum_1^{f-1} - \sum_1^{j-1} \end{aligned}$$

$$(3) \quad f = \lfloor \log_b n \rfloor; b \geq 2$$

$$\begin{aligned} S_n &= \sum_{k=1}^{n-1} k \quad [k+1 \text{ is a power of } b] \\ &= \sum_{k=1}^f (b^k - 1) \\ &= \sum_{k=1}^f b^k - \sum_{k=1}^f 1 \\ &= \sum_{k=1}^f b^k - f \\ &= \left[f b^f - \sum_{k=1}^{f-1} k(b^{k+1} - b^k) \right] - f \\ &= f(b^f - 1) - \sum_{k=1}^{f-1} k(b^{k+1} - b^k) \\ &= f(b^f - 1) - (b-1) \sum_{k=1}^{f-1} k b^k \\ &= f(b^f - 1) - (b-1) \sum_{j=1}^{f-1} \sum_{k=j}^{f-1} b^k \end{aligned}$$

(a)

$$\begin{aligned}
&= f(b^f - 1) - (b - 1) \sum_{j=1}^{f-1} \left(\sum_{k=1}^{f-1} b^k - \sum_{k=1}^{j-1} b^k \right) \quad (2) \\
&= f(b^f - 1) - (b - 1) \sum_{j=1}^{f-1} \left(\frac{b^f - 1}{b - 1} - \frac{b^j - 1}{b - 1} \right) \\
&= f(b^f - 1) - (b - 1) \sum_{j=1}^{f-1} \left(\frac{b^f}{b - 1} - \frac{b^j}{b - 1} \right) \\
&= f(b^f - 1) - (b - 1) \left[\frac{1}{b - 1} \left(\sum_{j=1}^{f-1} b^f - \sum_{j=1}^{f-1} b^j \right) \right] \\
&= f(b^f - 1) - \left(\sum_{j=1}^{f-1} b^f - \sum_{j=1}^{f-1} b^j \right) \\
&= f(b^f - 1) - \left(b^f(f - 1) - \sum_{j=1}^{f-1} b^j \right) \\
&= f(b^f - 1) - \left(b^f(f - 1) - \left(\frac{b^f - 1}{b - 1} \right) \right) \\
&= f(b^f - 1) - b^f(f - 1) + \left(\frac{b^f - 1}{b - 1} \right) \\
&= fb^f - f - fb^f + b^f + \left(\frac{b^f - 1}{b - 1} \right) \\
&= b^f - f + \left(\frac{b^f - 1}{b - 1} \right) \\
(b - 1)S_n &= (b - 1) \left(b^f - f \right) + \left(b^f - 1 \right) \\
&= b^{f+1} - fb + f - b^f + b^f - 1 \\
&= b^{f+1} - 1 - fb + f \\
&= b^{f+1} - 1 - f(b - 1) \\
S_n &= \left(\frac{b^{f+1} - 1}{b - 1} \right) - f \\
\Rightarrow S_n &= \frac{b^{\lfloor \log_b n \rfloor + 1} - 1}{b - 1} - \lfloor \log_b n \rfloor \quad **
\end{aligned}$$

e. o. a.

Solution:

$$\begin{aligned}
\sum_{k=1}^n \lfloor \log_b k \rfloor &= (n + 1) \lfloor \log_b n \rfloor - \frac{b^{\lfloor \log_b n \rfloor + 1} - b}{b - 1} \\
&= n \lfloor \log_b n \rfloor - \frac{b^{\lfloor \log_b n \rfloor + 1} - b}{b - 1} + \lfloor \log_b n \rfloor \\
&= \dots
\end{aligned}$$

§ 6.2.1

25. re: ... (p. 425)

Suppose that a binary tree has a_k internal nodes and b_k external nodes on level k , for $k = 0, 1, \dots$ (The root is at level zero.) Thus in Fig. 8 we have $(a_0, a_1, \dots, a_5) = (1, 2, 4, 4, 1, 0)$ and $(b_0, b_1, \dots, b_5) = (0, 0, 0, 4, 7, 2)$. *needs figure*