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|---|--------------|---|----|
| MPI_CART_COORDS(comm, rank, maxdims, coords) | | | 1 |
| IN | comm | communicator with Cartesian structure (handle) | 2 |
| IN | rank | rank of a process within group of comm (integer) | 3 |
| IN | maxdims | length of vector coords in the calling program (integer) | 4 |
| OUT | coords | integer array (of size ndims) containing the Cartesian coordinates of specified process (array of integers) | 5 |
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| | | | 7 |
| | | | 8 |
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| int MPI_Cart_coords(MPI_Comm comm, int rank, int maxdims, int *coords) | | | 10 |
| MPI_CART_COORDS(COMM, RANK, MAXDIMS, COORDS, IERROR) | | | 11 |
| INTEGER COMM, RANK, MAXDIMS, COORDS(*), IERROR | | | 12 |
| void MPI::Cartcomm::Get_coords(int rank, int maxdims, int coords[]) const | | | 13 |
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| The inverse mapping, rank-to-coordinates translation is provided by | | | 15 |
| MPI_CART_COORDS. | | | 16 |
| If comm is associated with a zero-dimensional Cartesian topology, | | | 17 |
| coords will be unchanged. | | | 18 |
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| MPI_GRAPH_NEIGHBORS_COUNT(comm, rank, nneighbors) | | | 22 |
| IN | comm | communicator with graph topology (handle) | 23 |
| IN | rank | rank of process in group of comm (integer) | 24 |
| OUT | nneighbors | number of neighbors of specified process (integer) | 25 |
| | | | 26 |
| | | | 27 |
| int MPI_Graph_neighbors_count(MPI_Comm comm, int rank, int *nneighbors) | | | 28 |
| MPI_GRAPH_NEIGHBORS_COUNT(COMM, RANK, NNEIGHBORS, IERROR) | | | 29 |
| INTEGER COMM, RANK, NNEIGHBORS, IERROR | | | 30 |
| int MPI::Graphcomm::Get_neighbors_count(int rank) const | | | 31 |
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| MPI_GRAPH_NEIGHBORS(comm, rank, maxneighbors, neighbors) | | | 36 |
| IN | comm | communicator with graph topology (handle) | 37 |
| IN | rank | rank of process in group of comm (integer) | 38 |
| IN | maxneighbors | size of array neighbors (integer) | 39 |
| OUT | neighbors | ranks of processes that are neighbors to specified process (array of integer) | 40 |
| | | | 41 |
| | | | 42 |
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| int MPI_Graph_neighbors(MPI_Comm comm, int rank, int maxneighbors, | | | 44 |
| int *neighbors) | | | 45 |
| MPI_GRAPH_NEIGHBORS(COMM, RANK, MAXNEIGHBORS, NEIGHBORS, IERROR) | | | 46 |
| INTEGER COMM, RANK, MAXNEIGHBORS, NEIGHBORS(*), IERROR | | | 47 |
| | | | 48 |

```

1 void MPI::Graphcomm::Get_neighbors(int rank, int maxneighbors, int
2     neighbors[]) const
3

```

MPI_GRAPH_NEIGHBORS_COUNT and MPI_GRAPH_NEIGHBORS provide adjacency information for a general graph topology. The returned count and array of neighbors for the queried rank will both include *all* neighbors and reflect the same edge ordering as was specified by the original call to MPI_GRAPH_CREATE. Specifically, MPI_GRAPH_NEIGHBORS_COUNT and MPI_GRAPH_NEIGHBORS will return values based on the original `index` and `edges` array passed to MPI_GRAPH_CREATE (assuming that `index[-1]` effectively equals zero):

- The count returned from MPI_GRAPH_NEIGHBORS_COUNT will be `(index[rank] - index[rank-1])`.
- The `neighbors` array returned from MPI_GRAPH_NEIGHBORS will be `edges[index[rank-1]]` through `edges[index[rank]]`.

Example 7.3 Assume there are four processes 0, 1, 2, 3 with the following adjacency matrix (note that some neighbors are listed multiple times):

| process | neighbors |
|---------|-----------|
| 0 | 1, 1, 3 |
| 1 | 0, 0 |
| 2 | 3 |
| 3 | 0, 2, 2 |

Thus, the input arguments to MPI_GRAPH_CREATE are:

```

nnodes = 4
index = 3, 5, 6, 9
edges = 1, 1, 3, 0, 0, 3, 0, 2, 2

```

Therefore, calling MPI_GRAPH_NEIGHBORS_COUNT and MPI_GRAPH_NEIGHBORS for each of the 4 processes will return:

| Input rank | Count | Neighbors |
|------------|-------|-----------|
| 0 | 3 | 1, 1, 3 |
| 1 | 2 | 0, 0 |
| 2 | 1 | 3 |
| 3 | 3 | 0, 2, 2 |

Example 7.4 Suppose that `comm` is a communicator with a shuffle-exchange topology. The group has 2^n members. Each process is labeled by a_1, \dots, a_n with $a_i \in \{0, 1\}$, and has three neighbors: $\text{exchange}(a_1, \dots, a_n) = a_1, \dots, a_{n-1}, \bar{a}_n$ ($\bar{a} = 1 - a$), $\text{shuffle}(a_1, \dots, a_n) = a_2, \dots, a_n, a_1$, and $\text{unshuffle}(a_1, \dots, a_n) = a_n, a_1, \dots, a_{n-1}$. The graph adjacency list is illustrated below for $n = 3$.