CS553: Cryptography

Assignment 2: Solutions

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### 1. How Many Keys?

r = 23 (formed from 2nd and 3rd digits of roll number)

$$\mathbb{Z}_{23}^+ = \{0, 1, ..., 22\}$$

Let K be the set of possible keys, such that  $\forall k \in K$ ,  $\gcd(r, k) = 1$ .

So, given r = 23,  $K = \{1, 2, ..., 22\}$ .

$$|\mathbf{K}| = 22.$$

#### 2. Euler Phi Function

Euler Phi Function, also known as Euler's totient function, returns the count of integers less than and relatively prime to n. It is denoted by  $\phi(n)$ .

For prime positive integers,  $\phi(n) = (n-1)$ . For other positive integers, if n = pq, where  $p,q \in \mathbb{Z}^+$  and prime,  $\phi(n) = (p-1)(q-1)$ .

The number of keys, |K|, for an Affine Cipher defined over  $\mathbb{Z}_r^+$  can be easily obtained by simply using Euler's totient function  $\phi(r)$ .

Since our r is prime,  $\phi(r) = (r-1) = 22$ .

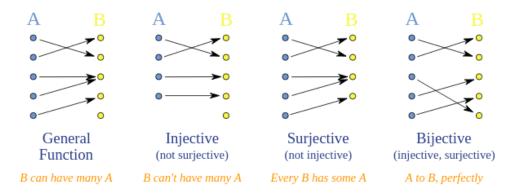
# 3. Bijection

Given  $X = \{x_1, x_2, ..., x_n\}, Y = \{y_1, y_2, ..., y_n\}$  and |X| = |Y|.

Also given,  $f: X \to Y$  is injective, i.e.  $\forall x \in X$ , f(x) is unique, or f(x) is a one-to-one function. So naturally,  $|X| \leq |Y|$ .

But, we know |X| = |Y|, meaning,  $\forall y \in Y, \exists x \in X$ , s.t. f(x) = y. Hence, f is also surjective, i.e. onto.

 $\therefore$  f is a bijection (Proved).



#### 4. Euclidean GCD

Python code (Python3): Euclidean\_gcd.py

```
def E_gcd(a,b): # function to calculate Euclidean GCD
    a,b = (b,a) if a < b else (a,b) # swaps a and b if a < b
    while True:
        r = a % b
        if (r == 0):
            break
        a = b
        b = r
    return b

try: # to allow only integer input
    a,b = [int(i)for i in(input("Enter two integers: ").split(" "))]
except ValueError:
    print("Please enter two integers!")
    exit(0)

print(E_gcd(a,b))</pre>
```

## 5. Involutory Key

### 5..1 Proof of involutory key

Given an Affine Cipher over  $\mathbb{Z}_m$  with K = (a, b).

Let  $e_K(x) = (ax + b) \mod m$ ,  $d_K(x) = a^{-1}(x - b) \mod m$  be the encryption and decryption function respectively.

<u>Proof 1</u>: Assuming key K is involutory, i.e.  $d_K(x) = e_K(x)$ ......(1), to prove  $a^{-1} \mod m = a$ , and  $b(a+1) \equiv 0 \mod m$ .

From (1),  $(ax + b) \mod m = a^{-1}(x - b) \mod m$ ,

 $\implies (ax + b) \mod \mathbf{m} = (a^{-1}x - a^{-1}b) \mod \mathbf{m}$ 

 $\therefore a \equiv a^{-1} \mod m...(2)$ , and

 $b \equiv -a^{-1}b \mod \text{m...}(3)$  [: two functions ax + b and cx + d are equal if their coefficients are equal, i.e. a = c and b = d].

From (2), we get  $a \equiv a^{-1} \mod m$  (Proved).

From (3), we get  $ab \equiv -b \mod m$ ,

$$\implies b(a+1) \equiv 0 \mod m \text{ (Proved)}$$

<u>Proof 2</u>: Assuming  $a \equiv a^{-1} \mod \text{m...}(4)$ , and  $b(a+1) \equiv 0 \mod \text{m}$ , to prove K is involutory, i.e.  $e_K(x) = d_K(x)$ 

 $b(a+1) \equiv 0 \mod m$ 

$$\implies ba \equiv -b \mod m, \implies b \equiv -a^{-1}b \mod m...(5)$$

Putting (4) and (5) in  $e_K(x)$ , we get

$$e_K(x) = (ax + b) \mod m = \{a^{-1}x \mod m + (-a^{-1}b) \mod m\} \mod m$$

$$= a^{-1}(x - b) \mod m = d_K(x) \implies \text{Key K is involutory}$$

... The above two proofs prove that Key K is involutory iff  $a \equiv a^{-1} \mod m$ , and  $b(a+1) \equiv 0 \mod m$ .(Proved)

## 5..2 Number of involutory keys

For an Affine Cipher defined over  $\mathbb{Z}_{15}$ , possible values for a, such that  $a = a^{-1}$  mod m are  $\{1, 4, 11, 14\}$  (using  $a^2$  mod 15). For each value of a, possible values of b such that  $b(a+1) \equiv 0$  mod n follow:

For a = 1, b(1 + 1) = 2b, meaning, only possible value for  $b \in \mathbb{Z}_{15}$  is 0. |b| = 1

For 
$$a = 4$$
,  $b(4 + 1) = 5b$ . So,  $b = \{0, 3, 6, 9, 12\}$ .  $|b| = 5$ .

For 
$$a = 11$$
,  $b(11 + 1) = 12b$ .  $b = \{0, 5, 10\}$ .  $|b| = 3$ .

For a = 14, b (14 + 1) = 15b. b can take any value from  $\mathbb{Z}_{15}$ . |b| = 15.

... Total no. of involutory keys = 1 + 5 + 3 + 15 = 24.