

# KUET\_Xtinction Team

## Notebook

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int	const N = 10000, mod = 12345;	
int	C[N][N];	
void	pre() {	
C[0][0] = 1;		
for(int n = 1; n < N; n++) {		
C[n][0] = 1;		
for (int k = 1; k <= n; k++) {		
C[n][k] = (C[n - 1][k - 1] + C[n - 1][k]) % mod;		
}		
}		
1.2	ncr modinv [23 lines] - eb48f866fa	
int	f[N], invf[N];	
void	pre() {	
f[0] = 1;		
for (int i = 1; i < N; i++) {		
f[i] = 1LL * i * f[i - 1] % mod;		
}		
invf[N - 1] = power(f[N - 1], mod - 2);		
for (int i = N - 2; i >= 0; i--) {		
invf[i] = 1LL * invf[i + 1] * (i + 1) % mod;		
}		
}		
int	nCr(int n, int r) {	
if	(n < r    n < 0) return 0;	
return	1LL * f[n] * invf[r] % mod * invf[n - r] % mod;	
}		
int	nCr(ll n, int r) {	
if	(n < r    n < 0) return 0;	
return	(1LL * ans * invf[r]) % mod;	
}		

```

        maxlen = k;           // Update
        LIS length
    }

    // Reconstruct LIS
    for (int i = mem[maxlen]; i != -1; i =
        prev_idx[i])
        res.push_back(num[i]);

    reverse(res.begin(), res.end()); // To get correct order
}

return res;
}

2.2 all node max dis [39 lines] - cf57f4a333
void solve() {
    int n;
    cin >> n;
    vector<vector<int>> g(n + 1);
    for(int i = 1; i < n; i++) {
        int u, v;
        cin >> u >> v;
        g[u].push_back(v);
        g[v].push_back(u);
    }

    vector<int> down(n + 1);
    vector<int> best1(n + 1, -1), best2(n
        + 1, -1);
    function<int(int, int)> dfs1 =
        [&](int u, int p) -> int {
            for(auto v : g[u]) if(v != p) {
                int cand = 1 + dfs1(v, u);

                if (cand > best1[u]) {
                    best2[u] = best1[u];
                    best1[u] = cand;
                } else if (cand > best2[u])
                    best2[u] = cand;
                down[u] = max(down[u],
                               cand);
            }
            return down[u];
        };
    dfs1(1, 0);

    vector<int> up(n + 1);
    function<void(int, int)> dfs2 =
        [&](int u, int p) -> void {
            for(auto v : g[u]) if(v != p) {
                int val = (down[v] + 1 ==
                           best1[u]) ? best2[u] :
                           best1[u];
                up[v] = 1 + max(up[u], val);
                dfs2(v, u);
            }
        };
    dfs2(1, 0);

    for(int i = 1; i <= n; i++)
        cout << max(up[i], down[i]) << "
            \n"[i == n];
}

```

**2.3 digit dp [14 lines] - d3853d7797**

---

```

ll dp[11][2][2];
string l,r; // l < r must
ll rec(ll ind, bool tightL, bool tightR){
    if(ind==l.size())return 0;
    if(dp[ind][tightL][tightR]!=-1)return
        dp[ind][tightL][tightR];
    char lo=tightL?l[ind]:'0',hi=_
        tightR?r[ind]:'9';
    ll ans=LLONG_MAX;
    for(char i=lo;i<=hi;i++){
        //Cost change kora lagbe
        //problemwise
        ll cost=(ll)(l[ind]==i)+(r[ind]==_
            =i);
        ans=min(ans,cost+rec(ind+1,tightL
            & l[ind]==i, tightR &
            r[ind]==i));
    }
    return dp[ind][tightL][tightR]=ans;
}

```

**2.4 longest path [27 lines] - b1bd570b6f**

---

```

queue<int> Q;
for(int i = 1; i <= n; i++) {
    if(indeg[i] == 0) Q.push(i);
}

vector<int> topo;
while(!Q.empty()) {
    int u = Q.front(); Q.pop();
    topo.push_back(u);
    for(auto v : g[u]) {
        if(--indeg[v] == 0) {
            Q.push(v);
        }
    }
}
vector<int> dp(n + 1, -1e9);
vector<int> par(n + 1, 0);
dp[1] = 1;
for(auto u : topo) {
    if(dp[u] == -1e9) continue;
    for(auto v : g[u]) {
        if(dp[u] + 1 > dp[v]) {
            dp[v] = dp[u] + 1;
            par[v] = u;
        }
    }
}

```

**2.5 longest zig zag path dp [92 lines] - 2ddaf1640b**

---

```

// find Longest zig zag path usnig BIT
// and DP
// red means peak ^
// white means vally v
// \/\/\/\/

struct BIT{ // stores (length, last-pos)
    using Node = pair<int,int>;
    int n; vector<Node> t; // (0,-1) = "nothing"
    BIT() {} // Node res{0,-1};
    void update(int index, int value) {
        for(int i = index; i < n; i += i & -i)
            t[i].first += value;
    }
    int query(int index) const {
        int sum = 0;
        for(int i = index; i >= 0; i -= i & -i)
            sum += t[i].first;
        return sum;
    }
}

```

```

BIT(int _n) : n(_n), t(n + 2, {0,-1}) {
}
Node query(int i) { // prefix-max
    Node res{0,-1};
    for (; i >= 1; i -= i & -i)
        if (t[i].first > res.first)
            res = t[i];
    return res;
}
void upd(int i, Node v) {
    for (; i <= n; i += i & -i)
        if (v.first > t[i].first)
            t[i] = v;
}

```

```

void solve() {
    int n;
    cin >> n;
    vector<int> v(n);
    for(auto &x : v) cin >> x;

    vector<int> tmp = v;
    sort(tmp.begin(), tmp.end());
    tmp.erase(unique(tmp.begin(),
                    tmp.end()), tmp.end());
    int m = tmp.size();
    vector<int> id(n), rid(n);
    for(int i = 0; i < n; i++) {
        id[i] = lower_bound(tmp.begin(),
                            tmp.end(), v[i]) -
            tmp.begin() + 1;
        rid[i] = m - id[i] + 1;
    }
}

```

```

BIT red(n), white(n);

vector<vector<int>> dp(2,
    vector<int>(n + 1, 0)); // 0 ->
    red 1 -> white
vector<vector<int>> pre(2,
    vector<int>(n + 1, -1)); // 0 ->
    red 1 -> white

```

```

for(int i = 0; i < n; i++) {
    // current is red
    auto [lenw, posw] =
        white.query(id[i] - 1);
    dp[0][i] = lenw + 1;
    pre[0][i] = posw;
    red.upd(rid[i], {dp[0][i], i});
    // current is white
    auto [lenr, posr] =
        red.query(rid[i] - 1);
    dp[1][i] = lenr + 1;
    pre[1][i] = posr;
    white.upd(id[i], {dp[1][i], i});
}

```

```

int mx = 0, pos = 0, f = 0;
for(int i = 0; i < n; i++) {
    if(mx < dp[0][i]) {
        mx = dp[0][i];
        pos = i;
        f = 0;
    }
}

```

```

}
if(mx < dp[1][i]) {
    mx = dp[1][i];
    pos = i;
    f = 1;
}
}

vector<int> ans;
while(pos >= 0) {
    ans.push_back(v[pos]);
    pos = pre[f][pos];
    f ^= 1;
}
reverse(ans.begin(), ans.end());
cout << ans.size() << "\n";
for(auto x : ans) {
    cout << x << " ";
}
cout << "\n";

int main() {
    ios_base::sync_with_stdio(false);
    cin.tie(0);
    int tc = 1;
    //cin >> tc;
    while (tc--) solve();
    return 0;
}

```

## 2.6 sibling dp [54 lines] - 473736aa00

```

/*
Problem:
Divide a rooted tree into the minimum
number of groups such that
the total cost (edge weights) within
each group does not exceed `k`.

```

**Definitions:**

- $dp[u][remk]$ : Minimum number of groups required to cover subtree rooted at node  $u$ , where  $remk$  is the remaining cost capacity of the current group.
- $adj[par][idx]$ : The  $idx$ -th child of node  $par$ , along with the edge cost.

```

ll n, k;
ll dp[mx][mx];
vector<pair<ll, ll>> adj[mx]; // adj[u] = {v, cost} edges in rooted tree
// Recursive function to compute
dp[par][remk]
// par = current parent node
// idx = index of child being processed
// remk = remaining budget in the current group
ll sibling_dp(ll par, ll idx, ll remk) {
    if (remk < 0) return inf; // invalid case, over budget
}

```

```
// No more children to process
if (adj[par].size() <= idx) return 0;

ll u = adj[par][idx].first; // current
child node

// Memoization check
if (dp[u][remk] != -1)
    return dp[u][remk];

ll ret = inf;
ll under = 0, sibling = 0;

// Option 1: Create a new group for
// this child (if not the root)
if (par != 0) {
    under = 1 + dfs(u, 0, k);
    // Entire subtree of `u` in a new
    group
    sibling = dfs(par, idx + 1, remk);
    // Process next siblings with
    same group
    ret = min(ret, under + sibling);
    // Total groups = groups under +
    groups from siblings
}

// Option 2: Try including this child
// in current group if cost allows
ll temp = remk - adj[par][idx].second;
// Remaining cost after adding this
edge

// Try all possible cost splits between
// `under` (child) and `sibling`
// (rest)
for (ll chk = temp; chk >= 0; chk--) {
    ll siblingk = temp - chk;
    under = dfs(u, 0, chk);
    // Child's subtree with budget
    // `chk`
    sibling = dfs(par, idx + 1,
    siblingk); // Remaining siblings
    // with remaining budget
    ret = min(ret, under + sibling);
}

return dp[u][remk] = ret; // Memoize
and return result
}
```

## 2.7 sos dp [19 lines] - e6cfca2350

```
// SUBSET SOS
// dp[mask] = sum of dp[submask] where
// submask ⊂ mask
void sos_subset(int M, vector<ll>& dp) {
    int N = 1 << M;
    for(int i = 0; i < M; i++)
        for(int mask = 0; mask < N;
            mask++)
            if(mask & (1 << i))
                dp[mask] += dp[mask ^ (1
                    << i)];
}
```

```
// SUPERSET SOS
// dp[mask] = sum of dp[supermask] where
// supermask ⊇ mask
void sos_superset(int M, vector<ll>& dp)
{
    int N = 1 << M;
    for(int i = 0; i < M; i++)
        for(int mask = N-1; mask >= 0;
            mask--)
            if(mask & (1 << i))
                dp[mask ^ (1 << i)] +=
                    dp[mask];
}
```

## 2.8 sum of dis of u [38 lines] - 6facf7b133

```
void solve() {
    int n;
    cin >> n;
    vector<vector<int>> g(n + 1);
    for(int i = 1; i < n; i++) {
        int u, v;
        cin >> u >> v;
        g[u].push_back(v);
        g[v].push_back(u);
    }
    vector<int> sz(n + 1);
    vector<ll> dp(n + 1);

    function<void(int, int)> dfs1 =
        [&](int u, int p) -> void {
        sz[u] = 1;
        for(auto v : g[u]) {
            if(v != p) {
                dfs1(v, u);
                sz[u] += sz[v];
                dp[1] += sz[v]; // root
                is 1
            }
        }
    };
    dfs1(1, 0);

    function<void(int, int)> dfs2 =
        [&](int u, int p) -> void {
        for(auto v : g[u]) if(v != p) {
            dp[v] = dp[u] - sz[v] + (n -
                sz[v]);
            dfs2(v, u);
        }
    };
    dfs2(1, 0);

    for(int i = 1; i <= n; i++) {
        cout << dp[i] << " \n"[i == n];
    }
}
```

## 3 game theory

### 3.1 note [14 lines] - 9730900ce0

```
>[First Write a Brute Force solution]
>Nim = all xor
>Misere Nim = Nim + corner case: if all
piles are 1, reverse(nim)
>Bogus Nim = Nim
```

>Staircase Nim = Odd indexed pile Nim  
(Even indexed pile doesn't matter, as one player can give bogus moves to drop all even piles to ground)

>Sprague Grundy: [Every impartial game under the normal play convention is equivalent to a one-heap game of nim]

Every tree = one nim pile = tree root value; tree leaf value = 0; tree node value = mex of all child nodes.

[Careful: one tree node can become multiple new tree roots (multiple elements in one node), then the value of that node = xor of all those root values]

>Hackenbush (Given a rooted tree; cut an edge in one move; subtree under that edge gets removed; last player to cut wins):

Colon: //G(u) =  
(G(v1) + 1) ⊕ (G(v2) + 1) ⊕ ... [v1, v2, ...  
are children of u]

For multiple trees ans is their xor

>Hackenbush on graph (instead of tree given an rooted graph):

fusion: All edges in a cycle can be fused to get a tree structure; build a super node, connect some single nodes with that super node, number of single nodes is the number of edges in the cycle.

Sol: [Bridge component tree] mark all bridges, a group of edges that are not bridges, becomes one component and contributes number of edges to the hackenbush. (even number of edges contributes 0, odd number of edges contributes 1)

## 4 strings

### 4.1 hash [50 lines] - 79b8bf4eea

```
struct SimpleHash {
    int len;
    long long base, mod;
    vector<int> P, H, R;
    SimpleHash() {}
    SimpleHash(string str, long long b,
    long long m) {
        base = b, mod = m, len =
            str.size();
        P.resize(len + 4, 1),
        H.resize(len + 3, 0),
        R.resize(len + 3, 0);
        for (int i = 1; i <= len + 3;
            i++)
            P[i] = (P[i - 1] * base) %
                mod;
        for (int i = 1; i <= len; i++)
            H[i] = (H[i - 1] * base +
                str[i - 1] + 1007) %
                mod;
        for (int i = len; i >= 1; i--)
```

```
R[i] = (R[i + 1] * base +
    str[i - 1] + 1007) %
    mod;
}
```

```
inline int range_hash(int l, int r) {
    int hashval = H[r + 1] - ((long
        long)P[r - 1 + 1] * H[l] %
        mod);
    return (hashval < 0 ? hashval +
        mod : hashval);
}
```

```
inline int reverse_hash(int l, int r) {
    int hashval = R[l + 1] - ((long
        long)P[r - 1 + 1] * R[r + 2] %
        mod);
    return (hashval < 0 ? hashval +
        mod : hashval);
}
```

```
}; struct DoubleHash {
    SimpleHash sh1, sh2;
    DoubleHash() {}
    DoubleHash(string str) {
        sh1 = SimpleHash(str,
            1949313259, 2091573227);
        sh2 = SimpleHash(str,
            1997293877, 2117566807);
    }
}
```

```
long long concate(DoubleHash& B, int
    l1, int r1, int l2, int r2) {
    int len1 = r1 - l1 + 1, len2 = r2 -
        l2 + 1;
    long long x1 =
        sh1.range_hash(l1, r1),
    x2 = B.sh1.range_hash(l2, r2),
    x1 = (x1 * B.sh1.P[len2]) %
        2091573227;
    long long newx1 = (x1 + x2) %
        2091573227;
    x1 = sh2.range_hash(l1, r1);
    x2 = B.sh2.range_hash(l2, r2),
    x1 = (x1 * B.sh2.P[len2]) %
        2117566807;
    long long newx2 = (x1 + x2) %
        2117566807;
    return (newx1 << 32) ^ newx2;
}
```

```
inline long long range_hash(int l,
    int r) {
    return ((long
        long)sh1.range_hash(l, r) <<
        32) ^ sh2.range_hash(l, r);
}
```

```
inline long long reverse_hash(int l,
    int r) {
    return ((long
        long)sh1.reverse_hash(l, r) <<
        32) ^ sh2.reverse_hash(l,
        r);
}
```

**4.2 kmp [13 lines] - 5a11b8b81b**

```
string t, p;
string s = p + "#" + t;
int n = s.size();
// pi[0...i] returns longest suffix which
// is prefix
vector<int> pi(n, 0);
for(int i = 1, j = 0; i < n; i++) {
    while(j >= 0 && s[i] != s[j]) {
        if(j >= 1) j = pi[j - 1];
        else j = -1;
    }
    j++;
    pi[i] = j;
}
```

**4.3 manachar [25 lines] - 934f373c7c**

```
struct Manacher {
    vector<int> p[2];
    // p[1][i] = (max odd length palindrome
    // centered at i) / 2 [floor division]
    // p[0][i] = same for even, it
    // considers the right center
    // e.g. for s = "abbabba", p[1][3] = 3,
    // p[0][2] = 2
    Manacher(string s) {
        int n = s.size();
        p[0].resize(n + 1);
        p[1].resize(n);
        for (int z = 0; z < 2; z++) {
            for (int i = 0, l = 0, r = 0; i <
                n; i++) {
                int t = r - i + !z;
                if (i < r) p[z][i] = min(t,
                    p[z][l + t]);
                int L = i - p[z][i], R = i +
                    p[z][i] - !z;
                while (L >= 1 && R + 1 < n && s[L -
                    1] == s[R + 1])
                    p[z][i]++;
                L--;
                R++;
                if (R > r) l = L, r = R;
            }
        }
        bool is_palindrome(int l, int r) {
            int mid = (l + r + 1) / 2, len = r -
                l + 1;
            return 2 * p[len % 2][mid] + len % 2
                >= len;
        }
    }
};
```

**4.4 palindromic tree [67 lines] - 8e5a153d34**

```
struct palindromic_tree {
    struct node {
        int len, link;
        int occ, occ1; // occ: total
        // in sts, occ1: only in first
        // n0
        array<int, 10> next;
        node(int l = 0) : len(l),
            link(0), occ(0), occ1(0) {
            next.fill(-1);
        }
    };
    string s;
```

```
vector<node> t;
int last;
int n0; // original
length (first-half cutoff)
```

```
palindromic_tree() {}
palindromic_tree(string _s, int _n0)
{
    s = _s; n0 = _n0;
    t.clear();
    t.reserve((int)s.size() + 3);
    t.push_back(node(-1));
    t.push_back(node(0));
    t[0].link = 0;
    t[1].link = 0;
    last = 1;
}
int getlink(int u, int i) {
    while (true) {
        int idx = i - 1 - t[u].len;
        if (idx >= 0 && s[idx] ==
            s[i]) return u;
        u = t[u].link;
    }
}
int add(int i) {
    int u = getlink(last, i);
    int ch = s[i] - '0';
    if (t[u].next[ch] != -1) {
        last = t[u].next[ch];
        t[last].occ++;
        if (i < n0) t[last].occ1++;
        // count only in first n0
        // for non-wrap
        return 0;
    }
    int llen = t[u].len + 2;
    node nu(llen);
    t.push_back(nu);
    int v = (int)t.size() - 1;
    t[u].next[ch] = v;
    if (t[v].len == 1) {
        t[v].link = 1;
    } else {
        int w = t[u].link;
        w = getlink(w, i);
        t[v].link = t[w].next[ch];
    }
    last = v;
    t[last].occ = 1;
    if (i < n0) t[last].occ1 = 1;
    return 1;
}
void pushOcc() {
    for (int v = (int)t.size() - 1; v
        >= 2; --v) {
        int p = t[v].link;
        t[p].occ += t[v].occ;
        t[p].occ1 += t[v].occ1;
    }
}
```

**4.5 suffix array [231 lines] - 1d9adcacccf**

```
struct SuffixArray {
```

```
string s;
int n;
vector<int> sa;
vector<int> lcp, lg;
vector<vector<int>> st;
```

```
// Constructor: builds SA, LCP, and
Sparse Table
```

```
SuffixArray(const string& str) :
s(str), n(str.size()) {
buildSA();
buildLCP();
buildSparse();
}
```

```
// Builds the suffix array using the
doubling method in O(n log n)
```

```
void buildSA() {
sa.resize(n);
vector<int> rank(n), tmp(n);
iota(sa.begin(), sa.end(), 0);
for (int i = 0; i < n; i++)
rank[i] = s[i];
```

```
for (int k = 1; k < n; k *= 2) {
auto cmp = [&](int i, int j)
```

```
{
if (rank[i] != rank[j])
return rank[i] <
rank[j];
int ri = (i + k < n) ?
rank[i + k] : -1;
int rj = (j + k < n) ?
rank[j + k] : -1;
return ri < rj;
};
```

```
sort(sa.begin(), sa.end(),
cmp);
```

```
tmp[sa[0]] = 0;
for (int i = 1; i < n; i++)
tmp[sa[i]] = tmp[sa[i - 1]] + (cmp(sa[i - 1], sa[i]) ? 1 : 0);
rank = tmp;
}
```

```
// Builds the LCP (Longest Common
Prefix) array in O(n)
```

```
void buildLCP() {
lcp.resize(n);
vector<int> inv(n);
for (int i = 0; i < n; i++)
inv[sa[i]] = i;
```

```
int k = 0;
lcp[n - 1] = 0;
for (int i = 0; i < n; i++) {
if (inv[i] == n - 1) {
k = 0;
continue;
}}
```

```
int j = sa[inv[i] + 1];
while (i + k < n && j + k < n
&& s[i + k] == s[j + k])
k++;
```

```
lcp[inv[i]] = k;
if (k) k--;
}

// Builds sparse table over sa for
rangeMin queries
void buildSparse() {
lg.resize(n + 1); lg[0] = -1;
for (int i = 1; i <= n; ++i)
lg[i] = lg[i >> 1] + 1;

int K = lg[n] + 1;
st.assign(K, vector<int>(n));
st[0] = sa;
for (int k = 1; k < K; ++k)
for (int i = 0; i + (1 << k) -
n; ++i)
st[k][i] = min(st[k - 1][i],
st[k - 1][i + (1 << (k - 1))]);
}

// Range minimum over sa[l..r] using
sparse table
int rangeMin(int l, int r) {
int k = lg[r - 1 + 1];
return min(st[k][l], st[k][r - (1
<< k) + 1]);
}

// Binary search for pattern range in
suffix array
pair<int, int> rangeSearch(const
string& pat, int &L, int &R, int
pos) {
//cout << L << " " << R << "\n";
int lo = L, hi = R, mid, ans =
-1;

while(lo <= hi) {
mid = (lo + hi) >> 1;
if(sa[mid] + pos >= n ||

s[sa[mid] + pos] <
pat[pos]) {
lo = mid + 1;
}
else {
ans = mid;
hi = mid - 1;
}
}

if(ans == -1) return {-1, -1};

int left = ans;
lo = ans, hi = R, ans = -1;
while(lo <= hi) {
mid = (lo + hi) >> 1;
if(sa[mid] + pos >= n ||
s[sa[mid] + pos] <=
pat[pos]) {
ans = mid;
lo = mid + 1;
}
else {
ans = mid;
hi = mid - 1;
}
}

if(ans == -1) return {-1, -1};
```

```
int left = ans;
lo = ans, hi = R, ans = -1;
while(lo <= hi) {
mid = (lo + hi) >> 1;
if(sa[mid] + pos >= n ||
s[sa[mid] + pos] <=
pat[pos]) {
ans = mid;
lo = mid + 1;
}
else {
ans = mid;
hi = mid - 1;
}
}
```

```

        }
    else {
        hi = mid - 1;
    }
}
if(ans == -1) return {-1, -1};
int right = ans;
return {left, right};
}

11 query(const string &p, int &L, int
&R, int pos) {
    auto [l, r] = rangeSearch(p, L,
    R, pos);
    if(l == -1) return -1;
    L = l, R = r;
    //cout << p << " " << l << " " << r <<
    "\n";
    return (r - l + 1);
}

11 occurrence_as_substring(const
    string &t) {
    string p;
    int ans = 0;
    int l = 0, r = n - 1;
    for(int i = 0; i < t.size(); i++) {
        p.push_back(t[i]);
        ans = query(p, l, r, i);
        if(ans == -1) break;
    }
    if(ans == -1) ans = 0;
    return ans;
}

// Returns the number of distinct
// substrings of each length
vector<int> countDistinctByLength() {
    vector<int> diff(n + 2, 0);
    for (int i = 0; i < n; i++) {
        int len = n - sa[i];
        int lcp_prev = (i == 0 ? 0 :
        lcp[i - 1]);
        int low = lcp_prev + 1;
        int high = len;
        if (low <= high) {
            diff[low]++;
            diff[high + 1]--;
        }
    }
    vector<int> result(n + 1);
    for (int i = 1; i <= n; i++)
        result[i] = result[i - 1] +
        diff[i];
    return result;
}

// Returns the k-th lexicographically
// smallest distinct substring
string getKthDistinctSubstring(long
long k) {
    for (int i = 0; i < n; i++) {
        int start = sa[i];
        int lcp_prev = (i == 0 ? 0 :
        lcp[i - 1]);
        int total = n - start -
        lcp_prev;
        if (k <= total)
            return s.substr(start,
            lcp_prev + k);
        k -= total;
    }
    return "";
}

// Returns the longest substring that
// appears more than once
string longestRepeatingSubstring() {
    int max_len = 0, index = -1;
    for (int i = 1; i < n; i++) {
        if (lcp[i] > max_len) {
            max_len = lcp[i];
            index = sa[i];
        }
    }
    if (max_len == 0) return "-1";
    return s.substr(index, max_len);
}

// Returns the longest common
// substring between s and another
// string s2
pair<int, string>
longestCommonSubstring(const
    string& s2) {
    string combined = s + '#' + s2;
    int len1 = s.size();
    SuffixArray combinedSA(combined);
    int max_len = 0, index = -1;
    for (int i = 1; i <
    combinedSA.n; i++) {
        int a = combinedSA.sa[i], b
        = combinedSA.sa[i - 1];
        if ((a < len1) != (b < len1))
        {
            if (combinedSA.lcp[i - 1]
            > max_len) {
                max_len =
                combinedSA.lcp[i
                - 1];
                index =
                combinedSA.sa[i];
            }
        }
    }
    if (max_len == 0) return {0, ""};
    return {max_len,
    combined.substr(index,
    max_len)};
}

// Returns the longest palindromic
// substring in s
string longestPalindromicSubstring()
{
    string rev = s;
    reverse(rev.begin(), rev.end());
    string joined = s + '#' + rev;
    SuffixArray pal(joined);
}

```

```

    int total = n - start -
    lcp_prev;
    if (k <= total)
        return s.substr(start,
        lcp_prev + k);
    k -= total;
}
return "";

// Returns the longest substring that
// appears more than once
string longestRepeatingSubstring() {
    int max_len = 0, index = -1;
    for (int i = 1; i < n; i++) {
        if (lcp[i] > max_len) {
            max_len = lcp[i];
            index = sa[i];
        }
    }
    if (max_len == 0) return "-1";
    return s.substr(index, max_len);
}

// Returns the longest common
// substring between s and another
// string s2
pair<int, string>
longestCommonSubstring(const
    string& s2) {
    string combined = s + '#' + s2;
    int len1 = s.size();
    SuffixArray combinedSA(combined);
    int max_len = 0, index = -1;
    for (int i = 1; i <
    combinedSA.n; i++) {
        int a = combinedSA.sa[i], b
        = combinedSA.sa[i - 1];
        if ((a < len1) != (b < len1))
        {
            if (combinedSA.lcp[i - 1]
            > max_len) {
                max_len =
                combinedSA.lcp[i
                - 1];
                index =
                combinedSA.sa[i];
            }
        }
    }
    if (max_len == 0) return {0, ""};
    return {max_len,
    combined.substr(index,
    max_len)};
}

// Returns the longest palindromic
// substring in s
string longestPalindromicSubstring()
{
    string rev = s;
    reverse(rev.begin(), rev.end());
    string joined = s + '#' + rev;
    SuffixArray pal(joined);
}

```

```

    int max_len = 0, index = -1;
    for (int i = 1; i < pal.n; i++) {
        int a = pal.sa[i], b =
        pal.sa[i - 1];
        bool in_s = (a < n), in_rev
        = (b > n);
        bool in_s2 = (b < n), in_rev2
        = (a > n);
        if ((in_s && in_rev) ||
        (in_s2 && in_rev2)) {
            int lcp_len = pal.lcp[i
            - 1];
            int pos = in_s ? a : b;
            if (lcp_len > max_len &&
            pos + lcp_len <= n) {
                max_len = lcp_len;
                index = pos;
            }
        }
    }
    if (max_len == 0) return "";
    return s.substr(index, max_len);
}

// Lexicographically compares
// s[l1..r1] and s[l2..r2]
int compareSubstrings(int l1, int
r1, int l2, int r2) {
    int len1 = r1 - l1 + 1, len2 = r2
    - l2 + 1;
    int min_len = min(len1, len2);
    for (int i = 0; i < min_len; i++)
    {
        if (s[l1 + i] != s[l2 + i])
            return (s[l1 + i] < s[l2
            + i]) ? -1 : 1;
    }
    if (len1 == len2) return 0;
    return (len1 < len2) ? -1 : 1;
}

// Finds the Lowest Common Ancestor
// (LCA) of x and y in the alternating
// tree
auto lca = [&](int x, int y) {
    timer++; // Increase visit timestamp
    for (; ; swap(x, y)) {
        if (x == -1) continue;
        if (aux[x] == timer) return x; // Found LCA
        aux[x] = timer;
        x = (mate[x] == -1 ? -1 :
        orig[parent[mate[x]]]); // Move up the alternating tree
    }
};

// Contracts a blossom (odd-length
// cycle) found during BFS
auto blossom = [&](int v, int w, int a)
{
    // v and w are two endpoints; a is
    // the LCA
    while (orig[v] != a) {
        parent[v] = w;
        w = mate[v];
        if (label[w] == 1) {
            label[w] = 0;
            q.push_back(w); // Re-add to BFS
            queue
        }
    }
    orig[v] = orig[w] = a; // Update
    base of blossom
}

```

## 4.6 z [14 lines] - 3133135f30

```

// z[i] returns the longest substring of
// length k from i which is also a
// prefix
function<vector<int>()> Zf = [&]()
{
    vector<int> z(n);
    for (int i = 1, l = 0, r = 0; i < n;
    ++i) {
        if (i <= r) {
            z[i] = min (r - i + 1, z[i -
            1]);
        }
        while (i + z[i] < n && s[z[i]] ==
        s[i + z[i]]) ++z[i];
        if (i + z[i] - 1 > r) {
            l = i, r = i + z[i] - 1;
        }
    }
    return z;
}

```

**5 flow**

**5.1 Blossom [86 lines] - 1c51297a78**

---

```

// Finds Maximum Matching in a General
Graph using Edmonds' Blossom
Algorithm
// Returns: vector<int> mate where
mate[i] = j means node i is matched
with node j
// Complexity: O(N * M), where N = number
of nodes, M = number of edges
vector<int> Blossom(vector<vector<int>>&
graph) {
    int n = graph.size(); // Number of nodes
    int timer = -1; // Used
    for tracking during LCA search
    // mate[i] = j means i is matched with
    j (or -1 if unmatched)
    vector<int> mate(n, -1);
    vector<int> label(n); // Labels: -1 = unused, 0 = even
    level, 1 = odd level
    vector<int> parent(n); // BFS
    tree parent
    vector<int> orig(n); // Original base of blossom for each
    node
    vector<int> aux(n, -1); // Aux
    array used in LCA finding
    vector<int> q; // BFS
    queue
    // Finds the Lowest Common Ancestor
    (LCA) of x and y in the alternating
    tree
    auto lca = [&](int x, int y) {
        timer++; // Increase visit timestamp
        for (; ; swap(x, y)) {
            if (x == -1) continue;
            if (aux[x] == timer) return x; // Found LCA
            aux[x] = timer;
            x = (mate[x] == -1 ? -1 :
            orig[parent[mate[x]]]); // Move up the alternating tree
        }
    };
    // Contracts a blossom (odd-length
    cycle) found during BFS
    auto blossom = [&](int v, int w, int a)
    {
        // v and w are two endpoints; a is
        the LCA
        while (orig[v] != a) {
            parent[v] = w;
            w = mate[v];
            if (label[w] == 1) {
                label[w] = 0;
                q.push_back(w); // Re-add to BFS
                queue
            }
        }
        orig[v] = orig[w] = a; // Update
        base of blossom
    };

```

```

        v = parent[w];
    }

    // Augments along the alternating path
    // ending at v to increase the
    // matching
    auto augment = [&](int v) {
        while (v != -1) {
            int pv = parent[v], nv = mate[pv];
            mate[v] = pv;           // Match v with
            // its parent
            mate[pv] = v;
            v = nv;                // Continue on
            // the unmatched end
        }

        // Breadth-First Search to find an
        // augmenting path starting from root
        auto bfs = [&](int root) {
            fill(label.begin(), label.end(),
                 -1); // Reset labels
            iota(orig.begin(), orig.end(), 0);
            // Each node is its own blossom
            // initially
            q.clear();
            // Reset BFS queue

            label[root] = 0;
            // Start root at even level
            q.push_back(root);

            for (int i = 0; i < (int)q.size();
                 ++i) {
                int v = q[i];
                for (auto x : graph[v]) {
                    if (label[x] == -1) {
                        // If x is unvisited
                        label[x] = 1;           // Mark as
                        // odd level
                        parent[x] = v;
                        if (mate[x] == -1) // Found an
                            // augmenting path
                            return augment(x), 1;
                        label[mate[x]] = 0; // Next
                        // level (even)
                        q.push_back(mate[x]);
                    }
                    else if (label[x] == 0 && orig[v]
                           != orig[x]) {
                        // Found a blossom (x and v
                        // have the same label and
                        // different bases)
                        int a = lca(orig[v], orig[x]);
                        blossom(x, v, a);
                        blossom(v, x, a);
                    }
                }
            }
            return 0; // No augmenting path
                      // found
        };

        // Try to find augmenting paths
        // starting from unmatched nodes
    };

```

**5.2 Hopcroft [69 lines] - 1867a462ee**

---

```

struct HopcroftKarp {
    vector<vector<int>> g;
    vector<int> pairU, pairV, dis;
    int n, m;
}

HopcroftKarp(int n, int m) : n(n),
    m(m) {
    g.resize(n + 1);
    dis.resize(n + 1);
    pairU.assign(n + 1, -1);
    pairV.assign(m + 1, -1);
}

void add_edge(int u, int v) {
    g[u].push_back(v);
}

bool bfs() {
    queue<int> Q;
    for(int u = 1; u <= n; u++) {
        if(pairU[u] == -1) {
            dis[u] = 0;
            Q.push(u);
        }
        else {
            dis[u] = -1;
        }
    }
    bool flag = false;
    while(not Q.empty()) {
        int u = Q.front(); Q.pop();
        for(auto v : g[u]) {
            int u2 = pairV[v];
            if(u2 == -1) {
                flag = true;
            }
            else if(dis[u2] == -1) {
                dis[u2] = dis[u] + 1;
                Q.push(u2);
            }
        }
    }
    return flag;
}

bool dfs(int u) {
    for(auto v : g[u]) {
        int u2 = pairV[v];
        if(u2 == -1 || (dis[u2] ==
                         dis[u] + 1 && dfs(u2))) {
            pairU[u] = v;
            pairV[v] = u;
            return true;
        }
    }
    dis[u] = -1;
    return false;
}

int run() {
    int match = 0;
    while(bfs()) {

```

---

```

        for(int u = 1; u <= n; u++) {
            if(pairU[u] == -1 &&
                dfs(u)) {
                match++;
            }
        }
    }
    return match;
}

5.3 MCMF [150 lines] - a17481cbb4
/* Minimum Cost Maximum Flow (MCMF)
 * Works for both directed and undirected
 * graphs, and supports negative edge
 * costs.
 * Does NOT work with negative cost
 * cycles.
 * For undirected edges, set `directed =
 * false`.
 * Complexity: O(min(E^2 * V log V, E log
 * V * flow))
 * Usage:
 *   MCMF mcmf(n); // Initialize with n nodes (0-indexed)
 *   mcmf.add_edge(u, v, cap, cost); // Add edge from u to v with capacity
 *   and cost
 *   auto [maxFlow, minCost] =
 *   mcmf.solve(source, sink);
 */

using T = long long;
const T inf = 1LL << 61;

struct MCMF {
    struct edge {
        int u, v;
        T cap, cost;
        int id;
        edge(int _u, int _v, T _cap, T
              _cost, int _id)
            : u(_u), v(_v), cap(_cap),
              cost(_cost), id(_id) {};
    };

    int n, s, t, mxid;
    T flow, cost;
    bool neg;

    vector<vector<int>> g; // Adjacency
    list: g[u] = indices of edges
    vector<edge> e; // Edge list
    vector<T> d, potential; // d =
    distances in Dijkstra; potential =
    reduced costs
    vector<T> flow_through; // Stores
    flow pushed through original edges
    vector<int> par; // Parent
    edge index in path

    // Default constructor
    MCMF() {}

```

---

```

    // Initialize graph with n nodes
    // (0-based indexing)
    MCMF(int _n) {
        n = _n + 10;
        g.assign(n, vector<int>());
        neg = false;
        mxid = 0;
    }

    // Adds an edge from u to v with given
    // capacity and cost
    // If directed = false, adds edge in
    // both directions
    void add_edge(int u, int v, T cap, T
                  cost, int id = -1, bool directed =
                  true) {
        if (cost < 0) neg = true;
        g[u].push_back(e.size());
        e.emplace_back(u, v, cap, cost, id);
        g[v].push_back(e.size());
        e.emplace_back(v, u, 0, -cost, -1);
        // Reverse edge
        mxid = max(mxid, id);
        if (!directed) add_edge(v, u, cap,
                               cost, -1, true); // Add reverse
                               // only once
    }

    // Dijkstra with potentials (Johnson's
    // algorithm) to find shortest path
    bool dijkstra() {
        par.assign(n, -1);
        d.assign(n, inf);
        priority_queue<pair<T, T>,
                      vector<pair<T, T>>, greater<>>
        q;

        d[s] = 0;
        q.emplace(0, s);

        while (!q.empty()) {
            int u = q.top().second;
            T nw = q.top().first;
            q.pop();
            if (nw != d[u]) continue;

            for (int i : g[u]) {
                int v = e[i].v;
                T cap = e[i].cap;
                T w = e[i].cost + potential[u] -
                potential[v]; // Reduced cost
                if (cap > 0 && d[u] + w < d[v]) {
                    d[v] = d[u] + w;
                    par[v] = i;
                    q.emplace(d[v], v);
                }
            }
        }

        for (int i = 0; i < n; i++) {
            if (d[i] < inf) potential[i] += d[i]; // Update potentials
        }
    }

```

```

        return d[t] != inf; // True if sink
    }

    // DFS-style recursive flow sending
    // function
    T send_flow(int v, T cur) {
        if (par[v] == -1) return cur;
        int id = par[v];
        int u = e[id].u;
        T w = e[id].cost;
        T f = send_flow(u, min(cur,
            e[id].cap));
        cost += f * w;
        e[id].cap -= f;
        e[id ^ 1].cap += f; // Update
        // reverse edge
        return f;
    }

    // Main function to compute minimum
    // cost max flow from s to t
    // Optionally limit total flow with
    // `goal`
    // Returns {maximum flow, minimum cost}
    pair<T, T> solve(int _s, int _t, T goal
        = inf) {
        s = _s; t = _t;
        flow = 0; cost = 0;
        potential.assign(n, 0);

        // If negative costs exist, use
        // Bellman-Ford to initialize
        // potentials
        if (neg) {
            d.assign(n, inf);
            d[s] = 0;
            for (int i = 0, relax = true; i < n
                && relax; ++i) {
                relax = false;
                for (int u = 0; u < n; ++u) {
                    if (d[u] == inf) continue;
                    for (int k : g[u]) {
                        int v = e[k].v;
                        T cap = e[k].cap, w =
                            e[k].cost;
                        if (cap > 0 && d[v] > d[u] +
                            w) {
                            d[v] = d[u] + w;
                            relax = true;
                        }
                    }
                }
                for (int i = 0; i < n; ++i) if
                    (d[i] < inf) potential[i] =
                        d[i];
            }
        }

        // Repeatedly find augmenting paths
        // using Dijkstra + potential
        while (flow < goal && dijkstra()) {
            flow += send_flow(t, goal - flow);
        }
    }
}

// Track flow through original edges
// (by ID)
flow_through.assign(mxid + 10, 0);
for (int u = 0; u < n; ++u) {
    for (int idx : g[u]) {
        if (e[idx].id >= 0) {
            flow_through[e[idx].id] = e[idx
                ^ 1].cap;
        }
    }
}

return {flow, cost};
}

5.4 dinic [91 lines] - 1ea72f97f9
// Dinic's Algorithm for Maximum Flow
// Complexity: O(V^2 * E) in general
// graphs
// Usage:
// 1. Initialize with Dinic
//    dinic(num_nodes);
// 2. Add edges with addEdge(u, v,
//    capacity);
// 3. Compute maximum flow with
//    maxFlow(source, sink);

#define eb emplace_back

struct Dinic {
    // Structure to represent an edge in
    // the flow graph
    struct Edge {
        int u, v; // from node u to
        // node v
        ll cap, flow = 0; // capacity and
        // current flow

        Edge() {}
        Edge(int u, int v, ll cap) : u(u),
            v(v), cap(cap) {}
    };

    int N; // Number of nodes in the graph
    vector<Edge> edge; // List
    // of all edges
    vector<vector<int>> adj; // Adjacency
    // list of edge indices
    vector<int> d, pt; // Distance (level) and pointer arrays

    // Constructor: Initializes flow
    // network for N nodes
    Dinic(int N) : N(N), edge(0), adj(N),
        d(N), pt(N) {}

    // Adds a directed edge from u to v
    // with given capacity
    // Adds reverse edge with 0 capacity
    // for residual graph
    void addEdge(int u, int v, ll cap) {
        if (u == v) return; // No self-loops
        edge.eb(u, v, cap);
        adj[u].eb(edge.size() - 1);
        edge.eb(v, u, 0); // Reverse edge
    }
}

```

```

    adj[v].eb(edge.size() - 1);
}

// Constructs level graph using BFS
// from source s
// Returns true if sink t is reachable
// from s
bool bfs(int s, int t) {
    queue<int> q({s});
    fill(pt.begin(), pt.end(), 0); // Reset
    // current edge pointers
    while (ll flow = dfs(s, t)) { // Push as
        // much flow as possible
        total += flow;
    }
    return total;
}

5.5 flow [6 lines] - 071f5a23a9
Covering Problems:
> Maximum Independent Set(Bipartite):
    Largest set of nodes which do not
    have any edge between them. sol:
    V-(MaxMatching)
> Minimum Vertex Cover(Bipartite):
    -Smallest set of nodes to cover all
    the edges -sol: MaxMatching
> Minimum Edge Cover(General graph):
    -Smallest set of edges to cover all
    the nodes -sol: V-(MaxMatching) (if
    edge cover exists, does not exit for
    isolated nodes)
> Minimum Path Cover(Vertex disjoint)
    DAG: -Minimum number of vertex
    disjoint paths that visit all the
    nodes -sol: make a bipartite graph
    using same nodes in two sides, one
    side is "from" other is "to", add
    edges from "from" to "to", then ans
    is V-(MaxMatching)
> Minimum Path Cover(Vertex Not Disjoint)
    General graph: -Minimum number of
    paths that visit all the nodes -sol:
    consider cycles as nodes then it will
    become a path cover problem with
    vertex disjoint on DAG

5.6 hungarian [60 lines] - 699d538345
struct Hungarian {
    vector<vector<int>> cost;
    vector<int> worker, job, parentjob,
        matchjob, assign;
    int n;

    Hungarian(int n) : n(n) {
        cost.resize(n, vector<int>(n));
        worker.resize(n + 1);
        job.resize(n + 1);
        parentjob.resize(n + 1);
        matchjob.resize(n + 1);
    }

    int solve() {
        for(int w = 1; w <= n; w++) {
            matchjob[0] = w;
            vector<int> minslack(n + 1,
                INT_MAX);
            vector<bool> used(n + 1, 0);
            int job0 = 0;

```

```

            adj[v].eb(edge.size() - 1);
}

// Constructs level graph using BFS
// from source s
// Returns true if sink t is reachable
// from s
bool bfs(int s, int t) {
    queue<int> q({s});
    fill(pt.begin(), pt.end(), 0); // Reset
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        total += flow;
    }
    return total;
}

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        worker.resize(n + 1);
        job.resize(n + 1);
        parentjob.resize(n + 1);
        matchjob.resize(n + 1);
    }

    int solve() {
        for(int w = 1; w <= n; w++) {
            matchjob[0] = w;
            vector<int> minslack(n + 1,
                INT_MAX);
            vector<bool> used(n + 1, 0);
            int job0 = 0;

```

```

            adj[v].eb(edge.size() - 1);
}

// Constructs level graph using BFS
// from source s
// Returns true if sink t is reachable
// from s
bool bfs(int s, int t) {
    queue<int> q({s});
    fill(pt.begin(), pt.end(), 0); // Reset
    // current edge pointers
    while (ll flow = dfs(s, t)) { // Push as
        // much flow as possible
        total += flow;
    }
    return total;
}

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    int n;

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        cost.resize(n, vector<int>(n));
        worker.resize(n + 1);
        job.resize(n + 1);
        parentjob.resize(n + 1);
        matchjob.resize(n + 1);
    }

    int solve() {
        for(int w = 1; w <= n; w++) {
            matchjob[0] = w;
            vector<int> minslack(n + 1,
                INT_MAX);
            vector<bool> used(n + 1, 0);
            int job0 = 0;

```

```

do {
    used[job0] = true;
    int worker0 =
        matchjob[job0], delta
        = INT_MAX, job1 = 0;
    for(int j = 1; j <= n;
        j++) {
        if(!used[j]) {
            int curcost =
                cost[worker0
                    - 1][j - 1]
                -
                worker[worker0
                    - job[j]];
            if(curocost <
                minslack[j]) {
                minslack[j]
                    =
                    curcost;
                parentjob[j]
                    = job0;
            }
            if(minslack[j] <
                delta) {
                delta =
                    minslack[j];
            }
            job1 = j;
        }
    }
    for(int j = 0; j <= n;
        j++) {
        if(used[j]) {
            worker[matchjob[j]] +=
                delta;
            job[j] -= delta;
        }
        else {
            minslack[j] -=
                delta;
        }
    }
    job0 = job1;
} while(matchjob[job0] != 0);

do {
    int job1 =
        parentjob[job0];
    matchjob[job0] =
        matchjob[job1];
    job0 = job1;
} while(job0);

assign.assign(n + 1, 0);
for(int j = 1; j <= n; j++) {
    assign[matchjob[j]] = j;
}

return -job[0];
}

```

## 6 geometry

### 6.1 basic geo [101 lines] - 187dfc5049

```

// basic geometry
// be careful! [th in radians]
Chord Length of a Circle
L = 2r * sin(th / 2)
Arc Length of a Circle
L = r * th
Sector Area of a Circle
A = (1/2) * r^2 * th
Area of Segment = Area of Sector - Area
of Triangle
A = (1/2) * r^2 * (th - sin(th))
Length of a Common Chord of Two Circles
L = 2 * sqrt( R^2 - ((d^2 + R^2 - r^2) /
(2d))^2 )

kissing circles
The plus sign gives the curvature of the
inner circle tangent to the three.
The minus sign gives the curvature of the
outer circle enclosing the three.
k4 = k1 + k2 + k3 ± 2 * sqrt(k1*k2 +
k2*k3 + k3*k1)
// where ki = 1 / ri (curvature), ri =
radius of circle i

// Given medians ma, mb, mc
s = (ma + mb + mc) / 2
D = sqrt(s * (s - ma) * (s - mb) * (s -
mc))
A = (4 / 3) * D // Area of original
triangle

// common formula is Area
= (1/2) * perimeter * apothem

// General Polygon
Each_Interior_Angle = ((n - 2) * pi) / n
Number_of_Diagonals = n * (n - 3) / 2

// Regular Polygon
// a = length of one side
// R = circumradius (radius of
circumscribed circle)
// r = apothem (inradius, perpendicular
from center to a side)
Side, a = 2 * R * sin(pi / n)
Apothem, r = a / (2 * tan(pi / n))
Circumradius = a / (2 * sin(pi / n))
Central_Angle = 360 / n
Area_from_side = (n * a^2) / (4 * tan(pi /
n))
Area_from_apothem = (1/2) * n * a * r
Area_from_R = (1/2) * n * R^2 * sin(2pi /
n)

// Volume of a spherical cap of height h
V_cap = (pi * h^2 * (3*R - h)) / 3
// Total volume of sphere
V_total = (4.0 / 3.0) * pi * R^3

// s = semi-perimeter = (a + b + c) / 2
// R = circumradius (radius of
circumcircle)

```

```

// r = inradius (radius of incircle)
// Ra, Rb, Rc = exradii opposite to a, b,
c respectively

```

```

// Incircle (circle tangent to all sides
inside triangle)
Inradius = r = A / s // onto
// Circumcircle (passes through all
triangle vertices)
Circumradius = R = (a * b * c) / (4 * A)
// pori
// Excircle (touches one side and
extensions of other two)
Exradius_Opposite_a = Ra = A / (s - a)

```

```

// Triangle Area Formulas
// Equilateral Triangle (all sides equal)
Area = (sqrt(3) / 4) * a^2
// Isosceles Triangle (two sides equal)
Area = (b / 4) * sqrt(4a^2 - b^2)

```

```

// 3D Volume Formulas
// Symbol Definitions:
// l = length, w = width, h = height
// a = edge length
// r = radius, d = diameter
// R = base radius (for cone/cylinder), H
= height
// Cylinder
Volume = pi * r^2 * h
// Cone
Volume = (1/3) * pi * r^2 * h
// Sphere
Volume = (4/3) * pi * r^3
// Hemisphere
Volume = (2/3) * pi * r^3
// Triangular Prism
Volume = (1/2) * b * h * l
// where b = triangle base, h = triangle
height, l = length/depth
// Pyramid (square or rectangular base)
Volume = (1/3) * base_area * height
// Tetrahedron (regular, edge = a)
Volume = (a^3) / (6 * sqrt(2))
// Frustum (of cone or pyramid)
Volume = (1/3) * pi * h * (R^2 + r^2 +
R*r)
// R = larger base radius, r = smaller
radius, h = vertical height
// Torus (donut shape)
Volume = 2 * pi^2 * R * r^2
// R = distance from center of tube to
center of torus
// r = radius of tube

```

### 6.2 custom geo [404 lines] - aa41681d08

```

#include <bits/stdc++.h>
using namespace std;
typedef long long ll;

```

```

typedef long double ld;
const int N = 1e7 + 9, mod = 1e9 + 7;
const ld inf = 1e100;
const ld eps = 1e-9;
const ld PI = acos((ld)-1.0);
int sign(ld x) { return (x > eps) - (x <
-eps); }
struct PT {
    ll x, y;
    PT() : x = 0, y = 0 {} 
    PT(ll x, ll y) : x(x), y(y) {}
    PT operator - (const PT&a) const {
        return PT(x - a.x, y - a.y); }
    PT operator * (const ll a) const {
        return PT(x * a, y * a); }
    friend PT operator * (const ll &a,
        const PT &b) { return PT(a *
        b.x, a * b.y); }
    PT operator / (const ll a) const {
        return PT(x / a, y / a); }
    bool operator == (PT a) const {
        return sign(a.x - x) == 0 &&
        sign(a.y - y) == 0; }
    bool operator != (PT a) const {
        return !(*this == a); }
    bool operator < (PT a) const { return
        sign(a.x - x) == 0 ? y < a.y : x
        < a.x; }
    bool operator > (PT a) const { return
        sign(a.x - x) == 0 ? y > a.y : x
        > a.x; }
};

istream &operator >> (istream &in, PT &p)
{ return in >> p.x >> p.y; }
ostream &operator << (ostream &out, PT
&p) { return out << "(" << p.x << ","
<< p.y << ")"; }

inline ll dot(PT a, PT b) { return a.x *
b.x + a.y * b.y; }
inline ll dist2(PT a, PT b) { return
dot(a - b, a - b); }
inline double dist(PT a, PT b) { return
sqrt(dot(a - b, a - b)); }
inline double cross(PT a, PT b) { return
a.x * b.y - a.y * b.x; }
inline double cross2(PT a, PT b, PT c) {
    return cross(b - a, c - a); }
inline int orientation(PT a, PT b, PT c) {
    return sign(cross(b - a, c - a)); }

// intersection point between segment ab
and segment cd assuming unique
intersection exists
bool seg_seg_intersection(PT a, PT b, PT
c, PT d, PT &ans) {
    ld oa = cross2(c, d, a), ob =
    cross2(c, d, b);
    ld oc = cross2(a, b, c), od =
    cross2(a, b, d);
    if (oa * ob < 0 && oc * od < 0){
        ans = (a * ob - b * oa) / (ob -
        oa);
        return 1;
    }
}

```

```

        else return 0;
    }
    // returns true if point p is on line
    // segment ab
    bool is_point_on_seg(PT a, PT b, PT p) {
        if (fabs(cross(p - b, a - b)) < eps)
            {
                if (p.x < min(a.x, b.x) - eps || p.x > max(a.x, b.x) + eps)
                    return false;
                if (p.y < min(a.y, b.y) - eps || p.y > max(a.y, b.y) + eps)
                    return false;
            }
        return true;
    }
    return false;
}

// intersection point between segment ab
// and segment cd assuming unique
// intersection exists
set<PT> seg_seg_intersection_inside(PT
    a, PT b, PT c, PT d) {
    PT ans;
    if (seg_seg_intersection(a, b, c, d,
        ans)) return {ans};
    set<PT> se;
    if (is_point_on_seg(c, d, a))
        se.insert(a);
    if (is_point_on_seg(c, d, b))
        se.insert(b);
    if (is_point_on_seg(a, b, c))
        se.insert(c);
    if (is_point_on_seg(a, b, d))
        se.insert(d);
    return se;
}

ld area(vector<PT> &p) {
    ld ans = 0; int n = p.size();
    for (int i = 0; i < n; i++) ans += cross(p[i], p[(i + 1) % n]);
    return fabs(ans) * 0.5;
}

ll area2(const vector<PT>& p) {
    ll ans = 0;
    int n = p.size();
    for (int i = 0; i < n; i++) {
        ans += cross(p[i], p[(i + 1) %
            n]);
    }
    return abs(ans); // not devided by 2
    for integers
}

bool is_point_on_polygon(vector<PT> &p,
    const PT& z) {
    int n = p.size();
    for (int i = 0; i < n; i++) {
        if (is_point_on_seg(p[i], p[(i +
            1) % n], z)) return 1;
    }
    return 0;
}

```

```

// returns 1e9 if the point is on the
// polygon
int winding_number(vector<PT> &p, const
    PT& z) { // O(n)
    if (is_point_on_polygon(p, z)) return
        1e9;
    int n = p.size(), ans = 0;
    for (int i = 0; i < n; ++i) {
        int j = (i + 1) % n;
        bool below = p[i].y < z.y;
        if (below != (p[j].y < z.y)) {
            auto orient = orientation(z,
                p[j], p[i]);
            if (orient == 0) return 0;
            if (below == (orient > 0))
                ans += below ? 1 : -1;
        }
    }
    return ans;
}

// -1 if strictly inside, 0 if on the
// polygon, 1 if strictly outside
int is_point_in_polygon(vector<PT> &p,
    const PT& z) { // O(n)
    int k = winding_number(p, z);
    return k == 1e9 ? 0 : k == 0 ? 1 :
        -1;
}

// -1 if strictly inside, 0 if on the
// polygon, 1 if strictly outside
// it must be strictly convex, otherwise
// make it strictly convex first
int is_point_in_convex(vector<PT> &p,
    const PT& x) { // O(log n)
    int n = p.size(); assert(n >= 3);
    int a = orientation(p[0], p[1], x), b
        = orientation(p[0], p[n - 1],
            x);
    if (a < 0 || b > 0) return 1;
    int l = 1, r = n - 1;
    while (l + 1 < r) {
        int mid = l + r >> 1;
        if (orientation(p[0], p[mid], x)
            >= 0) l = mid;
        else r = mid;
    }
    int k = orientation(p[l], p[r], x);
    if (k <= 0) return -k;
    if (l == 1 && a == 0) return 0;
    if (r == n - 1 && b == 0) return 0;
    return -1;
}

// Count boundary lattice points using
// GCD on each edge
ll boundary_lattice_points(const
    vector<PT>& p) {
    ll b = 0;
    int n = p.size();
    for (int i = 0; i < n; i++) {
        PT a = p[i], c = p[(i + 1) % n];
        b += __gcd(abs(a.x - c.x),
            abs(a.y - c.y));
    }
    return b;
}

```

```

return b;
}

// Computes (interior, boundary) lattice
// points using Pick's Theorem
pair<ll, ll> polygon_lattice_points(const
    vector<PT>& p) {
    // Pick's Theorem:
    // Area = I + B/2 - 1
    // I = the number of interior lattice
    // points
    // B = the number of boundary lattice
    // points
    ll A2 = area2(p); // Twice the area
    ll B = boundary_lattice_points(p);
    ll I = (A2 - B + 2) / 2;
    return {I, B};
}

// distance, p1, p2
pair<ll, pair<PT, PT>>
closest_pair_recursive(vector<PT>&
    pts, vector<PT>& tmp, int l, int r) {
    if (r - l <= 3) {
        pair<ll, pair<PT, PT>> res =
            {LLONG_MAX, {PT(), PT()}};
        for (int i = l; i < r; ++i) {
            for (int j = i + 1; j < r;
                ++j) {
                ll d = dist2(pts[i],
                    pts[j]);
                if (d < res.first) {
                    res = {d, {pts[i],
                        pts[j]}};
                }
            }
        }
        sort(pts.begin() + 1, pts.begin()
            + r, [](const PT& a, const
                PT& b) {
            return a.y < b.y;
        });
        return res;
    }

    int m = (l + r) / 2;
    ll midx = pts[m].x;

    auto left =
        closest_pair_recursive(pts, tmp,
            l, m);
    auto right =
        closest_pair_recursive(pts, tmp,
            m, r);
    auto res = min(left, right);

    merge(pts.begin() + 1, pts.begin() +
        m,
        pts.begin() + m, pts.begin() +
        r,
        tmp.begin() + 1, [](const PT&
            a, const PT& b) {
            return a.y < b.y;
        });
    return res;
}

// returns the boundary points of the
// convex hull.
vector<PT> ConvexHull(vector<PT>&p, int
    n) {
    int sz = 0;
    vector<PT> hull(n + n);
    sort(p.begin(), p.end());
    for (int i = 0; i < n; ++i) {
        while (sz > 1 and cross2(hull[sz -
            2], hull[sz - 1], p[i]) <
            0) --sz;
        hull[sz++] = p[i];
    }
    for (int i = n - 2, j = sz + 1; i >=
        0; --i) {
        while (sz >= j and cross2(hull[sz -
            2], hull[sz - 1], p[i]) <
            0) --sz;
        hull[sz++] = p[i];
    }
    hull.resize(sz - 1);
    return hull;
}

```

```

// keep only the corners that form a
// strictly-convex polygon
vector<PT> make_strict_convex(const
vector<PT>& P) {
int n = (int)P.size();
vector<PT> R;
for (int i = 0; i < n; ++i) {
    PT prv = P[(i + n - 1) % n], cur
        = P[i], nxt = P[(i + 1) %
n];
    if (orientation(prv, cur, nxt)
!= 0) R.push_back(cur);
}
return R;
}

// return true if a polygon strictly
// insider other
bool
polygon_polygon_intersect(vector<PT>
v1, vector<PT> v2) {
auto H1 = ConvexHull(v1, v1.size());
auto H2 = ConvexHull(v2, v2.size());
vector<PT> all = v1;
for(auto it : v2) {
    all.push_back(it);
}
auto H = ConvexHull(all, all.size());
bool ok = (H == H1 || H == H2);
return ok;
/*v1 = make_strict_convex(v1);
reverse(v1.begin(), v1.end());
v2 = make_strict_convex(v2);
reverse(v2.begin(), v2.end());
if (v1.size() < 3 || v2.size() < 3) {
    return false;
}
bool ok = true;
for (const auto &pt : v2)
    if (is_point_in_convex(v1, pt) !=
-1) { ok = false; break; }
if (!ok) {
    ok = true;
    for (const auto &pt : v1)
        if (is_point_in_convex(v2,
pt) != -1) { ok = false;
            break; }
}
return ok;*/
}

// Projection of point c onto line ab
PT project_from_point_to_line(PT a, PT
b, PT c) {
    return a + (b - a) * dot(c - a, b -
a) / dot(b - a, b - a);
}

// Reflection of point c across line ab
PT reflection_from_point_to_line(PT a, PT
b, PT c) {
}

```

```

PT p = project_from_point_to_line(a,
    b, c);
return p + (p - c);

// Minimum distance from point c to line
ab
ld dist_from_point_to_line(PT a, PT b, PT
c) {
    return fabs(cross(b - a, c - a)) /
        sqrtl(dist2(a, b));
}

// Minimum distance from point c to
// segment ab
ld dist_from_point_to_seg(PT a, PT b, PT
c) {
    ld r = dist2(a, b);
    if (r == 0) return dist(c, a);
    ld t = max(0.0L, min(1.0L, (ld)dot(c
        - a, b - a) / r));
    PT proj = a + (b - a) * t;
    return dist(c, proj);
}

// Minimum distance between two segments
ab and cd
ld dist_from_seg_to_seg(PT a, PT b, PT
c, PT d) {
    PT dummy;
    if (seg_seg_intersection(a, b, c, d,
        dummy)) return 0;
    return min({
        dist_from_point_to_seg(a, b, c),
        dist_from_point_to_seg(a, b, d),
        dist_from_point_to_seg(c, d, a),
        dist_from_point_to_seg(c, d, b)
    });
}

// Checks if a polygon is strictly convex
bool is_convex(const vector<PT>& p) {
    int n = p.size();
    bool has_pos = false, has_neg =
        false;
    for (int i = 0; i < n; i++) {
        ld z = cross2(p[i], p[(i + 1) %
n], p[(i + 2) % n]);
        has_pos |= (z > 0);
        has_neg |= (z < 0);
    }
    return !(has_pos && has_neg);
}

// Centroid (center of mass) of a polygon
PT centroid(const vector<PT>& p) {
    ld A = 0;
    PT c(0, 0);
    int n = p.size();
    for (int i = 0; i < n; i++) {
        ld cross_val = cross(p[i], p[(i +
1) % n]);
        A += cross_val;
        c.x += (p[i].x + p[(i + 1) %
n].x) * cross_val;
    }
    c.y += (p[i].y + p[(i + 1) %
n].y) * cross_val;
}

// Rotate a point counter-clockwise by
angle t (radians) around origin
PT rotateccw(PT a, ld t) {
    return PT(a.x * cos(t) - a.y *
        sin(t), a.x * sin(t) + a.y *
        cos(t));
}

// Rotate a point clockwise by angle t
(radians) around origin
PT rotatecw(PT a, ld t) {
    return PT(a.x * cos(t) + a.y *
        sin(t), -a.x * sin(t) + a.y *
        cos(t));
}

// 90-degree counter-clockwise rotation
PT rotateccw90(PT a) { return PT(-a.y,
    a.x); }

// 90-degree clockwise rotation
PT rotatecw90(PT a) { return PT(a.y,
    -a.x); }

// Check point in triangle (a, b, c)
// -1 = inside, 0 = on edge, 1 = outside
int is_point_in_triangle(PT a, PT b, PT
c, PT p) {
    if (sign(cross(b - a, c - a)) < 0)
        swap(b, c);
    int c1 = sign(cross(b - a, p - a));
    int c2 = sign(cross(c - b, p - b));
    int c3 = sign(cross(a - c, p - c));
    if (c1 < 0 || c2 < 0 || c3 < 0)
        return 1;
    if (c1 + c2 + c3 != 3) return 0;
    return -1;
}

```

### 6.3 my geo [77 lines] - 0869646f06

---

```

struct pt {
    double x, y;
    pt() {}
    pt(double _x, double _y) {
        x = _x;
        y = _y;
    }
    pt operator + (const pt &o) const {
        return pt(x + o.x, y + o.y); }
    pt operator - (const pt &o) const {
        return pt(x - o.x, y - o.y); }
    ld operator | (const pt &o) const {
        return x * o.x + y * o.y; }
    ld operator * (const pt &o) const {
        return x * o.y - y * o.x; }
    ld dis() { return hypotl(x, y); }
    bool operator < (const pt &o) const {
        return (x < o.x || (x == o.x && y
        < o.y)); }
    bool operator == (const pt &o) const
    {

```

```

        return (x == o.x && y == o.y);
    }
ld pt_to_seg_dis(pt p, pt a, pt b) {
    if(((p - a) | (b - a)) < 0) {
        return (p - a).dis();
    }
    else {
        if(((p - b) | (a - b)) < 0) {
            return (p - b).dis();
        }
        else {
            return fabs((p - b) * (b - a)) / (b - a).dis();
        }
    }
}
ld dis_bet_seg(pt a, pt b, pt c, pt d) {
ld dot1 = ((b - a) * (d - a)) * ((b - a) * (c - a));
ld dot2 = ((d - c) * (a - c)) * ((d - c) * (b - c));
if(dot1 <= 0 && dot2 <= 0) {
    cout << 0 << "\n";
    return;
}
ld ans = min(pt_to_seg_dis(c, a, b),
            pt_to_seg_dis(d, a, b));
ans = min(ans, pt_to_seg_dis(a, c,
                             d));
ans = min(ans, pt_to_seg_dis(b, c,
                             d));
return ans;
}
int orientation(pt a, pt b, pt c) {
ld val = (b - a) * (c - a);
if (val > 0) return 1;
if (val < 0) return -1;
return 0;
}
vector<pt> convex_hull(vector<pt> v) {
    sort(v.begin(), v.end());
    v.erase(unique(v.begin(), v.end()),
           v.end());
    int n = v.size();
    if (n <= 2) return v;
    vector<pt> lower, upper;
    for (int i = 0; i < n; i++) {
        int j = lower.size();
        while (j >= 2 &&
               orientation(lower[j-2],
                           lower.back(), v[i]) != 1) {
            lower.pop_back();
            j--;
        }
        lower.push_back(v[i]);
    }
    for (int i = n-1; i >= 0; i--) {
        int j = upper.size();
        while (j >= 2 &&
               orientation(upper[j-2],
                           upper.back(), v[i]) != 1) {
            upper.pop_back();
            j--;
        }
        upper.push_back(v[i]);
    }
}

```

```

    }
    lower.pop_back();
    upper.pop_back();
    lower.insert(lower.end(),
                 upper.begin(), upper.end());
    return lower;
}

7 graphs
7.1 articulations [22 lines] - 075d74c576
const int N=1e5+5;
vector<int> g[N];
int vis[N],dis[N],lo[N],isAP[N];
int timer;
vector<pair<int,int>> ans;
void dfs(int src, int par) {
    int child = 0;
    vis[src] = true;
    dis[src] = lo[src] = ++timer;
    for (auto v : g[src]) {
        if (!vis[v]) {
            child++;
            dfs(v,src);
            lo[src] = min(lo[src],
                          lo[v]);
            if (par != -1 && lo[v] >=
                dis[src]) isAP[src] =
                true;
            if(lo[v]>dis[src])
                ans.push_back({min(v, j
                                     src), max(v,src)} );
        }
        else if (v != par) lo[src] =
            min(lo[src], dis[v]);
    }
    if (par == -1 && child > 1) isAP[src] =
        true;
}

```

```

7.2 bellman full [82 lines] - 646df1f54a
// if reach can possible from 1 to n
vector<array<int, 3>> edges;
vector<vector<int>> g(n + 1);
edges.push_back({u, v, -w});
g[v].push_back(u);
vector<ll> dis(n + 1, inf);
dis[1] = 0;
// n - 1 relaxation
for(int c = 1; c < n; c++) {
    bool relax = false;
    for(auto [u, v, w] : edges) {
        if(dis[u] != inf && dis[u] + w <
           dis[v]) {
            dis[v] = dis[u] + w;
            relax = true;
        }
    }
    if(!relax) false;
}
vector<bool> bad(n);
for(auto [u, v, w] : edges) {
    if(dis[u] != inf && dis[u] + w <
       dis[v]) {
        bad[v] = 1;
    }
}

```

```

}
bool flag = 1;
queue<int> Q;
Q.push(n);
vector<bool> vis(n + 1);
while(not Q.empty()) {
    int u = Q.front(); Q.pop();
    if(bad[u]) {
        flag = 0;
        break;
    }
    if(vis[u]) continue;
    vis[u] = 1;
    for(auto v : g[u]) {
        Q.push(v);
    }
}
// negatative cycle print
edges.push_back({u, v, w});
vector<ll> dis(n + 1, inf);
vector<int> par(n + 1, -1);
int node = -1;
for(int i = 1; i <= n; i++) {
    if(dis[i] != inf) continue;
    dis[i] = 0;
    for(int c = 1; c <= n; c++) {
        bool relax = 1;
        for(auto [u, v, w] : edges) {
            if(dis[u] != inf && dis[u] +
               w < dis[v]) {
                dis[v] = dis[u] + w;
                par[v] = u;
                if(c == n) node = v;
                relax = 0;
            }
        }
        if(relax) break;
    }
    if(node != -1) break;
}
if(node == -1) {
    cout << "NO\n";
}
else {
    // Get a node in the cycle by walking
    // back n times
    for(int i = 0; i < n; i++) node =
        par[node];
    vector<int> path;
    int cur = node;
    do {
        path.push_back(cur);
        cur = par[cur];
    } while(cur != node);
    path.push_back(node);
    reverse(path.begin(), path.end());
    cout << "YES\n";
    for(int x : path) cout << x << " ";
    cout << "\n";
}

```

7.3 centroid [124 lines] - 6d12aab948

```

/* ----- Centroid-decomposition
with fast depth counting -----
*/
struct CentroidTree {
    int n;
    vector<vector<int>> adj, ctree;
    vector<int> subsize, cpar;
    vector<char> blocked;
    // faster than vector<bool>
    /* reusable frequency table */
    vector<int> cnt;
    // size = K+1, filled on demand
    vector<int> touched;
    // indices of cnt we modified
    CentroidTree(int nodes) : n(nodes) {
        adj.resize(n + 1);
        ctree.resize(n + 1);
        subsize.resize(n + 1);
        cpar.assign(n + 1, -1);
        blocked.assign(n + 1, 0);
    }
    void add_edge(int u, int v) {
        adj[u].push_back(v);
        adj[v].push_back(u);
    }
    /* ----- centroid
decomposition boiler-plate -----
*/
    void calculate_size(int u, int p) {
        subsize[u] = 1;
        for (int v : adj[u])
            if (v != p && !blocked[v]) {
                calculate_size(v, u);
                subsize[u] += subsize[v];
            }
    }
    int get_centroid(int u, int p, int sz) {
        for (int v : adj[u])
            if (v != p && !blocked[v] &&
                subsize[v] > sz / 2)
                return get_centroid(v,
                                    u, sz);
        return u;
    }
    /* ----- fast
depth-multiset helpers -----
*/
    inline void add_depth(int d) {
        if (cnt[d] == 0)
            touched.push_back(d);
        ++cnt[d];
    }
    inline void clear_touched() {
        for (int d : touched) cnt[d] = 0;
        touched.clear();
    }
}

```

```

void collect_depths(int u, int p, int d, vector<int>& depths, int K) {
    if (d > K) return;
    depths.push_back(d);
    for (int v : adj[u])
        if (v != p && !blocked[v])
            collect_depths(v, u, d + 1, depths, K);
}

/* count all paths of exact length K
   that go through centroid c */
void process_centroid(int c, int K, long long& ans) {
    add_depth(0);
    // centroid itself
    for (int v : adj[c]) if
        (!blocked[v]) {
            vector<int> depths;
            collect_depths(v, c, 1,
                           depths, K);

            for (int d : depths)
                // pairs: one in this
                // subtree,
                if (d <= K)
                    // other in previous
                    // ones / root
                    ans += cnt[K - d];

            for (int d : depths)
                // merge current subtree
                depths
                if (d <= K) add_depth(d);
        }
    clear_touched();
    // reset counter
}

void decompose(int u, int p, int K, long long& ans) {
    calculate_size(u, -1);
    int c = get_centroid(u, -1,
                         subsize[u]);
    blocked[c] = 1;
    cpar[c] = p;

    process_centroid(c, K, ans);

    if (p != -1) {
        ctree[p].push_back(c);
        ctree[c].push_back(p);
    }
    for (int v : adj[c])
        if (!blocked[v])
            decompose(v, c, K, ans);
}

long long count_paths_of_length_k(int K) {
    cnt.assign(K + 1, 0);
    // initialise once
    touched.clear();
    long long ans = 0;
    decompose(1, -1, K, ans);
}

```

```

        return ans;
    }

    /* optional helpers kept intact */
    int get_centroid_parent(int u) {
        return cpar[u];
    }
    vector<int>&
    get_centroid_children(int u) {
        return ctree[u];
    }

    /* ----- main ----- */
    int main() {
        ios::sync_with_stdio(false);
        cin.tie(nullptr);

        int n, k;
        cin >> n >> k;

        CentroidTree ct(n);
        for (int i = 0; i < n - 1; ++i) {
            int u, v; cin >> u >> v;
            ct.add_edge(u, v);
        }
        cout << ct.count_paths_of_length_k(k)
            << '\n';
        return 0;
    }
}

7.4 dsu on tree [97 lines] - 94c0401be0


---


struct DSUTree {
    int n; // Number of nodes in the tree

    // Graph structure and metadata
    vector<vector<int>> adj; // Original tree (adjacency list)
    vector<int> color, sz, res; // color[i]: color of node i, sz[i]: subtree size, res[i]: result for node i
    vector<map<int, int*>> cnt; // cnt[i]: pointer to map storing color frequencies in subtree of node i

    // Constructor: initialize vectors and process input + solve
    DSUTree(int n) : n(n) {
        adj.resize(n + 1);
        color.resize(n + 1);
        sz.resize(n + 1);
        res.resize(n + 1);
        cnt.resize(n + 1, nullptr);
        take_input(); // Read input
        build(); // Run DSU on tree
    }

    // Read color input and edges
    void take_input() {
        for (int i = 1; i <= n; i++) {
            cin >> color[i];
        }
        for (int i = 1; i < n; i++) {
            int u, v;
            cin >> u >> v;
        }
    }
}

```

```

    adj[u].push_back(v);
    adj[v].push_back(u);
}

// Step 1: Compute sizes of all subtrees
void compute_subtree_sizes(int u, int p) {
    sz[u] = 1;
    for (int v : adj[u]) {
        if (v != p) {
            compute_subtree_sizes(v, u);
            sz[u] += sz[v];
        }
    }
}

// Step 2: Perform DSU on Tree to count distinct colors
void dfs(int u, int p) {
    int big = -1, max_size = -1;

    // Find the heavy child (largest subtree)
    for (int v : adj[u]) {
        if (v != p && sz[v] > max_size) {
            max_size = sz[v];
            big = v;
        }
    }

    // Process all light children first (their data will be deleted)
    for (int v : adj[u]) {
        if (v != p && v != big)
            dfs(v, u);
    }

    // Process heavy child and reuse its color map
    if (big != -1) {
        dfs(big, u);
        cnt[u] = cnt[big]; // Inherit heavy child's map
    } else {
        cnt[u] = new map<int, int>(); // Create new map for leaf
    }

    // Add current node's color
    (*cnt[u])[color[u]]++;

    // Merge all light child maps into current node's map
    for (int v : adj[u]) {
        if (v == p || v == big)
            continue;
        for (auto &[col, freq] :
            *cnt[v]) {
            (*cnt[u])[col] += freq;
        }
    }
}

void take_input() {
    for (int i = 1; i <= m; i++) {
        int u, v;
        cin >> u >> v;
        u++, v++;
    }
}

// Store the number of distinct colors in the subtree of u
res[u] = cnt[u]->size();
}

// Entry point: preprocess and call DFS
void build() {
    compute_subtree_sizes(1, 0); // Root the tree at node 1
    dfs(1, 0); // Start DFS from root
}

// Output the result
void print() {
    for (int i = 1; i <= n; i++) {
        cout << res[i] << " "; // Output result for node i
    }
    cout << '\n';
}
};

7.5 euler graph [99 lines] - 9ea81ac57c


---


/*
all the edges should be in the same connected component
#undirected graph: euler path/trail: all degrees are even or exactly two of them are odd.
#undirected graph: euler circuit/cycle[start == end]: all degrees are even
#directed graph: euler path: for all -> indeg = outdeg or in - out == 1 && out - in == 1
#directed graph: euler circuit: for all -> indeg = outdeg
*/
struct EulerGraph {
    vector<vector<pair<int, int>>> g;
    vector<int> done, path, edges, indeg, outdeg;
    vector<bool> visedge;
    bool isdirected;
    int n, m;
    EulerGraph(int n, int m, bool isd = 1) : n(n), m(m) {
        g.resize(n + 1);
        done.assign(n + 1, 0);
        indeg.assign(n + 1, 0);
        outdeg.assign(n + 1, 0);
        visedge.assign(m + 1, 0);
        isdirected = isd;
        take_input();
    }

    void take_input() {
        for (int i = 1; i <= m; i++) {
            int u, v;
            cin >> u >> v;
            u++, v++;
        }
    }
}

```

```

g[u].push_back({v, i});
indeg[v]++;
outdeg[u]++;
if(!isdirected) {
    g[v].push_back({u, i});
    indeg[u]++;
    outdeg[v]++;
}
}

void dfs(int u) {
    while(done[u] < g[u].size()) {
        auto [v, eid] =
            g[u][done[u]++];
        if(!isdirected) {
            if(vistedge[eid])
                continue;
            vistedge[eid] = 1;
        }
        dfs(v);
        edges.push_back(eid);
    }
    path.push_back(u);
}

pair<int, bool> find_start() { // true means any otherwise fixed
if(isdirected) {
    bool allsame = 1;
    for(int i = 1; i <= n; i++) {
        if(indeg[i] != outdeg[i])
            allsame = 0;
        break;
    }
}
if(allsame) return {-1, 1};
// any can be root

int incnt = 0, outcnt = 0,
node = -1; // must be 1 1 and remaining same
for(int i = 1; i <= n; i++) {
    if((outdeg[i] - indeg[i]) > 1 || (indeg[i] - outdeg[i]) > 1) {
        return {-1, 0};
    }
    if((outdeg[i] - indeg[i]) == 1) {
        node = i;
        outcnt++;
    }
    if((outdeg[i] - indeg[i]) == -1) incnt++;
}
if((incnt == 1 && outcnt == 1)) return {node, 1};
return {-1, 0};
}
else {
    bool alleven = 1;
    for(int i = 1; i <= n; i++) {
        if(indeg[i] & 1) {
            alleven = 0;
            break;
        }
    }
}
}

7.6 floyed warshall with negative cycle [29 lines] - 6e9cf92edd
vector<vector<ll>> dis(n, vector<ll>(n, inf));
for (int i = 0; i < n; i++) dis[i][i] = 0;
for(int i = 1; i <= m; i++) {
    int u, v, w; cin >> u >> v >> w;
    dis[u][v] = min(dis[u][v], 1LL * w);
}
for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < n; j++) {
            if (dis[i][k] < inf && dis[k][j] < inf) {
                dis[i][j] = min(dis[i][j],
                                dis[i][k] +
                                dis[k][j]);
            }
        }
    }
}
/* if (d[i][k] + d[k][j] < d[i][j] - EPS)
d[i][j] = d[i][k] + d[k][j]; */
for (int k = 0; k < n; k++) {
    if(dis[k][k] >= 0) continue;
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < n; j++) {
            if (dis[i][k] < inf && dis[k][j] < inf) {
                dis[i][j] = -inf; // negative cycle found
            }
        }
    }
}
// dis[u][v] == inf -> not path, -inf -> negative cycle

```

7.7 khun algo [31 lines] - 9cf33812ac

```

int n, k;
vector<vector<int>> g;
vector<int> mt;
vector<bool> used;

if(alleven) return {-1, 1};
// any can be root

int odd = 0, node = -1;
for(int i = 1; i <= n; i++) {
    if(indeg[i] & 1) {
        odd++;
        node = i;
    }
}
if(odd == 2) return {node, 1};
return {-1, 0};

void solve() {
    auto [root, flag] = find_start();
}

7.6 floyed warshall with negative cycle [29 lines] - 6e9cf92edd
vector<vector<ll>> dis(n, vector<ll>(n, inf));
for (int i = 0; i < n; i++) dis[i][i] = 0;
for(int i = 1; i <= m; i++) {
    int u, v, w; cin >> u >> v >> w;
    dis[u][v] = min(dis[u][v], 1LL * w);
}
for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < n; j++) {
            if (dis[i][k] < inf && dis[k][j] < inf) {
                dis[i][j] = min(dis[i][j],
                                dis[i][k] +
                                dis[k][j]);
            }
        }
    }
}
/* if (d[i][k] + d[k][j] < d[i][j] - EPS)
d[i][j] = d[i][k] + d[k][j]; */
for (int k = 0; k < n; k++) {
    if(dis[k][k] >= 0) continue;
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < n; j++) {
            if (dis[i][k] < inf && dis[k][j] < inf) {
                dis[i][j] = -inf; // negative cycle found
            }
        }
    }
}
// dis[u][v] == inf -> not path, -inf -> negative cycle

```

```

bool try_kuhn(int v) {
    if (used[v])
        return false;
    used[v] = true;
    for (int to : g[v]) {
        if (mt[to] == -1 || try_kuhn(mt[to])) {
            mt[to] = v;
            return true;
        }
    }
    return false;
}

int main() {
    //... reading the graph ...
}

mt.assign(k, -1);
for (int v = 0; v < n; ++v) {
    used.assign(n, false);
    try_kuhn(v);
}

for (int i = 0; i < k; ++i)
    if (mt[i] != -1)
        printf("%d %d\n", mt[i] + 1, i + 1);
}

7.8 kth sortest path [49 lines] - 6449754e4d
Time Complexity = O(m * K * log(n * K))

int m, n, deg[MM], source, sink, K;
int val[MM][12]; // val[u][k] = k-th shortest path distance to node u (0-based k)

// Edge representation: destination node and weight
struct edge {
    int v, w;
} adj[MM][500]; // Adjacency list as a 2D array: adj[u][i] = i-th neighbor of u

// Info stored in priority queue for Dijkstra
struct info {
    int v, w, k; // v = current node, w = total weight, k = current path rank
    bool operator<(const info &b) const {
        return w > b.w; // Min-heap by weight
    }
};

priority_queue<info, vector<info>> Q;

void kthBestShortestPath() {
    // Step 1: Initialize all distances to INF
    for (int i = 0; i < n; i++)
        for (int j = 0; j < K; j++)
            val[i][j] = inf;
}

void dfs(int u = 1, int p = 0) {
    par[u][0] = p;
    dep[u] = dep[p] + 1;
    sz[u] = 1;
}

```

```

// Step 2: Start from source
Q.push({source, 0, 0});
val[source][0] = 0;

// Step 3: Modified Dijkstra
while (!Q.empty()) {
    info u = Q.top(); Q.pop();

    // For each neighbor of u
    for (int i = 0; i < deg[u.v]; i++) {
        int to = adj[u.v][i].v;
        int cost = adj[u.v][i].w + u.w;

        // Try to insert this new cost into val[to][0...K-1]
        for (int k = 0; k < K; k++) {
            if (cost < val[to][k]) {
                swap(cost, val[to][k]); // Insert and shift worse value forward
                Q.push({to, val[to][k], k});
            }
        }
    }
}

7.9 lca [74 lines] - 41fcde7445
class LCA {
public:
    int n;
    vector<vector<int>> g, par;
    vector<int> dep, sz;
    int lg;
    LCA() {}
    LCA(int n) {
        this->n = n;
        g.assign(n + 1, vector<int>());
        lg = log2(n) + 1;
        par.assign(n + 1, vector<int>(lg + 1, 0));
        dep.assign(n + 1, 0);
        sz.assign(n + 1, 0);
        take_input();
    }
    void take_input() {
        for(int i = 1; i < n; i++) {
            int u, v;
            cin >> u >> v;
            u++, v++;
            g[u].push_back(v);
            g[v].push_back(u);
        }
    }
    void dfs(int u = 1, int p = 0) {
        par[u][0] = p;
        dep[u] = dep[p] + 1;
        sz[u] = 1;
    }
}
```

```

        for(int i = 1; i <= lg; i++) {
            par[u][i] = par[par[u][i - 1]][i - 1];
        }
        for(auto v : g[u]) {
            if(v != p) {
                dfs(v, u);
                sz[u] += sz[v];
            }
        }
    }
    int lca(int u, int v) {
        if(dep[u] < dep[v]) swap(u, v);
        // make sure that dep[u] > dep[v]
        for(int k = lg; k >= 0; k--) {
            if(dep[par[u][k]] >= dep[v])
                u = par[u][k];
        }
        if(u == v) return v;
        for(int k = lg; k >= 0; k--) {
            if(par[u][k] != par[v][k]) {
                u = par[u][k];
                v = par[v][k];
            }
        }
        return par[u][0];
    }
    int kth(int u, int k) {
        for(int i = 0; i <= lg; i++) {
            if(k & (1 << i)) {
                u = par[u][i];
            }
        }
        return u;
    }
    int dist(int u, int v) {
        int l = lca(u, v);
        return dep[u] + dep[v] - (2 * dep[l]);
    }
    // kth node from u to v, 0th node is u
    // k < dist(u, v)
    int go(int u, int v, int k) {
        int l = lca(u, v);
        if(depl + k <= dep[u]) return kth(u, k);
        k -= dep[u] - dep[l];
        return kth(v, dep[v] - dep[l] - k);
    }
} lca_ob;

```

## 7.10 reachability tree [92 lines] - db920f5986

/\*  
- Find the minimal/maximal weight of the edges  
when traversing from vertex u to vertex v  
- merge the nodes with new nodes using DSU  
- max/min weight will be the root of the final tree  
\*/  
**struct Reachability\_tree {**  
 vector<vector<int>> g, par;

```

        vector<int> dep, dsupar, cost,
                    subtree;
        vector<ll> tot;
        int id, n, m, lg, timer;
        Reachability_tree(int n, int m) :
            id(n), n(n), m(m), g(n + m + 1),
            dsupar(n + m + 1), dep(n + m
                + 1), cost(n + m + 1),
            tot(n + m + 1), subtree(n + m
                + 1) {
                iota(dsupar.begin(),
                    dsupar.end(), 0);
                lg = log2(n + m + 1) + 1;
                par.assign(n + m + 1,
                    vector<int>(lg + 1));
                timer = 0;
                build();
            }
            int find(int v) {
                if(dsupar[v] == v) return v;
                return dsupar[v] =
                    find(dsupar[v]);
            }
            void join(int u, int v, int w) {
                u = find(u);
                v = find(v);
                if(u == v) return;
                int new_node = ++id;
                g[new_node].push_back(u);
                g[u].push_back(new_node);
                g[new_node].push_back(v);
                g[v].push_back(new_node);
                cost[new_node] = max(cost[u],
                    cost[v]);
                dsupar[u] = new_node;
                dsupar[v] = new_node;
            }
            void dfs(int u, int p) {
                par[u][0] = p;
                dep[u] = dep[p] + 1;
                subtree[u] = (u <= n);
                for(int i = 1; i <= lg; i++) {
                    par[u][i] = par[par[u][i - 1]][i - 1];
                }
                for(auto v : g[u]) if(v != p) {
                    dfs(v, u);
                    subtree[u] += subtree[v];
                }
            }
            int lca(int u, int v) {
                if(dep[u] < dep[v]) swap(u, v);
                // make sure that dep[u] > dep[v]
                for(int k = lg; k >= 0; k--) {
                    if(dep[par[u][k]] >= dep[v])
                        u = par[u][k];
                }
                if(u == v) return v;
                for(int k = lg; k >= 0; k--) {
                    if(par[u][k] != par[v][k]) {
                        u = par[u][k];
                        v = par[v][k];
                    }
                }
                return par[u][0];
            }

```

```

        }
        void f(int u, int p) {
            for(auto v : g[u]) if(v != p) {
                tot[v] = tot[u] + 1LL *
                    (subtree[u] - subtree[v])
                    * cost[u];
                f(v, u);
            }
        }
        void build() {
            for(int i = 1; i <= n; i++) {
                cin >> cost[i];
            }
            vector<array<int, 3>> edges;
            for(int i = 1; i <= m; i++) {
                int u, v;
                cin >> u >> v;
                edges.push_back({max(cost[u],
                    cost[v]), u, v});
            }
            sort(edges.begin(), edges.end());
            for(auto [w, u, v] : edges) {
                if(find(u) != find(v)) {
                    join(u, v, w);
                }
            }
            dfs(id, 0);
            f(id, 0);
            for(int i = 1; i <= n; i++) {
                tot[i] += cost[i];
            }
        }
    };

```

## 7.11 scc [44 lines] - 928d336921

/\*components: number of SCC.  
sz: size of each SCC.  
comp: component number of each node.  
Create reverse graph.  
Run find\_scc() to find SCC.  
Might need to create condensation graph by create\_condensed().  
Think about indeg/outdeg  
for multiple test cases- clear adj/radj \*/  
 /comp/vis/sz/topo/condensed.\*/  
 vector<int> adj[mx], radj[mx];  
 int comp[mx], vis[mx], sz[mx],  
 components;  
 vector<int> topo;  
 void dfs(int u) {
 vis[u] = 1;
 for(int v : adj[u])
 if(!vis[v]) dfs(v);
 topo.push\_back(u);
 }
 void dfs2(int u, int val) {
 comp[u] = val;
 sz[val]++;
 for(int v : radj[u])
 if(comp[v] == -1)
 dfs2(v, val);
 }
 void find\_scc(int n) {
 memset(vis, 0, sizeof vis);

```

        memset(comp, -1, sizeof comp);
        for(int i = 1; i <= n; i++) {
            if(!vis[i])
                dfs(i);
        }
        reverse(topo.begin(), topo.end());
        for(int u : topo)
            if(comp[u] == -1)
                dfs2(u, ++components);
    }
    vector<int> condensed[mx];
    void create_condensed(int n) {
        for(int i = 1; i <= n; i++) {
            for(int v : adj[i])
                if(comp[i] != comp[v])
                    condensed[comp[i]].push_back(comp[v]);
        }
    }
}
```

## 8 number theory

### 8.1 linear sieve [16 lines] - 5bf8f49c51

```

vector<int> pri;
vector<int> lp; // lowest prime factor
void sieve(int n) {
    lp.assign(n + 1, 0);
    pri.clear();
    for(int i = 2; i <= n; i++) {
        if(lp[i] == 0) {
            lp[i] = i;
            pri.push_back(i);
        }
        for(int p : pri) {
            if(p > lp[i] || 1LL * p * i
                > n) break;
            lp[p * i] = p;
        }
    }
}
```

### 8.2 mobius [9 lines] - bc823ec9cd

```

//mobius O(nlogn)
int mob[N];
void mobius() {
    for(int i=0;i<N;i++)mob[i]=0;
    mob[1]=1;
    for(int i=1;i<N;i++)
        for(int j=i+i;j<N;j+=i)mob[j]-=mob[i];
}
```

### 8.3 pollard rho [94 lines] - 3bab19be72

```

namespace PollardRho {
mt19937 rnd(chrono::steady_clock::now().time_since_epoch().count());
const int P = 1e6 + 9;
ll seq[P];
int primes[P], spf[P];
inline ll add_mod(ll x, ll y, ll m) {
    return (x += y) < m ? x : x - m;
}
inline ll mul_mod(ll x, ll y, ll m) {
    ll res = __int128(x) * y % m;
    return res;
}
```

```

// ll res = x * y - (ll)((long
// double)x * y / m + 0.5) * m;
// return res < 0 ? res + m : res;
}

inline ll pow_mod(ll x, ll n, ll m) {
    ll res = 1 % m;
    for (; n; n >>= 1) {
        if (n & 1) res = mul_mod(res, x,
            m);
        x = mul_mod(x, x, m);
    }
    return res;
}

// O(it * (log n)^3), it = number of
// rounds performed
inline bool miller_rabin(ll n) {
    if (n <= 2 || (n & 1 ^ 1)) return (n
        == 2);
    if (n < P) return spf[n] == n;
    ll c, d, s = 0, r = n - 1;
    for (; !(r & 1); r >>= 1, s++) {}
    // each iteration is a round
    for (int i = 0; primes[i] < n &&
        primes[i] < 32; i++) {
        c = pow_mod(primes[i], r, n);
        for (int j = 0; j < s; j++) {
            d = mul_mod(c, c, n);
            if (d == 1 && c != 1 && c != (n
                - 1)) return false;
            c = d;
        }
        if (c != 1) return false;
    }
    return true;
}

void init() {
    int cnt = 0;
    for (int i = 2; i < P; i++) {
        if (!spf[i]) primes[cnt++] = spf[i]
            = i;
        for (int j = 0, k; (k = i *
            primes[j]) < P; j++) {
            spf[k] = primes[j];
            if (spf[i] == spf[k]) break;
        }
    }
}

// returns O(n^(1/4))
ll pollard_rho(ll n) {
    while (1) {
        ll x = rnd() % n, y = x, c = rnd()
            % n, u = 1, v, t = 0;
        ll *px = seq, *py = seq;
        while (1) {
            *py++ = y = add_mod(mul_mod(y,
                y, n), c, n);
            *py++ = y = add_mod(mul_mod(y,
                y, n), c, n);
            if ((x = *px++) == y) break;
            v = u;
            u = mul_mod(u, abs(y - x), n);
            if (!u) return __gcd(v, n);
            if (++t == 32) {
                t = 0;
                if ((u = __gcd(u, n)) > 1 && u
                    < n) return u;
            }
        }
    }
}

```

```

    }

    if (t && (u = __gcd(u, n)) > 1 && u
        < n) return u;
}

vector<ll> factorize(ll n) {
    if (n == 1) return vector<ll>();
    if (miller_rabin(n)) return
        vector<ll>{n};
    vector<ll> v, w;
    while (n > 1 && n < P) {
        v.push_back(spf[n]);
        n /= spf[n];
    }
    if (n >= P) {
        ll x = pollard_rho(n);
        v = factorize(x);
        w = factorize(n / x);
        v.insert(v.end(), w.begin(),
            w.end());
    }
    return v;
}

void solve() {
    PollardRho::init();
    ll p; cin >> p;
    auto get1 = PollardRho::factorize(p);
}

8.4 segmented sieve [23 lines] - 2314fd45fd
vector<ll> segmented_sieve(ll l, ll r) {
    vector<ll> segpr;
    vector<bool> pr(r - l + 5, 1);
    if (l == 1) {
        pr[0] = false;
    }
    for (ll i = 0; svp[i] * svp[i] <= r;
        i++) {
        ll cur = svp[i];
        ll base = cur * cur;
        if (base < 1) {
            base = ((l + cur - 1) / cur)
                * cur;
        }
        for (ll j = base; j <= r; j +=
            cur) {
            pr[j - 1] = false;
        }
    }
    for (ll i = 0; i <= r - 1; i++) {
        if (pr[i]) {
            segpr.push_back(l + i);
        }
    }
    return segpr;
}

8.5 totient phi [28 lines] - c41619b4ad
// all of (1-10^6) -> O(n log n)
int phi[N];
void totient() {
    for (int i = 0; i < N; i++) phi[i] =
        i;
    for (int i = 2; i < N; i++) {
        if (phi[i] != i) continue;

```

```

        for (int j = i; j < N; j += i)
            phi[j] -= phi[j] / i;
    }
}

// 10^16 range->O(sqrt(n))
int phiValue(int n)
{
    int ans=1;
    int q=sqrt(n);
    for(int i=2;i<=q;i++)
    {
        if(n%i==0)
        {
            int tem=1;
            while(n%i==0) tem*=i, n/=i;
            ans=ans*tem/i*(i-1);
            q=sqrt(n);
        }
    }
    if(n>1)ans=ans*(n-1);
    return ans;
}

9 math
9.1 FFT [59 lines] - db73a70650
//Multiply returns ans[k]= $\sum ai*bj \% MOD$ 
//such that i+j==k
const double PI = acos(-1);
struct cd {
    double x, y;
    cd(double x = 0, double y = 0):
        x(x), y(y) {}
    cd operator+(cd o) const { return
        cd(x + o.x, y + o.y); }
    cd operator-(cd o) const { return
        cd(x - o.x, y - o.y); }
    cd operator*(cd o) const { return
        cd(x*x - y*y, x*y + y*x); }
    cd operator*(double d) const { return
        cd(x * d, y * d); }
    cd conj() const { return cd(x, -y); }
};

void fft(vector<cd> &a, bool invert) {
    int n = a.size();
    for (int i = 1, j = 0; i < n; ++i) {
        int bit = n >> 1;
        while (j & bit) {
            j ^= bit;
            bit >>= 1;
        }
        j ^= bit;
        if (i < j)
            swap(a[i], a[j]);
    }
    for (int len = 2; len <= n; len <=
        1) {
        double ang = 2 * PI / len *
            (invert ? -1 : 1);
        cd wlen(cos(ang), sin(ang));
        for (int i = 0; i < n; i += len)
        {
            cd w(1);
            for (int j = 0; j < len / 2;
                ++j) {

```

```

                cd u = a[i + j];
                cd v = a[i + j + len/2];
                *w;
                a[i + j] = u + v;
                a[i + j + len/2] = u - v;
                w = w * wlen;
            }
        }
        if (invert) {
            for (cd &x : a)
                x = x * (1.0 / n);
        }
    }
}

vector<long long> multiply(const
    vector<int> &a, const vector<int> &b)
{
    vector<cd> fa(a.begin(), a.end()),
        fb(b.begin(), b.end());
    int n = 1;
    while (n < a.size() + b.size())
        n <= 1;
    fa.resize(n);
    fb.resize(n);
    fft(fa, false);
    fft(fb, false);
    for (int i = 0; i < n; ++i)
        fa[i] = fa[i] * fb[i];
    fft(fa, true);
    vector<long long> result(n);
    for (int i = 0; i < n; ++i)
        result[i] = round(fa[i].x);
    return result;
}

9.2 NTT [55 lines] - 80b38224a2
//Multiply returns ans[k]= $\sum ai*bj \% MOD$ 
//such that i+j==k
//Frequency arrays are often used
const int MOD=998244353, ROOT=3;
int lim, inv_lim;
vector<int> rev, wn, w;
int modexp(int x,int e){
    int r=1;
    while(e){ 
        if(e&1) r=(long long)r*x%MOD;
        x=(long long)x*x%MOD;
        e>>=1;
    }
    return r;
}

void precompute(int n){
    lim=1; int L=0;
    while(lim<n) lim<=1, ++L;
    rev.assign(lim,0);
    for(int i=0;i<lim;i++)
        rev[i]=(rev[i]>>1)|((i&1)<<(L-1));
    wn.assign(lim,1);
    int g=modexp(ROOT,(MOD-1)/lim);
    for(int i=1;i<lim;i++) wn[i]=(long
        long)wn[i-1]*g%MOD;
    inv_lim=modexp(lim,MOD-2);
    w.resize(lim);
}


```

```

void ntt(vector<int>&a, bool invert){
    for(int i=0;i<lim;i++) if(i<rev[i])
        swap(a[i],a[rev[i]]);
    for(int len=1;len<lim;len<=1){
        int step=lim/(len<1);
        for(int i=0;i<len;i++)
            w[i]=wn[i*step];
        for(int i=0;i<lim;i+=len<1)
            for(int j=0;j<len;j++){
                int u=a[i+j];
                int v=(long)a[i+j+len]*w[j]%MOD;
                a[i+j]=u+v<MOD?u+v:u+v-MOD;
                a[i+j+len]=u-v>=0?u-v:u-v+MOD;
            }
        if(invert){
            reverse(a.begin()+1,a.begin()+lim);
            for(int&i:a) i=(long)
                long)i*inv_lim%MOD;
        }
    }
    vector<int>
        multiply(vector<int>a,vector<int>b){
    if(a.empty()||b.empty()) return {};
    int need=a.size()+b.size()-1;
    precompute(need);
    a.resize(lim); b.resize(lim);
    ntt(a,false); ntt(b,false);
    for(int i=0;i<lim;i++) a[i]=(long)
        long)a[i]*b[i]%MOD;
    ntt(a,true);
    a.resize(need);
    return a;
}

```

### 9.3 No of Digits in $n!$ in base B [14 lines] - 94373fb11a

```

ll NoOfDigitInNFactInBaseB(ll N, ll B) {
    ll i;
    double ans = 0;

    // Sum of logarithms:  $\log(N!) = \log(1) + \log(2) + \dots + \log(N)$ 
    for (i = 1; i <= N; i++) ans += log(i);

    // Convert log from base e to base B,
    // then add 1 to get digit count
    ans = ans / log(B);
    ans = ans + 1;

    // Return the result as an integer
    // (floor value)
    return (ll)ans;
}

```

### 9.4 Xor basis [35 lines] - ad03dee2a0

```

const int N = 1505;
bitset<N> bit[N];           // input vectors
// 1-indexed
int basis[N];               // basis[i] =
vector<int> g[N];           // for
                            // reconstruction: which vectors XORed
                            // to form current

```

```

int pivot[N];                  // pivot[j] =
                                // leading bit index of row j

void build_linear_basis(int n, int
maxBit) {
    fill(pivot, pivot + n + 1, -1);
    fill(basis, basis + maxBit + 1, -1);

    for(int i = maxBit; i >= 0; i--) {
        int pivotRow = -1;
        for(int j = 1; j <= n; j++) {
            if(pivot[j] == -1) continue;
            if(pivotRow != -1 &&
bit[j][i]) {
                bit[j] ^= bit[pivotRow];
                // eliminate bit i
                g[j].push_back(pivotRow);
                // record operation
            }
            else if(bit[j][i]) {
                pivotRow = j;
                // new pivot
                pivot[j] = i;
                basis[i] = j;
            }
        }
    }
}

// Example: maximum XOR of any subset
int maxXOR(int maxBit) {
    bitset<N> res;
    for(int i = maxBit; i >= 0; i--) {
        if(basis[i] != -1 && !res[i]) res
            ^= bit[basis[i]];
    }
    return (int)(res.to_ullong()); // works if N <= 64, else handle
                                // manually
}

```

### 9.5 inverse and mul mat [94 lines] - 7ddc88f99f

```

struct Matrix {
    vector<vector<double>> a;
    int n, m;

    Matrix(int r = 0, int c = 0) : n(r),
        m(c), a(r, vector<double>(c)) {}

    static Matrix identity(int size) {
        Matrix I(size, size);
        for (int i = 0; i < size; ++i)
            I.a[i][i] = 1;
        return I;
    }

    Matrix operator*(const Matrix &B)
        const {
        assert(m == B.n);
        Matrix res(n, B.m);
        for (int i = 0; i < n; ++i)
            for (int j = 0; j < B.m; ++j)
                for (int k = 0; k < m;
++k)

```

```

                    res.a[i][j] +=
                        a[i][k] *
                        B.a[k][j];
    }
    return res;
}

Matrix pow(ll exp) const {
    assert(n == m);
    Matrix base = *this, res =
        identity(n);
    while (exp) {
        if (exp & 1) res = res *
            base;
        base = base * base;
        exp >>= 1;
    }
    return res;
}

Matrix transpose() const {
    Matrix res(m, n);
    for (int i = 0; i < n; ++i)
        for (int j = 0; j < m; ++j)
            res.a[j][i] = a[i][j];
    return res;
}

double determinant() const {
    assert(n == m);
    Matrix tmp = *this;
    double det = 1;
    for (int i = 0; i < n; ++i) {
        int pivot = i;
        for (int j = i + 1; j < n;
++j)
            if (fabs(tmp.a[j][i]) >
                fabs(tmp.a[pivot][i]))
                pivot = j;
        if (fabs(tmp.a[pivot][i]) <
            1e-9) return 0;
        if (i != pivot)
            swap(tmp.a[i],
                tmp.a[pivot]), det *=
                -1;
        det *= tmp.a[i][i];
        for (int j = i + 1; j < n;
++j) {
            double f = tmp.a[j][i] /
                tmp.a[i][i];
            for (int k = i; k < n;
++k)
                tmp.a[j][k] -= f *
                    tmp.a[i][k];
        }
    }
    return det;
}

bool inverse(Matrix &inv) const {
    if (n != m) return false;
    inv = identity(n);
    Matrix tmp = *this;

    for (int i = 0; i < n; ++i) {
        int pivot = i;

```

```

        for (int j = i + 1; j < n;
++j)
            if (fabs(tmp.a[j][i]) >
                fabs(tmp.a[pivot][i]))
                pivot = j;
        if (fabs(tmp.a[pivot][i]) <
            1e-9) return false;
        swap(tmp.a[i], tmp.a[pivot]);
        swap(inv.a[i], inv.a[pivot]);

        double f = tmp.a[i][i];
        for (int j = 0; j < n; ++j) {
            tmp.a[i][j] /= f;
            inv.a[i][j] /= f;
        }
    }
    for (int j = 0; j < n; ++j) {
        if (i == j) continue;
        double fac = tmp.a[j][i];
        for (int k = 0; k < n;
++k) {
            tmp.a[j][k] -= fac *
                tmp.a[i][k];
            inv.a[j][k] -= fac *
                inv.a[i][k];
        }
    }
    return true;
}

```

### 9.6 linear diophantine eqn [73 lines] - e732e8bcad

```

ll extended_euclid(ll a, ll b, ll &x, ll
&y) {
    ll xx = y = 0, yy = x = 1;
    while (b) {
        ll q = a / b;
        ll t = b; b = a % b; a = t;
        t = xx; xx = x - q * xx; x = t;
        t = yy; yy = y - q * yy; y = t;
    }
    return a;
}

// Solves  $a*x + b*y = c$ . Finds any
// solution (x0, y0)
bool find_any_solution(ll a, ll b, ll c,
ll &x0, ll &y0, ll &g) {
    if (a == 0 && b == 0) {
        if (c) return false;
        x0 = y0 = g = 0;
        return true;
    }
    g = extended_euclid(abs(a), abs(b),
        x0, y0);
    if (c % g != 0) return false;
    x0 *= c / g;
    y0 *= c / g;
    if (a < 0) x0 *= -1;
    if (b < 0) y0 *= -1;
    return true;
}

```

```

void shift_solution(ll &x, ll &y, ll a,
ll b, ll cnt) {
x += cnt * b;
y -= cnt * a;
}

// Counts the number of solutions to a*x
// + b*y = c with x in [minx, maxx] and
// y in [miny, maxy]
11 find_all_solutions(ll a, ll b, ll c,
ll minx, ll maxx, ll miny, ll maxy) {
ll x, y, g;
if (!find_any_solution(a, b, c, x,
y, g)) return 0;

if (a == 0 && b == 0) {
assert(c == 0);
return 1LL * (maxx - minx + 1) *
(maxy - miny + 1);
}
if (a == 0) return (maxx - minx + 1) *
(miny <= c / b && c / b <=
maxy);
if (b == 0) return (maxy - miny + 1) *
(minx <= c / a && c / a <=
maxx);

a /= g, b /= g;
ll sign_a = a > 0 ? +1 : -1;
ll sign_b = b > 0 ? +1 : -1;

shift_solution(x, y, a, b, (minx - x
/ b));
if (x < minx) shift_solution(x, y,
a, b, sign_b);
if (x > maxx) return 0;
ll lx1 = x;

shift_solution(x, y, a, b, (maxx - x
/ b));
if (x > maxx) shift_solution(x, y,
a, b, -sign_b);
ll rx1 = x;

shift_solution(x, y, a, b, -(miny -
y) / a);
if (y < miny) shift_solution(x, y,
a, b, -sign_a);
if (y > maxy) return 0;
ll lx2 = x;

shift_solution(x, y, a, b, -(maxy -
y) / a);
if (y > maxy) shift_solution(x, y,
a, b, sign_a);
ll rx2 = x;

if (lx2 > rx2) swap(lx2, rx2);
ll lx = max(lx1, lx2);
ll rx = min(rx1, rx2);
if (lx > rx) return 0;

return (rx - lx) / abs(b) + 1;
}

```

## 10 misc

### 10.1 2d pref sum [26 lines] - 0d9b46b715

```

struct PrefixSum2D {
vector<vector<int>> prefix;
int n, m;
PrefixSum2D(const
vector<vector<int>>& grid) {
n = grid.size();
m = grid[0].size();
prefix.assign(n, vector<int>(m,
0));

for (int i = 0; i < n; ++i) {
for (int j = 0; j < m; ++j) {
prefix[i][j] =
grid[i][j];

if (i > 0) prefix[i][j] +=
prefix[i - 1][j];
// Top
if (j > 0) prefix[i][j] +=
prefix[i][j - 1];
// Left
if (i > 0 && j > 0)
prefix[i][j] -=
prefix[i - 1][j - 1]; // Top-left
overlap
}
}
int getSum(int x1, int y1, int x2,
int y2) {
int sum = prefix[x2][y2];
if (x1 > 0) sum -= prefix[x1 - 1][y2];
if (y1 > 0) sum -= prefix[x2][y1 - 1];
if (x1 > 0 && y1 > 0) sum +=
prefix[x1 - 1][y1 - 1];
return sum;
}

```

### 10.2 bit hacks [17 lines] - 13d6b87d3b

```

#define ckbit(n, k) (((n) >> (k)) & 1)
#define toggle(n, k) ((n) ^ (1LL << (k)))
#define setbit(n, k) ((n) |= (1LL << (k)))
#define unsetbit(n, k) ((n) &= ~(1LL <<
(k)))
#define lowbit(n) ((n) & -(n))
#define highbit(n) (63 -
__builtin_clzll(n)) // = floor
log2(n)
// a/b = a ~ b + a & b
// a ^ (a & b) = b ^ (a/b)
// (a & b) ^ (a/b) = a ~ b
// a+b = a/b + a & b
// a+b = a ~ b + 2(a & b)
// a-b = (a ^ (a & b)) - ((a/b) ^ a)
// a-b = ((a/b) ^ b) - ((a/b) ^ a)
// a-b = (a ^ (a & b)) - (b ^ (a & b))
// a-b = ((a/b) ^ b) - (b ^ (a & b))

```

### 10.3 job with 2 deadline [62 lines] - 546d97e7c2

```

/*
given n workers for a company
for each worker given 3 parameter
k, l, r
we need to assign each worker for a day i
for all n days
and no two worker can be assigned on same
date
if worker i is assigned on jth day and j
<= k then the profit 'l' otherwise
'r'
assign such that maximize profit
*/

```

#### #include <bits/stdc++.h>

```

using namespace std;
typedef long long ll;
typedef long double ld;
const ld inf = 1e100;
const ld eps = 1e-18;
```

#### void solve()

```

int n;
cin >> n;
vector<vector<int>> front(n + 1),
back(n + 1);
ll tot = 0;
for(int i = 1; i <= n; i++) {
int k, l, r;
cin >> k >> l >> r;
tot += min(l, r);
if(l > r) {
front[k].push_back(l - r);
}
else if(l < r && (n - k) > 0) {
back[n - k].push_back(r - 1);
}
}
```

```

auto add = [&](vector<vector<int>>
&vec) {
priority_queue<int> pq;
```

```

for(int i = 1; i <= n; i++) {
for(auto g : vec[i]) {
pq.push(g);
if(pq.size() > i)
pq.pop();
}
}
ll res = 0;
while(not pq.empty()) {
res += -pq.top(); pq.pop();
}
return res;
};
```

```

cout << tot + add(front) + add(back)
<< "\n";
}
```

```

int main() {
ios_base::sync_with_stdio(false);
cin.tie(0);
int tc = 1, cs = 1;
```

```

cin >> tc;
while (tc--) {
//cout << "Case " << cs++ << ":" ;
solve();
}
return 0;
}
```

### 10.4 pairs rearrange made palindrome [24 lines] - 449d38aff6

```

// palindrome count
// odd <= 1
// so Xor = 0 means odd = 0
// Xor = 1 means odd = 1 for this
// i need to flip exactly one bit of a-z
map<int, 11> freq;
int mask = 0;
// before processing any characters, we
// have one occurrence of mask 0
freq[mask] = 1;
ll result = 0;
// iterate over each character in the
// string
for (char c : s) {
// Toggle the bit corresponding to
// the current character
mask ^= (1 << (c - 'a'));
// count substrings ending here that
// are already balanced (exact
// match)
result += freq[mask];
// count substrings ending here that
// differ by exactly one bit
for (int b = 0; b < 26; ++b) {
result += freq[mask ^ (1 << b)];
}
// update frequency for the current
// mask
freq[mask]++;
}
cout << result << "\n";
```

### 10.5 run [16 lines] - be27096b28

# Open & set C++ build system in Sublime (Linux)

- 1) Open Sublime Text
- 2) Menu → Tools → Build System → New Build System
- 3) Paste this code:

```
{
"cmd": ["bash", "-c", "g++ \"$file\""
"-std=gnu++17 -O2 -o a.out &&
./a.out < input.txt"],
"selector": "source.c, source.c++",
"working_dir": "$file_path"
}
```

- 4) Ctrl + S → save as: cpp.sublime-build
- 5) Menu → Tools → Build System → cpp
- 6) Open main.cpp → Ctrl + B
- # input from input.txt, output in terminal

**10.6 sliding windows min-max using deque [13 lines] - 52c44b3137**

```
deque<int> dq;
for(int i = 1; i <= n; i++) {
    while(!dq.empty() && v[dq.back()] < v[i]) {
        dq.pop_back();
    }
    dq.push_back(i);
    if(dq.front() <= i - k) {
        dq.pop_front();
    }
    if(i >= k) {
        cout << v[dq.front()] << " ";
    }
}
```

**10.7 unique OR all subarray [22 lines] - b574681b4e**

```
int ans = 0;
map<int, bool> has;
for(int i = 1; i <= n; i++) {
    int cur = 0, next = i, cnt = 0;
    while(next <= n) {
        cur |= tree.query(1, 1, n, i,
                          next).val;
        if(!has.count(cur)) cnt++;
        has[cur] = 1;
        int mn = n + 2;
        for(int j = 0; j < 30; j++) {
            if((!cur & (1 << j))) {
                if(bitv[j].size() > 0) {
                    auto it =
                        lower_bound(bitv[j].begin(),
                                    bitv[j].end(),
                                    i);
                    if(it != bitv[j].end()) {
                        mn = min(mn,
                                  *it);
                    }
                }
            }
        }
        next = mn;
    }
}
```

## 11 Data structure

**11.1 2dbit [97 lines] - 089f12d2d7**

```
struct _2dbit {
    struct BITmx {
        vector<int> v;
        int n;
        BITmx(){}
        BITmx(int n) : n(n) {
            v.resize(n + 1);
        }
        void upd(int i, int val) {
            for(; i <= n; i += (i & -i))
            {
                v[i] = max(v[i], val);
            }
        }
        int query(int i) {

```

```
            int ans = 0;
            for(; i > 0; i -= (i & -i)) {
                ans = max(ans, v[i]);
            }
            return ans;
        }
        vector<vector<int>> cords;
        vector<BITmx> cnt;
        int n;
        _2dbit(int n) : n(n) {
            cnt.resize(n + 1);
            cords.resize(n + 1);
        }
        void add(int i, int val) {
            for(; i <= n; i += (i & -i)) {
                cords[i].push_back(val);
            }
        }
        void build() {
            for(int i = 1; i <= n; i++) {
                auto &v = cords[i];
                sort(v.begin(), v.end());
                v.erase(unique(v.begin(),
                               v.end()), v.end());
                cnt[i] = BITmx(v.size());
            }
        }
        int getidx(int p, int val) {
            auto &v = cords[p];
            return (lower_bound(v.begin(),
                                v.end(), val) - v.begin() +
                    1);
        }
        void upd(int x, int y, int val) {
            for(int i = x; i <= n; i += (i & -i)) {
                int id = getidx(i, y);
                cnt[i].upd(id, val);
            }
        }
        int query(int x, int y) {
            int ans = 0;
            for(int i = x; i > 0; i -= (i & -i)) {
                int id = getidx(i, y);
                ans = max(ans,
                          cnt[i].query(id - 1));
            }
            return ans;
        }
    };
    void solve() {
        int n;
        cin >> n;
        vector<array<int, 2>> v(n + 1);
        vector<int> allx, ally;
        for(int i = 1; i <= n; i++) {
            cin >> v[i][0] >> v[i][1];
            allx.push_back(v[i][0]);
            ally.push_back(v[i][1]);
        }
        sort(allx.begin(), allx.end());
        allx.erase(unique(allx.begin(),
                         allx.end()), allx.end());
    }
}
```

```
    auto getidxX = [&](int x) {
        return (lower_bound(allx.begin(),
                           allx.end(), x) - allx.begin() +
                1);
    };
    sort(ally.begin(), ally.end());
    ally.erase(unique(ally.begin(),
                     ally.end()), ally.end());
    auto getidxY = [&](int x) {
        return (lower_bound(ally.begin(),
                           ally.end(), x) - ally.begin() +
                1);
    };
    int m = allx.size();
    _2dbit bit(m);
    for(int i = 1; i <= n; i++) {
        int x = getidxX(v[i][0]), y =
            getidxY(v[i][1]);
        bit.add(x, y);
    }
    bit.build();
    int ans = 0;
    for(int i = 1; i <= n; i++) {
        int x = getidxX(v[i][0]), y =
            getidxY(v[i][1]);
        int val = 1 + bit.query(x - 1,
                               y);
        ans = max(ans, val);
        bit.upd(x, y, val);
    }
    cout << ans << "\n";
}

```

**11.2 bit [24 lines] - a82760cf8**

```
template <class T>
struct BIT { //1-indexed
    int n; vector<T> t;
    BIT() {}
    BIT(int _n) {
        n = _n; t.assign(n + 1, 0);
    }
    T prefixSum(int i) {
        T ans = 0;
        for (; i >= 1; i -= (i & -i)) ans +=
            t[i];
        return ans;
    }
    void upd(int i, T val) {
        if (i <= 0) return;
        for (; i <= n; i += (i & -i)) t[i] +=
            val;
    }
};
// inversion
BIT<ll> bit(n);
for(int i = 1; i <= n; i++) {
    int x; cin >> x;
    ans[i] = bit.query(n) - bit.query(x -
        1);
    bit.upd(x, 1);
}

```

**11.3 dsu with rollback [33 lines] - aa5a0e5006**

```
struct DSU {
    vector<ll> par, rank, sz;
    int n;
    DSU(int n) : n(n), par(n + 1), rank(n +
        1, 0), sz(n + 1, 1) {
        for(int i = 0; i < n; i++)
            par[i] = i;
        componentCount = 0;
    }
    int find(int v) {
        if (parent[v] == v) return v;
        return find(parent[v]);
    }
    bool merge(int u, int v) {
        u = find(u), v = find(v);
        if (u == v) return false;
        if (size[u] < size[v]) swap(u, v);
        history.push_back({v, size[v]});
        history.push_back({u, size[u]});
        parent[v] = u, size[u] += size[v], --componentCount;
        return true;
    }
    void rollback() {
        auto [u, oldSizeU] =
            history.back();
        history.pop_back();
        auto [v, oldSizeV] =
            history.back();
        history.pop_back();
        parent[v] = v, size[u] =
            oldSizeU, size[v] =
            oldSizeV, componentCount++;
    }
    int getComponentCount(){return componentCount;}
    bool same(int u, int v){return find(u) == find(v);}
    int getSize(int v){return size[find(v)];}
};

11.4 dsu [97 lines] - 4c9e4f01e0
/*
sometimes i need to modify the find()
function
first of all add all expreinece as
team[u]
then when join() make self[v] = team[v]
- team[u]
so for v the ans still team[v] = self[v]
+ team[u]
and when we find some other parents
then the added points add to all players
of a team
so we do self[v] += self[p] so here
self[p]
*/
struct DSU {
    vector<ll> par, rank, sz;
    int n;
    DSU(int n) : n(n), par(n + 1), rank(n +
        1, 0), sz(n + 1, 1) {
        for(int i = 0; i < n; i++)
            par[i] = i;
        componentCount = 0;
    }
    int find(int v) {
        if (parent[v] == v) return v;
        return find(parent[v]);
    }
    bool merge(int u, int v) {
        u = find(u), v = find(v);
        if (u == v) return false;
        if (size[u] < size[v]) swap(u, v);
        history.push_back({v, size[v]});
        history.push_back({u, size[u]});
        parent[v] = u, size[u] += size[v], --componentCount;
        return true;
    }
    void rollback() {
        auto [u, oldSizeU] =
            history.back();
        history.pop_back();
        auto [v, oldSizeV] =
            history.back();
        history.pop_back();
        parent[v] = v, size[u] =
            oldSizeU, size[v] =
            oldSizeV, componentCount++;
    }
    int getComponentCount(){return componentCount;}
    bool same(int u, int v){return find(u) == find(v);}
    int getSize(int v){return size[find(v)];}
}
```

```

    iota(par.begin(), par.end(), 0);
int find(int v) {
    return (par[v] == v ? v : (par[v] = find(par[v])));
}
bool same(int u, int v) {
    return find(u) == find(v);
}
int get_size(int v) {
    return sz[find(v)];
}
int count() {
    return n;
}
void join(int u, int v) {
    u = find(u);
    v = find(v);
    n--;
    if (rank[v] > rank[u]) swap(u, v);
    par[v] = u;
    sz[u] += sz[v];
    if (rank[u] == rank[v])
        rank[u]++;
    // u is the parent;
}
struct DSU {
    vector<ll> par, rank, sz, sum,
    elemwhichnode;
    int n, nextId;
    DSU(int n, int m)
        : par(n + m + 1), rank(n + m +
            1, 0), sz(n + m + 1, 1),
        sum(n + m + 1), elemwhichnode(n +
            1), n(n), nextId(n) {
            iota(par.begin(), par.end(), 0);
            for (int i = 1; i <= n; ++i)
                elemwhichnode[i] = i;
                sum[i] = i;
    }
    void join(int u, int v) {
        u = find(u);
        v = find(v);
        if (u == v) return;
        n--;
        if (rank[v] > rank[u]) swap(u, v);
        par[v] = u;
        sz[u] += sz[v];
        sum[u] += sum[v];
        if (rank[u] == rank[v])
            rank[u]++;
        // u is the parent;
    }
    void erase(int v) {
        int nv = elemwhichnode[v];
        if (nv == 0) return;
        // already deleted
        int pv = find(nv);
        sz[pv] -= 1;
        sum[pv] -= v;
    }
};

```

```

    elemwhichnode[v] = 0;
    // mark deleted
}
/* move p -> set(q) */
void move(int u, int v) {
    if (elemwhichnode[u] == 0) return;
    int pu = find(elemwhichnode[u]);
    int pv = find(elemwhichnode[v]);
    if (pu == pv) return;
    erase(u);
    nextId++;
    par[nextId] = nextId;
    rank[nextId] = 0;
    sz[nextId] = 1;
    sum[nextId] = u;
    elemwhichnode[u] = nextId;
    join(nextId, pv);
}
pair<int, ll> query(int p) {
    int r = find(elemwhichnode[p]);
    return { (int)sz[r], sum[r] };
}

```

### 11.5 dynamic segtree [95 lines] - 651c502af4

```

class SegTlazy {
public:
    struct node {
        ll mn, mncnt, lazy;
        bool haslazy;
        node *l, *r;
        node() : mn(inf), mncnt(0),
            lazy(0), haslazy(0),
            l(nullptr), r(nullptr) {}
        node(ll mn, ll mncnt) : mn(mn),
            mncnt(mncnt), lazy(0),
            haslazy(0), l(nullptr),
            r(nullptr) {}
    };
    node *root;
    ll n, m; // max and min range
    // default 1 to max
    explicit SegTlazy(ll _n, ll _m = 1) :
        n(_n), m(_m) { root = new
        node(); }
    inline void merge(node *nd) {
        nd->mncnt = 0;
        if(nd->l && nd->r) {
            nd->mn = min(nd->l->mn,
                nd->r->mn);
            if(nd->mn == nd->l->mn)
                nd->mncnt +=
                nd->l->mncnt;
            if(nd->mn == nd->r->mn)
                nd->mncnt +=
                nd->r->mncnt;
        }
        else if(nd->l) {
            nd->mn = nd->l->mn;
        }
    }
    void upd(ll l, ll r, ll v) {
        upd(root, m, l, r, v);
    }
};

```

```

    if(nd->mn == nd->l->mn)
        nd->mncnt +=
        nd->l->mncnt;
}
else if(nd->r) {
    nd->mn = nd->r->mn;
    if(nd->mn == nd->r->mn)
        nd->mncnt +=
        nd->r->mncnt;
}
void apply(node *nd, ll b, ll e, ll
lazy) {
    nd->mn += lazy;
}
inline void push(node *nd, ll b, ll
e) {
    if (!nd->haslazy) return;
    apply(nd, b, e, nd->lazy);
    if (b != e) {
        if (!nd->l) nd->l = new
        node();
        if (!nd->r) nd->r = new
        node();
        nd->l->lazy += nd->lazy;
        nd->r->lazy += nd->lazy;
        nd->l->haslazy = 1;
        nd->r->haslazy = 1;
    }
    nd->lazy = 0;
    nd->haslazy = 0;
}
void build(node *nd, ll b, ll e) {
    if(b == e) {
        nd->mn = 0;
        nd->mncnt = 1;
        return;
    }
    ll mid = (b + e) >> 1;
    if (!nd->l) nd->l = new node();
    if (!nd->r) nd->r = new node();
    build(nd->l, b, mid);
    build(nd->r, mid + 1, e);
    merge(nd);
}
void build() { build(root, m, n); }
void upd(node *nd, ll b, ll e, ll i,
j, ll v) {
    push(nd, b, e);
    if (j < b || e < i) return;
    if (i <= b && e <= j) {
        nd->lazy += v;
        nd->haslazy = 1;
        push(nd, b, e);
        return;
    }
    ll mid = (b + e) >> 1;
    if (!nd->l) nd->l = new node();
    if (!nd->r) nd->r = new node();
    upd(nd->l, b, mid, i, j, v);
    upd(nd->r, mid + 1, e, i, j, v);
    merge(nd);
}
void upd(ll l, ll r, ll v) {
    upd(root, m, l, r, v);
}

```

```

node query(node *nd, ll b, ll e, ll
i, ll j) {
    if (j < b || e < i) return
    node();
    if (!nd) return node();
    if (i <= b && e <= j) return *nd;
    push(nd, b, e);
    ll mid = (b + e) >> 1;
    node left = query(nd->l, b,
    mid, i, j);
    node right = query(nd->r, mid +
    1, e, i, j);
    node res;
    res.mn = min(left.mn, right.mn);
    if(res.mn == left.mn) res.mncnt
        += left.mncnt;
    if(res.mn == right.mn) res.mncnt
        += right.mncnt;
    return res;
}
node query(ll l, ll r) { return
query(root, m, n, l, r); }

```

### 11.6 gp hash table [19 lines] - 5b4a28d281

```

#include<ext/pb_ds/assoc_container.hpp>
#include<ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
struct custom_hash {
    static uint64_t splitmix64(uint64_t x)
    {
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) *
            0xbff58476d1ce4e5b9;
        x = (x ^ (x >> 27)) *
            0x94d049bb133111eb;
        return x ^ (x >> 31);
    }
    size_t operator()(uint64_t x) const {
        static const uint64_t FIXED_RANDOM =
            chrono::steady_clock::now().time_since_
            epoch();
        return splitmix64(x + FIXED_RANDOM);
    }
};
```

//pair (a, b) er jonne a \* MOD + b

```

gp_hash_table<int, int, custom_hash> mp;
unordered_set<ll, custom_hash> mp;
mp.reserve(n);

```

### 11.7 hld [109 lines] - 14d49ef41

```

class HLD : public SegTlazy {
public:
    vector<vector<pair<int, int>>> g;
    vector<vector<int>> par;
    vector<int> heavy, head, pos, depth,
    val;
    int timer, n, height;
    HLD(int n) : n(n), SegTlazy(n) {
        timer = 0;
        height = 31 - __builtin_clz(n);
        heavy.assign(n + 1, 0);
        head.assign(n + 1, 0);
    }
    void upd(ll l, ll r, ll v) {
        upd(root, m, l, r, v);
    }
};

```

```

pos.assign(n + 1, 0);
par.assign(n + 1,
    vector<int>(height + 1));
depth.assign(n + 1, 0);
val.assign(n + 1, 0);
g.resize(n + 1);

void add_edge(int u, int v, int w) {
    g[u].push_back({v, w});
    g[v].push_back({u, w});
}

int dfs(int u = 1, int p = 0) {
    int sub = 1, big = 0;
    par[u][0] = p;
    for(int j = 1; j <= height; j++) {
        par[u][j] = par[par[u][j - 1]][j - 1];
    }
    for (auto [v, w] : g[u]) if(v != p) {
        depth[v] = depth[u] + 1;
        val[v] = w;
        int subsize = dfs(v, u);
        if (subsize > big) big =
            subsize, heavy[u] = v;
        sub += subsize;
    }
    return sub;
}
void decompose(int u, int h) {
    head[u] = h,
    pos[u] = ++timer;
    if (heavy[u])
        decompose(heavy[u], h);
    for (auto [v, w] : g[u]) {
        if (v != par[u][0] && v != heavy[u]) decompose(v, v);
    }
}
void makeHLD(int root = 1) {
    dfs(root);
    decompose(root, root);
    for(int i = 1; i <= n; i++) {
        SegTlazy::a[pos[i]] = val[i];
    }
    SegTlazy::build(1, 1, n);
}
// if value on edge then call isEdge
= 1
11 Query(int u, int v, bool isEdge =
1) {
    node ret;
    for (; head[u] != head[v]; v =
        par[head[v]][0]) {
        if (depth[head[u]] >
            depth[head[v]]) swap(u,
            v);
        node tmp =
            SegTlazy::query(1, 1, n,
                pos[head[v]], pos[v]);
        SegTlazy::merge(ret, ret,
            tmp);
    }
    if (isEdge)
        SegTlazy::query(1, 1, n,
            pos[head[v]], pos[v]);
    SegTlazy::merge(ret, ret,
        tmp);
}

```

```

    }
    if (depth[u] > depth[v]) swap(u,
        v);
    node tmp = SegTlazy::query(1, 1,
        n, pos[u] + isEdge, pos[v]);
    SegTlazy::merge(ret, ret, tmp);
    return ret.val;
}
void Update(int u, int v, int val,
bool isEdge = 0) {
    for (; head[u] != head[v]; v =
        par[head[v]][0]) {
        if (depth[head[u]] >
            depth[head[v]]) swap(u,
            v);
        // cout<<"Updating:"<<v<<
        //<<head[v]<<'<<val<<endl;
        SegTlazy::upd(1, 1, n,
            pos[head[v]], pos[v],
            val);
    }
    if (depth[u] > depth[v]) swap(u,
        v);
    //cout<<"Updating:"<<v<<
    //<<val<<endl;
    SegTlazy::upd(1, 1, n, pos[u] +
        isEdge, pos[v], val);
}
void solve() {
    int n;
    cin>>n;
    HLD hld(n);
    // for vertex cin >> v[i]
    vector<pair<int, int>> edge;
    for(int i=1;i<n;i++){
        int u, v, w;
        cin >> u >> v >> w;
        hld.add_edge(u, v, w);
        edge.push_back({u, v});
    }
    hld.makeHLD();
    int q; cin >> q;
    while(q--){
        int ty; cin >> ty;
        if(ty == 1) {
            int id, val;
            cin >> id >> val;
            auto [u, v] = edge[id - 1];
            hld.Update(u, v, val, 1);
        }
        else {
            int u, v;
            cin >> u >> v;
            cout << hld.Query(u, v, 1) <<
                "\n";
        }
    }
}

```

11.8 mex using trie in logmax [26 lines] - 8195ee80bc

```

struct node{
    node *ch[2]; int cnt;
}

```

```

node() {ch[0] = ch[1] = NULL, cnt =
0; }
} *root;
void insert(int x) {
    node *curr = root;
    for(int i = 20; i >= 0; i--) {
        int bit = (x >> i) & 1;
        if(curr->ch[bit] == NULL)
            curr->ch[bit] = new
                node();
        curr = curr->ch[bit];
    }
    int mex() {
        node *curr = root; int ret = 0;
        for(int i = 20; i >= 0; i--) {
            if(curr == NULL || curr->ch[0]
                == NULL) return ret;
            if(curr->ch[0] > cnt >= (1 <<
                i)) {
                curr = curr->ch[1];
                ret |= (1 << i);
            } else curr = curr->ch[0];
        }
        return ret;
    }
}

```

### 11.9 mo with update [76 lines] - 6f0115f479

```

const ll N = 1e6 + 9;
const ll B = 1000;

struct query {
    ll l, r, t, id;
    bool operator < (const query &x)
        const {
        if(l / B == x.l / B) {
            if(r / B == x.r / B) return t
                < x.t;
            return r / B < x.r / B;
        }
        return l / B < x.l / B;
    }
} Q[N];
struct upd {
    ll pos, old, cur;
} U[N];

ll a[N];
ll cnt[N], ans[N], l, r, t;
ll tot;
inline void add(int x) {
    ++cnt[x];
    if(cnt[x]==1)tot++;
}
inline void del(int x) {
    --cnt[x];
    if(cnt[x]==0)tot--;
}
inline void update(int pos, int x) {
    if (l <= pos && pos <= r) {
        add(x);
        del(a[pos]);
    }
    a[pos] = x;
}

```

```

}
map<ll, ll> mp;
ll nxt = 0;
ll get(ll x) {
    return mp.count(x) ? mp[x] : mp[x] =
        ++nxt;
}
void solve(){
    ll n, q;
    cin >> n >> q;
    for (ll i = 1; i <= n; i++) {
        cin >> a[i];
        a[i] = get(a[i]);
    }
    ll nq = 0, nu = 0;
    for (ll i = 1; i <= q; i++) {
        char ty;ll l, r;
        cin >> ty >> l >> r;
        if (ty == 'Q') {l++;++nq, Q[nq]
            = {l, r, nu, nq};}
        else{l++;++nu, U[nu].pos = 1,
            U[nu].old = a[l], a[l] =
            get(r), U[nu].cur = a[l];}
    }
    sort(Q + 1, Q + nq + 1);
    t = nu, l = 1, r = 0;
    for (ll i = 1; i <= nq; i++) {
        ll L = Q[i].l, R = Q[i].r, T =
        Q[i].t;
        while(t < T){ t++,
            update(U[t].pos, U[t].cur);}
        while(t > T) {update(U[t].pos,
            U[t].old), t--;}
        if(R < l) {
            while(l > L) add(a[--l]);
            while(l < L) del(a[l++]);
            while(r < R) add(a[++r]);
            while(r > R) del(a[r--]);
        } else {
            while(r < R) add(a[++r]);
            while(r > R) del(a[r--]);
            while(l > L) add(a[-l]);
            while(l < L) del(a[l++]);
        }
        ans[Q[i].id] = tot;
    }
    for (ll i = 1; i <= nq; i++) cout <<
        ans[i] << '\n';
}

```

### 11.10 pbds [23 lines] - c8bf03595b

```

#include<bits/stdc++.h>
#include<ext/pb_ds/assoc_container.hpp>
#include<ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
using namespace std;
template <typename T> using o_set =
    tree<T, null_type,
    less_equal<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
template <typename T, typename R> using
o_map = tree<T, R,
less<T>, rb_tree_tag,
tree_order_statistics_node_update>;

```

```

int main() {
    o_set<int>se;
    se.insert(1);
    se.insert(2);
    cout << *se.find_by_order(0) << endl;
    //k th element
    cout << se.order_of_key(2) << endl;
    //number of elements less than k
    o_map<int, int>mp;
    mp.insert({1, 10});
    mp.insert({2, 20});
    cout << mp.find_by_order(0)->second <<
        endl;
    ///k th element
    cout << mp.order_of_key(2) << endl;
    //number of first elements less than k
}

11.11 persistance seg tree [68 lines] -
c7856e9a5a


---


class SegTlazy {
public:
    struct node {
        int cursum, premn, premx;
        node *left, *right;
        node(int _cur = 0, int mn = 1e9,
              int mx = -1e9)
            : cursum(_cur), premn(mn),
              premx(mx),
              left(nullptr),
              right(nullptr) {}
    };
    vector<int> a;
    vector<node*> roots; // persistent
                           // versions
    int n;
    SegTlazy() {
        n = 0;
        a.clear();
        roots.clear();
    }
    void init(int _n) {
        n = _n;
        a.assign(n + 1, 1); // 1-based
                            // initially all 1
        roots.clear();
    }
    void merge(node* nd, node* l, node*
              r) {
        nd->cursum = l->cursum +
                    r->cursum;
        nd->premn = min(l->premn,
                          l->cursum + r->premn);
        nd->premx = max(l->premx,
                          l->cursum + r->premx);
    }
    node* build(int b, int e) {
        if (b == e) {
            return new node(a[b], a[b],
                            a[b]);
        }
        int mid = (b + e) >> 1;
        node* cur = new node();
        cur->left = build(b, mid);
        cur->right = build(mid + 1, e);
    }
}

```

```

    merge(cur, cur->left,
          cur->right);
    return cur;
}
node* update(node* prev, int b, int
             e, int pos, int val) {
    if (b == e) {
        return new node(val, val,
                        val);
    }
    int mid = (b + e) >> 1;
    node* cur = new node();
    if (pos <= mid) {
        cur->left =
            update(prev->left, b,
                   mid, pos, val);
        cur->right = prev->right;
    } else {
        cur->left = prev->left;
        cur->right =
            update(prev->right, mid +
                  1, e, pos, val);
    }
    merge(cur, cur->left,
          cur->right);
    return cur;
}
node query(node* nd, int b, int e,
           int i, int j) {
    if (!nd || j < b || e < i) {
        return node(0, 1e9, -1e9);
    }
    if (i <= b && e <= j) {
        return *nd;
    }
    int mid = (b + e) >> 1;
    node l = query(nd->left, b, mid,
                   i, j);
    node r = query(nd->right, mid +
                   1, e, i, j);
    node res;
    merge(&res, &l, &r);
    return res;
}


---


11.12 segtree [124 lines] - df44add3e


---


#include <bits/stdc++.h>
using namespace std;
using ll = long long;

struct LazySegTree {
    vector<ll> tree, lazyAdd, assignVal;
    vector<char> assignFlag; // 0 = no
                             // assign pending, 1 = assign
                             // pending
    ll n;

    LazySegTree(ll arr[], ll sz) {
        n = sz;
        tree.assign(4*n, 0);
        lazyAdd.assign(4*n, 0);
        assignVal.assign(4*n, 0);
        assignFlag.assign(4*n, 0);
        build(arr, 0, 0, n-1);
    }
}

```

```

void build(ll arr[], ll ind, ll st,
           ll ed) {
    if (st == ed) { tree[ind] =
        arr[st]; return; }
    ll mid = (st + ed) / 2;
    build(arr, 2*ind+1, st, mid);
    build(arr, 2*ind+2, mid+1, ed);
    tree[ind] = tree[2*ind+1] +
        tree[2*ind+2];
}

// Push pending operations at node
// ind to children and apply to
// node's tree
void push(ll ind, ll st, ll ed) {
    if (st > ed) return;
    // First handle assignment (it
    // overrides additions)
    if (assignFlag[ind]) {
        tree[ind] = assignVal[ind] *
            (ed - st + 1);
        if (st != ed) {
            ll L = 2*ind+1, R =
                2*ind+2;
            // set children to this
            // assignment, clear
            // their adds
            assignFlag[L] =
                assignFlag[R] = 1;
            assignVal[L] =
                assignVal[R] =
                    assignVal[ind];
            lazyAdd[L] = lazyAdd[R]
                = 0;
        }
        assignFlag[ind] = 0; // consumed at this node
                            // (tree already updated)
        assignVal[ind] = 0;
    }
    // Then handle addition
    if (lazyAdd[ind] != 0) {
        tree[ind] += lazyAdd[ind] *
            (ed - st + 1);
        if (st != ed) {
            ll L = 2*ind+1, R =
                2*ind+2;
            // if child has assign
            // pending, adding
            // modifies assignVal
            if (assignFlag[L])
                assignVal[L] +=
                    lazyAdd[ind];
            else lazyAdd[L] +=
                lazyAdd[ind];
            if (assignFlag[R])
                assignVal[R] +=
                    lazyAdd[ind];
            else lazyAdd[R] +=
                lazyAdd[ind];
        }
        lazyAdd[ind] = 0;
    }
}

// Range add: add val to every
// element in [l,r]
void updateAdd(ll val, ll ind, ll l,
               ll r, ll st, ll ed) {
    push(ind, st, ed);
    if (r < st || ed < l) return;
    if (l <= st && ed <= r) {
        lazyAdd[ind] += val;
        push(ind, st, ed);
        return;
    }
    ll mid = (st + ed) / 2;
    updateAdd(val, 2*ind+1, l, r,
              st, mid);
    updateAdd(val, 2*ind+2, l, r,
              mid+1, ed);
    tree[ind] = tree[2*ind+1] +
        tree[2*ind+2];
}

// Range assign: set every element in
// [l,r] to val
void updateAssign(ll val, ll ind, ll
                  l, ll r, ll st, ll ed) {
    push(ind, st, ed);
    if (r < st || ed < l) return;
    if (l <= st && ed <= r) {
        // mark assignment on this
        // node and apply
        // immediately
        assignFlag[ind] = 1;
        assignVal[ind] = val;
        lazyAdd[ind] = 0; //
                           // assignment overrides
                           // previous adds
        push(ind, st, ed);
        return;
    }
    ll mid = (st + ed) / 2;
    updateAssign(val, 2*ind+1, l, r,
                 st, mid);
    updateAssign(val, 2*ind+2, l, r,
                 mid+1, ed);
    tree[ind] = tree[2*ind+1] +
        tree[2*ind+2];
}

// Range sum query [l,r]
ll sum(ll ind, ll l, ll r, ll st, ll
       ed) {
    push(ind, st, ed);
    if (r < st || ed < l) return 0;
    if (l <= st && ed <= r) return
        tree[ind];
    ll mid = (st + ed) / 2;
    return sum(2*ind+1, l, r, st,
              mid) + sum(2*ind+2, l, r,
              mid+1, ed);
}

// convenient wrappers:
void rangeAdd(ll l, ll r, ll val) {
    updateAdd(val, 0, l, r, 0, n-1);
}

```

```

void rangeAssign(ll l, ll r, ll val)
{ updateAssign(val, 0, l, r, 0,
n-1); }
ll rangeSum(ll l, ll r) { return
sum(0, l, r, 0, n-1); }

// Example usage
int main() {
ll a[] = {1,2,3,4,5};
ll sz = 5;
LazySegTree seg(a, sz);

// add 10 to range [1,3]
seg.rangeAdd(1,3,10); // array ->
{1,12,13,14,5}

// assign range [2,4] to 7
seg.rangeAssign(2,4,7); // array ->
{1,12,7,7,7}

// add 5 to [0,2]
seg.rangeAdd(0,2,5); // array ->
{6,17,12,7,7}

cout << seg.rangeSum(0,4) << "\n"; // prints 6+17+12+7+7 = 49
cout << seg.rangeSum(1,2) << "\n"; // prints 17+12 = 29
return 0;
}

```

```

11.13 sparse [34 lines] - 60d93b1d24
struct Spares_table { // now 1-based
vector<int> a;
vector<vector<int>> t;
int n, q, lg;
Spares_table() {}
Spares_table(int n, int q) : n(n),
q(q) {
lg = __lg(max(1, n)) + 1;
a.assign(n + 1, 0);
t.assign(n + 2, vector<int>(lg +
1, 0));
take_input();
}
void take_input() {
for (int i = 1; i <= n; i++) cin
>> a[i];
build();
}
void build() {
if (n <= 1) return;
// Build on diffs d[i] = |a[i] -
a[i-1]| for i = 2..n
for (int i = 2; i <= n; i++)
t[i][0] = abs(a[i] - a[i -
1]);
for (int j = 1; (1 << j) <= n;
j++) {
for (int i = 2; i + (1 << j) -
1 <= n; i++) {
t[i][j] = gcd(t[i][j -
1], t[i + (1 << (j -
1))][j - 1]);
}
}
}

```

```

}
// Query on the original array [l..r]
// (1-based).
// Returns GCD of |a[i] - a[i-1]| for
i in (l..r), and 0 if l == r.
int query(int l, int r) {
if (l >= r) return 0;
int len = r - l;
int k = __lg(len);
return gcd(t[l + 1][k], t[r - (1
<< k) + 1][k]);
}
}

11.14 trie [190 lines] - 8b1049f1f2
/*
if you need dfs on trie
at first make the trie an unique tree
inserting an unique id on each node
then just treat as node of a tree
*/
class Trie {
public:
class node { // 1 based
public:
node *child[27];
int leaf, sz, id;
node() {
for(int i = 0; i < 27; i++) {
child[i] = nullptr;
}
sz = 0;
leaf = 0;
id = 0;
}
} *root;
vector<int> dis;
int unqid = 1;
Trie () {
root = new node();
}
void insert(string s) {
auto cur = root;
for(auto x : s) {
int i = x - 'a' + 1; // 1-based
if(!cur->child[i]) {
cur->child[i] = new
node(); // edge dilam
}
cur = cur->child[i];
cur->sz++; // new edge e
ashalam ebar increment
korbo
cur->id = unqid++;
}
cur->leaf++;
}
int set(node *cur) {

```

```

dis[cur->id] = cur->leaf;
for(int i = 1; i <= 26; i++) {
if(cur->child[i]) {
set(cur->child[i]);
dis[cur->id] =
max(dis[cur->id],
dis[cur->child[i]->_id]);
}
}
return dis[cur->id];
}
};

Trie mytrie;
for(int i = 1; i <= n; i++) {
string t; cin >> t;
mytrie.insert(t);
}
mytrie.dis.assign(mytrie.unqid + 1, 0);
mytrie.set(mytrie.root);
/*
count how many subarray Xor < k
*/
struct Trie {
const int B = 20;
struct node { // 1 based
node *child[2];
int cnt;
node() {
child[0] = child[1] = 0;
cnt = 0;
}
} *root;
Trie () {
root = new node();
}
void insert(int x) {
auto cur = root;
for(int i = B - 1; i >= 0; i--) {
int id = (x >> i) & 1;
if(!cur->child[id])
cur->child[id] = new
node();
cur = cur->child[id];
cur->cnt++;
}
}
ll countLess(int y, int k) { // to
count greater n - countless
auto cur = root;
ll ans = 0;
for(int i = B - 1; i >= 0 &&
cur; i--) {
int yb = (y >> i) & 1;
int kb = (k >> i) & 1;
if(kb == 1) {
if(cur->child[yb]) {
ans += cur->_
child[yb]->cnt;
}
cur = cur->child[!yb];
}
}
cur = cur->child[!yb];
}

```

```

else {
cur = cur->child[yb];
}
}
return ans;
}
int occurence(int x) {
auto cur = root;
for(int i = B - 1; i >= 0; i--) {
int id = (x >> i) & 1;
if(!cur->child[id]) return 0;
cur = cur->child[id];
}
return cur->leaf;
}

void del(int x) {
stack<pair<node *, int>> stck;
auto cur = root;
for(int i = B - 1; i >= 0; i--) {
int id = (x >> i) & 1;
stck.push({cur, id});
cur = cur->child[id];
}
while(not stck.empty()) {
auto [par, id] = stck.top();
stck.pop();
auto child = par->child[id];
if(!child->child[0] &&
!child->child[1]) {
delete child;
par->child[id] = nullptr;
}
else {
break;
}
}
}

void remove(int x) {
if(occurence(x) == 1) {
del(x);
return;
}
auto cur = root;
for(int i = B - 1; i >= 0; i--) {
int id = (x >> i) & 1;
cur = cur->child[id];
}
cur->leaf--;
}

int get_max(int x) {
auto cur = root;
int ans = 0;
for(int i = B - 1; i >= 0; i--) {
int id = (x >> i) & 1;
if(cur->child[!id]) {
//cout << i << " " << id << "
" << sz << "\n";
ans += (1 << i);
cur = cur->child[!id];
}
else if(cur->child[id]){
cur = cur->child[id];
}
}
}

```

```

    }
    else break;
}

void clear(node *cur) {
    if(!cur) return;
    for(int i = 0; i < 2; i++) {
        if(cur->child[i])
            clear(cur->child[i]);
    }
    delete(cur);
}
~Trie() {
    clear(root);
}

11 ans = 0;
int pref = 0;
Trie trie;
trie.insert(0);
for(auto x : v) {
    pref ^= x;
    ans += trie.countLess(pref, k);
    trie.insert(pref);
}

```

### 11.15 wavelet tree [77 lines] - 90f623bba4

```

// Wavelet Tree structure used for
// answering range queries efficiently
struct wavelet_tree {
    int lo, hi; // range of
    // values represented at this node
    wavelet_tree *l, *r; // left and
    right children
    int *b; // b: bitmap of
    counts going left, c: prefix sums
    int bsz, csz; // sizes of b
    and c arrays

    // Constructor: initializes members
    wavelet_tree() : lo(1), hi(0),
    l(nullptr), r(nullptr),
    b(nullptr), bsz(0),
    c(nullptr),
    csz(0) {}

    // Initializes wavelet tree on range
    // [from, to) with values in [x, y]
    void init(int *from, int *to, int x,
    int y) {
        lo = x, hi = y;
        if(from >= to) return;

        int mid = (lo + hi) >> 1;
        auto f = [mid](int x) {
            return x <= mid;
        };

        int n = to - from;
        b = (int*)malloc((n + 2) *
        sizeof(int));
        c = (ll*)malloc((n + 2) *
        sizeof(ll));
        bsz = csz = 0;
        b[bsz++] = 0;
    }
}
```

```

c[csz++] = 0;

for(auto it = from; it != to;
    it++) {
    b[bsz] = b[bsz - 1] +
        f(*it); // Count of
    // elements ≤ mid
    c[csz] = c[csz - 1] + (*it);
    // Prefix sum of elements
    bsz++;
    csz++;
}

if(hi == lo) return;

auto pivot =
stable_partition(from, to,
f); // Separate elements
l = new wavelet_tree();
l->init(from, pivot, lo,
mid); // Left subtree
r = new wavelet_tree();
r->init(pivot, to, mid + 1,
hi); // Right subtree
}

// Returns count of numbers ≤ k in
// range [l, r]
int LTE(int l, int r, int k) {
    if(l > r || k < lo) return 0;
    if(hi <= k) return r - l + 1;

    int lb = b[l - 1], rb = b[r];
    return this->l->LTE(lb + 1, rb,
    k) + this->r->LTE(r - lb, r -
    rb, k);
}

// Returns sum of elements ≤ k in
// range [l, r]
int sum(int l, int r, int k) {
    if(l > r || k < lo) return 0;
    if(hi <= k) return c[r] - c[l - 1];

    int lb = b[l - 1], rb = b[r];
    return this->l->sum(lb + 1, rb,
    k) + this->r->sum(r - lb, r -
    rb, k);
}

//kth smallest element in [l, r]
//for array [1,2,1,3,5] 2nd smallest
//is 1 and 3rd smallest is 2
int kth(int l, int r, int k) {
    if(l > r) return 0;
    if(lo == hi) return lo;
    int inLeft = b[r] - b[l - 1], lb
    = b[l - 1], rb = b[r];
    if(k <= inLeft) return
        this->l->kth(lb + 1, rb, k);
    return this->r->kth(r - lb, r -
    rb, k - inLeft);
}

// Destructor to free memory
~wavelet_tree() {

```

```

    delete l;
    delete r;
    if(b) free(b);
    if(c) free(c);
}

```

## 12 Random

### 12.1 Combinatorics

$$\bullet \sum_{k=0}^n \binom{n-k}{k} = Fib_{n+1}$$

$$\bullet \binom{n}{k} + \binom{n}{k+1} = \binom{n+1}{k+1}$$

$$\bullet k \binom{n}{k} = n \binom{n-1}{k-1}$$

• Number of binary sequences of length n such that no two 0's are adjacent =  $Fib_{n+1}$

• Number of non-negative solution of  $x_1 + x_2 + x_3 + \dots + x_k = n$  is  $\binom{n+k-1}{n}$

#### 12.1.1 Catalan Number

$$\bullet C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$\bullet C_0 = 1, C_1 = 1, C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$$

$$\bullet 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786$$

• Number of correct bracket sequences consisting of n opening brackets.

• Number of ways to completely parenthesize  $n+1$  factors.

• The number of triangulations of a convex polygon with  $+2$  sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).

• The number of ways to connect the  $2n$  points on a circle to form  $n$  disjoint i.e. non-intersecting chords.

• The number of monotonic lattice paths from point  $(0,0)$  to point  $(n,n)$  in a square lattice of size  $n \times n$ , which do not pass above the main diagonal

• Number of permutation of length n that can be stack sorted.

- The number of non-crossing partitions of a set of n elements.

- The number of rooted full binary tree with  $n+1$  leaves.

- The number of Dyck words of length  $2n$ . A string consisting of n X's and n Y's such that no string prefix has more Y's than X's.

- Number of permutation of length n with no three-term increasing subsequence.

- Number of ways to tile a staircase shape of height n with n rectangle.

- $C_n^k = \frac{k+1}{n+1} \binom{2n-k}{n-k}$  denote the number of bracket sequences of size  $2n$  with the first k elements being (.

$$\bullet N(n, k) = \frac{1}{n} \binom{n}{k} \binom{n}{k-1}$$

- The number of expressions containing n pairs of correct parentheses, which contain k distinct nestings.  $N(4, 2) = 6$

- The number of paths from  $(0,0)$  to  $(2n, 0)$  with steps only northeast and southeast, not staying below the x-axis with k peaks. And sum of all number of peaks is Catalan number.

#### 12.1.2 Stirling Number of the First Kind

- Count permutation according to their number of cycles.

- $S(n, k)$  count the number of permutation of n elements with k disjoint cycles.

$$\bullet S(n, k) = (n-1) \times S(n-1, k) + S(n-1, k-1), S(0, 0) = 1, S(n, 0) = S(0, n) = 0$$

$$\bullet S(n, 1) = (n-1)!$$

$$\bullet S(n, n-1) = \binom{n}{2}$$

$$\bullet \sum_{k=0}^n S(n, k) = n!$$

### 12.1.3 Stirling Numbers of the Second Kind

- Number of ways to partition a set of n objects into k non-empty subsets.
- $S(n, k) = k * S(n-1, k) + S(n-1, k-1)$ ,  $S(0, 0) = 1$ ,  $S(n, 0) = S(0, n) = 0$
- $S(n, 2) = 2^{n-1} - 1$
- $S(n, k) = \frac{1}{k!} \sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^n$
- $S(n, k) * k!$  = number of ways to color n nodes using colors from 1 to k such that each color is used at least once.

### 12.1.4 Bell Number

- Counts the number of partitions of a set.
- $B_{n+1} = \sum_{k=0}^n \binom{n}{k} * B_k$
- $B_n = \sum_{k=0}^n S(n, k)$ , where S is Stirling number of second kind.
- The number of multiplicative partitions of a square free number with i prime factors is the i-th Bell number.
- $B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$
- If a deck is shuffled by removing and reinserting the top card n times, there are  $n^n$  possible shuffles. The number of shuffles that return the deck to its original order is  $B_n$ , so the probability of returning to the original order is  $B_n/n^n$ .

### 12.1.5 Lucas Theorem

- If p is prime then  $\binom{p^a}{k} \equiv 0 \pmod{p}$
- For non-negative integers m and n and a prime p:  

$$\binom{m}{n} = \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p}$$
 where  
 $m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0$   
 $n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$   
are the base p expansion.

### 12.1.6 Derangement

- A permutation such that no element appears in its original position.
- $d(n) = (n-1) * (d(n-1) + d(n-2))$ ,  $d(0) = 1$ ,  $d(1) = 0$
- $d(n) = nd(n-1) + (-1)^n = \lfloor \frac{n!}{e} \rfloor$ ,  $n \geq 1$

### 12.1.7 Burnside Lemma

Given a group G of symmetries and a set X, the number of elements of X up to symmetry equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|$$

where  $X^g$  are the elements fixed by g ( $g.x = x$ ) If f(n) counts "configurations" of some sort of length n, we can ignore rotational symmetry using  $G = \mathbb{Z}_n$  to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n, k)) = \frac{1}{n} \sum_{k|n} f(k) \phi(n/k)$$

### 12.1.8 Eulerian Number

- $E(n, k)$  is the number of permutations of the numbers 1 to n in which exactly k elements are greater than the previous element.
- $E(n, k) = (n-k)E(n-1, k-1) + (k+1)E(n-1, k)$ ,  $E(n, 0) = E(n, n-1) = 1$
- $E(n, k) = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n$
- $E(n, k) = E(n, n-1-k)$
- $E(0, k) = [k=0]$
- $E(n, 1) = 2^n - n - 1$

## 12.2 Number Theory

### 12.2.1 Möbius Function and Inversion

- define  $\mu(n)$  as the sum of the primitive nth roots of unity depending on the factorization of n into prime factors:

$$\mu(x) = \begin{cases} 0 & \text{n is not square free} \\ 1 & \text{n has even number of prime factors} \\ -1 & \text{n has odd number of prime factors} \end{cases}$$

- Möbius Inversion:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu(d) g(n/d)$$

- $\sum_{d|n} \mu(d) = [n=1]$

- $\phi(n) = \sum_{d|n} \mu(d) \cdot \frac{n}{d} = n \sum_{d|n} \frac{\mu(d)}{d} = \sum_{d|n} d \cdot \mu\left(\frac{n}{d}\right)$

- $a|b \rightarrow \phi(a)|\phi(b)$

- $\phi(mn) = \phi(m) \cdot \phi(n) \cdot \frac{d}{\phi(d)}$  where  $d = \gcd(m, n)$

- $\phi(n^m) = n^{m-1} \phi(n)$

- $\sum_{i=1}^n [\gcd(i, n) = k] = \phi\left(\frac{n}{k}\right)$

- $\sum_{i=1}^n \gcd(i, n) = \sum_{d|n} d \cdot \phi\left(\frac{n}{d}\right)$

- $\sum_{i=1}^n \frac{1}{\gcd(i, n)} = \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$

- $\sum_{i=1}^n \frac{i}{\gcd(i, n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$

- $\sum_{i=1}^n \frac{n}{\gcd(i, n)} = 2 \cdot \sum_{i=1}^n \frac{i}{\gcd(i, n)} - 1$

### 12.2.2 GCD and LCM

- $\gcd(a, b) = \gcd(b, a \bmod b)$
- If  $a|b.c$ , and  $\gcd(a, b) = d$ , then  $(a/d)|c$ .
- GCD is a multiplicative function.
- $\gcd(a, \text{lcm}(b, c)) = \text{lcm}(\gcd(a, b), \gcd(a, c))$
- $\gcd(n^a - 1, n^b - 1) = n^{\gcd(a, b)} - 1$

### 12.2.3 Gauss Circle Theorem

- Determine the number of lattice points in a circle centered at the origin with radius r.
- number of pairs (m,n) such that  $m^2 + n^2 \leq r^2$
- $N(r) = 1 + 4 \sum_{i=0}^{\infty} (\lfloor \frac{r^2}{4i+1} \rfloor - \lfloor \frac{r^2}{4i+3} \rfloor)$

### 12.2.4 Pick's Theorem

According to Pick's Theorem We can calculate the area of any polygon by just counting the number of Interior and Boundary lattice points of that polygon. If number of interior points are I and number of boundary lattice points are B then Area (A) of polygon will be:

$$Area = I + B/2 - 1$$

where I is the number of points in the interior shape, B stands for the number of points on the boundary of the shape.

### 12.2.5 Formula Cheatsheet

- $\sum_{i=1}^n = \frac{1}{m+1} [(n+1)^{m+1} - 1 - \sum_{i=1}^n ((i+1)^{m+1} - i^{m+1} - (m+1)i^m)]$
- $\sum_{i=0}^n c^i = \frac{c^{n+1}-1}{c-1}, c \neq 1$
- $\sum_{i=0}^{\infty} c^i = \frac{1}{1-c}, \sum_{i=1}^{\infty} c^i = \frac{c}{1-c}, |c| < 1$
- $H_n = \sum_{i=1}^n \frac{1}{i}, \sum_{i=1}^n i H_i = \frac{n(n+1)}{2} H_n - \frac{n(n-1)}{4}$
- $\sum_{k=0}^n \binom{r+k}{k} = \binom{r+n+1}{n}$