# Distributed File System: – File Service Architecture

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- The sharing of resources is a key goal for distributed systems.
- The sharing of <u>stored information</u> is perhaps the most important aspect of distributed resource sharing.

➤ A distributed file system enables programs to store and access remote files exactly as they do local ones.

➤ Allowing users to access files from any computer on a network.

- Web servers provide a restricted form of data sharing in which
  - Files stored locally in file systems at the server or in servers on a local network, are made available to clients throughout the Internet.

 The <u>performance and reliability</u> experienced for access to files stored at a server should be comparable to that for files stored on local disks.

## Distributed File System - Challenges

- The design of large-scale wide area readwrite file storage systems poses problems of:
  - Load balancing,
  - Reliability,
  - Availability and
  - Security,

# Distributed File System - Challenges

- A file service enables programs to store and access remote files exactly as they do local ones,
  - ➤ Allowing users to access their files from any computer in an intranet.

- File systems are responsible for the organization, storage, retrieval, naming, sharing and protection of files.
- They provide a programming interface that characterizes the file abstraction, freeing programmers from concern with the details of storage allocation and layout.

- Files are stored on disks or other non-volatile storage media.
- > Files contain both data and attributes.
- The data consist of a sequence of data items (typically 8-bit bytes), accessible by operations to read and write any portion of the sequence.

- ➤ The attributes are held as a single record containing information such as:
- The length of the file,
- Timestamps,
- File type,
- Owner's identity and
- Access control lists.

File attribute record structure

File length
Creation timestamp
Read timestamp
Write timestamp
Attribute timestamp
Reference count
Owner
File type
Access control list

A typical attribute record structure is illustrated in Figure.

The shaded attributes are managed by the file system and are not normally updatable by user programs.

- ➤ File systems are designed to store and manage large numbers of files, with facilities for creating, naming and deleting files.
- ➤ The naming of files is supported by the use of directories.

- Directories may include the names of other directories.
- Leading to the familiar hierarchic filenaming scheme and the multi-part pathnames for files used in UNIX and other operating systems.

- > File systems also take responsibility for:
  - 1. The control of access to files,
  - 2.Restricting access to files according to users' authorizations and
  - 3. The type of access requested (reading, updating, executing and so on).

- ➤ The main operations on files(create, open, close, read, and write etc...) that are available to applications in UNIX systems.
- > These are the system calls implemented by the kernel.

- The file system is responsible for applying access control for files.
- ➤ In local file systems such as UNIX, it does so when each file is opened,
- Checking the rights allowed for the user's identity in the access control list against the mode of access requested in the open system call.

- > If the rights match the mode,
  - The file is opened and
  - The mode is recorded in the open file state information.

#### 1. Access transparency: -

- Client programs should be unaware of the distribution of files.
- A single set of operations is provided for access to local and remote files.
- Programs written to operate on local files are able to access remote files without modification.

- 2. Location transparency: Client programs should see a uniform file name space.
- Files or groups of files may be relocated without changing their pathnames, and user programs see the same name space wherever they are executed.

- 3. **Mobility transparency** Neither client programs nor system administration tables in client nodes need to be changed when files are moved.
- This allows file mobility files or, more commonly, sets or volumes of files may be moved, either by system administrators or automatically.

**4. Performance transparency**: - Client programs should continue to perform satisfactorily while the load on the service varies within a specified range.

**5. Scaling transparency**:- The service can be expanded by incremental growth to deal with a wide range of loads and network sizes.

6. Concurrent file updates: Changes to a file by one client should not interfere with the operation of other clients simultaneously accessing or changing the same file.

**7. File replication** In a file service that supports replication, a file may be represented by several copies of its contents at different locations.

File replication has two benefits : -

- it enables multiple servers to share the load of providing a service to clients accessing the same set of files and
- It enhances fault tolerance by enabling clients to locate another server that holds a copy of the file when one has failed.

8. Hardware and operating system
heterogeneity - The service interfaces should
be defined so that client and server software can
be implemented for different operating systems

- This requirement is an important aspect of openness.

and computers.

**9.Fault tolerance** - The central role of the file service in distributed systems makes it essential that the service continue to operate in the face of client and server failures.

**10. Consistency** - Conventional file systems such as that provided in UNIX offer one-copy update semantics.

#### One-copy update semantics

- ➤ When files are replicated or cached at different sites,
- ➤ There is an inevitable delay in the propagation of modifications made at one site to all of the other sites that hold copies,
- ➤ and this may result in some deviation from one\_copy semantics.

- **11. Security** Virtually all file systems provide access-control mechanisms based on the use of access control lists.
- In distributed file systems, there is a need to authenticate client requests.

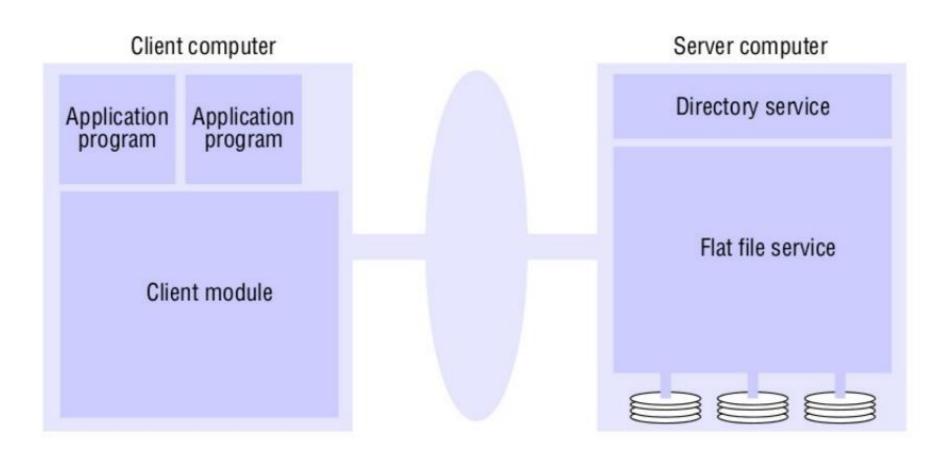
#### Security continues...-

- > So that access control at the server is based on correct user identities.
- ➤ Protect the contents of request and reply messages with digital signatures and (optionally) encryption of secret data.

- **11. Efficiency** A distributed file service should offer facilities that are of at least the same power.
- ➤ Generality as those found in conventional file systems.
- Should achieve a comparable level of performance.

- ➤ An architecture that offers a clear separation of the main concerns in providing access to files is obtained by structuring the file service as three components:-
- Flat file service
- Directory service
- Client module

File service architecture



The *flat file service and the directory service* each export an interface:

- > For use by client programs, and
- ➤ Their Request Procedure Call (RPC) interfaces,
- ➤ Taken together, provide a comprehensive set of operations for access to files.

The *client module* provides a single programming interface with operations on files similar to those found in conventional file systems.

The *design* is open in the sense that :-

- Different client modules can be used to implement different programming interfaces,
- Simulating the file operations of a variety of different operating systems.
- Optimizing the performance for different client and server hardware configurations.

#### Flat file service:

- ➤ The flat file service is concerned with implementing operations on the contents of files.
- ➤ Unique file identifiers (UFIDs) are used to refer to files in all requests for flat file service operations.
- ➤ The division of responsibilities between the file service and the directory service is based upon the use of UFIDs.

#### Flat file service:

- ➤ UFIDs are long sequences of bits and each file has a UFID that is unique among all of the files in a distributed system.
- ➤ When the flat file service receives a request to create a file, it generates a new UFID for it and returns the UFID to the requester.

### Directory service: -

➤ The <u>directory service provides a mapping</u> between text names for files and their UFIDs.

Clients may obtain the UFID of a file by quoting its text name to the directory service.

## Directory service: -

- ➤ The directory service provides the functions needed to generate directories, to add new file names to directories and to obtain UFIDs from directories.
- ➤ <u>It is a client of the flat file service</u>; its directory files are stored in files of the flat file service.

#### Client module: -

- > A client module runs in each client computer,
- ➤ Integrating and extending the operations of the flat file service and the directory service under a single application programming interface that is available to user-level programs in client computers.

#### Client module : -

- ➤ The client module also holds information about the network locations of the flat file server and directory server processes.
- ➤ The client module can play an important role in achieving satisfactory performance through the implementation of a cache of recently used file blocks at the client.

- > This is the RPC interface used by client modules.
- ➤ It is not normally used directly by user-level programs.

The information below contains a definition of the interface to a flat file service.

#### Flat file service operations

$Read(FileId, i, n) \rightarrow Data$ — throws $BadPosition$	If $1 \le i \le Length(File)$ : Reads a sequence of up to $n$ items from a file starting at item $i$ and returns it in $Data$ .
Write(FileId, i, Data) — throws BadPosition	If $1 \le i \le Length(File) + 1$ : Writes a sequence of <i>Data</i> to a file, starting at item $i$ , extending the file if necessary.
$Create() \rightarrow FileId$	Creates a new file of length 0 and delivers a UFID for it.
Delete(FileId)	Removes the file from the file store.
$GetAttributes(FileId) \rightarrow Attr$	Returns the file attributes for the file.
SetAttributes(FileId, Attr)	Sets the file attributes (only those attributes that are not shaded in Figure 12.3).

- ➤ The most important operations are those for reading and writing.
- ➤ Both the Read and the Write operation require a parameter 'i' specifying a position in the file.

#### Flat file service interface

➤ The Read operation copies the sequence of 'n' data items beginning at item 'i' from the specified file into Data, which is then returned to the client.

- ➤ The Write operation copies the sequence of data items in 'Data' into the specified file beginning at item 'i',
- ➤ Replacing the previous contents of the file at the corresponding position and extending the file if necessary.

- ➤ Create creates a new, empty file and returns the UFID that is generated.
- > Delete removes the specified file.

- ➤ GetAttributes and SetAttributes enable clients to access the attribute record.
- ➤ GetAttributes is normally available to any client that is allowed to read the file.
- ➤ Access to the SetAttributes operation would normally be restricted to the directory service that provides access to the file.

# THANK YOU!