# Computer Programming with Scala

Week 3: Functional Programming (FP)

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## Named and Default Parameters

## Refering paramters by name

#### Default Parameter values

### **Local Functions**

- ▶ Functions can be defined within other functions
- ► Functions are only visible in surrounding scope
- ▶ Inner function can access namespace of surrounding function

```
def filterEven(name: String, li:List[Int]):List[Int] = {
  def isEven(i:Int) = {
    println(name + " contains " + i)
    (i\%2 == 0)
  li match {
    case Nil => Nil
    case x::xs if (isEven(x)) => x::filterEven(name, xs)
                                     filterEven(name, xs)
    case x::xs
scala> filterEven("my list", List(1,2,3,4.5) )
mv list contains 1
my list contains 2
my list contains 3
my list contains 4
res0: List[Int] = List(2, 4)
```

# **Higher Order Functions**

#### First Class Functions → Functions are regular values

- ► Can be assigned to a variable
- Can be passed as arguments to functions
- Can be returned by other functions

## Higher Order Functions = Functions taking function as parameter

Powerful abstraction mechanism

```
def my_map (lst: List[Int] , fun: Int => Int) :List[Int] =
   for (l <- lst) yield fun (l)

val numbers = List (2 ,3 ,4 ,5)
def addone ( n : Int ) = n + 1

scala> my_map ( numbers , addone )
res0: List[Int] = List (3 , 4 , 5 , 6)
```

# Higher Order Functions on class List

### Filtering and Partitioning

► Functions as (named) values

```
val li = List(1, 2, 3, 4, 5)
def isEven (n: Int) = n%2 == 0
scala> li filter isEven
res0: List[Int] = List(2, 4)
```

► With an anonymous functions

```
scala> li filter (i => i%2 == 0)
res1: List[Int] = List(2, 4)
scala> li filter (_%2 == 0)
res2: List[Int] = List(2, 4)
```

```
scala> li partition (_%2 == 0)
res3: (List[Int], List[Int]) = (List(2, 4),List(1, 3, 5))
```

▶ Also defined: find, takeWhile, dropWhile and span. Check the doc

### Mapping over elements

```
scala> li map (_ + 1)
res4: List[Int] = List (2, 3 , 4 , 5 , 6)
scala> li foreach (x => print(x + ", "))
1, 2, 3, 4, 5,
```

## Folding List /: and \:

▶ Reduce all elements into a single value using the provided function

```
scala> def sum(xs: List[Int]): Int = (0 /: xs) (_ + _)
scala> sum( List(1,2,3,4) )
res0:Int = 10  # = 0 + 1 + 2 + 3 + 4
```

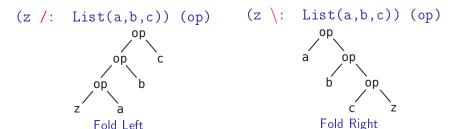
Variance

# Folding List /: and \:

▶ Reduce all elements into a single value using the provided function

```
scala> def sum(xs: List[Int]): Int = (0 /: xs) (_ + _)
scala> sum( List(1,2,3,4) )
res0:Int = 10  # = 0 + 1 + 2 + 3 + 4
scala> def sumRight(xs: List[Int]): Int = (0 \: xs) (_ + _)
scala> sum( List(1,2,3,4) )
res0:Int = 10  # = 0 + 4 + 3 + 2 + 1
```

▶ (z /: xs) (op) z: initial value, xs: list, op: operation to apply



Same result if op is associative; performance may vary

# Partially Applied Functions: Functions as Objects

- ▶ Passing in place of parameter list creates a partially applied function
- ► Function Object automatically built by the compiler

```
scala> def sum(a: Int, b: Int, c: Int) = a + b + c
sum: (a: Int, b: Int, c: Int)Int
scala > sum(1, 2, 3)
res0: Int = 6
                                             # This creates an object of type
scala> val a = sum _
                                             # <function3> (because sum takes
a: (Int, Int, Int) => Int = <function3>
                                             # 3 parameters)
scala > a(1.2.3)
                                             # Apply parameters to partially
res1: Int = 6
                                             # applied function => function call
scala > a.apply(1,2,3)
                                             # Exactly as before
res2: Int = 6
scala> val b = sum(1, _:Int , 3)
                                             # Here, only one parameter remains
b: Int => Int = <function1>
                                             # free. Thus the type <function1>
# Manual and bothersome definition (much simpler if it takes only one parameter)
scala> val f = \{case(a,b,c) \Rightarrow a + b + c\}: (Int,Int,Int) \Rightarrow Int
f: (Int, Int, Int) => Int = <function3>
```

# **Function Objects and Implicits**

▶ Underscore optional in contexts that require a function (and only there)

```
scala> someNumbers.foreach(print) # no need to write (print _) here
1234

scala> val c = sum
<console>:5: error: missing arguments for method sum...
follow this method with '_' if you want to treat it as a partially applied function
```

- ► Haskell doesn't require the for the partially applied function (implicit)
- ▶ But Scala targets Java developers ~ needs to detect missing parameters
- Thus the need for in general context
- is still optional where it can be no mistake

### Closures

▶ Free variable: variable without a value; Bound variable: variable with a value

#### Closure = when a function refers to an external free variable

```
scala> var more = 1
scala> val addMore = (x: Int) => x + more
addMore: (Int) => Int = <function1>
scala > addMore(10)
res0: Int = 11
```

- ▶ This function object is a closure, because it encloses (packs) the free variables
- ▶ In scala, captures the variables, not the values (Java captures constants)

```
scala> more = 3 ; addMore(10)
res1: Int = 13
```

#### Building Closures

```
scala> def makeIncreaser(more: Int) = (x: Int) => x + more
makeIncreaser: (more: Int)Int => Int
scala> val inc9999 = makeIncreaser(9999)
inc9999: (Int) => Int = <function1>
```

► Each closure gets its own copy of the captured elements

## Other Considerations

### Code Factorization with Higher Order Functions

```
def withOdd(nums: List[Int]): Boolean={
  var exists = false
  for (num <- nums)
    if (num \% 2 == 1)
      exists = true
  exists
}
```

```
def withOdd(nums: List[Int]): Boolean=
     nums.exists(_\%2 == 1)
```

- ▶ Q1: Implement List length with :/
- Q2: List reverse() with :/
- ▶ Q3: Type of |((x:Double) => x+1)|
- Q4: Write a function that adds 1 to every elements of a List[Int]
- ▶ Q5: Define  $S = \{a \times 2 \mid a \in [1, 100] \land a^2 < 99 \land a^3 > 9\}$
- ightharpoonup Q6: Explain | ((:Double)+2) | and | (:String).size

#### Tail Recusion Optimization

- Scala can optimize every tail recursive functions into a while loop
- Works only for basic forms (not mutually recursive, not partially applied)

# Lazy variables [lazy val ui = ...

Only evaluated on need (usually, scala values are evaluated when defined)

# Currying

Defining functions with multiple parameter lists

```
scala> def curriedSum(x: Int)(y: Int) = x + y
curriedSum: (x: Int)(y: Int)Int
scala> curriedSum(1)(2)
res5: Int = 3
```

You are actually defining two functions back to back

```
scala> def first(x: Int) = (y: Int) => x + y
first: (x: Int)(Int) => Int
scala> val second = first(1)
second: (Int) => Int = <function1>
```

Currying and Partially applied function

```
scala> curriedSum(1)
<console>:14: error: missing arguments for method curriedSum; follow this method wi
scala> curriedSum(1)_
res6: Int => Int = <function1>
```

This explains the :/ syntax

# **Function Composition**

```
def f(s: String) = "f(" + s + ")"
def g(s: String) = "g(" + s + ")"
compose makes a new function that composes its parameters: f(g(x))
scala> val FoG = f _ compose g _
FoG: String => String = <function1>
scala> FoG("yah")
res0: String = f(g(yah))
and Then does the same in the reverse order: g(f(x))
scala> val FthenG = f _ andThen g _
FthenG: String => String = <function1>
scala> FthenG("yah")
res1: String = g(f(yah))
```

### **PartialFunction**

- ▶ It's a function that is not defined for every parameter value
- It is not a Partially Applied Function

```
scala> val one: PartialFunction[Int, String] = { case 1 => "one" }
one: PartialFunction[Int,String] = <function1>
scala> one.isDefinedAt(1)
                                                scala> one.isDefinedAt(2)
res0: Boolean = true
                                                res1: Boolean = false
```

You can chain PartialFunctions with orElse

```
scala> val two: PartialFunction[Int, String] = { case 2 => "two" }
two: PartialFunction[Int,String] = <function1>
scala> val three: PartialFunction[Int, String] = { case 3 => "three" }
scala> val wildcard: PartialFunction[Int, String] = { case _ => "something else" }
scala> val partial = one orElse two orElse three orElse wildcard
partial: PartialFunction[Int,String] = <function1>
scala> partial(5)
                                      scala> partial(3)
res1: String = something else
                                      res2: String = three
scala> partial(2)
                                      scala> partial(1)
res3: String = two
                                      res4: String = two
```

# case class and Pattern Matching

### Defining a case class

```
abstract class Tree
case class Branch(left: Tree, right: Tree) extends Tree
case class Leaf(x: Int) extends Tree
```

#### Declaring a value

```
val t = Branch(Branch(Leaf(1), Leaf(2)), Branch(Leaf(3), Leaf(4)))
```

## Pattern Matching

```
def sumLeaves(t: Tree): Int = t match {
  case Branch(1, r) => sumLeaves(1) + sumLeaves(r)
  case Leaf(x) => x
}
```

### Matching on Variable Declaration

```
scala> val b = Branch(Leaf(1), Leaf(2))
b: Branch
scala> val Branch(_, 1) = b
```

1: Tree = Leaf(2)

Variance

# Parametrized types

### Defining a Tree[String] (without duplication)

```
abstract class Tree[A]
case class Branch[A](left: Tree[A], right: Tree[A]) extends Tree[A]
case class Leaf[A](x: A) extends Tree[A]

scala> val t = Branch(Branch(Leaf("a"), Leaf("b")), Branch(Leaf("c"), Leaf("d")))
t: Branch[String] = Branch(Branch(Leaf(a), Leaf(b)), Branch(Leaf(c), Leaf(d)))
```

► Tree is a trait while Tree[Int] is a type

# Parametrized types

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abstract class Tree[A]
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scala> val t = Branch(Branch(Leaf("a"), Leaf("b")), Branch(Leaf("c"), Leaf("d")))
t: Branch[String] = Branch(Branch(Leaf(a), Leaf(b)), Branch(Leaf(c), Leaf(d)))
```

Tree is a trait while Tree[Int] is a type

#### The Option type

- ▶ When you search for a value in a list, you don't know whether you'll find it
- ▶ An Option[A] can either be a Some (containing a value) or a None

```
val capitals = Map("France" -> "Paris", "Japan" -> "Tokyo")
scala> capitals get "France"
res0: Option[java.lang.String] = Some(Paris)
scala> capitals get "North Pole"
res1: Option[java.lang.String] = None
```

▶ Q: Define the follwing methods over Tree[A]: find, map, foreach, filter

### **Variance**

- ▶ Would you say that a Tree[Int] is-a Tree[Any]?
- ▶ Is it ok to provide a Tree[Int] where a Tree[Any] was expected?

## **V**ariance

- Would you say that a Tree[Int] is-a Tree[Any]?
- Is it ok to provide a Tree[Int] where a Tree[Any] was expected?
- ▶ Intuitively, yes, but by default, Scala generic types are nonvariant
- If your type Tree is covariant (flexible), just say so:

```
trait Tree[+T] { ... }  # a Tree[Int] is indeed a Tree[Any]
```

In some cases, you can tell that your type is contravariant

```
trait Tree[-T] { ... }  # WRONG! a Tree[Any] cannot be a Tree[Int]!
```

Purely functionnal types are often covariant

## Variance

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▶ Purely functionnal types are often covariant

#### Mutable Data is often not Covariant

```
def get = current
  def set(x: T) { current = x }
}
val c1 = new Cell[String]("abc")
val c2: Cell[Any] = c1
c2.set(1)
val s: String = c1.get # WOOOOPS
```

class Cell[+T](init: T) { # WRONG

private[this] var current = init

- ► This would sets the string to 1!
- Type system actually prevents this

```
Cell.scala:7: error: covariant type T
occurs in contravariant position in
type T of value x
def set(x: T) = current = x
```

# Variance and subtyping

```
class Animal { val sound = "rustle" }
class Bird extends Animal { override val sound = "call" }
class Chicken extends Bird { override val sound = "cluck" }
```

#### Specialization: You need a Bird and have a Chicken. That's OK.

► This is the Liskov Substitution Principle

But you cannot use an Animal in place of a Bird

# Variance and subtyping

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▶ But you cannot use an Animal in place of a Bird

#### Function parameters are contravariants

- ► Can't use a function that takes a Chicken for a function that takes a Bird
  - ▶ It would choke on a Duck; But a function that takes an Animal is OK

```
scala> val getTweet: (Bird => String) = ((a: Animal) => a.sound )
getTweet: Bird => String = <function1>
```

# Variance and subtyping

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- ► Can't use a function that takes a Chicken for a function that takes a Bird
  - ▶ It would choke on a Duck; But a function that takes an Animal is OK

```
scala> val getTweet: (Bird => String) = ((a: Animal) => a.sound )
getTweet: Bird => String = <function1>
```

#### Function return value are covariant

▶ Need a function that returns a Bird? A function returning a Chicken is OK

```
scala> val hatch: (() => Bird) = ((_) => new Chicken )
hatch: () => Bird = <function0>
```

CC

Variance

```
class Top
class Middle extends Top
class Bottom extends Middle
class Up {
  def cv(t:Top) = "Up"
  def inv(m:Middle) = "Up"
  def ctv(b:Bottom) = "Up"
class Down extends Up {
  def cv(m:Middle) = "Down"
  def inv(m:Middle) ="Down"
  def ctv(m:Middle) = "Down"
val u: Up = new Up
val d: Down= new Down
```

val ud:Up = new Down

	u.??()	d.??()	ud.??()
?.cv(Top)			
?.cv(Middle)			
? cv(Bottom)			
?.inv(Top)			
?.inv(Middle)			
? inv(Bottom)			
?.ctv(Top)			
?.ctv(Middle)			
?.ctv(Bottom)			

## Scala 2.x algorithm to select the Right Call

- ▶ Get signature from static types; Linearize receiver dynamic type to find it
- ▶ Other languages (and Scala 1.x) use other algorithms
- ▶ Don't do it in Real Projects

Courtesy of Antoine Beugnard (Telecom Bretagne)

```
class Top
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?.cv(Mi	ddle)	Up		
?.cv(Bo	ttom)	Up		
?.inv(To	p)			
?.inv(M	iddle)			
?.inv(Bo	ottom)			
?.ctv(To	op)			
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```

val ud:Up = new Down

	u.??()	d.??()	ud.??()
?.cv(Top)	Up		
?.cv(Middle)	Up		
?.cv(Bottom)	Up		
?.inv(Top)	Error		
?.inv(Middle)	Up		
?.inv(Bottom)	Up		
?.ctv(Top)			
?.ctv(Middle)			
?.ctv(Bottom)			

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```
u.??()
                         d.??()
                                  ud.??()
?.cv(Top)
                  Up
                          Up
?.cv(Middle)
                  Up
                         Down
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                  Up
                         Down
?.inv(Top)
                 Error
                         Error
?.inv(Middle)
                  Up
                         Down
?.inv(Bottom)
                  Up
                         Down
? ctv(Top)
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```

d.cv(Top)=Up because parameters are contravariant

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val u: Up = new Up
                             d.cv(Top)=Up because parameters are contravariant
```

val d: Down= new Down

```
u.??()
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                          Error
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                  Up
                         Down
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                  Up
                         Down
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                          Error
                 Error
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                          Error
```

d.ctv(Bot) ambiguous:  $Up.ctv(Bot) \approx Down.ctv(Mid)$ val ud:Up = new Down

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val ud:Up = new Down

```
u.??()
                         d.??()
                                 ud.??()
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                  Up
                          Up
                                    Up
?.cv(Middle)
                  Up
                         Down
                                    Up
?.cv(Bottom)
                  Up
                         Down
                                    Up
?.inv(Top)
                 Error
                         Error
?.inv(Middle)
                  Up
                         Down
?.inv(Bottom)
                  Up
                         Down
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                         Error
                 Error
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                 Error
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d.cv(Top)=Up because parameters are contravariant

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```

val ud:Up = new Down

```
u.??()
                         d.??()
                                 ud.??()
?.cv(Top)
                 Up
                          Up
                                   Up
?.cv(Middle)
                 Up
                         Down
                                   Up
?.cv(Bottom)
                  Up
                         Down
                                   Up
?.inv(Top)
                 Error
                         Error
                                   Error
?.inv(Middle)
                 Up
                         Down
                                  Down
?.inv(Bottom)
                  Up
                         Down
                                  Down
? ctv(Top)
                         Error
                 Error
?.ctv(Middle)
                 Error
                         Down
?.ctv(Bottom)
                  Up
                         Error
```

d.cv(Top)=Up because parameters are contravariant

d.ctv(Bot) ambiguous:  $Up.ctv(Bot) \approx Down.ctv(Mid)$ 

- ▶ Get signature from static types; Linearize receiver dynamic type to find it
- ▶ Other languages (and Scala 1.x) use other algorithms
- ► Don't do it in Real Projects

  Courtesy of Antoine Beugnard (Telecom Bretagne)

  http://public.enst-bretagne.fr/~beugnard/papiers/lb-sem.shtml

```
class Top
class Middle extends Top
class Bottom extends Middle
class Up {
  def cv(t:Top) = "Up"
  def inv(m:Middle) = "Up"
  def ctv(b:Bottom) = "Up"
class Down extends Up {
  def cv(m:Middle) = "Down"
  def inv(m:Middle) ="Down"
  def ctv(m:Middle) = "Down"
val u: Up = new Up
val d: Down= new Down
```

val ud:Up = new Down

```
u.??()
                         d.??()
                                  ud.??()
?.cv(Top)
                  Up
                          Up
                                    Up
?.cv(Middle)
                  Up
                         Down
                                    Up
?.cv(Bottom)
                  Up
                         Down
                                    Up
?.inv(Top)
                 Error
                         Error
                                   Error
?.inv(Middle)
                  Up
                                  Down
                         Down
?.inv(Bottom)
                  Up
                         Down
                                  Down
? ctv(Top)
                         Error
                                   Error
                 Error
?.ctv(Middle)
                 Error
                         Down
                                   Error
?.ctv(Bottom)
                  Up
                         Error
                                    Up
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Scala 2.x algorithm to select the Right Call

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18/21

# Polymorphism Bounds

### Refine your polymorphism

biophony takes any T that is-a Animal

# Polymorphism Bounds

### Refine your polymorphism

```
scala> def cacophony[T](things: Seq[T]) = things map (_.sound)
<console>:7: error: value sound is not a member of type parameter T
       def cacophony[T](things: Seq[T]) = things map (_.sound)
scala> def biophony[T <: Animal](things: Seq[T]) = things map (_.sound)</pre>
biophony: [T <: Animal](things: Seq[T])Seq[java.lang.String]
scala> biophony(Seg(new Chicken, new Bird))
res5: Seq[java.lang.String] = List(cluck, call)
```

biophony takes any T that is-a Animal

### Lower bound: List[T] defines ::(elem T) but also ::(U >: T)

```
scala> val flock = List(new Bird. new Bird)
flock: List[Bird] = List(Bird@7e1ec70e, Bird@169ea8d2)
scala> new Chicken :: flock
res6: List[Bird] = List(Chicken@56fbda05, Bird@7e1ec70e, Bird@169ea8d2)
scala> new Animal :: flock
res7: List[Animal] = List(Animal@56fbda05, Bird@7e1ec70e, Bird@169ea8d2)
```

# More on Polymorphism

Quantification: When one you don't care about one type

# More on Polymorphism

#### Quantification: When one you don't care about one type

```
scala> def count[A](1: List[A]) = 1.size
count: [A](List[A])Int
                                          # A is useless
scala> def count(1: List[_]) = 1.size
count: (List[])Int
                                          # It's thus omitted
```

### Structural Types: specify type requirements by interface structure

```
scala> def foo(x: { def get: Int }) = 123 + x.get
foo: (x: AnyRef{def get: Int})Int
scala> foo(new { def get = 10 })
res0: Int = 133
```

- x parameter must be provided with a get function to be concrete
- ▶ The second class makes an ad-hoc anonymous class

# More on Polymorphism

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- x parameter must be provided with a get function to be concrete
- ▶ The second class makes an ad-hoc anonymous class

#### Erasure

- ▶ Unfortunately, the JVM erases every type specialization
- From List[Int], only List[\_] remains at runtime

# Take Home Messages

### Functional Programming

- ► Avoid mutable values, prefer expressions over statements
- ► Higher Order: pass functions as parameters (to factorize behavior)
- Partially Applied Functions: Function objects as first-class citizens
- Closures: functions that encapsulate some external state
- ► Currying: functions with multiple parameter lists
- ► Parametrized types: containers such as Tree[A]
- Variance permits to refine what we expect (the type system to our rescue)
   But don't mess with Receiver and Parameters' variances at the same time

#### FP in Scala

- ▶ Having both OOP and FP is nice and funny, but that's a lot of tools
- ▶ Getting used to them requires a lot of practice
- ► Some Scala choices are very debatable (it targets the Java ecosystem)