Arithmetic expression geometry

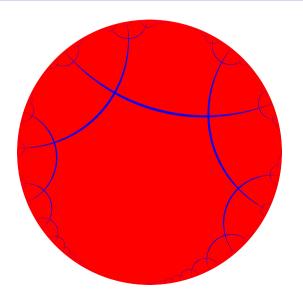
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The first glimpse



The beginning point

The famous example of word2vec

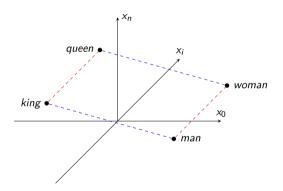


Figure: regulairty of word2vec

The case of numbers

$$(\alpha+1)\times 2\neq \alpha\times 2+1$$

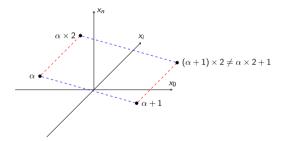
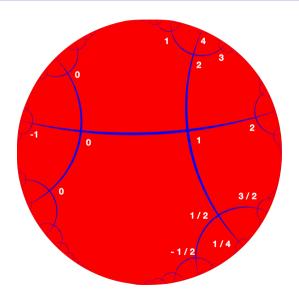


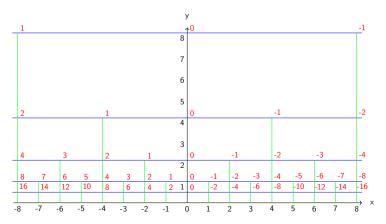
Figure: contradiction of numbers in Euclidean space

One arrangement in hyperbolic space



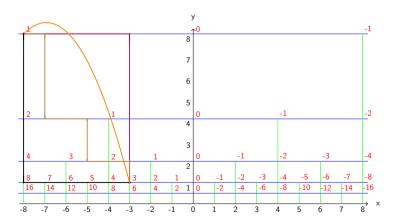
Another arrangement in hyperbolic space

$$a=-\frac{x}{y}$$



Encoding threadlike expressions as paths

• black line $1 \times 8 - 5 = 3$



The flow equation

Suppose we have a base point a_0 , and we step a small distance away from a_0 . Addition first

$$a_{\delta} = (a_0 + \mu \epsilon \cos \theta) e^{\lambda \epsilon \sin \theta}$$

Multiplication first

$$a_{\delta} = a_0 e^{\lambda \epsilon \sin \theta} + \mu \epsilon \cos \theta$$

Both formula can be simplified to the same result:

$$a_{\delta} = a_0 + \epsilon (a_0 \lambda \sin \theta + \mu \cos \theta)$$

Then, we have the following equation:

$$\frac{1}{\delta}(a_{\delta}-a_{0})=\frac{\epsilon}{\delta}(\mu\cos\theta+x_{0}\lambda\sin\theta)$$

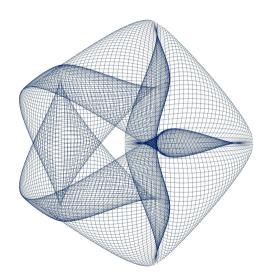
When both δ and ϵ are towards zero, we get da/dt, and hence

$$\frac{da}{dt} = u(\mu\cos\theta + a\lambda\sin\theta)$$

Or, we can change it to another form

$$\frac{da}{dc} = \mu \cos \theta + a\lambda \sin \theta \tag{1}$$

Fundamental problems



What is an arithmetic expression?

Giving an arithmetic expression, we can parse it into a syntax tree. For example, the expression

$$(((((1 \times 2) \times 2) - 1) \times (2 + 1)) - 6)$$
 (2)

and the parsed syntax tree

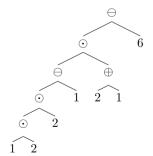


Figure: a tree representation of an arithmetic expression

A definition of arithmetic expression

Definition

An arithmetic expression a over \mathbb{Q} is a structure given by the following production rules:

$$a \leftarrow x$$

$$a \leftarrow (a+a)$$

$$a \leftarrow (a-a)$$

$$a \leftarrow (a \times a)$$

$$a \leftarrow (a \div a)$$

$$(3)$$

where $x \in \mathbb{Q}$, and we denote this as $a \in \mathbb{E}[\mathbb{Q}]$.

Evaluation of arithmetic expression

We can define evaluation $\nu(a)$ of a recursively as follows:

- Constant leaf: for any $x \in \mathbb{Q}$, $\nu(x) = x$.
- Compositional node by +: For any (a + b), $\nu((a + b)) = \nu(a) + \nu(b)$.
- Compositional node by -: For any (a-b), $\nu((a-b)) = \nu(a) \nu(b)$.
- Compositional node by \times : For any $(a \times b)$, $\nu((a \times b)) = \nu(a)\nu(b)$.
- Compositional node by \div : For any $(a \div b)$, if $\nu(b) \neq 0$, then $\nu((a \div b)) = \nu(a)/\nu(b)$.

Syntactic vs. Semantic

A careful reader may have noticed that the definition 1 is based on rational numbers $\mathbb Q$. Why can't we use real numbers $\mathbb R$ instead? The answer is that syntactically valid expressions may not be semantically valid. Dividing by zero can lead to invalid expressions, and the evaluation of the expression cannot be defined in this situation. Therefore, in real numbers, an expression may be syntactically valid but semantically not valid, and there is no algorithm that can decide whether an expression is semantically valid or not.

Topological arithmetic expression space

- 1 Convergence geometrically can lead to convergence of arithmetic evaluation
- 2 The arithmetic evaluation can be extended to the whole topological space continuously

Topological arithmetic expression space

We denote the set of all well-defined arithmetic expressions over the field of rational numbers $\mathbb Q$ as $\mathbb A$, where ν is a function from $\mathbb A$ to $\mathbb Q$.

Definition

There is a countable dense set G on the topological space A, and there exists an injection $\kappa:G\to\mathbb{A}$ between this dense set and the well-defined arithmetic expressions. We denote the image of this mapping as $\kappa(G)=\mathbb{K}$. If for any point $x\in A$, and any two sequences of points $y_i\in G$ and $w_j\in G$, when y_i converges to x and y_i also converges to y_i , the sequences of points over the rational numbers $v(\kappa(y_i))$ and $v(\kappa(w_i))$ both converge to the same real number; in this case, we can naturally make an extension:

- Extend G to a closed set \bar{G} on A;
- Extend ν to a mapping $\bar{\nu}$ from $\bar{\mathbb{K}}$ to the real numbers \mathbb{R} ;

If this extended valuation function $\bar{\nu}$ is a continuous function, then we call the topological space \mathcal{A} a topological arithmetic expression space. G is referred to as the grid on \mathcal{A} .



Classification problem

Local structure: decided by the flow equation 1. Classification of the global structure

Eigenfunction of Laplacian

On the hyperbolic plane

$$ds^2 = \frac{dx^2 + dy^2}{y^2}$$

the Laplacian is

$$\Delta = -y^2 \left(\frac{\partial^2}{\partial x^2} + \frac{\partial^2}{\partial y^2} \right)$$

Given

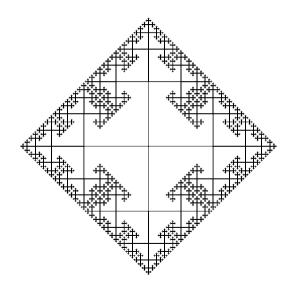
$$A = -\frac{x}{y} \tag{4}$$

We have

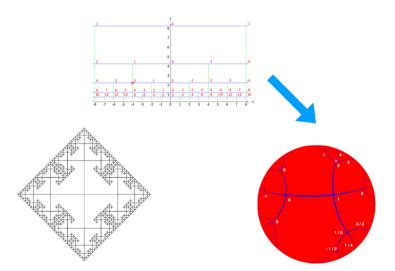
$$\Delta A = -y^2 \left(\frac{\partial^2}{\partial x^2} A + \frac{\partial^2}{\partial y^2} A \right) = y^2 \left(\frac{1}{\partial y} \left(\frac{1}{\partial y} \frac{x}{y} \right) \right) = 2A$$



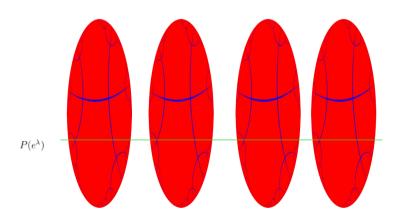
Further topics



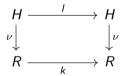
A jigsaw by Riemann mapping?



Tube structure?



Function as flow



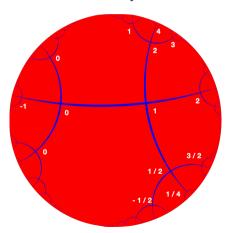


Extend calculus by mixing addition and multiplication?

Riemann sum is purely additional. Can we extend it by mixing addition and multiplication?

Calculus as a boundary algebra?

Infinitely small values and large values are appeared at the boundary. Limitation process can be treated as flow to the boundary.



Other analysis systems?

Formulate our systems as a tuple of

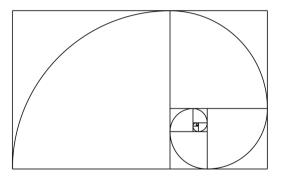
• *E*(*F*), (*H*, *a*), (*Path*, *Integ*)

Here we have

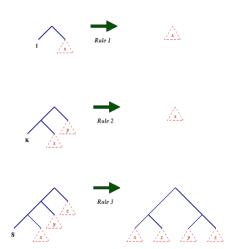
- E(F): Expressions over a field F
- (H, a): A scalar field "assignment" a on a space H
- (Path, Integ): all paths can be interpreted as an integral

Can we form a category? and complex analysis is an example?

Adventure in a wonderland



SKI combinator calculus



$$S(K(SI))K\alpha\beta \rightarrow$$
 $K(SI)\alpha(K\alpha)\beta \rightarrow$
 $SI(K\alpha)\beta \rightarrow$
 $I\beta(K\alpha\beta) \rightarrow$
 $I\beta\alpha \rightarrow$
 $\beta\alpha$

A space of SKI combinators

Any arithmetic expression can be represented by SKI combinators via Church numerals. So any arithmetic expression space can be encoded by SKI combinators. Program space!

Ancient Egyptian multiplication

1	1,1,1 U U U
II	$U_UU_UU_UU_U$
IIII	nn nn9
IIII IIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIII	1 ¹ 11 0 0 0
	ı ^{lıl} ınnnn ⁹

a example calculation

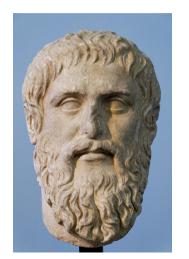
1*	35
2*	70
4	140
8*	280
1+2+8=11	35+70+280=385

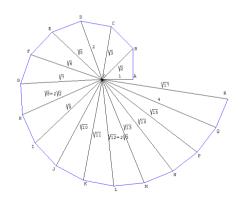


Ahmes Papyrus

A story of square root of 17

A story from Plato's Theaetetus: Theodorus of Cyrene, a young mathematician, was able to prove that $\sqrt{3}$, $\sqrt{5}$... are irrational, but not $\sqrt{17}$.

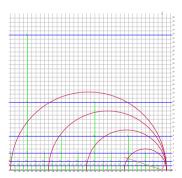




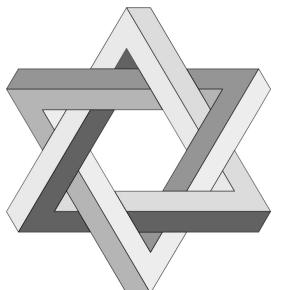
A logic system as a space?

The axioms by Victor Pambuccian and Celia Schacht are also fit into expressions, and then some terms of the system can be embedded into the expression space. Can we migrate the problem of irrationality of $\sqrt{17}$ from proof theory into a problem in group theory?

```
A.1. (x+y)+z=x+(y+z)
 A2. x + y = y + x
  A.3. (x \cdot y) \cdot z = x \cdot (y \cdot z)
 A 4. x \cdot y = y \cdot x
  A.S. x \cdot (y + z) = x \cdot y + x \cdot z
 A 6. x + 0 = x \wedge x \cdot 0 = 0
 A.7. \times 1 = x
 A.B. (x < y \land y < z) \rightarrow x < z
 A.9. \neg x < x
 A 10. x < y \lor x = y \lor y < x
 A II, x < y \rightarrow x + z < y + z
 A 12. (0 < z \land x < y) \rightarrow x \cdot z < y \cdot z
 A 13. x < y \rightarrow x + (y - x) = y
 A 14. 0 < 1 \land (r > 0 \rightarrow (r > 1 \lor r = 1))
A 16. m = \kappa(m, n) \cdot \mu(m, n) \wedge n = \kappa(m, n) \cdot \mu(n, m) \wedge (\mu(m, n) = 2 \left \lceil \frac{\mu(m, n)}{3} \right \rceil + 1 \vee \mu(n, m) = 2 \left \lceil \frac{\mu(m, n)}{3} \right \rceil + 1).
 A 17. x = \begin{bmatrix} 2t \end{bmatrix}.
 A 18, x = 2[5] \lor x = 2[5] + 1
 A 19. \pi_2(n) \land a \cdot b = \pi \land a > 1 \rightarrow a = 2 \begin{bmatrix} e \\ \end{bmatrix}
 A 20. 0 < n \rightarrow n = \tau(n) \cdot \omega(n) \wedge \pi_2(\tau(n)) \wedge \omega(n) = 2 \left[\frac{n(n)}{n}\right] + 1
 A 21. n < m \land \pi_2(m) \land \pi_2(n) \rightarrow \tau(m - n) = n.
```

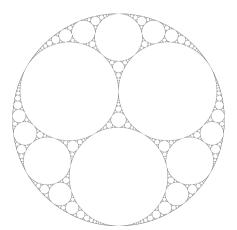


Final remarks



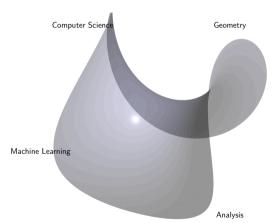
The unreasonable effectiveness of math

Math and even all human knowledge is also a geometrical object, just the same as the universe.



Knowledge geometry

A minimal surface of knowledge, every concept is a point, every relation is a line.



Thank you!