Optimised scheduling algorithm

Student Info:

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Optimising for:

Rental Cost

Links:

GitHub: https://github.com/mq45946396/comp3100-project

Introduction:

This project is a modified version of the stage one project code, with a new scheduling algorithm implemented in the project. The goal of the new algorithm is to improve the metrics over the all-to-largest (ATL) algorithm which was implemented in the stage one implementation (specifically, the rental cost of the algorithm).

In practice, compared to the previous algorithm it was able to optimize turnaround time and rental cost pretty consistently. However, it was not able to efficiently improve turnaround time or resource utilization compared to the three baseline algorithms (BF, FF and WF), and was only able to improve the rental cost in some configurations of the simulator compared to the baseline algorithms.

This is achieved because the algorithm attempts to allocate each job to the server which has the closest number of resources available for the job, which increases the efficiency of job cost by picking the most appropriate server based on the job's CPU requirements, and reduce the amount of time a job must wait before running on any given server.

Problem Definition

The primary issue with the current ATL algorithm is that it is very inefficient and does not use the server resources in the most effective manner, especially when compared to the three baseline algorithms. It performs the worst when looking at turnaround time. As shown below, when compared to the three baseline algorithms (FF, BF and WF respectively), ATL turnaround time performance is several times worse than the other baseline algorithms, meaning jobs will take much longer to complete when using the algorithm. This makes sense, as the largest server's cores will be quickly consumed and result in most jobs needing to queue for extensive periods of time before running.

Turnaround Time

Average	254086.33	1473.33	1462.83	6240.72

Despite this, ATL performs fairly well in terms of cost and resource utilisation. Since only one server is used, its cost per hour will be consistent, which would not happen with multiple servers being deployed. And since only one server is performing tasks, it is active at all times during the process, and therefore it's resource utilisation is always 100%.

Resource Util

Average	100.00	66.79	64.94	72.85
Cost				
Average	256.05	417.90	414.42	443.03

Algorithm Details

My new algorithm attempts to improve the scheduling efficiency by checking the core count of the servers which are available to be scheduled. It does this by calculating the difference between the number of cores required by the job and the number of cores which are available on the server. It then picks the server with the smallest difference and schedules the job to that. This means that each job goes to the server which is most appropriate for each given job. In an effort to promote preloading jobs while booting, the algorithm will also provide a bias to the difference value, halving it if the server is booting, ensuring that currently booting servers are more likely to get the jobs.

In pseudo code, the algorithm works as follows:

```
FOR EACH Job j:
GETS All AS servers
FOR EACH Server s:
let bootBias = 2 if s.state == booting, else 1
let difference = abs(s.cores - j.cores) / bootBias
SORT servers BY difference DESCENDING
SCHD j ON servers[0]
```

There may be situations where either:

- There are no servers with enough cores to complete the job
- There are no servers with sufficient memory or disk requirements

In this situation, the algorithm will schedule the job to the server which has the highest number of total cores (i.e. the largest server). However after some jobs have completed executing, the algorithm can continue allocating jobs to the most appropriately sized server.

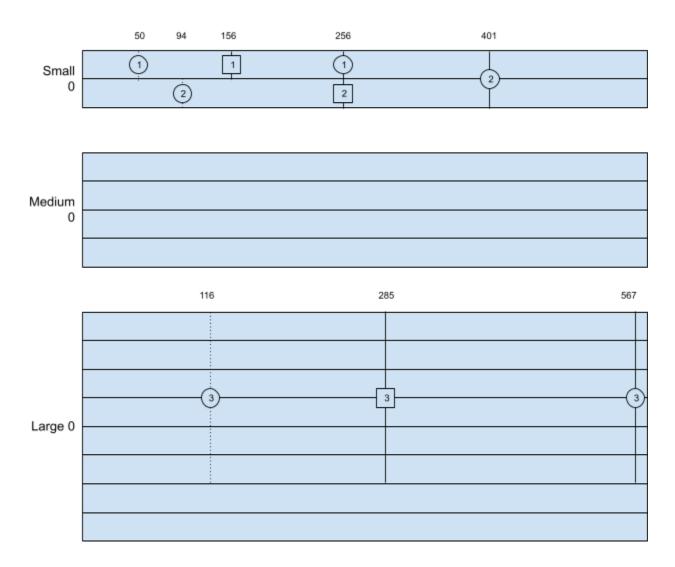
Scheduling Example

For this scenario, the following job and server configurations will be used:

Job	Cores Required	Server	# Cores
1	1	Small 0	2
2	2	Medium 0	4
3	6	Large 0	8

The resulting scheduling should result in the first job being allocated to small 0, since it's core count is the closest to job 1's requirements. The second job will be allocated to small 0, since it will be currently booting and has the smallest difference. Finally, job 3 will be allocated to large 0, since it is the only server large enough to handle job 3 and has the smallest core difference.

In a diagram form, the schedule may look like this.



Performance

When running with the stage 2 configurations, the performance is unfortunately not as good as I hoped. While it is effective in reducing the rental cost of scheduling compared to the baseline algorithms, ATL is still able to achieve a cheaper rental cost. I believe that this is mostly caused by ATL's design using only a single server to run all jobs.

While it dramatically improves turnaround time compared to ATL, the algorithm's turnaround time is at least twice as slow as BF and FF, but is half as slow as WF. It is quite similar in performance to the baseline algorithms in terms of utilization, but is still slightly worse compared to the baseline algorithms, and is significantly worse than ATL which achieves 100% utilization, since it only uses one server.

ATL	FF	BF	WF	Yours
672786	2428	2450	29714	8571
316359	2458	2458	2613	3674
679829	2356	2362	10244	5753
331382	1184	1198	12882	3853
283701	1205	1205	1245	2785
342754	1153	1154	4387	3993
244404	693	670	10424	2941
224174	673	673	746	1252
256797	645	644	5197	2706
240984	2852	2820	10768	6326
55746	2493	2494	2523	2757
139467	2491	2485	2803	3380
247673	1393	1254	8743	4953
52096	1209	1209	1230	1493
139670	1205	1205	1829	2242
145298	768	736	5403	3694
49299	665	665	704	763
151135	649	649	878	1248
254086.33	1473.33	1462.83	6240.72	3465.78
172.4568	1.0000	0.9929	4.2358	2.3523
173.6947	1.0072	1.0000	4.2662	2.3692
40.7143	0.2361	0.2344	1.0000	0.5553
	316359 679829 331382 283701 342754 244404 224174 256797 240984 55746 139467 247673 52096 139670 145298 49299 151135 254086.33 172.4568 173.6947	672786 2428 316359 2458 679829 2356 331382 1184 283701 1205 342754 1153 244404 693 224174 673 256797 645 240984 2852 55746 2493 139467 2491 247673 1393 52096 1209 145298 768 49299 665 151135 649 254086.33 1473.33 172.4568 1.0000 173.6947 1.0072	672786 2428 2458 316359 2458 2458 679829 2356 2362 331382 1184 1198 283701 1205 1205 342754 1153 1154 244404 693 670 224174 673 673 256797 645 644 240984 2852 2820 55746 2493 2494 139467 2491 2485 247673 1393 1254 52096 1209 1209 139670 1205 1205 145298 768 736 49299 665 665 151135 649 649 254086.33 1473.33 1462.83 172.4568 1.0000 0.9929 173.6947 1.0072 1.0000	672786 2428 2458 29714 316359 2458 2458 2613 679829 2356 2362 10244 331382 1184 1198 12882 283701 1205 1245 342754 1153 1154 4387 244404 693 670 10424 224174 673 673 746 256797 645 644 5197 240984 2852 2820 10768 55746 2493 2494 2523 139467 2491 2485 2803 247673 1393 1254 8743 52096 1209 1209 1230 145298 768 736 5403 49299 665 665 704 151135 649 649 878 254086.33 1473.33 1462.83 6240.72 172.4568 1.0000 0.9929 4.2358

Pros

- More cost effective than ATL, and is much more efficient in terms of turnaround time.
- Attempts to allocate jobs to the most appropriate server.

Cons

- In most situations, it is less efficient than the three baseline algorithms.
- When no suitable servers are present, it resorts to the inefficient ATL algorithm.

Conclusion

While the algorithm achieves its primary goal of reducing costs compared to the three baseline algorithms, it sacrifices its performance in turnaround time and resource utilization when compared to the other algorithms. That being said, it still performs significantly better than ATL in terms of turnaround time, making it a more efficient algorithm in terms of the time taken to run jobs.

In retrospect, I would have liked to better analyse the cause and effect of different scheduling techniques to derive an algorithm which performs better than all four existing algorithms.