

D R A F T

Document for a Standard Message-Passing Interface

Message Passing Interface Forum

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Chapter 6

Groups, Contexts, Communicators, and Caching

6.1 Introduction

This chapter introduces MPI features that support the development of parallel libraries. Parallel libraries are needed to encapsulate the distracting complications inherent in parallel implementations of key algorithms. They help to ensure consistent correctness of such procedures, and provide a “higher level” of portability than MPI itself can provide. As such, libraries prevent each programmer from repeating the work of defining consistent data structures, data layouts, and methods that implement key algorithms (such as matrix operations). Since the best libraries come with several variations on parallel systems (different data layouts, different strategies depending on the size of the system or problem, or type of floating point), this too needs to be hidden from the user.

We refer the reader to [7] and [1] for further information on writing libraries in MPI, using the features described in this chapter.

6.1.1 Features Needed to Support Libraries

The key features needed to support the creation of robust parallel libraries are as follows:

- Safe communication space, that guarantees that libraries can communicate as they need to, without conflicting with communication extraneous to the library,
- Group scope for collective operations, that allow libraries to avoid unnecessarily synchronizing uninvolved processes (potentially running unrelated code),
- Abstract process naming to allow libraries to describe their communication in terms suitable to their own data structures and algorithms,
- The ability to “adorn” a set of communicating processes with additional user-defined attributes, such as extra collective operations. This mechanism should provide a means for the user or library writer effectively to extend a message-passing notation.

In addition, a unified mechanism or object is needed for conveniently denoting communication context, the group of communicating processes, to house abstract process naming, and to store adornments.

6.1.2 MPI's Support for Libraries

The corresponding concepts that MPI provides, specifically to support robust libraries, are as follows:

- **Contexts** of communication,
- **Groups** of processes,
- **Virtual topologies**,
- **Attribute caching**,
- **Communicators**.

Communicators (see [4, 6, 8]) encapsulate all of these ideas in order to provide the appropriate scope for all communication operations in MPI. Communicators are divided into two kinds: intra-communicators for operations within a single group of processes and inter-communicators for operations between two groups of processes.

Caching. Communicators (see below) provide a “caching” mechanism that allows one to associate new attributes with communicators, on par with MPI built-in features. This can be used by advanced users to adorn communicators further, and by MPI to implement some communicator functions. For example, the virtual-topology functions described in Chapter 7 are likely to be supported this way.

Groups. Groups define an ordered collection of processes, each with a rank, and it is this group that defines the low-level names for inter-process communication (ranks are used for sending and receiving). Thus, groups define a scope for process names in point-to-point communication. In addition, groups define the scope of collective operations. Groups may be manipulated separately from communicators in MPI, but only communicators can be used in communication operations.

Intra-communicators. The most commonly used means for message passing in MPI is via intra-communicators. Intra-communicators contain an instance of a group, contexts of communication for both point-to-point and collective communication, and the ability to include virtual topology and other attributes. These features work as follows:

- **Contexts** provide the ability to have separate safe “universes” of message-passing in MPI. A context is akin to an additional tag that differentiates messages. The system manages this differentiation process. The use of separate communication contexts by distinct libraries (or distinct library invocations) insulates communication internal to the library execution from external communication. This allows the invocation of the library even if there are pending communications on “other” communicators, and avoids the need to synchronize entry or exit into library code. Pending point-to-point communications are also guaranteed not to interfere with collective communications within a single communicator.
- **Groups** define the participants in the communication (see above) of a communicator.

- A **virtual topology** defines a special mapping of the ranks in a group to and from a topology. Special constructors for communicators are defined in Chapter 7 to provide this feature. Intra-communicators as described in this chapter do not have topologies.
- **Attributes** define the local information that the user or library has added to a communicator for later reference.

Advice to users. The practice in many communication libraries is that there is a unique, predefined communication universe that includes all processes available when the parallel program is initiated; the processes are assigned consecutive ranks. Participants in a point-to-point communication are identified by their rank; a collective communication (such as broadcast) always involves all processes. This practice can be followed in MPI by using the predefined communicator `MPI_COMM_WORLD`. *Users who are satisfied with this practice can plug in `MPI_COMM_WORLD` wherever a communicator argument is required, and can consequently disregard the rest of this chapter. (End of advice to users.)*

Inter-communicators. The discussion has dealt so far with **intra-communication**: communication within a group. MPI also supports **inter-communication**: communication between two non-overlapping groups. When an application is built by composing several parallel modules, it is convenient to allow one module to communicate with another using local ranks for addressing within the second module. This is especially convenient in a client-server computing paradigm, where either client or server are parallel. The support of inter-communication also provides a mechanism for the extension of MPI to a dynamic model where not all processes are preallocated at initialization time. In such a situation, it becomes necessary to support communication across “universes.” Inter-communication is supported by objects called **inter-communicators**. These objects bind two groups together with communication contexts shared by both groups. For inter-communicators, these features work as follows:

- Contexts provide the ability to have a separate safe “universe” of message-passing between the two groups. A send in the local group is always a receive in the remote group, and vice versa. The system manages this differentiation process. The use of separate communication contexts by distinct libraries (or distinct library invocations) insulates communication internal to the library execution from external communication. This allows the invocation of the library even if there are pending communications on “other” communicators, and avoids the need to synchronize entry or exit into library code.
- A local and remote group specify the recipients and destinations for an inter-communicator.
- Virtual topology is undefined for an inter-communicator.
- As before, attributes cache defines the local information that the user or library has added to a communicator for later reference.

MPI provides mechanisms for creating and manipulating inter-communicators. They are used for point-to-point and collective communication in an related manner to intra-communicators. Users who do not need inter-communication in their applications can safely

ignore this extension. Users who require inter-communication between overlapping groups must layer this capability on top of MPI.

6.2 Basic Concepts

In this section, we turn to a more formal definition of the concepts introduced above.

6.2.1 Groups

A **group** is an ordered set of process identifiers (henceforth processes); processes are implementation-dependent objects. Each process in a group is associated with an integer **rank**. Ranks are contiguous and start from zero. Groups are represented by opaque **group objects**, and hence cannot be directly transferred from one process to another. A group is used within a communicator to describe the participants in a communication “universe” and to rank such participants (thus giving them unique names within that “universe” of communication).

There is a special pre-defined group: `MPI_GROUP_EMPTY`, which is a group with no members. The predefined constant `MPI_GROUP_NULL` is the value used for invalid group handles.

Advice to users. `MPI_GROUP_EMPTY`, which is a valid handle to an empty group, should not be confused with `MPI_GROUP_NULL`, which in turn is an invalid handle. The former may be used as an argument to group operations; the latter, which is returned when a group is freed, is not a valid argument. (*End of advice to users.*)

Advice to implementors. A group may be represented by a virtual-to-real process-address-translation table. Each communicator object (see below) would have a pointer to such a table.

Simple implementations of MPI will enumerate groups, such as in a table. However, more advanced data structures make sense in order to improve scalability and memory usage with large numbers of processes. Such implementations are possible with MPI. (*End of advice to implementors.*)

6.2.2 Contexts

A **context** is a property of communicators (defined next) that allows partitioning of the communication space. A message sent in one context cannot be received in another context. Furthermore, where permitted, collective operations are independent of pending point-to-point operations. Contexts are not explicit MPI objects; they appear only as part of the realization of communicators (below).

Advice to implementors. Distinct communicators in the same process have distinct contexts. A context is essentially a system-managed tag (or tags) needed to make a communicator safe for point-to-point and MPI-defined collective communication. Safety means that collective and point-to-point communication within one communicator do not interfere, and that communication over distinct communicators don’t interfere.

A possible implementation for a context is as a supplemental tag attached to messages on send and matched on receive. Each intra-communicator stores the value of its two tags (one for point-to-point and one for collective communication). Communicator-generating functions use a collective communication to agree on a new group-wide unique context.

Analogously, in inter-communication, two context tags are stored per communicator, one used by group A to send and group B to receive, and a second used by group B to send and for group A to receive.

Since contexts are not explicit objects, other implementations are also possible. (*End of advice to implementors.*)

6.2.3 Intra-Communicators

Intra-communicators bring together the concepts of group and context. To support implementation-specific optimizations, and application topologies (defined in the next chapter, Chapter 7), communicators may also “cache” additional information (see Section 6.7). MPI communication operations reference communicators to determine the scope and the “communication universe” in which a point-to-point or collective operation is to operate.

Each communicator contains a group of valid participants; this group always includes the local process. The source and destination of a message is identified by process rank within that group.

For collective communication, the intra-communicator specifies the set of processes that participate in the collective operation (and their order, when significant). Thus, the communicator restricts the “spatial” scope of communication, and provides machine-independent process addressing through ranks.

Intra-communicators are represented by opaque **intra-communicator objects**, and hence cannot be directly transferred from one process to another.

6.2.4 Predefined Intra-Communicators

An initial intra-communicator `MPI_COMM_WORLD` of all processes the local process can communicate with after initialization (itself included) is defined once `MPI_INIT` or `MPI_INIT_THREAD` has been called. In addition, the communicator `MPI_COMM_SELF` is provided, which includes only the process itself.

The predefined constant `MPI_COMM_NULL` is the value used for invalid communicator handles.

In a static-process-model implementation of MPI, all processes that participate in the computation are available after MPI is initialized. For this case, `MPI_COMM_WORLD` is a communicator of all processes available for the computation; this communicator has the same value in all processes. In an implementation of MPI where processes can dynamically join an MPI execution, it may be the case that a process starts an MPI computation without having access to all other processes. In such situations, `MPI_COMM_WORLD` is a communicator incorporating all processes with which the joining process can immediately communicate. Therefore, `MPI_COMM_WORLD` may simultaneously represent disjoint groups in different processes.

All MPI implementations are required to provide the `MPI_COMM_WORLD` communicator. It cannot be deallocated during the life of a process. The group corresponding to this communicator does not appear as a pre-defined constant, but it may be accessed using

MPI_COMM_GROUP (see below). MPI does not specify the correspondence between the process rank in MPI_COMM_WORLD and its (machine-dependent) absolute address. Neither does MPI specify the function of the host process, if any. Other implementation-dependent, predefined communicators may also be provided.

6.3 Group Management

This section describes the manipulation of process groups in MPI. These operations are local and their execution does not require interprocess communication.

6.3.1 Group Accessors

MPI_GROUP_SIZE(group, size)

IN	group	group (handle)
OUT	size	number of processes in the group (integer)

int MPI_Group_size(MPI_Group group, int *size)

MPI_Group_size(group, size, ierror)
 TYPE(MPI_Group), INTENT(IN) :: group
 INTEGER, INTENT(OUT) :: size
 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GROUP_SIZE(GROUP, SIZE, IERROR)
 INTEGER GROUP, SIZE, IERROR

MPI_GROUP_RANK(group, rank)

IN	group	group (handle)
OUT	rank	rank of the calling process in group, or MPI_UNDEFINED if the process is not a member (integer)

int MPI_Group_rank(MPI_Group group, int *rank)

MPI_Group_rank(group, rank, ierror)
 TYPE(MPI_Group), INTENT(IN) :: group
 INTEGER, INTENT(OUT) :: rank
 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GROUP_RANK(GROUP, RANK, IERROR)
 INTEGER GROUP, RANK, IERROR

MPI_GROUP_TRANSLATE_RANKS(group1, n, ranks1, group2, ranks2)			1
IN	group1	group1 (handle)	2
IN	n	number of ranks in ranks1 and ranks2 arrays (integer)	3
IN	ranks1	array of zero or more valid ranks in group1	4
IN	group2	group2 (handle)	5
OUT	ranks2	array of corresponding ranks in group2, MPI_UNDEFINED when no correspondence exists.	6
			7
			8
			9
			10
int MPI_Group_translate_ranks(MPI_Group group1, int n, const int ranks1[], MPI_Group group2, int ranks2[])			11
			12
MPI_Group_translate_ranks(group1, n, ranks1, group2, ranks2, ierror)			13
TYPE(MPI_Group), INTENT(IN) :: group1, group2			14
INTEGER, INTENT(IN) :: n, ranks1(n)			15
INTEGER, INTENT(OUT) :: ranks2(n)			16
INTEGER, OPTIONAL, INTENT(OUT) :: ierror			17
			18
MPI_GROUP_TRANSLATE_RANKS(GROUP1, N, RANKS1, GROUP2, RANKS2, IERROR)			19
INTEGER GROUP1, N, RANKS1(*), GROUP2, RANKS2(*), IERROR			20
			21
This function is important for determining the relative numbering of the same processes in two different groups. For instance, if one knows the ranks of certain processes in the group of MPI_COMM_WORLD, one might want to know their ranks in a subset of that group.			22
			23
MPI_PROC_NULL is a valid rank for input to MPI_GROUP_TRANSLATE_RANKS, which returns MPI_PROC_NULL as the translated rank.			24
			25
			26
			27
MPI_GROUP_COMPARE(group1, group2, result)			28
IN	group1	first group (handle)	29
IN	group2	second group (handle)	30
OUT	result	result (integer)	31
			32
			33
			34
int MPI_Group_compare(MPI_Group group1, MPI_Group group2, int *result)			35
			36
MPI_Group_compare(group1, group2, result, ierror)			37
TYPE(MPI_Group), INTENT(IN) :: group1, group2			38
INTEGER, INTENT(OUT) :: result			39
INTEGER, OPTIONAL, INTENT(OUT) :: ierror			40
			41
MPI_GROUP_COMPARE(GROUP1, GROUP2, RESULT, IERROR)			42
INTEGER GROUP1, GROUP2, RESULT, IERROR			43
			44
MPI_IDENT results if the group members and group order is exactly the same in both groups. This happens for instance if group1 and group2 are the same handle. MPI_SIMILAR results if the group members are the same but the order is different. MPI_UNEQUAL results otherwise.			45
			46
			47
			48

6.3.2 Group Constructors

Group constructors are used to subset and superset existing groups. These constructors construct new groups from existing groups. These are local operations, and distinct groups may be defined on different processes; a process may also define a group that does not include itself. Consistent definitions are required when groups are used as arguments in communicator-building functions. MPI does not provide a mechanism to build a group from scratch, but only from other, previously defined groups. The base group, upon which all other groups are defined, is the group associated with the initial communicator `MPI_COMM_WORLD` (accessible through the function `MPI_COMM_GROUP`).

Rationale. In what follows, there is no group duplication function analogous to `MPI_COMM_DUP`, defined later in this chapter. There is no need for a group duplicator. A group, once created, can have several references to it by making copies of the handle. The following constructors address the need for subsets and supersets of existing groups. (*End of rationale.*)

Advice to implementors. Each group constructor behaves as if it returned a new group object. When this new group is a copy of an existing group, then one can avoid creating such new objects, using a reference-count mechanism. (*End of advice to implementors.*)

`MPI_COMM_GROUP(comm, group)`

IN	comm	communicator (handle)
OUT	group	group corresponding to comm (handle)

`int MPI_Comm_group(MPI_Comm comm, MPI_Group *group)`

`MPI_Comm_group(comm, group, ierror)`
 TYPE(MPI_Comm), INTENT(IN) :: comm
 TYPE(MPI_Group), INTENT(OUT) :: group
 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

`MPI_COMM_GROUP(COMM, GROUP, IERROR)`
 INTEGER COMM, GROUP, IERROR

`MPI_COMM_GROUP` returns in `group` a handle to the group of `comm`.

`MPI_GROUP_UNION(group1, group2, newgroup)`

IN	group1	first group (handle)
IN	group2	second group (handle)
OUT	newgroup	union group (handle)

`int MPI_Group_union(MPI_Group group1, MPI_Group group2,
 MPI_Group *newgroup)`

```

MPI_Group_union(group1, group2, newgroup, ierror)
    TYPE(MPI_Group), INTENT(IN) :: group1, group2
    TYPE(MPI_Group), INTENT(OUT) :: newgroup
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GROUP_UNION(GROUP1, GROUP2, NEWGROUP, IERROR)
    INTEGER GROUP1, GROUP2, NEWGROUP, IERROR

MPI_GROUP_INTERSECTION(group1, group2, newgroup)
    IN      group1      first group (handle)
    IN      group2      second group (handle)
    OUT     newgroup     intersection group (handle)

int MPI_Group_intersection(MPI_Group group1, MPI_Group group2,
    MPI_Group *newgroup)

MPI_Group_intersection(group1, group2, newgroup, ierror)
    TYPE(MPI_Group), INTENT(IN) :: group1, group2
    TYPE(MPI_Group), INTENT(OUT) :: newgroup
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GROUP_INTERSECTION(GROUP1, GROUP2, NEWGROUP, IERROR)
    INTEGER GROUP1, GROUP2, NEWGROUP, IERROR

MPI_GROUP_DIFFERENCE(group1, group2, newgroup)
    IN      group1      first group (handle)
    IN      group2      second group (handle)
    OUT     newgroup     difference group (handle)

int MPI_Group_difference(MPI_Group group1, MPI_Group group2,
    MPI_Group *newgroup)

MPI_Group_difference(group1, group2, newgroup, ierror)
    TYPE(MPI_Group), INTENT(IN) :: group1, group2
    TYPE(MPI_Group), INTENT(OUT) :: newgroup
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GROUP_DIFFERENCE(GROUP1, GROUP2, NEWGROUP, IERROR)
    INTEGER GROUP1, GROUP2, NEWGROUP, IERROR

The set-like operations are defined as follows:

union All elements of the first group (group1), followed by all elements of second group
    (group2) not in the first group.

intersect all elements of the first group that are also in the second group, ordered as in
    the first group.

```

difference all elements of the first group that are not in the second group, ordered as in the first group.

Note that for these operations the order of processes in the output group is determined primarily by order in the first group (if possible) and then, if necessary, by order in the second group. Neither union nor intersection are commutative, but both are associative.

The new group can be empty, that is, equal to `MPI_GROUP_EMPTY`.

`MPI_GROUP_INCL(group, n, ranks, newgroup)`

IN	group	group (handle)
IN	n	number of elements in array ranks (and size of newgroup) (integer)
IN	ranks	ranks of processes in group to appear in newgroup (array of integers)
OUT	newgroup	new group derived from above, in the order defined by ranks (handle)

```
int MPI_Group_incl(MPI_Group group, int n, const int ranks[],
                  MPI_Group *newgroup)
```

```
MPI_Group_incl(group, n, ranks, newgroup, ierror)
```

```
    TYPE(MPI_Group), INTENT(IN) :: group
```

```
    INTEGER, INTENT(IN) :: n, ranks(n)
```

```
    TYPE(MPI_Group), INTENT(OUT) :: newgroup
```

```
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GROUP_INCL(GROUP, N, RANKS, NEWGROUP, IERROR)
```

```
    INTEGER GROUP, N, RANKS(*), NEWGROUP, IERROR
```

The function `MPI_GROUP_INCL` creates a group `newgroup` that consists of the `n` processes in `group` with ranks `ranks[0], ..., ranks[n-1]`; the process with rank `i` in `newgroup` is the process with rank `ranks[i]` in `group`. Each of the `n` elements of `ranks` must be a valid rank in `group` and all elements must be distinct, or else the program is erroneous. If `n = 0`, then `newgroup` is `MPI_GROUP_EMPTY`. This function can, for instance, be used to reorder the elements of a group. See also `MPI_GROUP_COMPARE`.

`MPI_GROUP_EXCL(group, n, ranks, newgroup)`

IN	group	group (handle)
IN	n	number of elements in array ranks (integer)
IN	ranks	array of integer ranks in group not to appear in newgroup
OUT	newgroup	new group derived from above, preserving the order defined by group (handle)

```
int MPI_Group_excl(MPI_Group group, int n, const int ranks[],
                  MPI_Group *newgroup)
```

```
MPI_Group_excl(group, n, ranks, newgroup, ierror)
    TYPE(MPI_Group), INTENT(IN) :: group
    INTEGER, INTENT(IN) :: n, ranks(n)
    TYPE(MPI_Group), INTENT(OUT) :: newgroup
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GROUP_EXCL(GROUP, N, RANKS, NEWGROUP, IERROR)
    INTEGER GROUP, N, RANKS(*), NEWGROUP, IERROR
```

The function `MPI_GROUP_EXCL` creates a group of processes `newgroup` that is obtained by deleting from `group` those processes with ranks `ranks[0]` ... `ranks[n-1]`. The ordering of processes in `newgroup` is identical to the ordering in `group`. Each of the `n` elements of `ranks` must be a valid rank in `group` and all elements must be distinct; otherwise, the program is erroneous. If `n = 0`, then `newgroup` is identical to `group`.

```
MPI_GROUP_RANGE_INCL(group, n, ranges, newgroup)
```

IN	group	group (handle)
IN	n	number of triplets in array <code>ranges</code> (integer)
IN	ranges	a one-dimensional array of integer triplets, of the form (first rank, last rank, stride) indicating ranks in <code>group</code> of processes to be included in <code>newgroup</code>
OUT	newgroup	new group derived from above, in the order defined by <code>ranges</code> (handle)

```
int MPI_Group_range_incl(MPI_Group group, int n, int ranges[][3],
                        MPI_Group *newgroup)
```

```
MPI_Group_range_incl(group, n, ranges, newgroup, ierror)
    TYPE(MPI_Group), INTENT(IN) :: group
    INTEGER, INTENT(IN) :: n, ranges(3,n)
    TYPE(MPI_Group), INTENT(OUT) :: newgroup
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GROUP_RANGE_INCL(GROUP, N, RANGES, NEWGROUP, IERROR)
    INTEGER GROUP, N, RANGES(3,*), NEWGROUP, IERROR
```

If `ranges` consists of the triplets

$$(first_1, last_1, stride_1), \dots, (first_n, last_n, stride_n)$$

then `newgroup` consists of the sequence of processes in `group` with ranks

$$first_1, first_1 + stride_1, \dots, first_1 + \left\lfloor \frac{last_1 - first_1}{stride_1} \right\rfloor stride_1, \dots,$$

$$first_n, first_n + stride_n, \dots, first_n + \left\lfloor \frac{last_n - first_n}{stride_n} \right\rfloor stride_n.$$

Each computed rank must be a valid rank in `group` and all computed ranks must be distinct, or else the program is erroneous. Note that we may have $first_i > last_i$, and $stride_i$ may be negative, but cannot be zero.

The functionality of this routine is specified to be equivalent to expanding the array of ranges to an array of the included ranks and passing the resulting array of ranks and other arguments to `MPI_GROUP_INCL`. A call to `MPI_GROUP_INCL` is equivalent to a call to `MPI_GROUP_RANGE_INCL` with each rank i in `ranges` replaced by the triplet $(i,i,1)$ in the argument `ranges`.

`MPI_GROUP_RANGE_EXCL(group, n, ranges, newgroup)`

IN	group	group (handle)
IN	n	number of elements in array ranges (integer)
IN	ranges	a one-dimensional array of integer triplets of the form (first rank, last rank, stride), indicating the ranks in group of processes to be excluded from the output group newgroup.
OUT	newgroup	new group derived from above, preserving the order in group (handle)

```
int MPI_Group_range_excl(MPI_Group group, int n, int ranges[][3],
                        MPI_Group *newgroup)
```

```
MPI_Group_range_excl(group, n, ranges, newgroup, ierror)
```

```
TYPE(MPI_Group), INTENT(IN) :: group
INTEGER, INTENT(IN) :: n, ranges(3,n)
TYPE(MPI_Group), INTENT(OUT) :: newgroup
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GROUP_RANGE_EXCL(GROUP, N, RANGES, NEWGROUP, IERROR)
```

```
INTEGER GROUP, N, RANGES(3,*), NEWGROUP, IERROR
```

Each computed rank must be a valid rank in `group` and all computed ranks must be distinct, or else the program is erroneous.

The functionality of this routine is specified to be equivalent to expanding the array of ranges to an array of the excluded ranks and passing the resulting array of ranks and other arguments to `MPI_GROUP_EXCL`. A call to `MPI_GROUP_EXCL` is equivalent to a call to `MPI_GROUP_RANGE_EXCL` with each rank i in `ranges` replaced by the triplet $(i,i,1)$ in the argument `ranges`.

Advice to users. The range operations do not explicitly enumerate ranks, and therefore are more scalable if implemented efficiently. Hence, we recommend MPI programmers to use them whenever possible, as high-quality implementations will take advantage of this fact. (*End of advice to users.*)

Advice to implementors. The range operations should be implemented, if possible, without enumerating the group members, in order to obtain better scalability (time and space). (*End of advice to implementors.*)

6.3.3 Group Destructors

MPI_GROUP_FREE(group)

INOUT	group	group (handle)
-------	-------	----------------

int MPI_Group_free(MPI_Group *group)

MPI_Group_free(group, ierror)

TYPE(MPI_Group), INTENT(INOUT) ::	group
-----------------------------------	-------

INTEGER, OPTIONAL, INTENT(OUT) ::	ierror
-----------------------------------	--------

MPI_GROUP_FREE(GROUP, IERROR)

INTEGER GROUP, IERROR

This operation marks a group object for deallocation. The handle `group` is set to `MPI_GROUP_NULL` by the call. Any on-going operation using this group will complete normally.

Advice to implementors. One can keep a reference count that is incremented for each call to `MPI_COMM_GROUP`, `MPI_COMM_CREATE`, `MPI_COMM_DUP`, and `MPI_COMM_IDUP`, and decremented for each call to `MPI_GROUP_FREE` or `MPI_COMM_FREE`; the group object is ultimately deallocated when the reference count drops to zero. (*End of advice to implementors.*)

6.4 Communicator Management

This section describes the manipulation of communicators in MPI. Operations that access communicators are local and their execution does not require interprocess communication. Operations that create communicators are collective and may require interprocess communication.

Advice to implementors. High-quality implementations should amortize the overheads associated with the creation of communicators (for the same group, or subsets thereof) over several calls, by allocating multiple contexts with one collective communication. (*End of advice to implementors.*)

6.4.1 Communicator Accessors

The following are all local operations.

MPI_COMM_SIZE(comm, size)

IN	comm	communicator (handle)
----	------	-----------------------

OUT	size	number of processes in the group of <code>comm</code> (integer)
-----	------	---

int MPI_Comm_size(MPI_Comm comm, int *size)

MPI_Comm_size(comm, size, ierror)

```

1  TYPE(MPI_Comm), INTENT(IN) :: comm
2  INTEGER, INTENT(OUT) :: size
3  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
4
5  MPI_COMM_SIZE(COMM, SIZE, IERROR)
6  INTEGER COMM, SIZE, IERROR
7

```

Rationale. This function is equivalent to accessing the communicator’s group with `MPI_COMM_GROUP` (see above), computing the size using `MPI_GROUP_SIZE`, and then freeing the temporary group via `MPI_GROUP_FREE`. However, this function is so commonly used that this shortcut was introduced. (*End of rationale.*)

Advice to users. This function indicates the number of processes involved in a communicator. For `MPI_COMM_WORLD`, it indicates the total number of processes available unless the number of processes has been changed by using the functions described in Chapter 10; note that the number of processes in `MPI_COMM_WORLD` does not change during the life of an MPI program.

This call is often used with the next call to determine the amount of concurrency available for a specific library or program. The following call, `MPI_COMM_RANK` indicates the rank of the process that calls it in the range from $0 \dots \text{size}-1$, where `size` is the return value of `MPI_COMM_SIZE`. (*End of advice to users.*)

```

24 MPI_COMM_RANK(comm, rank)
25
26 IN      comm      communicator (handle)
27 OUT     rank      rank of the calling process in group of comm (integer)
28
29 int MPI_Comm_rank(MPI_Comm comm, int *rank)
30
31 MPI_Comm_rank(comm, rank, ierror)
32 TYPE(MPI_Comm), INTENT(IN) :: comm
33 INTEGER, INTENT(OUT) :: rank
34 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
35
36 MPI_COMM_RANK(COMM, RANK, IERROR)
37 INTEGER COMM, RANK, IERROR
38

```

Rationale. This function is equivalent to accessing the communicator’s group with `MPI_COMM_GROUP` (see above), computing the rank using `MPI_GROUP_RANK`, and then freeing the temporary group via `MPI_GROUP_FREE`. However, this function is so commonly used that this shortcut was introduced. (*End of rationale.*)

Advice to users. This function gives the rank of the process in the particular communicator’s group. It is useful, as noted above, in conjunction with `MPI_COMM_SIZE`.

Many programs will be written with the master-slave model, where one process (such as the rank-zero process) will play a supervisory role, and the other processes will serve as compute nodes. In this framework, the two preceding calls are useful for

determining the roles of the various processes of a communicator. (*End of advice to users.*)

`MPI_COMM_COMPARE(comm1, comm2, result)`

IN	<code>comm1</code>	first communicator (handle)
IN	<code>comm2</code>	second communicator (handle)
OUT	<code>result</code>	result (integer)

`int MPI_Comm_compare(MPI_Comm comm1, MPI_Comm comm2, int *result)`

```

MPI_Comm_compare(comm1, comm2, result, ierror)
    TYPE(MPI_Comm), INTENT(IN) :: comm1, comm2
    INTEGER, INTENT(OUT) :: result
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_COMM_COMPARE(COMM1, COMM2, RESULT, IERROR)
    INTEGER COMM1, COMM2, RESULT, IERROR

```

`MPI_IDENT` results if and only if `comm1` and `comm2` are handles for the same object (identical groups and same contexts). `MPI_CONGRUENT` results if the underlying groups are identical in constituents and rank order; these communicators differ only by context. `MPI_SIMILAR` results if the group members of both communicators are the same but the rank order differs. `MPI_UNEQUAL` results otherwise.

6.4.2 Communicator Constructors

The following are collective functions that are invoked by all processes in the group or groups associated with `comm`, with the exception of `MPI_COMM_CREATE_GROUP`, which is invoked only by the processes in the group of the new communicator being constructed.

Rationale. Note that there is a chicken-and-egg aspect to MPI in that a communicator is needed to create a new communicator. The base communicator for all MPI communicators is predefined outside of MPI, and is `MPI_COMM_WORLD`. This model was arrived at after considerable debate, and was chosen to increase “safety” of programs written in MPI. (*End of rationale.*)

This chapter presents the following communicator construction routines:

`MPI_COMM_CREATE`, `MPI_COMM_DUP`, `MPI_COMM_IDUP`, `MPI_COMM_DUP_WITH_INFO`, `MPI_COMM_IDUP_WITH_INFO` and `MPI_COMM_SPLIT` can be used to create both intracommunicators and intercommunicators; `MPI_COMM_CREATE_GROUP` and `MPI_INTERCOMM_MERGE` (see Section 6.6.2) can be used to create intracommunicators; and `MPI_INTERCOMM_CREATE` (see Section 6.6.2) can be used to create intercommunicators.

An intracommunicator involves a single group while an intercommunicator involves two groups. Where the following discussions address intercommunicator semantics, the two groups in an intercommunicator are called the *left* and *right* groups. A process in an intercommunicator is a member of either the left or the right group. From the point of view

of that process, the group that the process is a member of is called the *local group*; the other group (relative to that process) is the *remote group*. The left and right group labels give us a way to describe the two groups in an intercommunicator that is not relative to any particular process (as the local and remote groups are).

```

MPI_COMM_DUP(comm, newcomm)
    IN      comm      communicator (handle)
    OUT     newcomm    copy of comm (handle)

int MPI_Comm_dup(MPI_Comm comm, MPI_Comm *newcomm)
MPI_Comm_dup(comm, newcomm, ierror)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Comm), INTENT(OUT) :: newcomm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_COMM_DUP(COMM, NEWCOMM, IERROR)
    INTEGER COMM, NEWCOMM, IERROR

```

`MPI_COMM_DUP` duplicates the existing communicator `comm` with associated key values and topology information. For each key value, the respective copy callback function determines the attribute value associated with this key in the new communicator; one particular action that a copy callback may take is to delete the attribute from the new communicator. `MPI_COMM_DUP` returns in `newcomm` a new communicator with the same group or groups, same topology, and any copied cached information, but a new context (see Section 6.7.1).

Advice to users. This operation is used to provide a parallel library with a duplicate communication space that has the same properties as the original communicator. This includes any attributes (see below) and topologies (see Chapter 7). This call is valid even if there are pending point-to-point communications involving the communicator `comm`. A typical call might involve a `MPI_COMM_DUP` at the beginning of the parallel call, and an `MPI_COMM_FREE` of that duplicated communicator at the end of the call. Other models of communicator management are also possible.

This call applies to both intra- and inter-communicators. (*End of advice to users.*)

Advice to implementors. One need not actually copy the group information, but only add a new reference and increment the reference count. Copy on write can be used for the cached information. (*End of advice to implementors.*)

```

MPI_COMM_DUP_WITH_INFO(comm, info, newcomm)
    IN      comm      communicator (handle)
    IN      info       info object (handle)
    OUT     newcomm    copy of comm (handle)

```

```

int MPI_Comm_dup_with_info(MPI_Comm comm, MPI_Info info, MPI_Comm *newcomm)
MPI_Comm_dup_with_info(comm, info, newcomm, ierror)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Info), INTENT(IN) :: info
    TYPE(MPI_Comm), INTENT(OUT) :: newcomm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_COMM_DUP_WITH_INFO(COMM, INFO, NEWCOMM, IERROR)
    INTEGER COMM, INFO, NEWCOMM, IERROR

```

MPI_COMM_DUP_WITH_INFO behaves exactly as MPI_COMM_DUP except that the hints provided by the argument `info` are associated with the output communicator `newcomm`.

Rationale. It is expected that some hints will only be valid at communicator creation time. However, for legacy reasons, most communicator creation calls do not provide an `info` argument. One may associate `info` hints with a duplicate of any communicator at creation time through a call to MPI_COMM_DUP_WITH_INFO. (*End of rationale.*)

```

MPI_COMM_IDUP(comm, newcomm, request)

```

IN	comm	communicator (handle)
OUT	newcomm	copy of comm (handle)
OUT	request	communication request (handle)

```

int MPI_Comm_idup(MPI_Comm comm, MPI_Comm *newcomm, MPI_Request *request)
MPI_Comm_idup(comm, newcomm, request, ierror)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Comm), INTENT(OUT), ASYNCHRONOUS :: newcomm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_COMM_IDUP(COMM, NEWCOMM, REQUEST, IERROR)
    INTEGER COMM, NEWCOMM, REQUEST, IERROR

```

MPI_COMM_IDUP is a nonblocking variant of MPI_COMM_DUP. With the exception of its nonblocking behavior, the semantics of MPI_COMM_IDUP are as if MPI_COMM_DUP was executed at the time that MPI_COMM_IDUP is called. For example, attributes changed after MPI_COMM_IDUP will not be copied to the new communicator. All restrictions and assumptions for nonblocking collective operations (see Section 5.12) apply to MPI_COMM_IDUP and the returned request.

It is erroneous to use the communicator `newcomm` as an input argument to other MPI functions before the MPI_COMM_IDUP operation completes.

```

1 MPI_COMM_IDUP_WITH_INFO(comm, info, newcomm, request)
2     IN      comm      communicator (handle)
3     IN      info      info object (handle)
4     OUT     newcomm    copy of comm (handle)
5     OUT     request    communication request (handle)
6
7

```

```

8
9 int MPI_Comm_idup_with_info(MPI_Comm comm, MPI_Info info,
10     MPI_Comm *newcomm, MPI_Request *request)
11
12 MPI_Comm_idup_with_info(comm, info, newcomm, request, ierror)
13     TYPE(MPI_Comm), INTENT(IN) :: comm
14     TYPE(MPI_Info), INTENT(IN) :: info
15     TYPE(MPI_Comm), INTENT(OUT), ASYNCHRONOUS :: newcomm
16     TYPE(MPI_Request), INTENT(OUT) :: request
17     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
18
19 MPI_COMM_IDUP_WITH_INFO(COMM, INFO, NEWCOMM, REQUEST, IERROR)
20     INTEGER COMM, INFO, NEWCOMM, REQUEST, IERROR
21

```

MPI_COMM_IDUP_WITH_INFO is a nonblocking variant of MPI_COMM_DUP_WITH_INFO. With the exception of its nonblocking behavior, the semantics of MPI_COMM_IDUP_WITH_INFO are as if MPI_COMM_DUP_WITH_INFO was executed at the time that MPI_COMM_IDUP_WITH_INFO is called. For example, attributes or info hints changed after MPI_COMM_IDUP_WITH_INFO will not be copied to the new communicator. All restrictions and assumptions for nonblocking collective operations (see Section 5.12) apply to MPI_COMM_IDUP_WITH_INFO and the returned request.

It is erroneous to use the communicator newcomm as an input argument to other MPI functions before the MPI_COMM_IDUP_WITH_INFO operation completes.

Rationale. The MPI_COMM_IDUP and MPI_COMM_IDUP_WITH_INFO functions are crucial for the development of purely nonblocking libraries (see [5]). (*End of rationale.*)

```

35 MPI_COMM_CREATE(comm, group, newcomm)
36
37     IN      comm      communicator (handle)
38     IN      group      group, which is a subset of the group of comm (handle)
39     OUT     newcomm    new communicator (handle)
40
41

```

```

42 int MPI_Comm_create(MPI_Comm comm, MPI_Group group, MPI_Comm *newcomm)
43
44 MPI_Comm_create(comm, group, newcomm, ierror)
45     TYPE(MPI_Comm), INTENT(IN) :: comm
46     TYPE(MPI_Group), INTENT(IN) :: group
47     TYPE(MPI_Comm), INTENT(OUT) :: newcomm
48     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
49

```

```

MPI_COMM_CREATE(COMM, GROUP, NEWCOMM, IERROR)
    INTEGER COMM, GROUP, NEWCOMM, IERROR

```

If `comm` is an intracommunicator, this function returns a new communicator `newcomm` with communication group defined by the `group` argument. No cached information propagates from `comm` to `newcomm`. Each process must call `MPI_COMM_CREATE` with a `group` argument that is a subgroup of the `group` associated with `comm`; this could be `MPI_GROUP_EMPTY`. The processes may specify different values for the `group` argument. If a process calls with a non-empty `group` then all processes in that `group` must call the function with the same `group` as argument, that is the same processes in the same order. Otherwise, the call is erroneous. This implies that the set of groups specified across the processes must be disjoint. If the calling process is a member of the group given as `group` argument, then `newcomm` is a communicator with `group` as its associated group. In the case that a process calls with a `group` to which it does not belong, e.g., `MPI_GROUP_EMPTY`, then `MPI_COMM_NULL` is returned as `newcomm`. The function is collective and must be called by all processes in the group of `comm`.

Rationale. The interface supports the original mechanism from MPI-1.1, which required the same `group` in all processes of `comm`. It was extended in MPI-2.2 to allow the use of disjoint subgroups in order to allow implementations to eliminate unnecessary communication that `MPI_COMM_SPLIT` would incur when the user already knows the membership of the disjoint subgroups. (*End of rationale.*)

Rationale. The requirement that the entire group of `comm` participate in the call stems from the following considerations:

- It allows the implementation to layer `MPI_COMM_CREATE` on top of regular collective communications.
- It provides additional safety, in particular in the case where partially overlapping groups are used to create new communicators.
- It permits implementations to sometimes avoid communication related to context creation.

(*End of rationale.*)

Advice to users. `MPI_COMM_CREATE` provides a means to subset a group of processes for the purpose of separate MIMD computation, with separate communication space. `newcomm`, which emerges from `MPI_COMM_CREATE`, can be used in subsequent calls to `MPI_COMM_CREATE` (or other communicator constructors) to further subdivide a computation into parallel sub-computations. A more general service is provided by `MPI_COMM_SPLIT`, below. (*End of advice to users.*)

Advice to implementors. When calling `MPI_COMM_DUP`, all processes call with the same `group` (the `group` associated with the communicator). When calling `MPI_COMM_CREATE`, the processes provide the same `group` or disjoint subgroups. For both calls, it is theoretically possible to agree on a group-wide unique context with no communication. However, local execution of these functions requires use of a larger context name space and reduces error checking. Implementations may strike various compromises between these conflicting goals, such as bulk allocation of multiple contexts in one collective operation.

Important: If new communicators are created without synchronizing the processes involved then the communication system must be able to cope with messages arriving in a context that has not yet been allocated at the receiving process. (*End of advice to implementors.*)

If `comm` is an intercommunicator, then the output communicator is also an intercommunicator where the local group consists only of those processes contained in `group` (see Figure 6.1). The `group` argument should only contain those processes in the local group of the input intercommunicator that are to be a part of `newcomm`. All processes in the same local group of `comm` must specify the same value for `group`, i.e., the same members in the same order. If either `group` does not specify at least one process in the local group of the intercommunicator, or if the calling process is not included in the `group`, `MPI_COMM_NULL` is returned.

Rationale. In the case where either the left or right group is empty, a null communicator is returned instead of an intercommunicator with `MPI_GROUP_EMPTY` because the side with the empty group must return `MPI_COMM_NULL`. (*End of rationale.*)

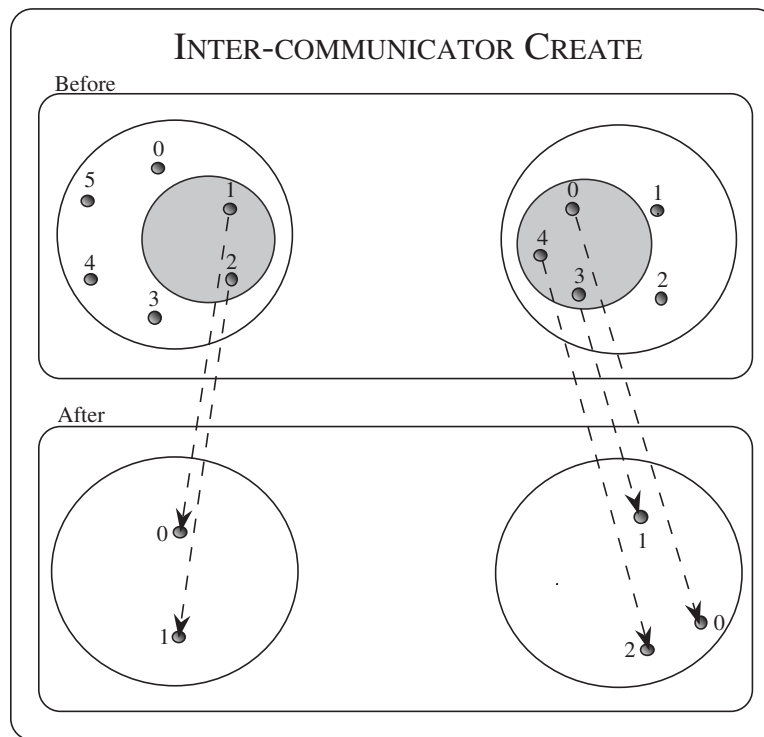


Figure 6.1: Intercommunicator creation using `MPI_COMM_CREATE` extended to intercommunicators. The input groups are those in the grey circle.

Example 6.1 The following example illustrates how the first node in the left side of an intercommunicator could be joined with all members on the right side of an intercommunicator to form a new intercommunicator.

```

MPI_Comm inter_comm, new_inter_comm;
MPI_Group local_group, group;
int      rank = 0; /* rank on left side to include in
                    new inter-comm */

/* Construct the original intercommunicator: "inter_comm" */
...

/* Construct the group of processes to be in new
   intercommunicator */
if (/* I'm on the left side of the intercommunicator */) {
    MPI_Comm_group(inter_comm, &local_group);
    MPI_Group_incl(local_group, 1, &rank, &group);
    MPI_Group_free(&local_group);
}
else
    MPI_Comm_group(inter_comm, &group);

MPI_Comm_create(inter_comm, group, &new_inter_comm);
MPI_Group_free(&group);

MPI_COMM_CREATE_GROUP(comm, group, tag, newcomm)

IN      comm      intracommunicator (handle)
IN      group     group, which is a subset of the group of comm (handle)
IN      tag       tag (integer)
OUT     newcomm   new communicator (handle)

int MPI_Comm_create_group(MPI_Comm comm, MPI_Group group, int tag,
                        MPI_Comm *newcomm)

MPI_Comm_create_group(comm, group, tag, newcomm, ierror)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Group), INTENT(IN) :: group
    INTEGER, INTENT(IN) :: tag
    TYPE(MPI_Comm), INTENT(OUT) :: newcomm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_COMM_CREATE_GROUP(COMM, GROUP, TAG, NEWCOMM, IERROR)
    INTEGER COMM, GROUP, TAG, NEWCOMM, IERROR

MPI_COMM_CREATE_GROUP is similar to MPI_COMM_CREATE; however,
MPI_COMM_CREATE must be called by all processes in the group of
comm, whereas MPI_COMM_CREATE_GROUP must be called by all processes in group,
which is a subgroup of the group of comm. In addition, MPI_COMM_CREATE_GROUP
requires that comm is an intracommunicator. MPI_COMM_CREATE_GROUP returns a new
intracommunicator, newcomm, for which the group argument defines the communication

```

group. No cached information propagates from `comm` to `newcomm`. Each process must provide a group argument that is a subgroup of the group associated with `comm`; this could be `MPI_GROUP_EMPTY`. If a non-empty group is specified, then all processes in that group must call the function, and each of these processes must provide the same arguments, including a group that contains the same members with the same ordering. Otherwise the call is erroneous. If the calling process is a member of the group given as the `group` argument, then `newcomm` is a communicator with `group` as its associated group. If the calling process is not a member of `group`, e.g., `group` is `MPI_GROUP_EMPTY`, then the call is a local operation and `MPI_COMM_NULL` is returned as `newcomm`.

Rationale. Functionality similar to `MPI_COMM_CREATE_GROUP` can be implemented through repeated `MPI_INTERCOMM_CREATE` and `MPI_INTERCOMM_MERGE` calls that start with the `MPI_COMM_SELF` communicators at each process in `group` and build up an intracommunicator with `group` `group` [3]. Such an algorithm requires the creation of many intermediate communicators; `MPI_COMM_CREATE_GROUP` can provide a more efficient implementation that avoids this overhead. (*End of rationale.*)

Advice to users. An intercommunicator can be created collectively over processes in the union of the local and remote groups by creating the local communicator using `MPI_COMM_CREATE_GROUP` and using that communicator as the local communicator argument to `MPI_INTERCOMM_CREATE`. (*End of advice to users.*)

The `tag` argument does not conflict with tags used in point-to-point communication and is not permitted to be a wildcard. If multiple threads at a given process perform concurrent `MPI_COMM_CREATE_GROUP` operations, the user must distinguish these operations by providing different `tag` or `comm` arguments.

Advice to users. `MPI_COMM_CREATE` may provide lower overhead than `MPI_COMM_CREATE_GROUP` because it can take advantage of collective communication on `comm` when constructing `newcomm`. (*End of advice to users.*)

`MPI_COMM_SPLIT(comm, color, key, newcomm)`

IN	<code>comm</code>	communicator (handle)
IN	<code>color</code>	control of subset assignment (integer)
IN	<code>key</code>	control of rank assignment (integer)
OUT	<code>newcomm</code>	new communicator (handle)

`int MPI_Comm_split(MPI_Comm comm, int color, int key, MPI_Comm *newcomm)`

`MPI_Comm_split(comm, color, key, newcomm, ierror)`

`TYPE(MPI_Comm), INTENT(IN) :: comm`

`INTEGER, INTENT(IN) :: color, key`

`TYPE(MPI_Comm), INTENT(OUT) :: newcomm`

`INTEGER, OPTIONAL, INTENT(OUT) :: ierror`

`MPI_COMM_SPLIT(COMM, COLOR, KEY, NEWCOMM, IERROR)`

INTEGER COMM, COLOR, KEY, NEWCOMM, IERROR

This function partitions the group associated with `comm` into disjoint subgroups, one for each value of `color`. Each subgroup contains all processes of the same color. Within each subgroup, the processes are ranked in the order defined by the value of the argument `key`, with ties broken according to their rank in the old group. A new communicator is created for each subgroup and returned in `newcomm`. A process may supply the color value `MPI_UNDEFINED`, in which case `newcomm` returns `MPI_COMM_NULL`. This is a collective call, but each process is permitted to provide different values for `color` and `key`.

With an intracommunicator `comm`, a call to `MPI_COMM_CREATE(comm, group, newcomm)` is equivalent to a call to `MPI_COMM_SPLIT(comm, color, key, newcomm)`, where processes that are members of their `group` argument provide `color` = number of the `group` (based on a unique numbering of all disjoint groups) and `key` = rank in `group`, and all processes that are not members of their `group` argument provide `color` = `MPI_UNDEFINED`.

The value of `color` must be non-negative or `MPI_UNDEFINED`.

Advice to users. This is an extremely powerful mechanism for dividing a single communicating group of processes into k subgroups, with k chosen implicitly by the user (by the number of colors asserted over all the processes). Each resulting communicator will be non-overlapping. Such a division could be useful for defining a hierarchy of computations, such as for multigrid, or linear algebra. For intracommunicators, `MPI_COMM_SPLIT` provides similar capability as `MPI_COMM_CREATE` to split a communicating group into disjoint subgroups. `MPI_COMM_SPLIT` is useful when some processes do not have complete information of the other members in their group, but all processes know (the color of) the group to which they belong. In this case, the MPI implementation discovers the other group members via communication. `MPI_COMM_CREATE` is useful when all processes have complete information of the members of their group. In this case, MPI can avoid the extra communication required to discover group membership. `MPI_COMM_CREATE_GROUP` is useful when all processes in a given group have complete information of the members of their group and synchronization with processes outside the group can be avoided.

Multiple calls to `MPI_COMM_SPLIT` can be used to overcome the requirement that any call have no overlap of the resulting communicators (each process is of only one color per call). In this way, multiple overlapping communication structures can be created. Creative use of the `color` and `key` in such splitting operations is encouraged.

Note that, for a fixed color, the keys need not be unique. It is `MPI_COMM_SPLIT`'s responsibility to sort processes in ascending order according to this key, and to break ties in a consistent way. If all the keys are specified in the same way, then all the processes in a given color will have the relative rank order as they did in their parent group.

Essentially, making the key value zero for all processes of a given color means that one does not really care about the rank-order of the processes in the new communicator. (*End of advice to users.*)

Rationale. `color` is restricted to be non-negative, so as not to conflict with the value assigned to `MPI_UNDEFINED`. (*End of rationale.*)

The result of `MPI_COMM_SPLIT` on an intercommunicator is that those processes on the left with the same color as those processes on the right combine to create a new intercom-

municator. The key argument describes the relative rank of processes on each side of the intercommunicator (see Figure 6.2). For those colors that are specified only on one side of the intercommunicator, MPI_COMM_NULL is returned. MPI_COMM_NULL is also returned to those processes that specify MPI_UNDEFINED as the color.

Advice to users. For intercommunicators, MPI_COMM_SPLIT is more general than MPI_COMM_CREATE. A single call to MPI_COMM_SPLIT can create a set of disjoint intercommunicators, while a call to MPI_COMM_CREATE creates only one. (*End of advice to users.*)

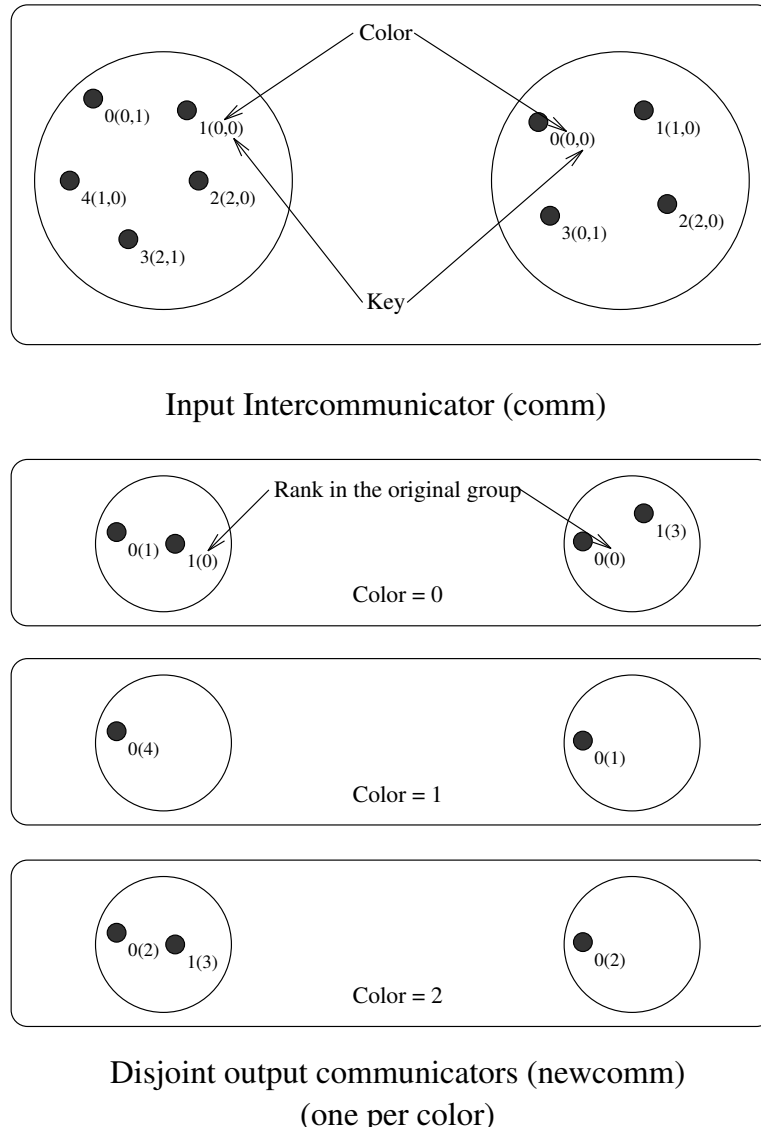


Figure 6.2: Intercommunicator construction achieved by splitting an existing intercommunicator with MPI_COMM_SPLIT extended to intercommunicators.

Example 6.2 (Parallel client-server model). The following client code illustrates how clients on the left side of an intercommunicator could be assigned to a single server from a pool of servers on the right side of an intercommunicator.

```

/* Client code */
MPI_Comm  multiple_server_comm;
MPI_Comm  single_server_comm;
int        color, rank, num_servers;

/* Create intercommunicator with clients and servers:
   multiple_server_comm */
...

/* Find out the number of servers available */
MPI_Comm_remote_size(multiple_server_comm, &num_servers);

/* Determine my color */
MPI_Comm_rank(multiple_server_comm, &rank);
color = rank % num_servers;

/* Split the intercommunicator */
MPI_Comm_split(multiple_server_comm, color, rank,
               &single_server_comm);

```

The following is the corresponding server code:

```

/* Server code */
MPI_Comm  multiple_client_comm;
MPI_Comm  single_server_comm;
int        rank;

/* Create intercommunicator with clients and servers:
   multiple_client_comm */
...

/* Split the intercommunicator for a single server per group
   of clients */
MPI_Comm_rank(multiple_client_comm, &rank);
MPI_Comm_split(multiple_client_comm, rank, 0,
               &single_server_comm);

```

MPI_COMM_SPLIT_TYPE(comm, split_type, key, info, newcomm)

IN	comm	communicator (handle)
IN	split_type	type of processes to be grouped together (integer)
IN	key	control of rank assignment (integer)
IN	info	info argument (handle)
OUT	newcomm	new communicator (handle)

```

int MPI_Comm_split_type(MPI_Comm comm, int split_type, int key,
                        MPI_Info info, MPI_Comm *newcomm)

```

```

1 MPI_Comm_split_type(comm, split_type, key, info, newcomm, ierror)
2     TYPE(MPI_Comm), INTENT(IN) :: comm
3     INTEGER, INTENT(IN) :: split_type, key
4     TYPE(MPI_Info), INTENT(IN) :: info
5     TYPE(MPI_Comm), INTENT(OUT) :: newcomm
6     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
7
8 MPI_COMM_SPLIT_TYPE(COMM, SPLIT_TYPE, KEY, INFO, NEWCOMM, IERROR)
9     INTEGER COMM, SPLIT_TYPE, KEY, INFO, NEWCOMM, IERROR

```

This function partitions the group associated with `comm` into disjoint subgroups, based on the type specified by `split_type`. Each subgroup contains all processes of the same type. Within each subgroup, the processes are ranked in the order defined by the value of the argument `key`, with ties broken according to their rank in the old group. A new communicator is created for each subgroup and returned in `newcomm`. This is a collective call; all processes must provide the same `split_type`, but each process is permitted to provide different values for `key`. An exception to this rule is that a process may supply the type value `MPI_UNDEFINED`, in which case `newcomm` returns `MPI_COMM_NULL`.

The following types are predefined by MPI:

`MPI_COMM_TYPE_SHARED` — this type splits the communicator into subcommunicators, each of which can create a shared memory region.

`MPI_COMM_TYPE_HW_SUBDOMAIN` — all MPI processes in `newcomm` share a **hardware resource** (e.g., a network switch, a computing core, an L3 cache, a GPU).

When this value is used, the MPI implementation considers the hardware as hierarchically structured and presents it to the application as a hierarchy of resources. More specifically, such a hierarchy is composed of a series of levels where each level comprises a finite number of non-overlapping sets of homogeneous resources. Moreover, each set of resources located at a certain level in the hierarchy is also fully encompassed in a set of resources located in a higher (previous) level, except for the top level. The potential connectivity between resources located at a certain level in the hierarchy are not considered, nor presented to the application when using the value `MPI_COMM_TYPE_HW_SUBDOMAIN`.

Advice to users. The **process binding**, (i.e. the set of resources an MPI process is restricted to) influences the result of the splitting operation. In particular, when an MPI process is bound to a resource located at a specific level in the hardware hierarchy, such level is considered for the MPI process as its *bottom (last) level*. That is, the extent of the hardware hierarchy that a particular MPI process has access to will change according to the binding of such MPI process. If the MPI process binding is modified during an application execution, the communicators created the value with `MPI_COMM_TYPE_HW_SUBDOMAIN` may not reflect the actual status of the hardware locality of the MPI process. (*End of advice to users.*)

Advice to implementors. A high quality implementation will return in the group of the output communicator `newcomm` the largest subset of MPI processes that fit the splitting criterion. (*End of advice to implementors.*)

There are two possible modes for `MPI_COMM_TYPE_HW_SUBDOMAIN`: the *guided* mode and the *unguided* mode.

With the *guided* mode, the user can further constrain with the `info` argument the splitting of the input communicator. For this, the `info` key `mpi_hw_domain_type` is reserved and its value is an implementation-defined string designating the requested hardware resource, for instance: “NUMANode”, “Package” or “L2”. As an exception, the value `mpi_shared_memory` is defined and reserved in order to produce the same communicators as the ones that would be created if the `MPI_COMM_TYPE_SHARED` value was used for the `split_type` parameter.

Rationale. The value `mpi_shared_memory` is defined in order to ensure consistency between the use of `MPI_COMM_TYPE_SHARED` and the use of `MPI_COMM_TYPE_HW_SUBDOMAIN` with the *guided* mode. (*End of rationale.*)

If the requested hardware resource does not exist in the hierarchy or is unavailable (e.g., a CPU that is shut down during the application execution) or if the requested resource is located at the *bottom level* in the hierarchy for the calling MPI process, then `MPI_COMM_NULL` is returned in `newcomm` for such process.

The *unguided* mode is enabled when the `info` argument is not used for the splitting of the input communicator (i.e., the key `mpi_hw_domain_type` must not be set in the `info` object provided). The *unguided* mode allows `MPI_COMM_SPLIT_TYPE` to split an input communicator for which the MPI process associated with its group all share a hardware resource A into n subcommunicators for which the MPI process associated with their respective subgroups all share a hardware resource B_i with $i \in \{0, \dots, n-1\}$ and where all B_i are included into A (see Figure 6.3 for an example).

It is then possible to expose the hierarchical nature of the underlying hardware by recursively splitting the communicators with a series of calls to `MPI_COMM_SPLIT_TYPE`, as shown in Example 6.3. The set of hierarchy levels used for the unguided splitting is implementation dependent.

Advice to implementors. This set should reflect the hardware of the actual system on which the application is started. (*End of advice to implementors.*)

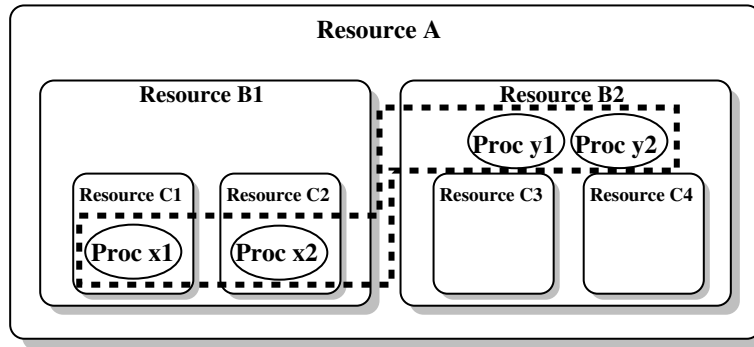
Example 6.3 (Recursive splitting of a communicator).

```
#define MAX_NUM_DOMAINS 32

MPI_Comm hcomm[MAX_NUM_DOMAINS];
int rank, domain_num = 0;

hcomm[domain_num] = MPI_COMM_WORLD;

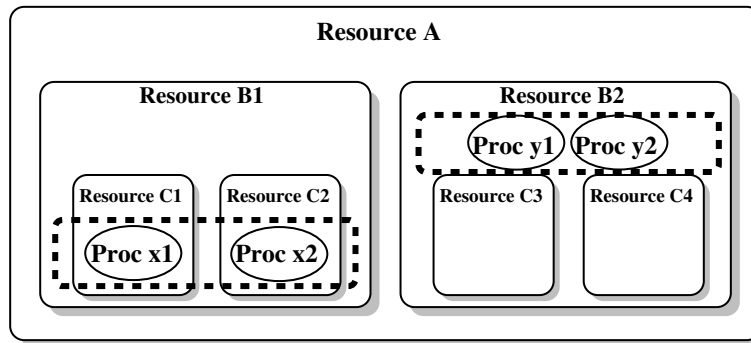
while((hcomm[domain_num] != MPI_COMM_NULL)
    && (domain_num < MAX_NUM_DOMAINS-1)){
    MPI_Comm_rank(hcomm[domain_num], &rank);
    MPI_Comm_split_type(hcomm[domain_num],
                        MPI_COMM_TYPE_HW_SUBDOMAIN,
```



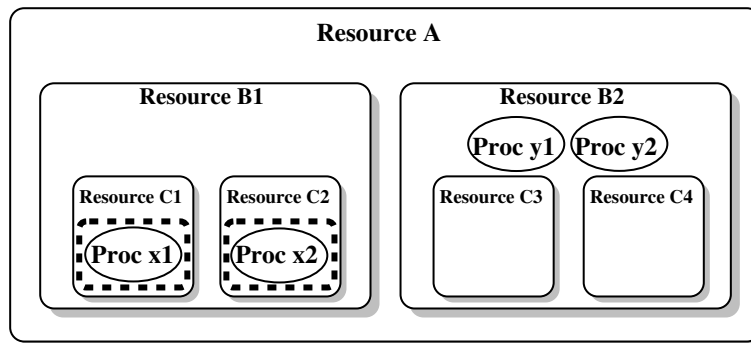
Resource A (original) communicator is in dashed line.

Process x1 is restricted to Resource C1, x2 to C2.

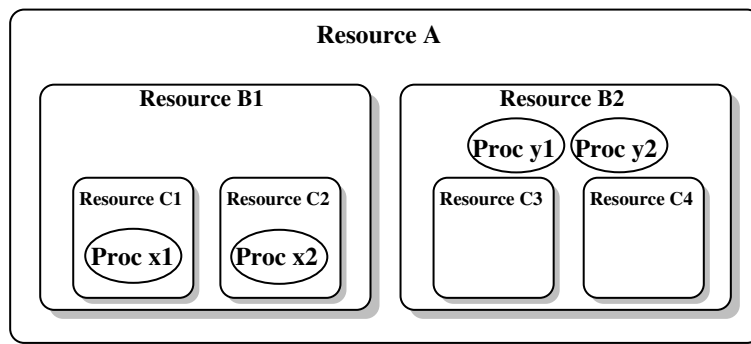
Processes y1 and y2 are restricted to B2.



Step 1: splitting of Resource A communicator
into Resource B1 and Resource B2 communicators.



Step 2: splitting of Resource B1 and B2 communicators
into Resource C1, C2 and NULL communicators



Step 3: splitting of Resource C1 and C2 communicators
into NULL communicators

Figure 6.3: Recursive splitting of an existing communicator with `MPI_COMM_SPLIT_TYPE` and `MPI_COMM_TYPE_HW_SUBDOMAIN` in *unguided* mode.

```

        rank,
        MPI_INFO_NULL,
        &hcomm[domain_num+1]);
    domain_num++;
}

```

To avoid infinite recursion in *unguided* mode, the group of MPI processes associated with the newly created communicator `newcomm` must be a *strict* subset of the group associated with the input (parent) communicator `comm`. More specifically, a call to `MPI_COMM_COMPARE(comm, newcomm, result)` shall return `result = MPI_UNEQUAL` when the input communicator `comm` is valid. In the case where further splitting of the input communicator is not possible because the *bottom level* in the hierarchy has already been reached then `MPI_COMM_NULL` is returned in `newcomm` in such process.

Advice to implementors. The sets of resources used for the splitting operation are implementation dependent. Moreover, in the case where the newly created communicator corresponds to several different hardware resources (for instance, a resource A featuring a single resource B featuring in its turn a single resource C), the selection of the resource (A , B or C) is implementation defined. (*End of advice to implementors.*)

Advice to implementors. Implementations can define their own `split_type` values, or use the `info` argument, to assist in creating communicators based on their own criteria, other than hardware. (*End of advice to implementors.*)

6.4.3 Communicator Destructors

MPI_COMM_FREE(comm)

INOUT	comm	communicator to be destroyed (handle)
-------	------	---------------------------------------

```
int MPI_Comm_free(MPI_Comm *comm)
```

```
MPI_Comm_free(comm, ierror)
```

```
TYPE(MPI_Comm), INTENT(INOUT) :: comm
```

```
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

MPI_COMM_FREE(COMM, IERROR)

INTEGER COMM, IERROR

This collective operation marks the communication object for deallocation. The handle is set to `MPI_COMM_NULL`. Any pending operations that use this communicator will complete normally; the object is actually deallocated only if there are no other active references to it. This call applies to intra- and inter-communicators. The delete callback functions for all cached attributes (see Section 6.7) are called in arbitrary order.

Advice to implementors. Though collective, it is anticipated that this operation will normally be implemented to be local, though a debugging version of an MPI library might choose to synchronize. (*End of advice to implementors.*)

6.4.4 Communicator Info

Hints specified via info (see Chapter 9) allow a user to provide information to direct optimization. Providing hints may enable an implementation to deliver increased performance or minimize use of system resources. An implementation is free to ignore all hints; however, applications must comply with any info hints they provide that are used by the MPI implementation (i.e., are returned by a call to `MPI_COMM_GET_INFO`) and that place a restriction on the behavior of the application. Hints are specified on a per communicator basis, in `MPI_COMM_DUP_WITH_INFO`, `MPI_COMM_IDUP_WITH_INFO`, `MPI_COMM_SET_INFO`, `MPI_COMM_SPLIT_TYPE`, `MPI_DIST_GRAPH_CREATE`, and `MPI_DIST_GRAPH_CREATE_ADJACENT`, via the opaque info object. When an info object that specifies a subset of valid hints is passed to `MPI_COMM_SET_INFO`, there will be no effect on previously set or defaulted hints that the info does not specify.

Advice to implementors. It may happen that a program is coded with hints for one system, and later executes on another system that does not support these hints. In general, unsupported hints should simply be ignored. Needless to say, no hint can be mandatory. However, for each hint used by a specific implementation, a default value must be provided when the user does not specify a value for this hint. (*End of advice to implementors.*)

Info hints are not propagated by MPI from one communicator to another. The following info keys are valid for all communicators.

`mpi_assert_no_any_tag` (**boolean, default: false**): If set to true, then the implementation may assume that the process will not use the `MPI_ANY_TAG` wildcard on the given communicator.

`mpi_assert_no_any_source` (**boolean, default: false**): If set to true, then the implementation may assume that the process will not use the `MPI_ANY_SOURCE` wildcard on the given communicator.

`mpi_assert_exact_length` (**boolean, default: false**): If set to true, then the implementation may assume that the lengths of messages received by the process are equal to the lengths of the corresponding receive buffers, for point-to-point communication operations on the given communicator.

`mpi_assert_allow_overtaking` (**boolean, default: false**): If set to true, then the implementation may assume that point-to-point communications on the given communicator do not rely on the non-overtaking rule specified in Section 3.5. In other words, the application asserts that send operations are not required to be matched at the receiver in the order in which the send operations were posted by the sender, and receive operations are not required to be matched in the order in which they were posted by the receiver.

Advice to users. Use of the `mpi_assert_allow_overtaking` info key can result in non-determinism in the message matching order. (*End of advice to users.*)

Advice to users. Some optimizations may only be possible when all processes in the group of the communicator provide a given info key with the same value. (*End of advice to users.*)

MPI_COMM_SET_INFO(comm, info)

INOUT comm communicator (handle)
IN info info object (handle)

int MPI_Comm_set_info(MPI_Comm comm, MPI_Info info)

MPI_Comm_set_info(comm, info, ierror)
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Info), INTENT(IN) :: info
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_COMM_SET_INFO(COMM, INFO, IERROR)
INTEGER COMM, INFO, IERROR

MPI_COMM_SET_INFO updates the hints of the communicator associated with `comm` using the hints provided in `info`. This operation has no effect on previously set or defaulted hints that are not specified by `info`. It also has no effect on previously set or defaulted hints that are specified by `info`, but are ignored by the MPI implementation in this call to MPI_COMM_SET_INFO. MPI_COMM_SET_INFO is a collective routine. The `info` object may be different on each process, but any `info` entries that an implementation requires to be the same on all processes must appear with the same value in each process's `info` object.

Advice to users. Some `info` items that an implementation can use when it creates a communicator cannot easily be changed once the communicator has been created. Thus, an implementation may ignore hints issued in this call that it would have accepted in a creation call. An implementation may also be unable to update certain `info` hints in a call to MPI_COMM_SET_INFO. MPI_COMM_GET_INFO can be used to determine whether updates to existing `info` hints were ignored by the implementation. (*End of advice to users.*)

Advice to users. Setting `info` hints on the predefined communicators MPI_COMM_WORLD and MPI_COMM_SELF may have unintended effects, as changes to these global objects may affect all components of the application, including libraries and tools. Users must ensure that all components of the application that use a given communicator, including libraries and tools, can comply with any `info` hints associated with that communicator. (*End of advice to users.*)

MPI_COMM_GET_INFO(comm, info_used)

IN comm communicator object (handle)
OUT info_used new `info` object (handle)

int MPI_Comm_get_info(MPI_Comm comm, MPI_Info *info_used)

MPI_Comm_get_info(comm, info_used, ierror)
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Info), INTENT(OUT) :: info_used
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

1 MPI_COMM_GET_INFO(COMM, INFO_USED, IERROR)
2     INTEGER COMM, INFO_USED, IERROR

```

MPI_COMM_GET_INFO returns a new info object containing the hints of the communicator associated with `comm`. The current setting of all hints related to this communicator is returned in `info_used`. An MPI implementation is required to return all hints that are supported by the implementation and have default values specified; any user-supplied hints that were not ignored by the implementation; and any additional hints that were set by the implementation. If no such hints exist, a handle to a newly created info object is returned that contains no key/value pair. The user is responsible for freeing `info_used` via `MPI_INFO_FREE`.

6.5 Motivating Examples

6.5.1 Current Practice #1

Example #1a:

```

18 int main(int argc, char *argv[])
19 {
20     int me, size;
21     ...
22     MPI_Init(&argc, &argv);
23     MPI_Comm_rank(MPI_COMM_WORLD, &me);
24     MPI_Comm_size(MPI_COMM_WORLD, &size);
25
26     (void)printf("Process %d size %d\n", me, size);
27     ...
28     MPI_Finalize();
29     return 0;
30 }

```

Example #1a is a do-nothing program that initializes itself, and refers to the “all” communicator, and prints a message. It terminates itself too. This example does not imply that MPI supports `printf`-like communication itself.

Example #1b (supposing that `size` is even):

```

36 int main(int argc, char *argv[])
37 {
38     int me, size;
39     int SOME_TAG = 0;
40     ...
41     MPI_Init(&argc, &argv);
42
43     MPI_Comm_rank(MPI_COMM_WORLD, &me); /* local */
44     MPI_Comm_size(MPI_COMM_WORLD, &size); /* local */
45
46     if((me % 2) == 0)
47     {
48         /* send unless highest-numbered process */

```

```

        if((me + 1) < size)
            MPI_Send(..., me + 1, SOME_TAG, MPI_COMM_WORLD);
    }
    else
        MPI_Recv(..., me - 1, SOME_TAG, MPI_COMM_WORLD, &status);

    ...
    MPI_Finalize();
    return 0;
}

```

Example #1b schematically illustrates message exchanges between “even” and “odd” processes in the “all” communicator.

6.5.2 Current Practice #2

```

int main(int argc, char *argv[])
{
    int me, count;
    void *data;
    ...

    MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &me);

    if(me == 0)
    {
        /* get input, create buffer ‘data’ */
        ...
    }

    MPI_Bcast(data, count, MPI_BYTE, 0, MPI_COMM_WORLD);

    ...
    MPI_Finalize();
    return 0;
}

```

This example illustrates the use of a collective communication.

6.5.3 (Approximate) Current Practice #3

```

int main(int argc, char *argv[])
{
    int me, count, count2;
    void *send_buf, *recv_buf, *send_buf2, *recv_buf2;
    MPI_Group group_world, grprem;
    MPI_Comm commslave;
    static int ranks[] = {0};

```

```

1      ...
2      MPI_Init(&argc, &argv);
3      MPI_Comm_group(MPI_COMM_WORLD, &group_world);
4      MPI_Comm_rank(MPI_COMM_WORLD, &me); /* local */
5
6      MPI_Group_excl(group_world, 1, ranks, &grpem); /* local */
7      MPI_Comm_create(MPI_COMM_WORLD, grpem, &commslave);
8
9      if(me != 0)
10     {
11         /* compute on slave */
12         ...
13         MPI_Reduce(send_buf,recv_buf,count, MPI_INT, MPI_SUM, 1, commslave);
14         ...
15         MPI_Comm_free(&commslave);
16     }
17     /* zero falls through immediately to this reduce, others do later... */
18     MPI_Reduce(send_buf2, recv_buf2, count2,
19               MPI_INT, MPI_SUM, 0, MPI_COMM_WORLD);
20
21     MPI_Group_free(&group_world);
22     MPI_Group_free(&grpem);
23     MPI_Finalize();
24     return 0;
25 }

```

This example illustrates how a group consisting of all but the zeroth process of the “all” group is created, and then how a communicator is formed (`commslave`) for that new group. The new communicator is used in a collective call, and all processes execute a collective call in the `MPI_COMM_WORLD` context. This example illustrates how the two communicators (that inherently possess distinct contexts) protect communication. That is, communication in `MPI_COMM_WORLD` is insulated from communication in `commslave`, and vice versa.

In summary, “group safety” is achieved via communicators because distinct contexts within communicators are enforced to be unique on any process.

6.5.4 Example #4

The following example is meant to illustrate “safety” between point-to-point and collective communication. MPI guarantees that a single communicator can do safe point-to-point and collective communication.

```

41     #define TAG_ARBITRARY 12345
42     #define SOME_COUNT    50
43
44     int main(int argc, char *argv[])
45     {
46         int me;
47         MPI_Request request[2];
48         MPI_Status status[2];

```

```

MPI_Group group_world, subgroup;
int ranks[] = {2, 4, 6, 8};
MPI_Comm the_comm;
...
MPI_Init(&argc, &argv);
MPI_Comm_group(MPI_COMM_WORLD, &group_world);

MPI_Group_incl(group_world, 4, ranks, &subgroup); /* local */
MPI_Group_rank(subgroup, &me); /* local */

MPI_Comm_create(MPI_COMM_WORLD, subgroup, &the_comm);

if(me != MPI_UNDEFINED)
{
    MPI_Irecv(buff1, count, MPI_DOUBLE, MPI_ANY_SOURCE, TAG_ARBITRARY,
              the_comm, request);
    MPI_Isend(buff2, count, MPI_DOUBLE, (me+1)%4, TAG_ARBITRARY,
              the_comm, request+1);
    for(i = 0; i < SOME_COUNT; i++)
        MPI_Reduce(..., the_comm);
    MPI_Waitall(2, request, status);

    MPI_Comm_free(&the_comm);
}

MPI_Group_free(&group_world);
MPI_Group_free(&subgroup);
MPI_Finalize();
return 0;
}

```

6.5.5 Library Example #1

The main program:

```

int main(int argc, char *argv[])
{
    int done = 0;
    user_lib_t *libh_a, *libh_b;
    void *dataset1, *dataset2;
    ...
    MPI_Init(&argc, &argv);
    ...
    init_user_lib(MPI_COMM_WORLD, &libh_a);
    init_user_lib(MPI_COMM_WORLD, &libh_b);
    ...
    user_start_op(libh_a, dataset1);
    user_start_op(libh_b, dataset2);
}

```

```

1      ...
2      while(!done)
3      {
4          /* work */
5          ...
6          MPI_Reduce(..., MPI_COMM_WORLD);
7          ...
8          /* see if done */
9          ...
10     }
11     user_end_op(libh_a);
12     user_end_op(libh_b);
13
14     uninit_user_lib(libh_a);
15     uninit_user_lib(libh_b);
16     MPI_Finalize();
17     return 0;
18 }

```

The user library initialization code:

```

21     void init_user_lib(MPI_Comm comm, user_lib_t **handle)
22     {
23         user_lib_t *save;
24
25         user_lib_initsave(&save); /* local */
26         MPI_Comm_dup(comm, &(save->comm));
27
28         /* other inits */
29         ...
30
31         *handle = save;
32     }

```

User start-up code:

```

35     void user_start_op(user_lib_t *handle, void *data)
36     {
37         MPI_Irecv( ..., handle->comm, &(handle->irecv_handle) );
38         MPI_Isend( ..., handle->comm, &(handle->isend_handle) );
39     }

```

User communication clean-up code:

```

42     void user_end_op(user_lib_t *handle)
43     {
44         MPI_Status status;
45         MPI_Wait(&handle->isend_handle, &status);
46         MPI_Wait(&handle->irecv_handle, &status);
47     }

```

User object clean-up code:

```

void uninit_user_lib(user_lib_t *handle)
{
    MPI_Comm_free(&(handle->comm));
    free(handle);
}

```

6.5.6 Library Example #2

The main program:

```

int main(int argc, char *argv[])
{
    int ma, mb;
    MPI_Group group_world, group_a, group_b;
    MPI_Comm comm_a, comm_b;

    static int list_a[] = {0, 1};
    #if defined(EXAMPLE_2B) || defined(EXAMPLE_2C)
        static int list_b[] = {0, 2, 3};
    #else /* EXAMPLE_2A */
        static int list_b[] = {0, 2};
    #endif

    int size_list_a = sizeof(list_a)/sizeof(int);
    int size_list_b = sizeof(list_b)/sizeof(int);

    ...
    MPI_Init(&argc, &argv);
    MPI_Comm_group(MPI_COMM_WORLD, &group_world);

    MPI_Group_incl(group_world, size_list_a, list_a, &group_a);
    MPI_Group_incl(group_world, size_list_b, list_b, &group_b);

    MPI_Comm_create(MPI_COMM_WORLD, group_a, &comm_a);
    MPI_Comm_create(MPI_COMM_WORLD, group_b, &comm_b);

    if(comm_a != MPI_COMM_NULL)
        MPI_Comm_rank(comm_a, &ma);
    if(comm_b != MPI_COMM_NULL)
        MPI_Comm_rank(comm_b, &mb);

    if(comm_a != MPI_COMM_NULL)
        lib_call(comm_a);

    if(comm_b != MPI_COMM_NULL)
    {
        lib_call(comm_b);
        lib_call(comm_b);
    }
}

```

```

1      }
2
3      if(comm_a != MPI_COMM_NULL)
4          MPI_Comm_free(&comm_a);
5      if(comm_b != MPI_COMM_NULL)
6          MPI_Comm_free(&comm_b);
7      MPI_Group_free(&group_a);
8      MPI_Group_free(&group_b);
9      MPI_Group_free(&group_world);
10     MPI_Finalize();
11     return 0;
12 }
13
14 The library:
15
16 void lib_call(MPI_Comm comm)
17 {
18     int me, done = 0;
19     MPI_Status status;
20     MPI_Comm_rank(comm, &me);
21     if(me == 0)
22         while(!done)
23         {
24             MPI_Recv(..., MPI_ANY_SOURCE, MPI_ANY_TAG, comm, &status);
25             ...
26         }
27     else
28     {
29         /* work */
30         MPI_Send(..., 0, ARBITRARY_TAG, comm);
31         ...
32     }
33 #ifdef EXAMPLE_2C
34     /* include (resp, exclude) for safety (resp, no safety): */
35     MPI_Barrier(comm);
36 #endif
37 }

```

The above example is really three examples, depending on whether or not one includes rank 3 in `list_b`, and whether or not a synchronize is included in `lib_call`. This example illustrates that, despite contexts, subsequent calls to `lib_call` with the same context need not be safe from one another (colloquially, “back-masking”). Safety is realized if the `MPI_Barrier` is added. What this demonstrates is that libraries have to be written carefully, even with contexts. When rank 3 is excluded, then the synchronize is not needed to get safety from back-masking.

Algorithms like “reduce” and “allreduce” have strong enough source selectivity properties so that they are inherently okay (no back-masking), provided that MPI provides basic guarantees. So are multiple calls to a typical tree-broadcast algorithm with the same root or different roots (see [8]). Here we rely on two guarantees of MPI: pairwise ordering of

messages between processes in the same context, and source selectivity — deleting either feature removes the guarantee that back-masking cannot be required.

Algorithms that try to do non-deterministic broadcasts or other calls that include wildcard operations will not generally have the good properties of the deterministic implementations of “reduce,” “allreduce,” and “broadcast.” Such algorithms would have to utilize the monotonically increasing tags (within a communicator scope) to keep things straight.

All of the foregoing is a supposition of “collective calls” implemented with point-to-point operations. MPI implementations may or may not implement collective calls using point-to-point operations. These algorithms are used to illustrate the issues of correctness and safety, independent of how MPI implements its collective calls. See also Section 6.9.

6.6 Inter-Communication

This section introduces the concept of inter-communication and describes the portions of MPI that support it. It describes support for writing programs that contain user-level servers.

All communication described thus far has involved communication between processes that are members of the same group. This type of communication is called “**intra-communication**” and the communicator used is called an “**intra-communicator**,” as we have noted earlier in the chapter.

In modular and multi-disciplinary applications, different process groups execute distinct modules and processes within different modules communicate with one another in a pipeline or a more general module graph. In these applications, the most natural way for a process to specify a target process is by the rank of the target process within the target group. In applications that contain internal user-level servers, each server may be a process group that provides services to one or more clients, and each client may be a process group that uses the services of one or more servers. It is again most natural to specify the target process by rank within the target group in these applications. This type of communication is called “**inter-communication**” and the communicator used is called an “**inter-communicator**,” as introduced earlier.

An **inter-communication** is a point-to-point communication between processes in different groups. The group containing a process that initiates an inter-communication operation is called the “local group,” that is, the sender in a send and the receiver in a receive. The group containing the target process is called the “remote group,” that is, the receiver in a send and the sender in a receive. As in intra-communication, the target process is specified using a (communicator, rank) pair. Unlike intra-communication, the rank is relative to a second, remote group.

All inter-communicator constructors are blocking except for `MPI_COMM_IDUP` and require that the local and remote groups be disjoint.

Advice to users. The groups must be disjoint for several reasons. Primarily, this is the intent of the intercommunicators — to provide a communicator for communication between disjoint groups. This is reflected in the definition of `MPI_INTERCOMM_MERGE`, which allows the user to control the ranking of the processes in the created intracommunicator; this ranking makes little sense if the groups are not disjoint. In addition, the natural extension of collective operations to inter-communicators makes the most sense when the groups are disjoint. (*End of advice to users.*)

Here is a summary of the properties of inter-communication and inter-communicators:

- The syntax of point-to-point and collective communication is the same for both inter- and intra-communication. The same communicator can be used both for send and for receive operations.
- A target process is addressed by its rank in the remote group, both for sends and for receives.
- Communications using an inter-communicator are guaranteed not to conflict with any communications that use a different communicator.
- A communicator will provide either intra- or inter-communication, never both.

The routine `MPI_COMM_TEST_INTER` may be used to determine if a communicator is an inter- or intra-communicator. Inter-communicators can be used as arguments to some of the other communicator access routines. Inter-communicators cannot be used as input to some of the constructor routines for intra-communicators (for instance, `MPI_CART_CREATE`).

Advice to implementors. For the purpose of point-to-point communication, communicators can be represented in each process by a tuple consisting of:

```

group
send_context
receive_context
source

```

For inter-communicators, *group* describes the remote group, and *source* is the rank of the process in the local group. For intra-communicators, *group* is the communicator group (remote=local), *source* is the rank of the process in this group, and *send context* and *receive context* are identical. A group can be represented by a rank-to-absolute-address translation table.

The inter-communicator cannot be discussed sensibly without considering processes in both the local and remote groups. Imagine a process **P** in group \mathcal{P} , which has an inter-communicator $\mathbf{C}_{\mathcal{P}}$, and a process **Q** in group \mathcal{Q} , which has an inter-communicator $\mathbf{C}_{\mathcal{Q}}$. Then

- $\mathbf{C}_{\mathcal{P}}.\mathbf{group}$ describes the group \mathcal{Q} and $\mathbf{C}_{\mathcal{Q}}.\mathbf{group}$ describes the group \mathcal{P} .
- $\mathbf{C}_{\mathcal{P}}.\mathbf{send_context} = \mathbf{C}_{\mathcal{Q}}.\mathbf{receive_context}$ and the context is unique in \mathcal{Q} ;
 $\mathbf{C}_{\mathcal{P}}.\mathbf{receive_context} = \mathbf{C}_{\mathcal{Q}}.\mathbf{send_context}$ and this context is unique in \mathcal{P} .
- $\mathbf{C}_{\mathcal{P}}.\mathbf{source}$ is rank of **P** in \mathcal{P} and $\mathbf{C}_{\mathcal{Q}}.\mathbf{source}$ is rank of **Q** in \mathcal{Q} .

Assume that **P** sends a message to **Q** using the inter-communicator. Then **P** uses the **group** table to find the absolute address of **Q**; **source** and **send_context** are appended to the message.

Assume that **Q** posts a receive with an explicit source argument using the inter-communicator. Then **Q** matches **receive_context** to the message context and source argument to the message source.

The same algorithm is appropriate for intra-communicators as well.

In order to support inter-communicator accessors and constructors, it is necessary to supplement this model with additional structures, that store information about the local communication group, and additional safe contexts. (*End of advice to implementors.*)

6.6.1 Inter-communicator Accessors

`MPI_COMM_TEST_INTER(comm, flag)`

IN comm communicator (handle)
OUT flag (logical)

`int MPI_Comm_test_inter(MPI_Comm comm, int *flag)`

`MPI_Comm_test_inter(comm, flag, ierror)`
 TYPE(MPI_Comm), INTENT(IN) :: comm
 LOGICAL, INTENT(OUT) :: flag
 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

`MPI_COMM_TEST_INTER(COMM, FLAG, IERROR)`
 INTEGER COMM, IERROR
 LOGICAL FLAG

This local routine allows the calling process to determine if a communicator is an inter-communicator or an intra-communicator. It returns `true` if it is an inter-communicator, otherwise false.

When an inter-communicator is used as an input argument to the communicator accessors described above under intra-communication, the following table describes behavior.

<code>MPI_COMM_SIZE</code>	returns the size of the local group.
<code>MPI_COMM_GROUP</code>	returns the local group.
<code>MPI_COMM_RANK</code>	returns the rank in the local group

Table 6.1: `MPI_COMM_*` Function Behavior (in Inter-Communication Mode)

Furthermore, the operation `MPI_COMM_COMPARE` is valid for inter-communicators. Both communicators must be either intra- or inter-communicators, or else `MPI_UNEQUAL` results. Both corresponding local and remote groups must compare correctly to get the results `MPI_CONGRUENT` or `MPI_SIMILAR`. In particular, it is possible for `MPI_SIMILAR` to result because either the local or remote groups were similar but not identical.

The following accessors provide consistent access to the remote group of an inter-communicator. The following are all local operations.

```

1 MPI_COMM_REMOTE_SIZE(comm, size)
2     IN      comm      inter-communicator (handle)
3     OUT     size      number of processes in the remote group of comm
4                      (integer)
5
6
7 int MPI_Comm_remote_size(MPI_Comm comm, int *size)
8
9 MPI_Comm_remote_size(comm, size, ierror)
10    TYPE(MPI_Comm), INTENT(IN) :: comm
11    INTEGER, INTENT(OUT) :: size
12    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
13
14 MPI_COMM_REMOTE_SIZE(COMM, SIZE, IERROR)
15    INTEGER COMM, SIZE, IERROR
16
17 MPI_COMM_REMOTE_GROUP(comm, group)
18     IN      comm      inter-communicator (handle)
19     OUT     group     remote group corresponding to comm (handle)
20
21
22 int MPI_Comm_remote_group(MPI_Comm comm, MPI_Group *group)
23
24 MPI_Comm_remote_group(comm, group, ierror)
25    TYPE(MPI_Comm), INTENT(IN) :: comm
26    TYPE(MPI_Group), INTENT(OUT) :: group
27    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
28
29 MPI_COMM_REMOTE_GROUP(COMM, GROUP, IERROR)
30    INTEGER COMM, GROUP, IERROR

```

Rationale. Symmetric access to both the local and remote groups of an inter-communicator is important, so this function, as well as MPI_COMM_REMOTE_SIZE have been provided. (*End of rationale.*)

6.6.2 Inter-communicator Operations

This section introduces four blocking inter-communicator operations.

MPI_INTERCOMM_CREATE is used to bind two intra-communicators into an inter-communicator; the function MPI_INTERCOMM_MERGE creates an intra-communicator by merging the local and remote groups of an inter-communicator. The functions MPI_COMM_DUP and MPI_COMM_FREE, introduced previously, duplicate and free an inter-communicator, respectively.

Overlap of local and remote groups that are bound into an inter-communicator is prohibited. If there is overlap, then the program is erroneous and is likely to deadlock. (If a process is multithreaded, and MPI calls block only a thread, rather than a process, then “dual membership” can be supported. It is then the user’s responsibility to make sure that calls on behalf of the two “roles” of a process are executed by two independent threads.)

The function `MPI_INTERCOMM_CREATE` can be used to create an inter-communicator from two existing intra-communicators, in the following situation: At least one selected member from each group (the “group leader”) has the ability to communicate with the selected member from the other group; that is, a “peer” communicator exists to which both leaders belong, and each leader knows the rank of the other leader in this peer communicator. Furthermore, members of each group know the rank of their leader.

Construction of an inter-communicator from two intra-communicators requires separate collective operations in the local group and in the remote group, as well as a point-to-point communication between a process in the local group and a process in the remote group.

In standard MPI implementations (with static process allocation at initialization), the `MPI_COMM_WORLD` communicator (or preferably a dedicated duplicate thereof) can be this peer communicator. For applications that have used `spawn` or `join`, it may be necessary to first create an intracommunicator to be used as peer.

The application topology functions described in Chapter 7 do not apply to inter-communicators. Users that require this capability should utilize `MPI_INTERCOMM_MERGE` to build an intra-communicator, then apply the graph or cartesian topology capabilities to that intra-communicator, creating an appropriate topology-oriented intra-communicator. Alternatively, it may be reasonable to devise one’s own application topology mechanisms for this case, without loss of generality.

`MPI_INTERCOMM_CREATE(local_comm, local_leader, peer_comm, remote_leader, tag, newintercomm)`

IN	<code>local_comm</code>	local intra-communicator (handle)
IN	<code>local_leader</code>	rank of local group leader in <code>local_comm</code> (integer)
IN	<code>peer_comm</code>	“peer” communicator; significant only at the <code>local_leader</code> (handle)
IN	<code>remote_leader</code>	rank of remote group leader in <code>peer_comm</code> ; significant only at the <code>local_leader</code> (integer)
IN	<code>tag</code>	tag (integer)
OUT	<code>newintercomm</code>	new inter-communicator (handle)

```
int MPI_Intercomm_create(MPI_Comm local_comm, int local_leader,
                        MPI_Comm peer_comm, int remote_leader, int tag,
                        MPI_Comm *newintercomm)
```

```
MPI_Intercomm_create(local_comm, local_leader, peer_comm, remote_leader,
                    tag, newintercomm, ierror)
```

```
TYPE(MPI_Comm), INTENT(IN) :: local_comm, peer_comm
INTEGER, INTENT(IN) :: local_leader, remote_leader, tag
TYPE(MPI_Comm), INTENT(OUT) :: newintercomm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_INTERCOMM_CREATE(LOCAL_COMM, LOCAL_LEADER, PEER_COMM, REMOTE_LEADER,
                    TAG, NEWINTERCOMM, IERROR)
```

```
INTEGER LOCAL_COMM, LOCAL_LEADER, PEER_COMM, REMOTE_LEADER, TAG,
NEWINTERCOMM, IERROR
```

This call creates an inter-communicator. It is collective over the union of the local and remote groups. Processes should provide identical `local_comm` and `local_leader` arguments within each group. Wildcards are not permitted for `remote_leader`, `local_leader`, and `tag`.

`MPI_INTERCOMM_MERGE(intercomm, high, newintracomm)`

IN	<code>intercomm</code>	Inter-Communicator (handle)
IN	<code>high</code>	(logical)
OUT	<code>newintracomm</code>	new intra-communicator (handle)

```
int MPI_Intercomm_merge(MPI_Comm intercomm, int high,
                        MPI_Comm *newintracomm)
```

```
MPI_Intercomm_merge(intercomm, high, newintracomm, ierror)
    TYPE(MPI_Comm), INTENT(IN) :: intercomm
    LOGICAL, INTENT(IN) :: high
    TYPE(MPI_Comm), INTENT(OUT) :: newintracomm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_INTERCOMM_MERGE(INTERCOMM, HIGH, NEWINTRACOMM, IERROR)
    INTEGER INTERCOMM, NEWINTRACOMM, IERROR
    LOGICAL HIGH
```

This function creates an intra-communicator from the union of the two groups that are associated with `intercomm`. All processes should provide the same `high` value within each of the two groups. If processes in one group provided the value `high = false` and processes in the other group provided the value `high = true` then the union orders the “low” group before the “high” group. If all processes provided the same `high` argument then the order of the union is arbitrary. This call is blocking and collective within the union of the two groups.

The error handler on the new intercommunicator in each process is inherited from the communicator that contributes the local group. Note that this can result in different processes in the same communicator having different error handlers.

Advice to implementors. The implementation of `MPI_INTERCOMM_MERGE`, `MPI_COMM_FREE`, and `MPI_COMM_DUP` are similar to the implementation of `MPI_INTERCOMM_CREATE`, except that contexts private to the input inter-communicator are used for communication between group leaders rather than contexts inside a bridge communicator. (*End of advice to implementors.*)

6.6.3 Inter-Communication Examples

Example 1: Three-Group “Pipeline”

Groups 0 and 1 communicate. Groups 1 and 2 communicate. Therefore, group 0 requires one inter-communicator, group 1 requires two inter-communicators, and group 2 requires 1 inter-communicator.

```
int main(int argc, char *argv[])
```

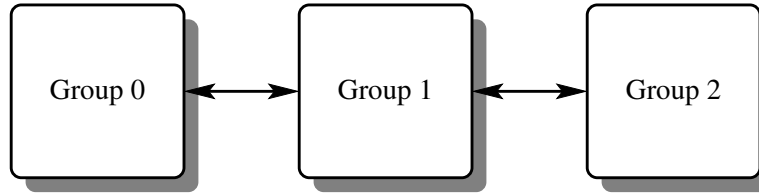


Figure 6.4: Three-group pipeline

```

{
  MPI_Comm  myComm;          /* intra-communicator of local sub-group */
  MPI_Comm  myFirstComm;    /* inter-communicator */
  MPI_Comm  mySecondComm;   /* second inter-communicator (group 1 only) */
  int membershipKey;
  int rank;

  MPI_Init(&argc, &argv);
  MPI_Comm_rank(MPI_COMM_WORLD, &rank);

  /* User code must generate membershipKey in the range [0, 1, 2] */
  membershipKey = rank % 3;

  /* Build intra-communicator for local sub-group */
  MPI_Comm_split(MPI_COMM_WORLD, membershipKey, rank, &myComm);

  /* Build inter-communicators. Tags are hard-coded. */
  if (membershipKey == 0)
  {
      /* Group 0 communicates with group 1. */
      MPI_Intercomm_create(myComm, 0, MPI_COMM_WORLD, 1,
                          1, &myFirstComm);
  }
  else if (membershipKey == 1)
  {
      /* Group 1 communicates with groups 0 and 2. */
      MPI_Intercomm_create(myComm, 0, MPI_COMM_WORLD, 0,
                          1, &myFirstComm);
      MPI_Intercomm_create(myComm, 0, MPI_COMM_WORLD, 2,
                          12, &mySecondComm);
  }
  else if (membershipKey == 2)
  {
      /* Group 2 communicates with group 1. */
      MPI_Intercomm_create(myComm, 0, MPI_COMM_WORLD, 1,
                          12, &myFirstComm);
  }

  /* Do work ... */

  switch(membershipKey) /* free communicators appropriately */
  {

```

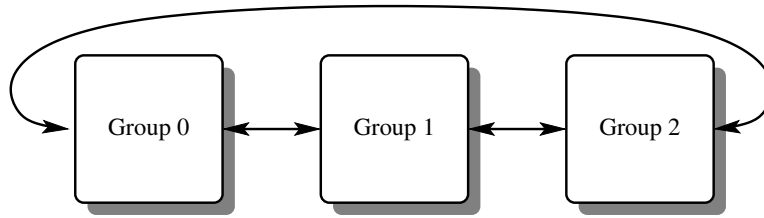


Figure 6.5: Three-group ring

```

10 case 1:
11     MPI_Comm_free(&mySecondComm);
12 case 0:
13 case 2:
14     MPI_Comm_free(&myFirstComm);
15     break;
16 }
17
18 MPI_Finalize();
19 return 0;
20 }
  
```

Example 2: Three-Group “Ring”

Groups 0 and 1 communicate. Groups 1 and 2 communicate. Groups 0 and 2 communicate. Therefore, each requires two inter-communicators.

```

26 int main(int argc, char *argv[])
27 {
28     MPI_Comm myComm; /* intra-communicator of local sub-group */
29     MPI_Comm myFirstComm; /* inter-communicators */
30     MPI_Comm mySecondComm;
31     int membershipKey;
32     int rank;
33
34     MPI_Init(&argc, &argv);
35     MPI_Comm_rank(MPI_COMM_WORLD, &rank);
36     ...
37
38     /* User code must generate membershipKey in the range [0, 1, 2] */
39     membershipKey = rank % 3;
40
41     /* Build intra-communicator for local sub-group */
42     MPI_Comm_split(MPI_COMM_WORLD, membershipKey, rank, &myComm);
43
44     /* Build inter-communicators. Tags are hard-coded. */
45     if (membershipKey == 0)
46     {
47         /* Group 0 communicates with groups 1 and 2. */
48         MPI_Intercomm_create(myComm, 0, MPI_COMM_WORLD, 1,
  
```



```

        1, &myFirstComm);
MPI_Intercomm_create(myComm, 0, MPI_COMM_WORLD, 2,
        2, &mySecondComm);
}
else if (membershipKey == 1)
{
    /* Group 1 communicates with groups 0 and 2. */
    MPI_Intercomm_create(myComm, 0, MPI_COMM_WORLD, 0,
        1, &myFirstComm);
    MPI_Intercomm_create(myComm, 0, MPI_COMM_WORLD, 2,
        12, &mySecondComm);
}
else if (membershipKey == 2)
{
    /* Group 2 communicates with groups 0 and 1. */
    MPI_Intercomm_create(myComm, 0, MPI_COMM_WORLD, 0,
        2, &myFirstComm);
    MPI_Intercomm_create(myComm, 0, MPI_COMM_WORLD, 1,
        12, &mySecondComm);
}

/* Do some work ... */

/* Then free communicators before terminating... */
MPI_Comm_free(&myFirstComm);
MPI_Comm_free(&mySecondComm);
MPI_Comm_free(&myComm);
MPI_Finalize();
return 0;
}

```

6.7 Caching

MPI provides a “caching” facility that allows an application to attach arbitrary pieces of information, called **attributes**, to three kinds of MPI objects, communicators, windows, and datatypes. More precisely, the caching facility allows a portable library to do the following:

- pass information between calls by associating it with an MPI intra- or inter-communicator, window, or datatype,
- quickly retrieve that information, and
- be guaranteed that out-of-date information is never retrieved, even if the object is freed and its handle subsequently reused by MPI.

The caching capabilities, in some form, are required by built-in MPI routines such as collective communication and application topology. Defining an interface to these capabilities as part of the MPI standard is valuable because it permits routines like collective communication and application topologies to be implemented as portable code, and also because it makes MPI more extensible by allowing user-written routines to use standard MPI calling sequences.

Advice to users. The communicator `MPI_COMM_SELF` is a suitable choice for posting process-local attributes, via this attribute-caching mechanism. (*End of advice to users.*)

Rationale. In one extreme one can allow caching on all opaque handles. The other extreme is to only allow it on communicators. Caching has a cost associated with it and should only be allowed when it is clearly needed and the increased cost is modest. This is the reason that windows and datatypes were added but not other handles. (*End of rationale.*)

One difficulty is the potential for size differences between Fortran integers and C pointers. For this reason, the Fortran versions of these routines use integers of kind `MPI_ADDRESS_KIND`.

Advice to implementors. High-quality implementations should raise an error when a keyval that was created by a call to `MPI_XXX_CREATE_KEYVAL` is used with an object of the wrong type with a call to `MPI_YYY_GET_ATTR`, `MPI_YYY_SET_ATTR`, `MPI_YYY_DELETE_ATTR`, or `MPI_YYY_FREE_KEYVAL`. To do so, it is necessary to maintain, with each keyval, information on the type of the associated user function. (*End of advice to implementors.*)

6.7.1 Functionality

Attributes can be attached to communicators, windows, and datatypes. Attributes are local to the process and specific to the communicator to which they are attached. Attributes are not propagated by MPI from one communicator to another except when the communicator is duplicated using `MPI_COMM_DUP` or `MPI_COMM_IDUP` (and even then the application must give specific permission through callback functions for the attribute to be copied).

Advice to users. Attributes in C are of type `void *`. Typically, such an attribute will be a pointer to a structure that contains further information, or a handle to an MPI object. In Fortran, attributes are of type `INTEGER`. Such attribute can be a handle to an MPI object, or just an integer-valued attribute. (*End of advice to users.*)

Advice to implementors. Attributes are scalar values, equal in size to, or larger than a C-language pointer. Attributes can always hold an MPI handle. (*End of advice to implementors.*)

The caching interface defined here requires that attributes be stored by MPI opaquely within a communicator, window, and datatype. Accessor functions include the following:

- obtain a key value (used to identify an attribute); the user specifies “callback” functions by which MPI informs the application when the communicator is destroyed or copied.
- store and retrieve the value of an attribute;

Advice to implementors. Caching and callback functions are only called synchronously, in response to explicit application requests. This avoids problems that result from repeated crossings between user and system space. (This synchronous calling rule is a general property of MPI.)

The choice of key values is under control of MPI. This allows MPI to optimize its implementation of attribute sets. It also avoids conflict between independent modules caching information on the same communicators.

A much smaller interface, consisting of just a callback facility, would allow the entire caching facility to be implemented by portable code. However, with the minimal callback interface, some form of table searching is implied by the need to handle arbitrary communicators. In contrast, the more complete interface defined here permits rapid access to attributes through the use of pointers in communicators (to find the attribute table) and cleverly chosen key values (to retrieve individual attributes). In light of the efficiency “hit” inherent in the minimal interface, the more complete interface defined here is seen to be superior. (*End of advice to implementors.*)

MPI provides the following services related to caching. They are all process local.

6.7.2 Communicators

Functions for caching on communicators are:

`MPI_COMM_CREATE_KEYVAL(comm_copy_attr_fn, comm_delete_attr_fn, comm_keyval, extra_state)`

IN	<code>comm_copy_attr_fn</code>	copy callback function for <code>comm_keyval</code> (function)
IN	<code>comm_delete_attr_fn</code>	delete callback function for <code>comm_keyval</code> (function)
OUT	<code>comm_keyval</code>	key value for future access (integer)
IN	<code>extra_state</code>	extra state for callback functions

```
int MPI_Comm_create_keyval(MPI_Comm_copy_attr_function *comm_copy_attr_fn,
                           MPI_Comm_delete_attr_function *comm_delete_attr_fn,
                           int *comm_keyval, void *extra_state)
```

```
MPI_Comm_create_keyval(comm_copy_attr_fn, comm_delete_attr_fn, comm_keyval,
                       extra_state, ierror)
```

```
PROCEDURE(MPI_Comm_copy_attr_function) :: comm_copy_attr_fn
PROCEDURE(MPI_Comm_delete_attr_function) :: comm_delete_attr_fn
INTEGER, INTENT(OUT) :: comm_keyval
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_COMM_CREATE_KEYVAL(COMM_COPY_ATTR_FN, COMM_DELETE_ATTR_FN, COMM_KEYVAL,
                       EXTRA_STATE, IERROR)
```

```
EXTERNAL COMM_COPY_ATTR_FN, COMM_DELETE_ATTR_FN
INTEGER COMM_KEYVAL, IERROR
INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
```

Generates a new attribute key. Keys are locally unique in a process, and opaque to user, though they are explicitly stored in integers. Once allocated, the key value can be used to associate attributes and access them on any locally defined communicator.

The C callback functions are:

```

1  typedef int MPI_Comm_copy_attr_function(MPI_Comm oldcomm, int comm_keyval,
2      void *extra_state, void *attribute_val_in,
3      void *attribute_val_out, int *flag);

```

4 and

```

5  typedef int MPI_Comm_delete_attr_function(MPI_Comm comm, int comm_keyval,
6      void *attribute_val, void *extra_state);
7

```

8 which are the same as the MPI-1.1 calls but with a new name. The old names are deprecated.
9 With the `mpi_f08` module, the Fortran callback functions are:

10 ABSTRACT INTERFACE

```

11     SUBROUTINE MPI_Comm_copy_attr_function(oldcomm, comm_keyval, extra_state,
12         attribute_val_in, attribute_val_out, flag, ierror)
13         TYPE(MPI_Comm) :: oldcomm
14         INTEGER :: comm_keyval, ierror
15         INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state, attribute_val_in,
16             attribute_val_out
17         LOGICAL :: flag

```

18 and

19 ABSTRACT INTERFACE

```

20     SUBROUTINE MPI_Comm_delete_attr_function(comm, comm_keyval,
21         attribute_val, extra_state, ierror)
22         TYPE(MPI_Comm) :: comm
23         INTEGER :: comm_keyval, ierror
24         INTEGER(KIND=MPI_ADDRESS_KIND) :: attribute_val, extra_state
25

```

26 With the `mpi` module and `mpif.h`, the Fortran callback functions are:

```

27     SUBROUTINE COMM_COPY_ATTR_FUNCTION(OLDCOMM, COMM_KEYVAL, EXTRA_STATE,
28         ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT, FLAG, IERROR)
29         INTEGER OLDCOMM, COMM_KEYVAL, IERROR
30         INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
31             ATTRIBUTE_VAL_OUT
32         LOGICAL FLAG

```

33 and

```

34     SUBROUTINE COMM_DELETE_ATTR_FUNCTION(COMM, COMM_KEYVAL, ATTRIBUTE_VAL,
35         EXTRA_STATE, IERROR)
36         INTEGER COMM, COMM_KEYVAL, IERROR
37         INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE
38

```

39 The `comm_copy_attr_fn` function is invoked when a communicator is duplicated by
40 `MPI_COMM_DUP` or `MPI_COMM_IDUP`. `comm_copy_attr_fn` should be of type
41 `MPI_Comm_copy_attr_function`. The copy callback function is invoked for each key value in
42 `oldcomm` in arbitrary order. Each call to the copy callback is made with a key value and its
43 corresponding attribute. If it returns `flag = 0` or `.FALSE.`, then the attribute is deleted in the
44 duplicated communicator. Otherwise (`flag = 1` or `.TRUE.`), the new attribute value is set to
45 the value returned in `attribute_val_out`. The function returns `MPI_SUCCESS` on success and
46 an error code on failure (in which case `MPI_COMM_DUP` or `MPI_COMM_IDUP` will fail).

47 The argument `comm_copy_attr_fn` may be specified as `MPI_COMM_NULL_COPY_FN`

or `MPI_COMM_DUP_FN` from either C or Fortran. `MPI_COMM_NULL_COPY_FN` is a function that does nothing other than returning `flag = 0` or `.FALSE.` (depending on whether the keyval was created with a C or Fortran binding to `MPI_COMM_CREATE_KEYVAL`) and `MPI_SUCCESS`. `MPI_COMM_DUP_FN` is a simple-minded copy function that sets `flag = 1` or `.TRUE.`, returns the value of `attribute_val_in` in `attribute_val_out`, and returns `MPI_SUCCESS`. These replace the MPI-1 predefined callbacks `MPI_NULL_COPY_FN` and `MPI_DUP_FN`, whose use is deprecated.

Advice to users. Even though both formal arguments `attribute_val_in` and `attribute_val_out` are of type `void *`, their usage differs. The C copy function is passed by MPI in `attribute_val_in` the *value* of the attribute, and in `attribute_val_out` the *address* of the attribute, so as to allow the function to return the (new) attribute value. The use of type `void *` for both is to avoid messy type casts.

A valid copy function is one that completely duplicates the information by making a full duplicate copy of the data structures implied by an attribute; another might just make another reference to that data structure, while using a reference-count mechanism. Other types of attributes might not copy at all (they might be specific to oldcomm only). (*End of advice to users.*)

Advice to implementors. A C interface should be assumed for copy and delete functions associated with key values created in C; a Fortran calling interface should be assumed for key values created in Fortran. (*End of advice to implementors.*)

Analogous to `comm_copy_attr_fn` is a callback deletion function, defined as follows. The `comm_delete_attr_fn` function is invoked when a communicator is deleted by `MPI_COMM_FREE` or when a call is made explicitly to `MPI_COMM_DELETE_ATTR`. `comm_delete_attr_fn` should be of type `MPI_Comm_delete_attr_function`.

This function is called by `MPI_COMM_FREE`, `MPI_COMM_DELETE_ATTR`, and `MPI_COMM_SET_ATTR` to do whatever is needed to remove an attribute. The function returns `MPI_SUCCESS` on success and an error code on failure (in which case `MPI_COMM_FREE` will fail).

The argument `comm_delete_attr_fn` may be specified as `MPI_COMM_NULL_DELETE_FN` from either C or Fortran. `MPI_COMM_NULL_DELETE_FN` is a function that does nothing, other than returning `MPI_SUCCESS`. `MPI_COMM_NULL_DELETE_FN` replaces `MPI_NULL_DELETE_FN`, whose use is deprecated.

If an attribute copy function or attribute delete function returns other than `MPI_SUCCESS`, then the call that caused it to be invoked (for example, `MPI_COMM_FREE`), is erroneous.

The special key value `MPI_KEYVAL_INVALID` is never returned by `MPI_COMM_CREATE_KEYVAL`. Therefore, it can be used for static initialization of key values.

Advice to implementors. The predefined Fortran functions `MPI_COMM_NULL_COPY_FN`, `MPI_COMM_DUP_FN`, and `MPI_COMM_NULL_DELETE_FN` are defined in the `mpi` module (and `mpif.h`) and the `mpi_f08` module with the same name, but with different interfaces. Each function can coexist twice with the same name in the same MPI library, one routine as an

implicit interface outside of the `mpi` module, i.e., declared as `EXTERNAL`, and the other routine within `mpi_f08` declared with `CONTAINS`. These routines have different link names, which are also different to the link names used for the routines used in C. (*End of advice to implementors.*)

Advice to users. Callbacks, including the predefined Fortran functions `MPI_COMM_NULL_COPY_FN`, `MPI_COMM_DUP_FN`, and `MPI_COMM_NULL_DELETE_FN` should not be passed from one application routine that uses the `mpi_f08` module to another application routine that uses the `mpi` module or `mpif.h`, and vice versa; see also the advice to users on page ?? (*End of advice to users.*)

`MPI_COMM_FREE_KEYVAL(comm_keyval)`

INOUT `comm_keyval` key value (integer)

`int MPI_Comm_free_keyval(int *comm_keyval)`

`MPI_Comm_free_keyval(comm_keyval, ierror)`

INTEGER, INTENT(INOUT) :: `comm_keyval`

INTEGER, OPTIONAL, INTENT(OUT) :: `ierror`

`MPI_COMM_FREE_KEYVAL(COMM_KEYVAL, IERROR)`

INTEGER `COMM_KEYVAL`, `IERROR`

Frees an extant attribute key. This function sets the value of `keyval` to `MPI_KEYVAL_INVALID`. Note that it is not erroneous to free an attribute key that is in use, because the actual free does not transpire until after all references (in other communicators on the process) to the key have been freed. These references need to be explicitly freed by the program, either via calls to `MPI_COMM_DELETE_ATTR` that free one attribute instance, or by calls to `MPI_COMM_FREE` that free all attribute instances associated with the freed communicator.

`MPI_COMM_SET_ATTR(comm, comm_keyval, attribute_val)`

INOUT `comm` communicator from which attribute will be attached
(handle)

IN `comm_keyval` key value (integer)

IN `attribute_val` attribute value

`int MPI_Comm_set_attr(MPI_Comm comm, int comm_keyval, void *attribute_val)`

`MPI_Comm_set_attr(comm, comm_keyval, attribute_val, ierror)`

TYPE(MPI_Comm), INTENT(IN) :: `comm`

INTEGER, INTENT(IN) :: `comm_keyval`

INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: `attribute_val`

INTEGER, OPTIONAL, INTENT(OUT) :: `ierror`

`MPI_COMM_SET_ATTR(COMM, COMM_KEYVAL, ATTRIBUTE_VAL, IERROR)`

```

INTEGER COMM, COMM_KEYVAL, IERROR
INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL

```

This function stores the stipulated attribute value `attribute_val` for subsequent retrieval by `MPI_COMM_GET_ATTR`. If the value is already present, then the outcome is as if `MPI_COMM_DELETE_ATTR` was first called to delete the previous value (and the callback function `comm_delete_attr_fn` was executed), and a new value was next stored. The call is erroneous if there is no key with value `keyval`; in particular `MPI_KEYVAL_INVALID` is an erroneous key value. The call will fail if the `comm_delete_attr_fn` function returned an error code other than `MPI_SUCCESS`.

```

MPI_COMM_GET_ATTR(comm, comm_keyval, attribute_val, flag)

```

IN	comm	communicator to which the attribute is attached (handle)
IN	comm_keyval	key value (integer)
OUT	attribute_val	attribute value, unless <code>flag = false</code>
OUT	flag	false if no attribute is associated with the key (logical)

```

int MPI_Comm_get_attr(MPI_Comm comm, int comm_keyval, void *attribute_val,
                      int *flag)

```

```

MPI_Comm_get_attr(comm, comm_keyval, attribute_val, flag, ierror)
  TYPE(MPI_Comm), INTENT(IN) :: comm
  INTEGER, INTENT(IN) :: comm_keyval
  INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val
  LOGICAL, INTENT(OUT) :: flag
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_COMM_GET_ATTR(COMM, COMM_KEYVAL, ATTRIBUTE_VAL, FLAG, IERROR)
  INTEGER COMM, COMM_KEYVAL, IERROR
  INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
  LOGICAL FLAG

```

Retrieves attribute value by key. The call is erroneous if there is no key with value `keyval`. On the other hand, the call is correct if the key value exists, but no attribute is attached on `comm` for that key; in such case, the call returns `flag = false`. In particular `MPI_KEYVAL_INVALID` is an erroneous key value.

Advice to users. The call to `MPI_Comm_set_attr` passes in `attribute_val` the *value* of the attribute; the call to `MPI_Comm_get_attr` passes in `attribute_val` the *address* of the location where the attribute value is to be returned. Thus, if the attribute value itself is a pointer of type `void*`, then the actual `attribute_val` parameter to `MPI_Comm_set_attr` will be of type `void*` and the actual `attribute_val` parameter to `MPI_Comm_get_attr` will be of type `void**`. (*End of advice to users.*)

Rationale. The use of a formal parameter `attribute_val` of type `void*` (rather than `void**`) avoids the messy type casting that would be needed if the attribute value is declared with a type other than `void*`. (*End of rationale.*)

```
1 MPI_COMM_DELETE_ATTR(comm, comm_keyval)
```

```
2     INOUT    comm                communicator from which the attribute is deleted (han-
3                                     dle)
```

```
4     IN       comm_keyval        key value (integer)
```

```
6
7 int MPI_Comm_delete_attr(MPI_Comm comm, int comm_keyval)
```

```
8 MPI_Comm_delete_attr(comm, comm_keyval, ierror)
```

```
9     TYPE(MPI_Comm), INTENT(IN) :: comm
```

```
10    INTEGER, INTENT(IN) :: comm_keyval
```

```
11    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
12
13 MPI_COMM_DELETE_ATTR(COMM, COMM_KEYVAL, IERROR)
```

```
14    INTEGER COMM, COMM_KEYVAL, IERROR
```

```
15
16     Delete attribute from cache by key. This function invokes the attribute delete function
17 comm_delete_attr_fn specified when the keyval was created. The call will fail if the
18 comm_delete_attr_fn function returns an error code other than MPI_SUCCESS.
```

```
19     Whenever a communicator is replicated using the function MPI_COMM_DUP or
20 MPI_COMM_IDUP, all call-back copy functions for attributes that are currently set are
21 invoked (in arbitrary order). Whenever a communicator is deleted using the function
22 MPI_COMM_FREE all callback delete functions for attributes that are currently set are
23 invoked.
```

24 6.7.3 Windows

```
25
26 The functions for caching on windows are:
```

```
27
28 MPI_WIN_CREATE_KEYVAL(win_copy_attr_fn, win_delete_attr_fn, win_keyval, extra_state)
```

```
29
30
31     IN       win_copy_attr_fn    copy callback function for win_keyval (function)
```

```
32     IN       win_delete_attr_fn  delete callback function for win_keyval (function)
```

```
33     OUT      win_keyval          key value for future access (integer)
```

```
34     IN       extra_state         extra state for callback functions
```

```
35
36
37 int MPI_Win_create_keyval(MPI_Win_copy_attr_function *win_copy_attr_fn,
38     MPI_Win_delete_attr_function *win_delete_attr_fn,
39     int *win_keyval, void *extra_state)
```

```
40
41 MPI_Win_create_keyval(win_copy_attr_fn, win_delete_attr_fn, win_keyval,
42     extra_state, ierror)
```

```
43     PROCEDURE(MPI_Win_copy_attr_function) :: win_copy_attr_fn
```

```
44     PROCEDURE(MPI_Win_delete_attr_function) :: win_delete_attr_fn
```

```
45     INTEGER, INTENT(OUT) :: win_keyval
```

```
46     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
```

```
47     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```



```

MPI_WIN_CREATE_KEYVAL(WIN_COPY_ATTR_FN, WIN_DELETE_ATTR_FN, WIN_KEYVAL,
                      EXTRA_STATE, IERROR)
EXTERNAL WIN_COPY_ATTR_FN, WIN_DELETE_ATTR_FN
INTEGER WIN_KEYVAL, IERROR
INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE

```

The argument `win_copy_attr_fn` may be specified as `MPI_WIN_NULL_COPY_FN` or `MPI_WIN_DUP_FN` from either C or Fortran. `MPI_WIN_NULL_COPY_FN` is a function that does nothing other than returning `flag = 0` and `MPI_SUCCESS`. `MPI_WIN_DUP_FN` is a simple-minded copy function that sets `flag = 1`, returns the value of `attribute_val_in` in `attribute_val_out`, and returns `MPI_SUCCESS`.

The argument `win_delete_attr_fn` may be specified as `MPI_WIN_NULL_DELETE_FN` from either C or Fortran. `MPI_WIN_NULL_DELETE_FN` is a function that does nothing, other than returning `MPI_SUCCESS`.

The C callback functions are:

```

typedef int MPI_Win_copy_attr_function(MPI_Win oldwin, int win_keyval,
                                       void *extra_state, void *attribute_val_in,
                                       void *attribute_val_out, int *flag);

```

and

```

typedef int MPI_Win_delete_attr_function(MPI_Win win, int win_keyval,
                                       void *attribute_val, void *extra_state);

```

With the `mpi_f08` module, the Fortran callback functions are:

```

ABSTRACT INTERFACE
  SUBROUTINE MPI_Win_copy_attr_function(oldwin, win_keyval, extra_state,
    attribute_val_in, attribute_val_out, flag, ierror)
    TYPE(MPI_Win) :: oldwin
    INTEGER :: win_keyval, ierror
    INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state, attribute_val_in,
    attribute_val_out
    LOGICAL :: flag

```

and

```

ABSTRACT INTERFACE
  SUBROUTINE MPI_Win_delete_attr_function(win, win_keyval, attribute_val,
    extra_state, ierror)
    TYPE(MPI_Win) :: win
    INTEGER :: win_keyval, ierror
    INTEGER(KIND=MPI_ADDRESS_KIND) :: attribute_val, extra_state

```

With the `mpi` module and `mpif.h`, the Fortran callback functions are:

```

SUBROUTINE WIN_COPY_ATTR_FUNCTION(OLDWIN, WIN_KEYVAL, EXTRA_STATE,
    ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT, FLAG, IERROR)
  INTEGER OLDWIN, WIN_KEYVAL, IERROR
  INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
    ATTRIBUTE_VAL_OUT
  LOGICAL FLAG

```

and

```

1  SUBROUTINE WIN_DELETE_ATTR_FUNCTION(WIN, WIN_KEYVAL, ATTRIBUTE_VAL,
2      EXTRA_STATE, IERROR)
3      INTEGER WIN, WIN_KEYVAL, IERROR
4      INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE

```

If an attribute copy function or attribute delete function returns other than MPI_SUCCESS, then the call that caused it to be invoked (for example, MPI_WIN_FREE), is erroneous.

```

10 MPI_WIN_FREE_KEYVAL(win_keyval)

```

```

11     INOUT    win_keyval                key value (integer)

```

```

13 int MPI_Win_free_keyval(int *win_keyval)

```

```

15 MPI_Win_free_keyval(win_keyval, ierror)
16     INTEGER, INTENT(INOUT) :: win_keyval
17     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

19 MPI_WIN_FREE_KEYVAL(WIN_KEYVAL, IERROR)
20     INTEGER WIN_KEYVAL, IERROR

```

```

23 MPI_WIN_SET_ATTR(win, win_keyval, attribute_val)

```

```

24     INOUT    win                window to which attribute will be attached (handle)
25     IN       win_keyval        key value (integer)
26     IN       attribute_val     attribute value

```

```

29 int MPI_Win_set_attr(MPI_Win win, int win_keyval, void *attribute_val)

```

```

31 MPI_Win_set_attr(win, win_keyval, attribute_val, ierror)
32     TYPE(MPI_Win), INTENT(IN) :: win
33     INTEGER, INTENT(IN) :: win_keyval
34     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: attribute_val
35     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

36 MPI_WIN_SET_ATTR(WIN, WIN_KEYVAL, ATTRIBUTE_VAL, IERROR)
37     INTEGER WIN, WIN_KEYVAL, IERROR
38     INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL

```

```

41 MPI_WIN_GET_ATTR(win, win_keyval, attribute_val, flag)

```

```

43     IN       win                window to which the attribute is attached (handle)
44     IN       win_keyval        key value (integer)
45     OUT      attribute_val     attribute value, unless flag = false
46     OUT      flag              false if no attribute is associated with the key (logical)

```

```

int MPI_Win_get_attr(MPI_Win win, int win_keyval, void *attribute_val,
                    int *flag)
MPI_Win_get_attr(win, win_keyval, attribute_val, flag, ierror)
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, INTENT(IN) :: win_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val
    LOGICAL, INTENT(OUT) :: flag
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_WIN_GET_ATTR(WIN, WIN_KEYVAL, ATTRIBUTE_VAL, FLAG, IERROR)
    INTEGER WIN, WIN_KEYVAL, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
    LOGICAL FLAG

MPI_WIN_DELETE_ATTR(win, win_keyval)
    INOUT    win                window from which the attribute is deleted (handle)
    IN       win_keyval         key value (integer)

int MPI_Win_delete_attr(MPI_Win win, int win_keyval)
MPI_Win_delete_attr(win, win_keyval, ierror)
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, INTENT(IN) :: win_keyval
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_WIN_DELETE_ATTR(WIN, WIN_KEYVAL, IERROR)
    INTEGER WIN, WIN_KEYVAL, IERROR

```

6.7.4 Datatypes

The new functions for caching on datatypes are:

```

MPI_TYPE_CREATE_KEYVAL(type_copy_attr_fn, type_delete_attr_fn, type_keyval,
                      extra_state)
    IN      type_copy_attr_fn    copy callback function for type_keyval (function)
    IN      type_delete_attr_fn  delete callback function for type_keyval (function)
    OUT     type_keyval          key value for future access (integer)
    IN      extra_state          extra state for callback functions

int MPI_Type_create_keyval(MPI_Type_copy_attr_function *type_copy_attr_fn,
                          MPI_Type_delete_attr_function *type_delete_attr_fn,
                          int *type_keyval, void *extra_state)
MPI_Type_create_keyval(type_copy_attr_fn, type_delete_attr_fn, type_keyval,
                      extra_state, ierror)

```

```

1  PROCEDURE(MPI_Type_copy_attr_function) :: type_copy_attr_fn
2  PROCEDURE(MPI_Type_delete_attr_function) :: type_delete_attr_fn
3  INTEGER, INTENT(OUT) :: type_keyval
4  INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
5  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
6
7  MPI_TYPE_CREATE_KEYVAL(TYPE_COPY_ATTR_FN, TYPE_DELETE_ATTR_FN, TYPE_KEYVAL,
8      EXTRA_STATE, IERROR)
9  EXTERNAL TYPE_COPY_ATTR_FN, TYPE_DELETE_ATTR_FN
10 INTEGER TYPE_KEYVAL, IERROR
11 INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE

```

The argument `type_copy_attr_fn` may be specified as `MPI_TYPE_NULL_COPY_FN` or `MPI_TYPE_DUP_FN` from either C or Fortran. `MPI_TYPE_NULL_COPY_FN` is a function that does nothing other than returning `flag = 0` and `MPI_SUCCESS`. `MPI_TYPE_DUP_FN` is a simple-minded copy function that sets `flag = 1`, returns the value of `attribute_val_in` in `attribute_val_out`, and returns `MPI_SUCCESS`.

The argument `type_delete_attr_fn` may be specified as `MPI_TYPE_NULL_DELETE_FN` from either C or Fortran. `MPI_TYPE_NULL_DELETE_FN` is a function that does nothing, other than returning `MPI_SUCCESS`.

The C callback functions are:

```

21 typedef int MPI_Type_copy_attr_function(MPI_Datatype oldtype,
22     int type_keyval, void *extra_state, void *attribute_val_in,
23     void *attribute_val_out, int *flag);

```

and

```

25 typedef int MPI_Type_delete_attr_function(MPI_Datatype datatype,
26     int type_keyval, void *attribute_val, void *extra_state);

```

With the `mpi_f08` module, the Fortran callback functions are:

```

29 ABSTRACT INTERFACE
30   SUBROUTINE MPI_Type_copy_attr_function(oldtype, type_keyval, extra_state,
31       attribute_val_in, attribute_val_out, flag, ierror)
32     TYPE(MPI_Datatype) :: oldtype
33     INTEGER :: type_keyval, ierror
34     INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state, attribute_val_in,
35         attribute_val_out
36     LOGICAL :: flag

```

and

```

38 ABSTRACT INTERFACE
39   SUBROUTINE MPI_Type_delete_attr_function(datatype, type_keyval,
40       attribute_val, extra_state, ierror)
41     TYPE(MPI_Datatype) :: datatype
42     INTEGER :: type_keyval, ierror
43     INTEGER(KIND=MPI_ADDRESS_KIND) :: attribute_val, extra_state

```

With the `mpi` module and `mpif.h`, the Fortran callback functions are:

```

46 SUBROUTINE TYPE_COPY_ATTR_FUNCTION(OLDTYPE, TYPE_KEYVAL, EXTRA_STATE,
47     ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT, FLAG, IERROR)

```

```

    INTEGER OLDTYPE, TYPE_KEYVAL, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE,
        ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT
    LOGICAL FLAG
and
SUBROUTINE TYPE_DELETE_ATTR_FUNCTION(DATATYPE, TYPE_KEYVAL, ATTRIBUTE_VAL,
    EXTRA_STATE, IERROR)
    INTEGER DATATYPE, TYPE_KEYVAL, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE

    If an attribute copy function or attribute delete function returns other than
    MPI_SUCCESS, then the call that caused it to be invoked (for example, MPI_TYPE_FREE),
    is erroneous.

MPI_TYPE_FREE_KEYVAL(type_keyval)
    INOUT    type_keyval                key value (integer)

int MPI_Type_free_keyval(int *type_keyval)
MPI_Type_free_keyval(type_keyval, ierror)
    INTEGER, INTENT(INOUT) :: type_keyval
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_TYPE_FREE_KEYVAL(TYPE_KEYVAL, IERROR)
    INTEGER TYPE_KEYVAL, IERROR

MPI_TYPE_SET_ATTR(datatype, type_keyval, attribute_val)
    INOUT    datatype                datatype to which attribute will be attached (handle)
    IN       type_keyval                key value (integer)
    IN       attribute_val            attribute value

int MPI_Type_set_attr(MPI_Datatype datatype, int type_keyval,
    void *attribute_val)

MPI_Type_set_attr(datatype, type_keyval, attribute_val, ierror)
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    INTEGER, INTENT(IN) :: type_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: attribute_val
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_TYPE_SET_ATTR(DATATYPE, TYPE_KEYVAL, ATTRIBUTE_VAL, IERROR)
    INTEGER DATATYPE, TYPE_KEYVAL, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL

```

```

1 MPI_TYPE_GET_ATTR(datatype, type_keyval, attribute_val, flag)
2     IN      datatype          datatype to which the attribute is attached (handle)
3     IN      type_keyval       key value (integer)
4     OUT     attribute_val     attribute value, unless flag = false
5     OUT     flag              false if no attribute is associated with the key (logical)

```

```

8 int MPI_Type_get_attr(MPI_Datatype datatype, int type_keyval,
9                       void *attribute_val, int *flag)

```

```

10 MPI_Type_get_attr(datatype, type_keyval, attribute_val, flag, ierror)
11     TYPE(MPI_Datatype), INTENT(IN) :: datatype
12     INTEGER, INTENT(IN) :: type_keyval
13     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val
14     LOGICAL, INTENT(OUT) :: flag
15     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

16 MPI_TYPE_GET_ATTR(DATATYPE, TYPE_KEYVAL, ATTRIBUTE_VAL, FLAG, IERROR)
17     INTEGER DATATYPE, TYPE_KEYVAL, IERROR
18     INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
19     LOGICAL FLAG

```

```

20 MPI_TYPE_DELETE_ATTR(datatype, type_keyval)
21     INOUT   datatype          datatype from which the attribute is deleted (handle)
22     IN      type_keyval       key value (integer)

```

```

23 int MPI_Type_delete_attr(MPI_Datatype datatype, int type_keyval)

```

```

24 MPI_Type_delete_attr(datatype, type_keyval, ierror)
25     TYPE(MPI_Datatype), INTENT(IN) :: datatype
26     INTEGER, INTENT(IN) :: type_keyval
27     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

28 MPI_TYPE_DELETE_ATTR(DATATYPE, TYPE_KEYVAL, IERROR)
29     INTEGER DATATYPE, TYPE_KEYVAL, IERROR

```

6.7.5 Error Class for Invalid Keyval

Key values for attributes are system-allocated, by `MPI_{TYPE,COMM,WIN}_CREATE_KEYVAL`. Only such values can be passed to the functions that use key values as input arguments. In order to signal that an erroneous key value has been passed to one of these functions, there is a new MPI error class: `MPI_ERR_KEYVAL`. It can be returned by `MPI_ATTR_PUT`, `MPI_ATTR_GET`, `MPI_ATTR_DELETE`, `MPI_KEYVAL_FREE`, `MPI_{TYPE,COMM,WIN}_DELETE_ATTR`, `MPI_{TYPE,COMM,WIN}_SET_ATTR`, `MPI_{TYPE,COMM,WIN}_GET_ATTR`, `MPI_{TYPE,COMM,WIN}_FREE_KEYVAL`, `MPI_COMM_DUP`, `MPI_COMM_IDUP`,

MPI_COMM_DISCONNECT, and MPI_COMM_FREE. The last four are included because keyval is an argument to the copy and delete functions for attributes.

6.7.6 Attributes Example

Advice to users. This example shows how to write a collective communication operation that uses caching to be more efficient after the first call. (*End of advice to users.*)

```

/* key for this module's stuff: */
static int gop_key = MPI_KEYVAL_INVALID;

typedef struct
{
    int ref_count;          /* reference count */
    /* other stuff, whatever else we want */
} gop_stuff_type;

void Efficient_Collective_Op(MPI_Comm comm, ...)
{
    gop_stuff_type *gop_stuff;
    MPI_Group      group;
    int            foundflag;

    MPI_Comm_group(comm, &group);

    if (gop_key == MPI_KEYVAL_INVALID) /* get a key on first call ever */
    {
        if ( ! MPI_Comm_create_keyval(gop_stuff_copier,
                                      gop_stuff_destructor,
                                      &gop_key, (void *)0)) {
            /* get the key while assigning its copy and delete callback
               behavior. */
        } else
            MPI_Abort(comm, 99);
    }

    MPI_Comm_get_attr(comm, gop_key, &gop_stuff, &foundflag);
    if (foundflag)
    { /* This module has executed in this group before.
       We will use the cached information */
    }
    else
    { /* This is a group that we have not yet cached anything in.
       We will now do so.
       */

        /* First, allocate storage for the stuff we want,

```

```

1      and initialize the reference count */
2
3      gop_stuff = (gop_stuff_type *) malloc(sizeof(gop_stuff_type));
4      if (gop_stuff == NULL) { /* abort on out-of-memory error */ }
5
6      gop_stuff->ref_count = 1;
7
8      /* Second, fill in *gop_stuff with whatever we want.
9         This part isn't shown here */
10
11     /* Third, store gop_stuff as the attribute value */
12     MPI_Comm_set_attr(comm, gop_key, gop_stuff);
13 }
14 /* Then, in any case, use contents of *gop_stuff
15    to do the global op ... */
16 }
17
18 /* The following routine is called by MPI when a group is freed */
19
20 int gop_stuff_destructor(MPI_Comm comm, int keyval, void *gop_stuffP,
21                          void *extra)
22 {
23     gop_stuff_type *gop_stuff = (gop_stuff_type *)gop_stuffP;
24     if (keyval != gop_key) { /* abort -- programming error */ }
25
26     /* The group's being freed removes one reference to gop_stuff */
27     gop_stuff->ref_count -= 1;
28
29     /* If no references remain, then free the storage */
30     if (gop_stuff->ref_count == 0) {
31         free((void *)gop_stuff);
32     }
33     return MPI_SUCCESS;
34 }
35
36 /* The following routine is called by MPI when a group is copied */
37 int gop_stuff_copier(MPI_Comm comm, int keyval, void *extra,
38                      void *gop_stuff_inP, void *gop_stuff_outP, int *flag)
39 {
40     gop_stuff_type *gop_stuff_in = (gop_stuff_type *)gop_stuff_inP;
41     gop_stuff_type **gop_stuff_out = (gop_stuff_type **)gop_stuff_outP;
42     if (keyval != gop_key) { /* abort -- programming error */ }
43
44     /* The new group adds one reference to this gop_stuff */
45     gop_stuff_in->ref_count += 1;
46     *gop_stuff_out = gop_stuff_in;
47     return MPI_SUCCESS;
48 }

```


6.8 Naming Objects

There are many occasions on which it would be useful to allow a user to associate a printable identifier with an MPI communicator, window, or datatype, for instance error reporting, debugging, and profiling. The names attached to opaque objects do not propagate when the object is duplicated or copied by MPI routines. For communicators this can be achieved using the following two functions.

MPI_COMM_SET_NAME (comm, comm_name)

INOUT	comm	communicator whose identifier is to be set (handle)
IN	comm_name	the character string which is remembered as the name (string)

```
int MPI_Comm_set_name(MPI_Comm comm, const char *comm_name)
```

```
MPI_Comm_set_name(comm, comm_name, ierror)
  TYPE(MPI_Comm), INTENT(IN) :: comm
  CHARACTER(LEN=*), INTENT(IN) :: comm_name
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_COMM_SET_NAME(COMM, COMM_NAME, IERROR)
```

```
  INTEGER COMM, IERROR
  CHARACTER*(*) COMM_NAME
```

MPI_COMM_SET_NAME allows a user to associate a name string with a communicator. The character string which is passed to **MPI_COMM_SET_NAME** will be saved inside the MPI library (so it can be freed by the caller immediately after the call, or allocated on the stack). Leading spaces in name are significant but trailing ones are not.

MPI_COMM_SET_NAME is a local (non-collective) operation, which only affects the name of the communicator as seen in the process which made the **MPI_COMM_SET_NAME** call. There is no requirement that the same (or any) name be assigned to a communicator in every process where it exists.

Advice to users. Since **MPI_COMM_SET_NAME** is provided to help debug code, it is sensible to give the same name to a communicator in all of the processes where it exists, to avoid confusion. (*End of advice to users.*)

The length of the name which can be stored is limited to the value of **MPI_MAX_OBJECT_NAME** in Fortran and **MPI_MAX_OBJECT_NAME-1** in C to allow for the null terminator. Attempts to put names longer than this will result in truncation of the name. **MPI_MAX_OBJECT_NAME** must have a value of at least 64.

Advice to users. Under circumstances of store exhaustion an attempt to put a name of any length could fail, therefore the value of **MPI_MAX_OBJECT_NAME** should be viewed only as a strict upper bound on the name length, not a guarantee that setting names of less than this length will always succeed. (*End of advice to users.*)

Advice to implementors. Implementations which pre-allocate a fixed size space for a name should use the length of that allocation as the value of **MPI_MAX_OBJECT_NAME**.

Implementations which allocate space for the name from the heap should still define `MPI_MAX_OBJECT_NAME` to be a relatively small value, since the user has to allocate space for a string of up to this size when calling `MPI_COMM_GET_NAME`. (*End of advice to implementors.*)

`MPI_COMM_GET_NAME (comm, comm_name, resultlen)`

IN	comm	communicator whose name is to be returned (handle)
OUT	comm_name	the name previously stored on the communicator, or an empty string if no such name exists (string)
OUT	resultlen	length of returned name (integer)

`int MPI_Comm_get_name(MPI_Comm comm, char *comm_name, int *resultlen)`

`MPI_Comm_get_name(comm, comm_name, resultlen, ierror)`

TYPE(MPI_Comm), INTENT(IN) :: comm

CHARACTER(LEN=MPI_MAX_OBJECT_NAME), INTENT(OUT) :: comm_name

INTEGER, INTENT(OUT) :: resultlen

INTEGER, OPTIONAL, INTENT(OUT) :: ierror

`MPI_COMM_GET_NAME(COMM, COMM_NAME, RESULTLEN, IERROR)`

INTEGER COMM, RESULTLEN, IERROR

CHARACTER*(*) COMM_NAME

`MPI_COMM_GET_NAME` returns the last name which has previously been associated with the given communicator. The name may be set and retrieved from any language. The same name will be returned independent of the language used. `name` should be allocated so that it can hold a resulting string of length `MPI_MAX_OBJECT_NAME` characters.

`MPI_COMM_GET_NAME` returns a copy of the set name in `name`.

In C, a null character is additionally stored at `name[resultlen]`. The value of `resultlen` cannot be larger than `MPI_MAX_OBJECT_NAME-1`. In Fortran, `name` is padded on the right with blank characters. The value of `resultlen` cannot be larger than `MPI_MAX_OBJECT_NAME`.

If the user has not associated a name with a communicator, or an error occurs, `MPI_COMM_GET_NAME` will return an empty string (all spaces in Fortran, "" in C). The three predefined communicators will have predefined names associated with them. Thus, the names of `MPI_COMM_WORLD`, `MPI_COMM_SELF`, and the communicator returned by `MPI_COMM_GET_PARENT` (if not `MPI_COMM_NULL`) will have the default of `MPI_COMM_WORLD`, `MPI_COMM_SELF`, and `MPI_COMM_PARENT`. The fact that the system may have chosen to give a default name to a communicator does not prevent the user from setting a name on the same communicator; doing this removes the old name and assigns the new one.

Rationale. We provide separate functions for setting and getting the name of a communicator, rather than simply providing a predefined attribute key for the following reasons:

- It is not, in general, possible to store a string as an attribute from Fortran.

- It is not easy to set up the delete function for a string attribute unless it is known to have been allocated from the heap.
- To make the attribute key useful additional code to call `strdup` is necessary. If this is not standardized then users have to write it. This is extra unneeded work which we can easily eliminate.
- The Fortran binding is not trivial to write (it will depend on details of the Fortran compilation system), and will not be portable. Therefore it should be in the library rather than in user code.

(End of rationale.)

Advice to users. The above definition means that it is safe simply to print the string returned by `MPI_COMM_GET_NAME`, as it is always a valid string even if there was no name.

Note that associating a name with a communicator has no effect on the semantics of an MPI program, and will (necessarily) increase the store requirement of the program, since the names must be saved. Therefore there is no requirement that users use these functions to associate names with communicators. However debugging and profiling MPI applications may be made easier if names are associated with communicators, since the debugger or profiler should then be able to present information in a less cryptic manner. *(End of advice to users.)*

The following functions are used for setting and getting names of datatypes. The constant `MPI_MAX_OBJECT_NAME` also applies to these names.

`MPI_TYPE_SET_NAME (datatype, type_name)`

INOUT	datatype	datatype whose identifier is to be set (handle)
IN	type_name	the character string which is remembered as the name (string)

`int MPI_Type_set_name(MPI_Datatype datatype, const char *type_name)`

```

MPI_Type_set_name(datatype, type_name, ierror)
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  CHARACTER(LEN=*), INTENT(IN) :: type_name
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_TYPE_SET_NAME(DATATYPE, TYPE_NAME, IERROR)
  INTEGER DATATYPE, IERROR
  CHARACTER*(*) TYPE_NAME

```

1 `MPI_TYPE_GET_NAME (datatype, type_name, resultlen)`

2	IN	datatype	datatype whose name is to be returned (handle)
3			
4	OUT	type_name	the name previously stored on the datatype, or a empty string if no such name exists (string)
5			
6	OUT	resultlen	length of returned name (integer)

7

8 `int MPI_Type_get_name(MPI_Datatype datatype, char *type_name,`
 9 `int *resultlen)`

10

11 `MPI_Type_get_name(datatype, type_name, resultlen, ierror)`
 12 `TYPE(MPI_Datatype), INTENT(IN) :: datatype`
 13 `CHARACTER(LEN=MPI_MAX_OBJECT_NAME), INTENT(OUT) :: type_name`
 14 `INTEGER, INTENT(OUT) :: resultlen`
 15 `INTEGER, OPTIONAL, INTENT(OUT) :: ierror`

16 `MPI_TYPE_GET_NAME(DATATYPE, TYPE_NAME, RESULTLEN, IERROR)`
 17 `INTEGER DATATYPE, RESULTLEN, IERROR`
 18 `CHARACTER*(*) TYPE_NAME`

19

20 Named predefined datatypes have the default names of the datatype name. For exam-
 21 ple, `MPI_WCHAR` has the default name of `MPI_WCHAR`.

22 The following functions are used for setting and getting names of windows. The con-
 23 stant `MPI_MAX_OBJECT_NAME` also applies to these names.

24

25 `MPI_WIN_SET_NAME (win, win_name)`

26			
27	INOUT	win	window whose identifier is to be set (handle)
28	IN	win_name	the character string which is remembered as the name
29			(string)

30

31 `int MPI_Win_set_name(MPI_Win win, const char *win_name)`

32

33 `MPI_Win_set_name(win, win_name, ierror)`
 34 `TYPE(MPI_Win), INTENT(IN) :: win`
 35 `CHARACTER(LEN=*) , INTENT(IN) :: win_name`
 36 `INTEGER, OPTIONAL, INTENT(OUT) :: ierror`

37 `MPI_WIN_SET_NAME(WIN, WIN_NAME, IERROR)`
 38 `INTEGER WIN, IERROR`
 39 `CHARACTER*(*) WIN_NAME`

```

MPI_WIN_GET_NAME (win, win_name, resultlen)
    IN      win      window whose name is to be returned (handle)
    OUT     win_name  the name previously stored on the window, or a empty
                      string if no such name exists (string)
    OUT     resultlen length of returned name (integer)

int MPI_Win_get_name(MPI_Win win, char *win_name, int *resultlen)

MPI_Win_get_name(win, win_name, resultlen, ierror)
    TYPE(MPI_Win), INTENT(IN) :: win
    CHARACTER(LEN=MPI_MAX_OBJECT_NAME), INTENT(OUT) :: win_name
    INTEGER, INTENT(OUT) :: resultlen
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_WIN_GET_NAME(WIN, WIN_NAME, RESULTLEN, IERROR)
    INTEGER WIN, RESULTLEN, IERROR
    CHARACTER*(*) WIN_NAME

```

6.9 Formalizing the Loosely Synchronous Model

In this section, we make further statements about the loosely synchronous model, with particular attention to intra-communication.

6.9.1 Basic Statements

When a caller passes a communicator (that contains a context and group) to a callee, that communicator must be free of side effects throughout execution of the subprogram: there should be no active operations on that communicator that might involve the process. This provides one model in which libraries can be written, and work “safely.” For libraries so designated, the callee has permission to do whatever communication it likes with the communicator, and under the above guarantee knows that no other communications will interfere. Since we permit good implementations to create new communicators without synchronization (such as by preallocated contexts on communicators), this does not impose a significant overhead.

This form of safety is analogous to other common computer-science usages, such as passing a descriptor of an array to a library routine. The library routine has every right to expect such a descriptor to be valid and modifiable.

6.9.2 Models of Execution

In the loosely synchronous model, transfer of control to a **parallel procedure** is effected by having each executing process invoke the procedure. The invocation is a collective operation: it is executed by all processes in the execution group, and invocations are similarly ordered at all processes. However, the invocation need not be synchronized.

We say that a parallel procedure is *active* in a process if the process belongs to a group that may collectively execute the procedure, and some member of that group is currently executing the procedure code. If a parallel procedure is active in a process, then this process

1 may be receiving messages pertaining to this procedure, even if it does not currently execute
 2 the code of this procedure.

3 4 Static Communicator Allocation

5 This covers the case where, at any point in time, at most one invocation of a parallel
 6 procedure can be active at any process, and the group of executing processes is fixed. For
 7 example, all invocations of parallel procedures involve all processes, processes are single-
 8 threaded, and there are no recursive invocations.

9 In such a case, a communicator can be statically allocated to each procedure. The
 10 static allocation can be done in a preamble, as part of initialization code. If the parallel
 11 procedures can be organized into libraries, so that only one procedure of each library can
 12 be concurrently active in each processor, then it is sufficient to allocate one communicator
 13 per library.

15 16 Dynamic Communicator Allocation

17 Calls of parallel procedures are well-nested if a new parallel procedure is always invoked in
 18 a subset of a group executing the same parallel procedure. Thus, processes that execute
 19 the same parallel procedure have the same execution stack.

20 In such a case, a new communicator needs to be dynamically allocated for each new
 21 invocation of a parallel procedure. The allocation is done by the caller. A new communicator
 22 can be generated by a call to `MPI_COMM_DUP`, if the callee execution group is identical to
 23 the caller execution group, or by a call to `MPI_COMM_SPLIT` if the caller execution group
 24 is split into several subgroups executing distinct parallel routines. The new communicator
 25 is passed as an argument to the invoked routine.

26 The need for generating a new communicator at each invocation can be alleviated or
 27 avoided altogether in some cases: If the execution group is not split, then one can allocate
 28 a stack of communicators in a preamble, and next manage the stack in a way that mimics
 29 the stack of recursive calls.

30 One can also take advantage of the well-ordering property of communication to avoid
 31 confusing caller and callee communication, even if both use the same communicator. To do
 32 so, one needs to abide by the following two rules:

- 34 • messages sent before a procedure call (or before a return from the procedure) are also
 35 received before the matching call (or return) at the receiving end;
- 36 • messages are always selected by source (no use is made of `MPI_ANY_SOURCE`).

38 39 The General Case

40 In the general case, there may be multiple concurrently active invocations of the same
 41 parallel procedure within the same group; invocations may not be well-nested. A new
 42 communicator needs to be created for each invocation. It is the user's responsibility to make
 43 sure that, should two distinct parallel procedures be invoked concurrently on overlapping
 44 sets of processes, communicator creation is properly coordinated.

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