

Recall from last time.

Thm (recurrence state thm).

State  $i$  is recurrent  $\iff$

$$\sum_{n=1}^{+\infty} P_{ii}^{(n)} = +\infty.$$

(w.p.l., MC starting from  $i$  will return to  $i$  infinitely often)

$P_{ii}^{(n)}$  =  $n$ -step transition.

e.g. application, multidimensional RW.

$$\sum_{n=1}^{+\infty} P_{ii}^{(n)} = \sum_{n=1}^{+\infty} \left(\frac{c}{\sqrt{n}}\right)^d \begin{cases} < +\infty & d \geq 3 \\ = +\infty & d = 1, 2. \end{cases}$$

Proof of the theorem.

Detour: Borel-Cantelli Lemma.

General tool for reasoning infinite seq of events  
(finitely or infinitely often)

Let  $(E_n)_{n=1}^{+\infty}$  be a sequence of events.

If  $\sum_{n=1}^{+\infty} P(E_n) < +\infty$

then

$$P\left(\left(E_n\right)_{n=1}^{+\infty} \text{ happens only finitely often}\right) = 1.$$

Note: converse is false in general.

Application to recurrence state thm:

We let  $E_n = \{X_n = i\}$  (visit back  $i$  at time  $n$ ).

If  $\sum_{n=1}^{+\infty} P(E_n) = \sum_{n=1}^{+\infty} p_{ii}^{(n)} < +\infty$

then by BC lemma.  $(E_n)_{n=1}^{+\infty}$  happens only finitely many times, w.p. 1

i.e. we only visit back to  $i$  finitely many times, w.p. 1.

This implies transience and proves " $\Rightarrow$ " of thm.

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Back to the detour.

Proof of BC lemma.

$$\sum_{n=1}^{+\infty} P(E_n) = E\left[\sum_{n=1}^{+\infty} \mathbb{1}_{E_n}\right] < +\infty \quad (\text{by assumption}).$$

Fubini  
Tonelli!

↑  
# of times event happens.

So the r.v.  $\sum_{n=1}^{+\infty} \mathbb{1}_{E_n}$  has finite expectation.

and thus finite a.s.

From the proof, it's clear that converse cannot be guaranteed exists real-valued r.v. w/ infinite  $E[\cdot]$ .

However, for M.C.s.  $N(i) = \sum_{n=1}^{+\infty} \mathbb{1}_{\{X_n=i\}}$

is either infinite a.s. or

finite & geometrically distributed.

Using this observation, we can prove " $\Leftarrow$ " of recurrence state thm.

$$\sum_{n=1}^{+\infty} p_{il}^{(n)} = E_i \left[ \sum_{n=1}^{+\infty} \mathbb{1}_{X_n=i} \right] \quad (\text{Fubini-Tonelli})$$

$$= E_i [N(i)].$$

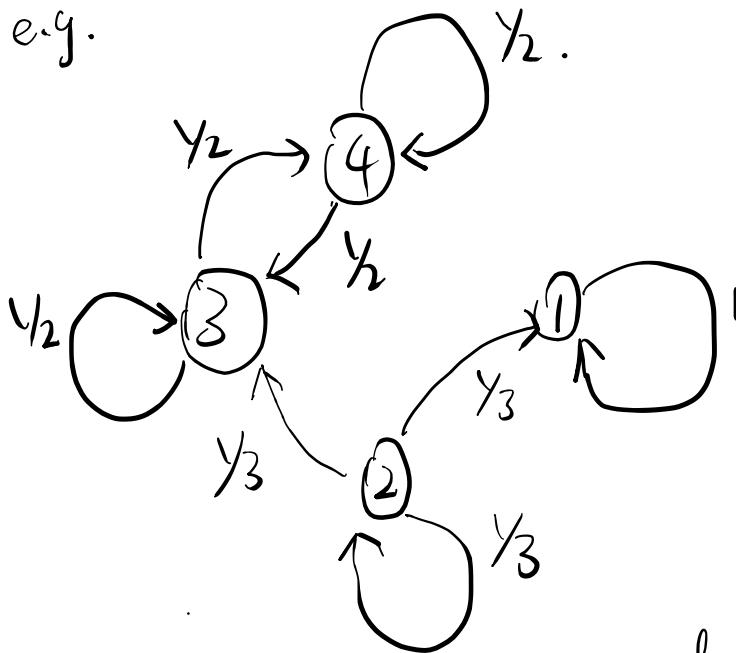
If  $i$  is transient, then  $f_{ii} < 1$

$$E_i [N(i)] = \frac{f_{ii}}{1-f_{ii}} < +\infty.$$

This proves " $\Leftarrow$ " of thm.

Next question: how to compute  $f_{ij}$ 's systematically?

Motivating e.g.



$$f_{11} = 1, \quad f_{1j} = 0 \quad \text{for } j \in \{2, 3, 4\}.$$

$$f_{22} = \frac{1}{3},$$

$$f_{21} = \frac{1}{3} + \frac{1}{3} \times \frac{1}{3} + \left(\frac{1}{3}\right)^3 + \dots$$

$$= \frac{y_3}{1 - y_3} = \frac{1}{2}.$$

$$\text{Similarly. } f_{23} = f_{24} = \frac{1}{2}$$

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In general, we use  $f$ -expansion.

i.e. solving out  $f_{ij}$ 's using linear systems.

$$f_{ij} = P_i(\text{visit } j).$$

(Expand based on the first step)

$$= \sum_{k \in S} P_i(X_1 = k) \cdot P_k(\text{visit } j \text{ starting from time } 0).$$

$\left( \neq \sum_{k \in S} P_k f_{kj} \right)$

$$= \sum_{\substack{k \in S \\ k \neq j}} P_k f_{kj} + P_j$$

If  $|S| < \infty$ ,  $f$ -expansion gives  $|S|^2$  equations

for  $|S|^2$  variables.

We can solve out these variables.

A more interesting example.

Gambler's ruin.

Suppose a gambler plays coin-tossing games.

- Initially, a  $\in N^+$  units of money.
- For each round, bet 1 unit

With winning prob  $p \in (0, 1)$ .

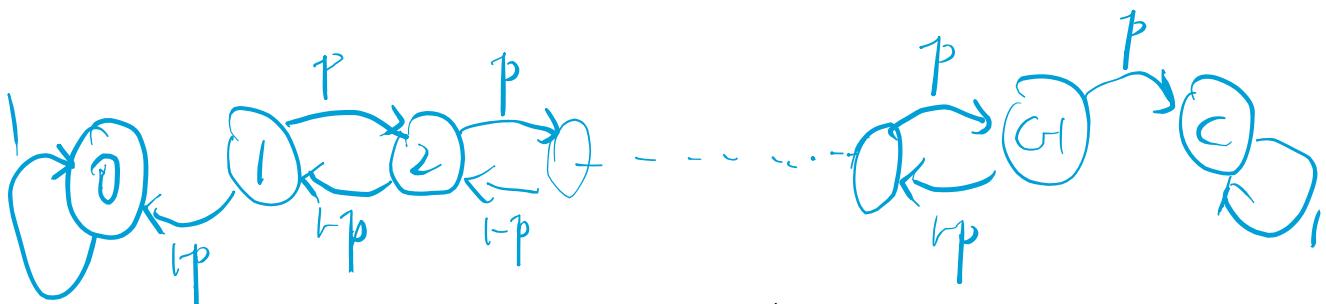
- Stop when
  - money = 0
  - money =  $c > a$ .

In MC language, let  $X_n$  be the amount of money in  $n$ -th round.

$$X_0 = a. \quad X_{n+1} = X_n + \varepsilon_{n+1} \text{ when } X_n \in \{0, c\}$$

$$\text{where } \varepsilon_{n+1} = \begin{cases} +1 & \text{w.p. } p \\ -1 & \text{w.p. } (1-p) \end{cases} \text{ indp.}$$

$$X_{n+1} = X_n \text{ when } X_n \notin \{0, c\}.$$



Interested in ruin probability

$$P_a(\text{visit } \varnothing) = f_{a0}.$$

Computation of  $f_{a0}$ :

$$f_{00} = 1 \quad f_{c0} = 0$$

By f-expansion.

$$f_{a0} = P f_{(a+1)0} + (1-P) f_{(a-1)0}$$

(when  $a \geq 2$ )

For  $a=1$

$$f_{10} = Pf_{20} + (1-P)$$
$$(= Pf_{20} + (1-P)f_{00} \quad \text{since } f_{00}=1)$$

$$\text{So } f_{a0} = Pf_{(a+1)0} + (1-P) f_{(a-1)0} \quad \text{for } a \in \{1, \dots, c\}$$

Solving the f-expansion eqs.

Case I:  $P = \frac{1}{2}$ .

$$f_{a0} = \frac{1}{2} \left( f_{(a+1)0} + f_{(a-1)0} \right),$$

So  $(f_{a0})_{a=0}^c$  is an arithmetic sequence, i.e.

$$f_{c0} - f_{(c-1)0} = f_{(c+1)0} - f_{c0} = f_{c0} - f_{(c-1)0} = \dots = f_{10} - f_{00}$$

$$f_{00} = 1, \quad f_{c0} = 0$$

$$-\frac{1}{c}$$

and therefore,

$$f_{a0} = 1 - \frac{a}{c}.$$

Similarly,  $f_{ac} = \frac{a}{c}.$

Remark: We cannot make money from it  
(in expectation).

$$\mathbb{E}[X_{final}] = c \cdot \frac{a}{c} + 0 \cdot \left(1 - \frac{a}{c}\right) = a.$$

(Indeed, this is not a coincidence).  
C.f. martingale part.

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Case II:  $p \neq \frac{1}{2}$ .

$$f_{a0} = p f_{(a+1)0} + (1-p) f_{(a-1)0}$$

$$p f_{(a+1)0} - p f_{a0} + (1-p) f_{(a-1)0} - (1-p) f_{a0} = 0.$$

So

$$f_{(a+1)0} - f_{a0} = \frac{1-p}{p} (f_{a0} - f_{(a-1)0}).$$

The increments form a geometric seq.

$$f_{(a+1)0} - f_{a0} = \left(\frac{1-p}{p}\right)^a (f_{10} - f_{00})$$

Summing up over  $a$ ,

$$\begin{aligned} -1 &= f_{c0} - f_{00} = \sum_{a=0}^{c-1} (f_{(a+1)0} - f_{a0}) \\ &= (f_{10} - f_{00}) \cdot \left[ \sum_{a=0}^{c-1} \left(\frac{1-p}{p}\right)^a \right]. \end{aligned}$$

Following some calculation,

$$f_{c0} = \frac{\left(\frac{1-p}{p}\right)^c - \left(\frac{1-p}{p}\right)^a}{\left(\frac{1-p}{p}\right)^c - 1} \quad (p \neq \frac{1}{2})$$

Remark: when  $p \neq \frac{1}{2}$ .

run prob is } exponentially small when  $p > \frac{1}{2}$ .  
} exponentially close to 1 when  $p < \frac{1}{2}$ .

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Relation between recurrence/transience of different states.

Recall notation :

$i \rightarrow j$  means  $j$  is reachable from  $i$ ,  $f_{ij} > 0$ .

$i \leftrightarrow j$  means  $i \rightarrow j$  and  $j \rightarrow i$ .

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Def. We call an MC irreducible when

$i \leftrightarrow j$  for  $i, j \in S$ .

(Implicitly, an MC not satisfying this  
can be "reduced" into smaller ones).

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Fact. If  $i \leftrightarrow j$ , then " $i$  is recurrent"



" $j$  is recurrent".

Corollary of the fact: "base theorem"

If a MC is irreducible, then one of the  
following is true:

- All states are recurrent. ("recurrent MC")
- All states are transient. ("transient MC").

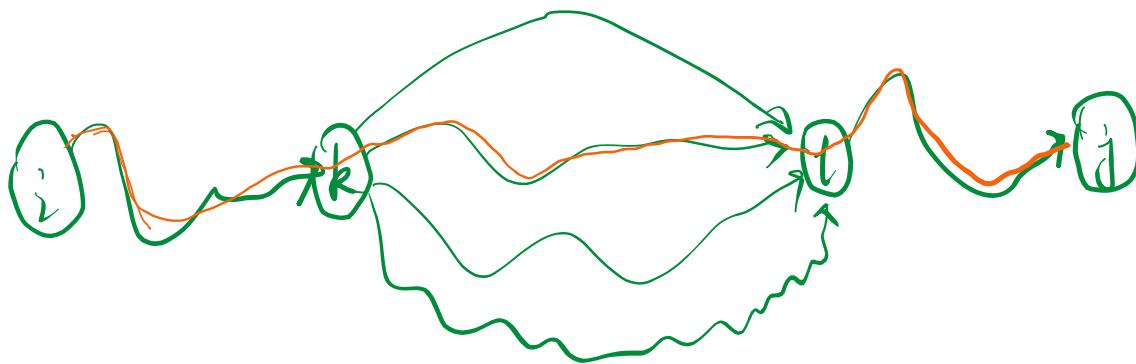
Proof of the fact:

We use the "sum lemma".

Lemma: If  $i, j, k, l \in S$  (allow repeating).

Suppose  $i \rightarrow k$  and  $l \rightarrow j$

If  $\sum_{n=1}^{+\infty} p_{kl}^{(n)} = +\infty$  then  $\sum_{n=1}^{+\infty} p_{ij}^{(n)} = +\infty$ .



Proof: Since  $i \rightarrow k$ ,  $\exists m > 0$ ,  $p_{ik}^{(m)} > 0$

$\hookrightarrow j$ ,  $\exists r > 0$ ,  $p_{lj}^{(r)} > 0$ .

Then (by Kolmogorov - Chapman equation)

$$p_{ij}^{(n+m+r)} \geq p_{ik}^{(m)} \cdot p_{kl}^{(n)} \cdot p_{lj}^{(r)}$$

um them together

$$\sum_{n=1}^{+\infty} P_{ij}^{(n)} \geq \underbrace{P_{ih}^{(m)}}_{>0} \cdot \underbrace{P_{hl}^{(r)}}_{>0} \cdot \underbrace{\sum_{k=1}^{+\infty} P_{kl}^{(n)}}_{= +\infty} = +\infty.$$

This is sum lemma

I'd like to prove of recurrence equivalence:

Apply sum lemma w/  $j=i$ ,  $l=k$ .

$$i \rightarrow k, l \rightarrow j \iff i \leftrightarrow l.$$

$$l \text{ is recurrent} \iff \sum_{n=1}^{+\infty} P_{il}^{(n)} = +\infty.$$

Sum lemma implies

$$\sum_{n=1}^{+\infty} P_{il}^{(n)} = +\infty.$$

which proves desired fact.

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Next question.

What do we know about  $\sum_{n=1}^{+\infty} P_{ij}^{(n)}$ ? (for  $i \neq j$ ).

Fact. If MC is Recurrent,  $\sum_{n=1}^{+\infty} P_{ij}^{(n)} = +\infty$

If MC is Transient,  $\sum_{n=1}^{+\infty} P_{ij}^{(n)} < +\infty$ .

(This is if-and-only-if: any ij pair  
is equivalent to MC recurrence/transience)

Proof.  $\sum_{n=1}^{+\infty} P_{ij}^{(n)} = \mathbb{E}_i[N(j)].$

- Recurrent,  $\exists m, P_{ij}^{(m)} > 0$  by irreducibility

$$\mathbb{E}_i[N(j)] \geq P_{ij}^{(m)} \cdot \underbrace{\mathbb{E}_j[N(j)]}_{=+\infty \text{ by recurrence}} = +\infty.$$

- Transient,

$$\mathbb{E}_i[N(j)] = f_{ij} \cdot (\underbrace{\mathbb{E}_j[N(j)] + 1}_{= \frac{f_{ii}}{1-f_{ij}}} + 1) = \frac{f_{ij}}{1-f_{ij}} < +\infty.$$

e.g. When  $|S| < +\infty$ , and irreducible,  
then the MC is recurrent

Proof. For any fixed  $i \in S$ .

$$\sum_{j \in S} \left( \sum_{n=1}^{+\infty} P_{ij}^{(n)} \right) = \sum_{n=1}^{+\infty} \left( \sum_{j \in S} P_{ij}^{(n)} \right) = \sum_{n=1}^{+\infty} 1 = +\infty$$

Sum of finitely many terms.

$\exists j_0$ , s.t.

$$\sum_{n=1}^{+\infty} P_{j_0}^{(n)} = +\infty.$$

and by irreducibility, this implies MC is recurrent.

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Another useful fact: "f-lemma".

Motivation: 3 characterization for recurrence/transience

$$\left\{ \begin{array}{l} \sum_{n=1}^{+\infty} P_{ij}^{(n)} \text{ equivalent} \\ N(j) \text{ the same} \\ f_{ij} ? \end{array} \right.$$

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f Lemma. If  $j \rightarrow i$ ,  $f_{jj} = 1$

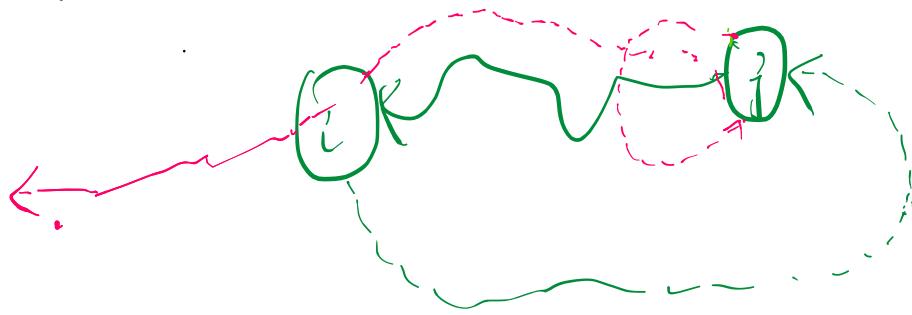
then  $f_{ij} = 1$ .

(Compute off-diagonal f for recurrent states).

Direct corollary:

For recurrent MC,  $f_{ij} = 1$  for  $i, j \in S$ .

# Proof of f-Lemma.



Technical tool: "hit lemma".

Define the event

$H_{ij} = \{ \text{MC hits } i \text{ before returning to } j \}$ .

$f_j \rightarrow i$ , then  $P_j(H_{ij}) > 0$ .

(This lemma removes the trouble of " $\rightarrow$ " path).

Proof of "Hit Lemma".

$j \rightarrow i$ . So  $\exists x_0 x_1 \dots x_m$  such that  $x_0 = j$  and  $x_m = i$ .

We have

$$P_{x_0 x_1} P_{x_1 x_2} \dots P_{x_{m-1} x_m} > 0.$$

This implies

$$P_{x_1 x_{t+1}} P_{x_{t+1} x_{t+2}} \cdots P_{x_m x_m} > 0.$$

Let  $t$  be the last time  $j$  appears  
in this sequence.

$$P_j(H_{ij}) \geq P_{x_1 x_{t+1}} \cdots P_{x_m x_m} > 0$$

which proves hit lemma.

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Proof of f-lemma using hit lemma.

$$D = P_j(\text{never return to } j)$$

$$\geq \underbrace{P_j(H_{ij})}_{>0} \cdot \underbrace{P_i(\text{never visit } j)}_{\text{must be } 0}.$$

So  $f_{ij} = 1$ , which proves f-lemma.