

logistics.

wenlong.mou @ uta.utoronto.ca

$(X_t : t \in T)$. T a period of time.

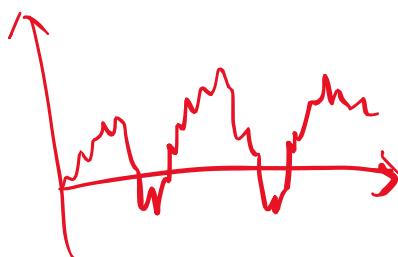
Temporal structures.

Markov process
(Discrete states).

$T = \mathbb{N}$ (non-neg ints).
(DT MC) gambling.
 $T = \mathbb{R}_+$ (non-neg real).
(end of semester) (CT MC)

martingales (discrete time). — gambling.

Brownian motion and stochastic calculus.



Def. Discrete-time, discrete-state, time-homogeneous MC.

(X_0, X_1, X_2, \dots) . $X_i \in S$.

(i) State space

S

(finite or countably infinite).

(ii) Initial distribution

$(v_i)_{i \in S}$

$$X_0 \sim v$$

($P(X_0 = i) = v_i$ for $i \in S$).

(iii). Transition probabilities.

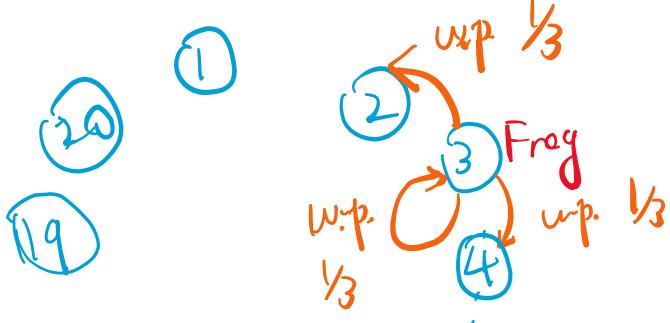
$$P = (P_{ij})_{i,j \in S}$$

Does not depend on time.
(In general, time-inhomogeneous M).

$$P_{ij} = P(X_{t+1} = j \mid X_t = i)$$

(for $i, j \in S$).

e.g.



X_i := index of the lily pad frog sitting on.

$$S = \{1, 2, 3, \dots, 20\}$$

$$v_i = \begin{cases} 1 & i=3, \\ 0 & i \neq 3. \end{cases}$$

$$P_{ij} = \begin{cases} \frac{1}{3} & \text{when } i=j \text{ or } \\ & i=(j+1) \bmod 20 \text{ or } \\ & i=(j-1) \bmod 20 \\ 0 & \text{otherwise.} \end{cases}$$

$$\begin{bmatrix} \frac{1}{3} & \frac{1}{3} & & & \\ \frac{1}{3} & \frac{1}{3} & & & \\ & & 0 & & \\ & & & & \frac{1}{3} \\ & & & & \frac{1}{3} \end{bmatrix}$$

"transition matrix".

e.g. Coin tossing.

At each time $t = 1, 2, 3, \dots$, toss a coin.

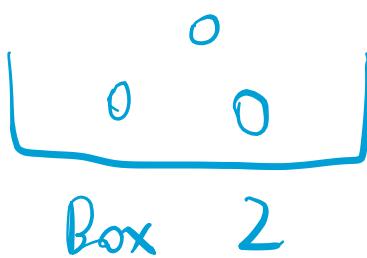
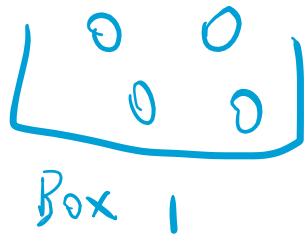
$X_t := \# \text{heads in first } t \text{ rounds.}$

$$S = \{0, 1, 2, \dots\}$$

$$V(0) = V_0 = 1, \quad V_i = 0 \text{ (for } i \geq 1).$$

$$P_{ij} = \begin{cases} \frac{1}{2} & (i=j \text{, or } i+1=j) \\ 0 & (\text{otherwise}). \end{cases}$$

e.g. Ehrenfest's Urn.



d balls

In total.

$$d = a+b$$

At each time.

— Randomly select a ball (uniformly)

— Move it to the opposite side.

Start w/ a balls in 1 and b balls in 2.

X_t := # of balls in box 1
at time t

$$= \{0, 1, 2, \dots, d\}$$

$$v_i = 1, \quad v_i = 0 \quad \text{for } (i \neq a)$$

$$\cdot P_{ij} \quad P(X_{t+1}=j \mid X_t=i)$$

$$X_{t+1} = \begin{cases} X_t + 1 & \text{w.p. } 1 - \frac{1}{d} \\ X_t - 1 & \text{w.p. } \frac{1}{d} \end{cases}$$

$$P_{ij} = \begin{cases} 1 - \frac{1}{d} & \text{when } j = i+1 \\ \frac{1}{d} & \text{when } j = i-1 \\ 0 & \text{otherwise.} \end{cases}$$

Non-example if the move depends on history before time t .

Important property of MC

"Markov property".

$$P(X_t = j \mid X_0 = i_0, X_1 = i_1, X_2 = i_2, \dots, X_{t-1} = i_{t-1})$$

$$= P(X_t = j \mid X_{t-1} = i_{t-1}) = P_{i_{t-1}, j}$$

Corollary.

$$P(X_0 = i_0, X_1 = i_1, \dots, X_n = i_n)$$

$$= P(X_0 = i_0) \cdot P(X_1 = i_1 \mid X_0 = i_0) \cdot P(X_2 = i_2 \mid X_0 = i_0, X_1 = i_1)$$

$$\dots \cdot P(X_n = i_n \mid X_0 = i_0, \dots, X_{n-1} = i_{n-1})$$

$$= v_{i_0} \cdot P_{i_0, i_1} \cdot P_{i_1, i_2} \cdot \dots \cdot P_{i_{n-1}, i_n}.$$

$$P(X_0 = i_0, X_1 = i_1, X_2 = i_2) = v_{i_0} \cdot P_{i_0, i_1} \cdot P_{i_1, i_2}$$

$$P(X_0 = i_0, X_2 = i_2) = v_{i_0} \cdot \sum_{i_1 \in S} P_{i_0, i_1} \cdot P_{i_1, i_2}$$

Matrix multiplication.

$$P(X_2 = i_2 \mid X_0 = i_0) = [P^2]_{i_0, i_2}.$$

} Detour

A, B are infinite-dim matrices (Countably inf S).

$$A = (a_{ij})_{i,j \in S} \quad B = (b_{ij})_{i,j \in S}$$

$$[A \cdot B]_{i,j} = \sum_{k \in S} a_{ik} b_{kj}$$

infinite sum convergence unclear in general
 but convergence is true
 when A, B are prob. transition matrices.

In general, for integer $k \geq 0$.

$$P(X_k=j | X_0=i) = [P^k]_{i,j}$$

$\overbrace{P(X_k=j | X_0=i)}$ $\rightarrow = P(X_{n+k}=j | X_n=i)$

$$P(X_k=j) = [\underbrace{v \cdot P^k}]_j.$$

(Row vector)

Def. $P_{i,j}^{(n)} := P(X_n=j | X_0=i)$

$(\Leftarrow P(X_{n+m}=j | X_m=i))$

for $i, j \in S$.

We have $P_{ij}^{(n)} = [P^n]_{:,j}$.

$$P_{ij}^{(m+n)} = [P^{m+n}]_{:,j}$$

$$= [P^m \cdot P^n]_{:,j}$$

$$= \sum_{k \in S} P_{ik}^{(m)} \cdot P_{kj}^{(n)}$$

"Chapman-Kolmogorov Eq"

Recurrence and transience.

Def. $N(i) :=$ total number of times that MC visits state i

$$= \sum_{t=1}^{+\infty} \mathbb{I}\{X_t = i\}$$

(r.v., may be infinite in some cases).

$$f_{ij} := P(N(j) \geq 1 \mid X_0 = i)$$

(Prob. reach j when starting from i).

f_{ii} : prob of returning to i
if we start from i .

Def. A state i is
 recurrent if $f_{ii} = 1$
 transient if $f_{ii} < 1$

Fact. $P_i(N(i) \geq k) = f_{ii}^k$

$$= \overbrace{P(N(i) \geq k | X_0 = i)}$$

(Not to confused w/ $P_X(X=x)$, not used
in my lec).

"Proof." Induction. $k=1$. By def.

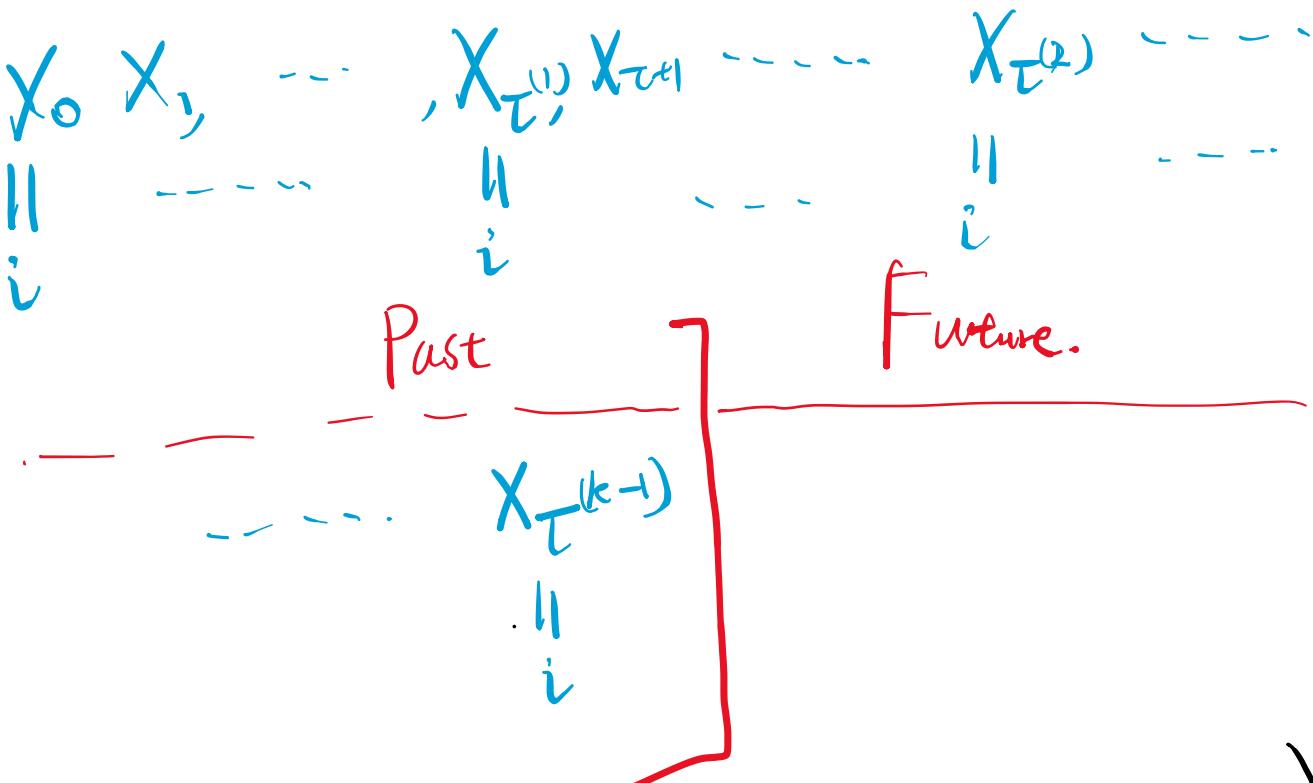
Assuming the claim is true for $k-1$.

$$P_i(N(i) \geq k)$$

$$= \boxed{P_i(N(i) \geq k | N(i) \geq k-1)} \cdot P_i(N(i) \geq k-1)$$

Need to show $= f_{ii}$.

$= f_{ii}^{k-1}$ by induction.



$$\begin{aligned}
 & P(\text{visit } i \text{ again} \mid X_0, X_1, \dots, X_{T^{(k-1)}}) \\
 &= P(\text{visit } i \text{ again} \mid X_{T^{(k-1)}} = i) \\
 &= P(\text{visit } i \text{ again} \mid X_0 = i) = f_i.
 \end{aligned}$$

a strong Markov property".
Let τ be the hitting time of state i ,
then we have
 $(X_\tau, X_{\tau+1}, \dots) \stackrel{d}{=} (X_0, X_1, \dots)$.

Also applies to

$\tau^{(k)} = k\text{-th hitting time of } i$.

By strong Markov property,

$$P_i(N(i) \geq k \mid N(i) \geq k-1) = f_{ii}^k.$$

Def. hitting time of i

$$\tau_i := \inf \{t \geq 1 : X_t = i\}$$

Corollary.

$$P_i(N(j) \geq k) = f_{ij} \cdot f_{jj}^{k-1}.$$

$$X_0 - \cdots \left[\begin{matrix} X_j \\ \parallel \\ j \end{matrix} \cdots \cdots \cdots \right]$$

Fresh new MC starting
from j .

$$P_i(N(j) \geq k)$$

$$\begin{aligned}
 &= P_i(N(j) \geq 1) \cdot \underbrace{P_i(N(j) \geq k \mid N(j) \geq 1)}_{f_{ij}} \\
 &= f_{ij} \cdot P_j(N(j) \geq k-1) \\
 &= f_{ij} \cdot f_{jj}^{k-1}
 \end{aligned}$$

Notation

$$P_i(N(j) \geq k \mid N(j) \geq 1)$$

$$= P(N(j) \geq k \mid N(j) \geq 1, X_0 = i).$$

$$= P(N(j) \geq k \mid N(j) \geq 1)$$

$$= P(N(j) \geq k-1 \mid X_0 = j).$$

Conclusion. (starting from i)

If $f_{ii} = 1$ then $N(i) = +\infty$
w.p. 1

If $f_{ii} < 1$ then

$$P(N(i)=k) = f_{ii}^k - f_{ii}^{k+1}$$

(Geometric distribution).

$$\text{e.g. } E_i[N(j)] = \begin{cases} \frac{f_{ij}}{1-f_{jj}} & (f_{jj} < 1) \\ +\infty & (f_{jj} = 1, f_{ij} > 0) \\ 0 & (f_{ij} = 0). \end{cases}$$

Question: can we determine
recurrence/transience by computing
sth. using $(p_{ij})_{i,j \in S}$.

Recurrence State Theorem.

State i is recurrent if and only if

$$\sum_{n=1}^{+\infty} P_{ii}^{(n)} = +\infty.$$

Proof. $\sum_{n=1}^{+\infty} P_{ii}^{(n)}$

$$= \sum_{n=1}^{+\infty} P_i(X_n=i)$$

$$= \sum_{n=1}^{+\infty} E_i[1_{X_n=i}]$$

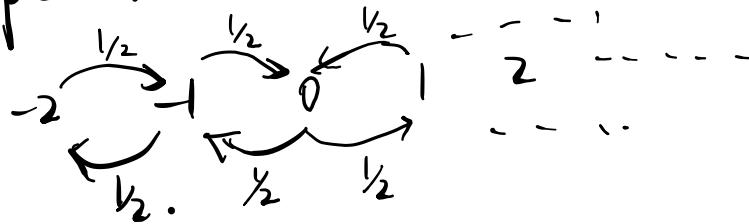
(Fubini - Tonelli)

$$= E_i\left[\sum_{n=1}^{+\infty} 1_{X_n=i}\right]$$

$$= E_i[N(i)] = \begin{cases} \frac{f_{ii}}{1-f_{ii}} & f_{ii} < 1 \\ +\infty & f_{ii} = 1. \end{cases}$$

(By-product, able to compute
 f_{ii} from $(P_{ij})_{i,j \in S}$).

e.g. Simple Random Walk.



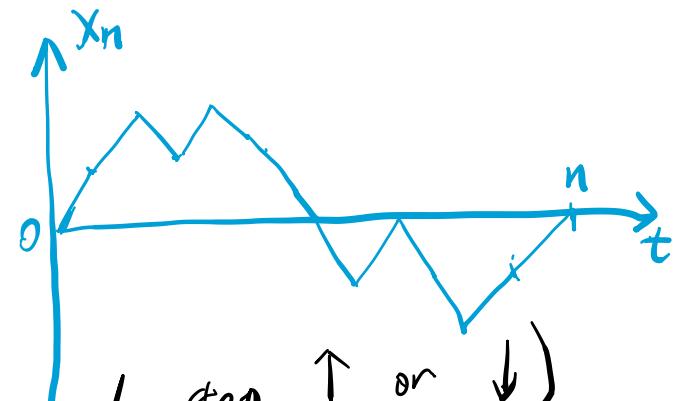
$$P_{i(i+1)} = P_{i(i-1)} = \frac{1}{2} \quad \text{for } i \in S = \mathbb{Z}.$$

Question: whether $f_{00} = 1$?

Suffices to check (by recurrent state thm)

$$\sum_{n=1}^{+\infty} P_{00}^{(n)}.$$

$$P_{00}^{(n)} = 0 \quad (\text{if } n \text{ is odd}).$$



Total # paths = 2^n (for each step, ↑ or ↓).

paths that go back to 0 at time n = $\binom{n}{n/2} = \frac{n!}{(n/2)!^2}$.

($\frac{n}{2}$ ↑'s and $\frac{n}{2}$ ↓'s for n steps)

Stirling's approx $n! \approx \sqrt{2\pi n} \cdot \left(\frac{n}{e}\right)^n$.

$$P_{00}^{(n)} \approx \sqrt{\frac{2}{\pi}} \cdot \frac{1}{\sqrt{n}}.$$

$$\sum_{n=1}^{+\infty} P_{00}^{(n)} = +\infty$$

So 0 is recurrent.