AVR Microcontroller

Microprocessor Course

Chapter 6

AVR ADVANCED ASSEMBLY LANGUAGE PROGRAMMING

Aban 1401

Arithmetic and logic expressions with constant values

The AVR Studio IDE supports logic operations between expressions as well.

Table 6-1: Arithmetic Operators

Symbol	Action	
+	Addition	
_	Subtraction	
*	Multiplication	
/	Division	
%	Modulo	

For example, in the following $\frac{}{\sqrt{}}$ program R21 is loaded with 0x14:

```
.EQU C1 = 0x50

.EQU C2 = 0x10

.EQU C3 = 0x04

LDI R21, (C1&C2) | C3 ; R21=(0x10&0x50) | 0x04 = 0x10 | 0x04 = 0x14
```

In Table 6-3 you see the shift operators, which are very useful. They shift left and right a constant value.

Table 6-3: Shift Operators

Symbol	Action	Example	_
<<	Shifts left the left expression by the number of places given by the right expression	LDI R20,0b101<<2 ;R20=0b10100)
>>	Shifts right the left expression by the number of places given by the right expression	LDI R20,0b100>>1 ;R20=0b010	_

For example, the following instruction loads the R20 register with 0b00001110:

LDI R16,0b00000111<<1 ; R16 =
$$0$$
b00001110

One of the uses of shift operators is for initializing the registers. For example suppose we want to set the Z and C bits of the SREG (Status Register) register and clear the others. Look at Figure 6-1.

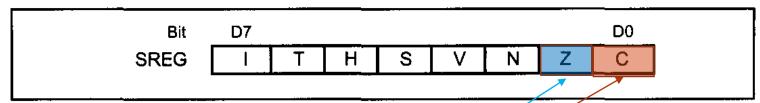


Figure 6-1. Bits of the Status Register

If we load 0b00000011 to SREG the task will be done:

LDI R20,
$$0b00000011$$
; $Z = 1$, $C = 1$
OUT SREG, R20

In this example, we calculated the 0b00000011 number by looking at Figure 6-1.

But imagine you are writing a program and you want to do the same task; you have to open the datasheet or a reference book to see the structure of the SREG register. To make the task simpler, the names of the register bits are defined in the header files of each AVR microcontroller. For example, in M32DEF.INC there are the following lines of code:

```
; SREG - Status Register
.equ SREG_C = 0 ; carry flag
.equ SREG_Z = 1 ; zero flag
.equ SREG_N = 2 ; negative flag
.equ SREG_V = 3 ; 2's complement overflow flag
.equ SREG_S = 4 ; sign bit
.equ SREG_H = 5 ; half carry flag
.equ SREG_T = 6 ; bit copy storage
.equ SREG_I = 7 ; global interrupt enable
```

So, we can use the names of the bits instead of remembering the structure of the registers or finding them in the datasheet. For example, the following program sets the Z flag of the SREG register and clears the other bits:

```
LDI R16, 1<<SREG_Z ;R16= 1 << 1 = 0b00000010
OUT SREG,R16 ;SREG = 0b00000010 (set Z and clear others)
```

As another example, the following program sets the V and S flags of SREG:

```
LDI R16, (1<<SREG_V) | (1<<SREG_S) ;R16=0b1000|0b10000=0b11000

OUT SREG,R16 ;SREG = 0b00011000 (set V and S, clear others)</pre>
```

In Example 6-1, you see the usage of the directives in I/O port programming.

Example 6-1

Write codes to set PB2 and PB4 of PORTB to 1 and clear the other pins

- (a) without the directives, and
- (b) using the directives.

Solution:

```
(a) LDI R20,0x14 ; R20 = 0x14
OUT PORTB, R20 ; PORTB = R20
```

To make the code more readable, we can write the number in binary as well:

```
LDI R20,0b00010100 ; R20 = 0x14
OUT PORTB, R20 ; PORTB = 0x14
```

```
(b) LDI R20, (1<<4) | (1<<2) ; R20 = (0b10000 | 0b00100) = 0b10100
OUT PORTB, R20 ; PORTB = R20
```

As we mentioned before, the names of the register bits are defined in the header files of each AVR microcontroller. PB2 and PB4 are defined equal to 2 and 4, as well. Therefore, we can write the code as shown below:

```
LDI R20, (1 << PB4) \mid (1 << PB2); set the PB4 and PB2 bits OUT PORTB, R20; PORTB = R20
```

Notice that when the assembler wants to convert a code to machine language it substitutes all of the assembler directives with their equivalent values. Thus, using the directives has no side effects on the performance of our code but rather makes our code more readable.

Example 6-2

What does the AVR assembler do while assembling the following program?

Solution:

.equ is an assembler directive. When assembling ".equ C1 = 2", the assembler assigns value 2 to C1. Similarly, while assembling the ".equ C2 = 3" instruction, it assigns the value 3 to C2.

When the assembler converts the "LDI R20,C1|(1<<C2)" instruction to machine language, it knows the values of C1 and C2. Thus it calculates the value of "C1|(1<<C2)", and then replaces the expression with its value. Therefore, "LDI R20,C1|(1<<C2)" will be converted to "LDI R20,0b00001010". Then the assembler converts the instruction to machine language.

Example 6-3

What does the AVR assembler do while assembling the following program?

```
.INCLUDE "M32DEF.INC"
```

LDI R20, $(1 << PB4) \mid (1 << PB2)$; set the PB4 and PB2 bits

OUT DDRB, R20 ; DDRB = R20

HERE: RJMP HERE

Solution:

Including a header file at the beginning of a program is similar to copying all the contents of the header file to the beginning of the program. Thus, the assembler, first assembles the contents of M32DEF.INC. The header file contains some ".equ" instructions, such as ".equ PB4 = 4". Thus, after reading the header file the assembler learns that PB4 is equal to 4, PB2 is equal to 2, and so on. Thus, when it wants to assemble instructions such as "LDI R20, (1<<PB4) | (1<<PB2)", it knows the values of PB2 and PB4. It calculates the value of "(1<<PB4) | (1<<PB2)" and substitutes it.

HIGH() and LOW() functions

The HIGH() and LOW() functions give the higher and the lower bytes of a 16-bit value. For example, in the following program 0x55 and 0x44 are

loaded into R16 and R17, respectively:

```
LDI R16, LOW (0x4455); R16 = 0x55
LDI R17, HIGH (0x4455); R17 = 0x44
```

In Chapter 2, we used the following instructions to make the stack pointer refer to the last location of the memory:

```
HIGH LOW

0x4455

R17 R16
```

```
LDI R16, HIGH (RAMEND); R16 = 0 \times 08 (for ATmega32)

OUT SPH, R16 ;SPH = the high byte of address

LDI R16, LOW (RAMEND); R16 = 0 \times 5f

OUT SPL, R16 ;SPL = the low byte of address
```

REGISTER AND DIRECT ADDRESSING MODES

The CPU can access data in various ways. The data could be in a register, or in memory, or provided as an immediate value. These various ways of accessing data are called *addressing modes*. These ways can be categorized into the following groups:

- 1. Single-Register (Immediate)
- 2. Register
- 3. Direct
- 4. Register indirect
- 5. Flash Direct
- 6. Flash Indirect

Single-register (immediate) addressing mode

In this addressing mode, the operand is a register.

```
NEG
      R18
                        ;negate the contents of R18
COM
      R19
                        ; complement the contents of R19
    R20
                        ;increment R20
INC
DEC
    R21
                        ;decrement R21
                        ;rotate right R22
ROR
    R22
```

In some of the instructions there is also a constant value with the register operand.

```
LDI R19,0x25
                      ;load 0x25 into R19
SUBI R19,0x6
                      ;subtract 0x6 from R19
ANDI R19,0b01000000
                       ;AND R19 with 0x40
```

The constant value is sometimes referred to as *immediate data* since the operand comes immediately after the opcode when the instruction is assembled; and the addressing mode is referred to as immediate addressing *mode* in some microcontrollers.

This addressing mode can be used to load data into any of the R16-R31 general purpose registers. The immediate addressing mode is also used for arithmetic and logic instructions. Note that the letter "I" in instructions such as LDI, ANDI, and SUBI means "Immediate."

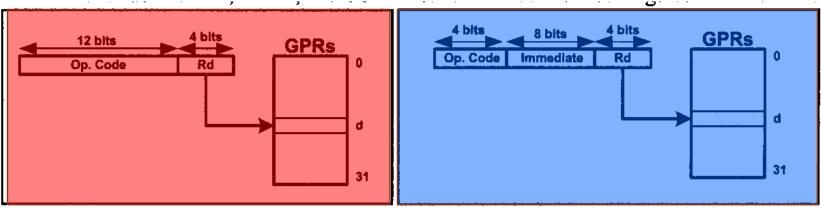


Figure 6-2a. Single-Register Addressing

Figure 6-2b. Single-Register (with immediate)

We can use the .EQU directive to access immediate data, as shown below.

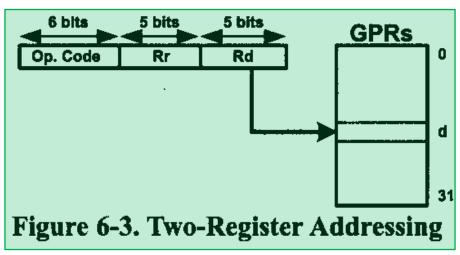
.EQU COUNT =
$$0x30$$

... ...
LDI R16, COUNT ; R16 = $0x30$

Two-register addressing mode

Two-register addressing mode involves the use of two registers to hold the

data to be manipulated.



Examples of two-register addressing mode are as follows:

```
ADD R20,R23 ; add R23 to R20

SUB R29,R20 ; subtract R20 from R29

AND R16,R17 ; AND R16 with 0x40

MOV R23,R19 ; copy the contents of R19 to R23
```

Direct addressing mode

The entire data memory can be accessed using either direct or register indirect addressing modes. In direct addressing mode, the operand data is in a RAM memory location whose address is known, and this address is given as a part of the instruction.

```
LDS R19,0x560 ; load R19 with the contents of memory loc $560 STS 0x40,R19 ; store R19 to data space location 0x40
```

The two instructions use direct addressing mode. If we dissect the opcode we see that the addresses are embedded in The instruction

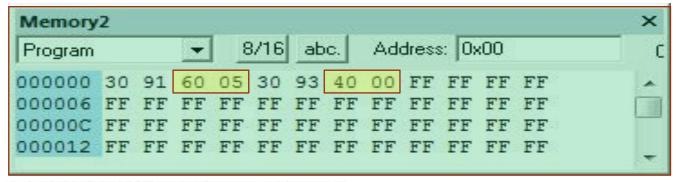
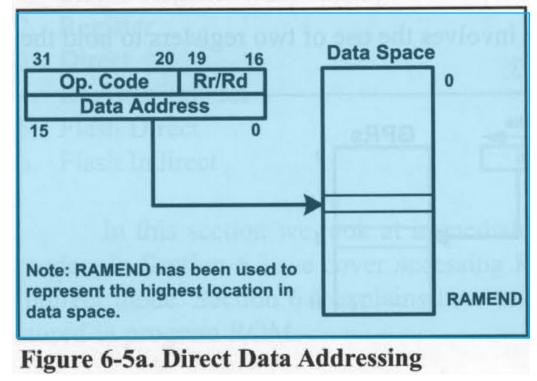


Figure 6-4. Direct Addressing Opcode

As shown in Figure 6-5a, the address field is a 16-bit address and can take values from \$0000-\$FFFF. Of course, it is much easier to use names instead of addresses in the program, and we have seen many examples of them in the last few chapters.



It must be noted that data memory does not support immediate addressing mode. In other words, to move data into internal RAM or to I/O registers, we must first move it to a GPR (R16-R31), and then move it from the GPR to the data memory space using the STS instruction.

LDI

R19,0x95 0x520,R19

I/O direct addressing mode

To access the I/O registers there is a special mode called I/O direct addressing mode. The I/O direct addressing mode can address only the standard I/O registers. The IN and OUT instructions use this addressing mode. Examine the following instruction, which copies the contents of PINB to PORTC:

IN OUT

R18,0x16 0x15,R18 ;PINB ;PORTC

As shown in Figure 6-5b, the address field is a 6-bit address and can take values from \$00 to \$3F, which is from 00 to 63 in decimal. So, it can address the entire standard I/O register memory space.

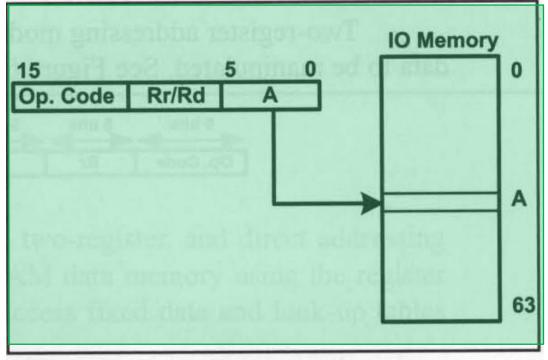


Figure 6-5b. I/O Direct Addressing

The AVR registers for Ports A, B, and so on are part of the group of registers commonly referred to as *I/O registers*. There are many *I/O* registers and they are widely used. The *I/O* registers can be accessed by their names or by their addresses. For example, PINB has address 0x16, and PORTC the address \$15, as shown in Table 6-4. Notice how the following pairs of instructions mean the same thing:

OUT	0x15,R19	
OUT	PORTC,R19	
IN	R19,0x16	
IN	R19,PINB	

Table 6-4: Selected ATmega32 I/O Register Addresses

Symbol	Name	I/O Address	Data Memory Addr.
PIND	Port D input pins	\$10	\$30
DDRD	Data Direction, Port D	\$11	\$31
PORTD_	Port D data register	\$12	\$32
PINC	Port C input pins	\$13	\$33
DDRC	Data Direction, Port C	\$14	\$34
PORTC	Port C data register	\$15	\$35
PINB	Port B input pins	\$16	\$36
DDRB	Data Direction, Port B	\$17	\$37
PORTB	Port B data register	\$18	\$38
PINA	Port A input pins	\$19	\$39
DDRA	Data Direction, Port A	\$1A	\$3A
PORTA	Port A data register	\$1B	\$3B
SPL	Stack Pointer, Low byte	\$3D	\$5D
SPH	Stack Pointer, High byte	\$3E	\$5E

Table 6-4 lists some of the AVR I/O registers and their addresses. The following points should be noted about the addresses of I/O registers:

- 1. As shown in Figures 2-3 and 2-4, the addresses between \$20 and \$5F of the data space have been assigned to standard I/O registers in all of the AVRs. These I/O registers have two addresses: I/O address and data memory address. The I/O address is used when we use the I/O direct addressing mode, while the data memory address is used when we use the direct addressing mode; in other words, the standard I/O registers can be accessed using both the direct addressing and the I/O addressing modes.
- 2. Some AVRs have less than 64 I/O registers. So, some locations of the standard I/O memory are not used by the I/O registers. The unused locations are reserved and must not be used by the AVR programmer.

3. Some AVRs have more than 64 I/O registers. The extra I/O registers are located above the data memory address \$5F. The data memory allocated to the extra I/O registers is called extended I/O memory.

To access the extended I/O registers we can use the direct addressing mode. For example, inATmega128, PORTF has the memory address of 0x62. So, the following instruction stores the contents of R20 in PORTF.

STS 0x62,R20

4. The I/O registers can have different addresses in different AVR microcontrollers. For example, the I/O address \$2 is assigned to TWAR in the ATmega32, while the same address is assigned to DDRE inATmega128. This can cause problems if you want to run programs written for one AVR on another AVR. The best way to solve this problem is to use the names of the registers instead of their addresses.

Register indirect addressing mode

In the register indirect addressing mode, a register is used as a pointer to the data memory location. In the AVR, three registers are used for this purpose: X, Y, and Z. These are 16-bit registers allowing access to the entire 65,536 bytes of data memory space in the AVR.

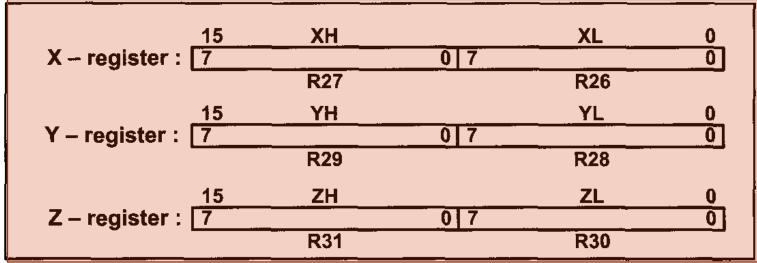


Figure 6-6. Registers X, Y, and Z

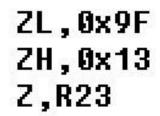
We can use them with the LD instruction to read the value of a location pointed to by these registers. For example, the following instruction reads the value of the location pointed to by the X pointer.

```
LD R24, X ;load into R24 from location pointed to by X
```

For instance, the following program loads the contents of location 0x130 into R18:

```
LDI XL, 0x30; load R26 (the low byte of X) with 0x30; LDI XH, 0x01; load R27 (the high byte of X) with 0x1; copy the contents of location 0x130 to R18
```

The ST instruction can be used to write a value to a location to which any of the X, Y, and Z registers points. For example, the following program stores the contents of R23 into location 0x139F:



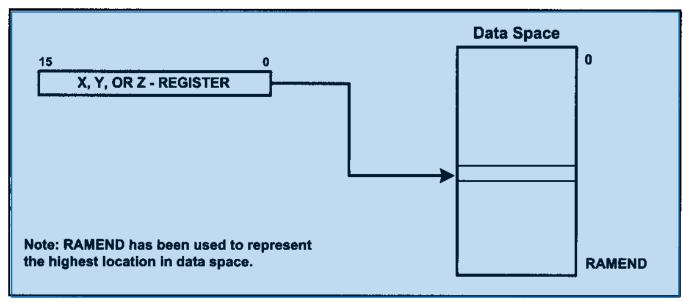


Figure 6-7. Register Indirect Addressing Mode

Advantages of register indirect addressing mode

```
Example 6-5
Write a program to copy the value $55 into memory locations $140 through $144 using
(a) direct addressing mode,
(b) register indirect addressing mode without a loop, and
(c) a loop.
Solution:
(a)
       LDI
             R17,0x55
                                ;load R17 with value 0x55
       STS
             0x140,R17
                                ;copy R17 to memory location 0x140
             0x141,R17
                                ; copy R17 to memory location 0x141
       STS
             0x142,R17
       STS
                                ; copy R17 to memory location 0x142
       STS
             0x143,R17
                                ; copy R17 to memory location 0x143
                                copy R17 to memory location 0x144
             0x144,R17
       STS
(b)
             R16,0x55
                          :load R16 with value 0x55
       LDI
             YL, 0x40
       LDI
                          ; load R28 with value 0x40 (low byte of addr.)
                          ;load R29 with value 0x1 (high byte of addr.)
       LDI
             YH, 0x1
             Y,R16
                          ; copy R16 to memory location 0x140
       ST
       INC
             YL
                          ;increment the low byte of Y
             Y, R16
                          ; copy R16 to memory location 0x141
       ST
                          ;increment the pointer
       INC
       ST
             Y.R16
                          ; copy R16 to memory location 0x142
                          ;increment the pointer
       INC
            YL
                          copy R16 to memory location 0x143
       ST
             Y. R16
                          ;increment the pointer
       INC
             YL
                          ; copy R16 to memory location 0x144
       ST
             Y,R16
```

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node.

```
Example 6-5 (Cont.)
(c)
                         :R16 = 5 (R16 \text{ for counter})
       LDI
           R16,0x5
           R20,0x55 ;load R20 with value 0x55 (value to be copied)
       LDÏ
                       :load YL with value 0x40
       LDI
            YL,0x40
                       ;load YH with value 0x1
            YH, 0x1
      LDI
                         ; copy R20 to memory pointed to by Y
           Y,R20
       ST
L1:
                         ; increment the pointer
            YL
       INC
                         :decrement the counter
            R16
       DEC
                         ;loop while counter is not zero
       BRNE L1
Use the AVR Studio simulator to examine memory contents after the above program is
run.
       $140 = (\$55) $141 = (\$55) $142 = (\$55) $143 = (\$55) 144 = (\$55)
```

Auto-increment and auto-decrement options for pointer registers

Because the pointer registers (X, Y, and Z) are 16-bit registers, they can go from \$0000 to \$FFFF, which covers the entire 64K memory space of the AVR. Using the "INC ZL" instruction to increment the pointer can cause a problem.

Table 6-5: AVR Auto-Increment/Decrement of Pointer Registers for LD Instruction

Instru	ıction	Function
LD	Rn,X	After loading location pointed to by X, the X stays the same.
LD	Rn,X+	After loading location pointed to by X, the X is incremented.
LD	Rn,-X	The X is decremented, then the location pointed to by X is loaded.
LD	Rn,Y	After loading location pointed to by Y, the Y stays the same.
LD	Rn,Y+	After loading location pointed to by Y, the Y is incremented.
LD	Rn,-Y	The Y is decremented, then the location pointed to by Y is loaded.
LDD	Rn,Y+q	After loading location pointed to by Y+q, the Y stays the same.
LD	Rn,Z	After loading location pointed to by Z, the Z stays the same.
LD	Rn,Z+	After loading location pointed to by Z, the Z is incremented.
LD	Rn,-Z	The Z is decremented, then the location pointed to by Z is loaded.
LDD	Rn,Z+q	After loading location pointed to by Z+q, the Z stays the same.

Example 6-6

Assume that RAM location \$90-\$94 have string of ASCII data, as shown below:

$$$90 = ('H')$$
 $$91 = ('E')$ $$92 = ('L')$ $$93 = ('L')$ $$94 = ('O')$

Write a program to get each character and send it to PORT B one byte at a time.

Show the program using:

- (a) Direct Accessing mode.
- (b) Register indirect addressing mode.

Solution:

(a) Using direct Addressing mode

	<i>J</i>
LDI	R20,0xFF
OUT	DDRB,R20
LDS	R20,0x90
OUT	PORTB, R20
LDS	R20,0x91
OUT	PORTB, R20
LDS	R20,0x92
OUT	PORTB, R20
LDS	R20,0x93
OUT	PORTB, R20
LDS	R20,0x94
OUT	PORTB, R20

(b) Using Register indirect addressing mode

	T LDI	R16,0x5
	LDI	R20,0xFF
	OUT	DDRB,R20
	LDI	ZL,0x90
	LDI	ZH,0x0
L1:	LD	R20,Z
	INC	ZL
	OUT	PORTB, R20
	DEC	R16
	BRNE	L1

When simulating the above program on the AVR Studio, make sure memory location \$90-\$94 have the message "HELLO"

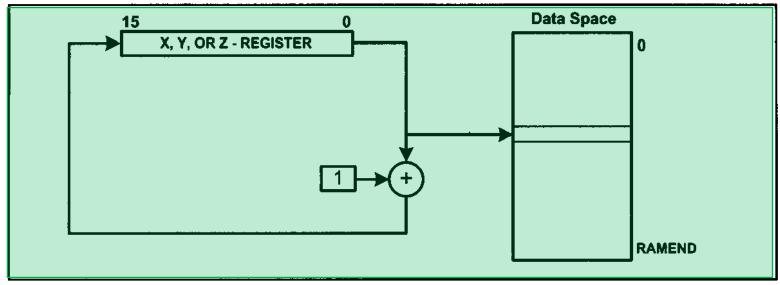


Figure 6-8. Register Indirect Addressing with Post-increment

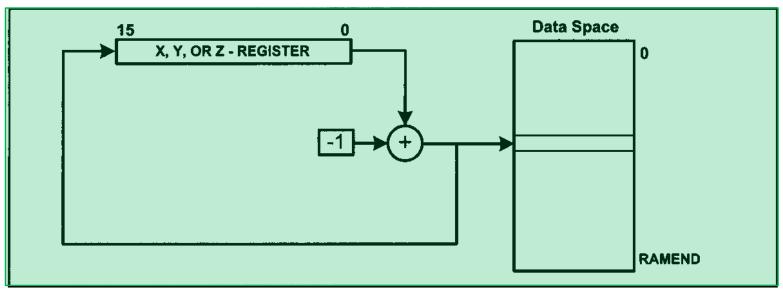


Figure 6-9. Register Indirect Addressing with Pre-decrement

Write a program to clear 16 memory location starting at data memory address \$60. Use the following:

- (a) INC Rn
- (b) Auto-increment

.inc	lude "m32	def.inc"
	LDI	R16,16
	LDI	XL,0x60
	LDI	XH,0x00
	LDI	R20,0x0
L1:	ST	X,R20
	INC	XL
	DEC	R16
	BRNE	L1

.INCLUDE "m32def.inc"		
	LDI	R16,16
	LDI	XL,0x60
	LDI	XH,0x00
0.00	LDI	R20,0x0
L1:	ST	X+,R20
	DEC	R16
	BRNE	L1

Example 6-8

Assume that data memory locations \$240-\$243 have the following hex data. Write a program to add them together and place the result in locations \$220 and \$221.

$$$240 = ($7D)$$
 $$241 = ($EB)$ $$242 = ($C5)$ $$243 = ($5B)$

Solution:

```
.INCLUDE "M32DEF.INC"
     .EQU L BYTE = 0x220 ; RAM loc for L Byte
     .EQU H BYTE = 0x221 ; RAM loc for H Byte
     LDI R16,4
     LDI R20,0
     LDI R21,0
     LDI XL, 0x40 ; the low byte of X = 0x40
     LDI XH, 0x02; the high byte of X = 02
    LD R22, X+ ; read contents of location where X points to
L1:
     ADD R20, R22
     BRCC L2
                     :branch if C = 0
     INC R21 ;increment R21
     DEC
     DEC R16
BRNE L1
                   ;decrement counter
L2:
                 ;loop until counter is zero
          L BYTE, R20; store the low byte of the result in $220
     ST
          H BYTE, R21; store the high byte of the result in $221
```

Example 6-9

Write a program to copy a block of 5 bytes of data from data memory locations starting at \$130 to RAM locations starting at \$60.

Solution:

```
LDI R16, 16 ;R16 = 16 (counter value)
    LDI XL, 0x30; the low byte of address
         XH, 0x01 ; the high byte of address
    LDI
    LDI YL, 0x60; the low byte of address
    LDI YH, 0x00; the high byte of address
    LD R20, X+ ; read where X points to
L1:
    ST Y+, R20 ;store R20 where Y points to
         R16 ;decrement counter
     DEC
    BRNE L1 ;loop until counter = zero
```

Before we run the above program.

```
130 = ('H') 131 = ('E') 132 = ('L') 133 = ('L') 134 = ('O')
```

After the program is run, the addresses \$60-\$64 have the same data as \$130-\$134.

```
130 = ('H') 131 = ('E') 132 = ('L') 133 = ('L') 134 = ('O')
60 = ('H') 61 = ('E') 62 = ('L') 63 = ('L') 64 = ('O')
```

AVR ADVANCED ASSEMBLY LANGUAGE PROGRAMMING 6.3 REGISTER INDIRECT ADDRESSING MODE

Example 6-10

Two multibyte numbers are stored in locations \$130-\$133 and \$150-\$153. Write a program to add the multibyte numbers and save the result in address \$160-\$163.

\$C7659812

\$2978742A

Solution:

```
.INCLUDE "M32DEF.INC"
                            :R16 = 4 (counter value)
     LDI R16, 4
     LDI XL, 0x30
     LDI XH, 0x1
                            ; load pointer. X = $130
     LDI YL, 0x50
     LDI YH, 0x1
                            ; load pointer. Y = $150
     LDI ZL, 0x60
                            ; load pointer. Z = $160
     LDI ZH, 0x1
                            ; clear carry
     CLC
                            ; copy memory to R18 and INC X
L1:
     LD R18, X+
     LD R19, Y+
                            ; copy memory to R19 and INC Y
     ADC R18,R19
                            ;R18 = R18 + R19 + carry
                            ;store R18 in memory and INC Z
     ST Z+, R18
                            :decrement R16 (counter)
     DEC
           R16
     BRNE L1
                            ;loop until counter = zero
```

AVR ADVANCED ASSEMBLY LANGUAGE PROGRAMMING 6.3 REGISTER INDIRECT ADDRESSING MODE

Before the addition we have:

```
MSByte
133 = (\$C7) \ 132 = (\$65) \ 131 = (\$98) \ 130 = (\$12)
153 = (\$29) \ 152 = (\$78) \ 151 = (\$74) \ 150 = (\$2A)
```

After the addition we have:

```
163 = (\$F0) \ 162 = (\$DE) \ 161 = (0C) \ 160 = (3C)
```

Notice that we are using the little endian convention of storing a low byte to a low address, and a high byte to a high address. Single-step the program in AVR Studio and examine the pointer registers and memory contents to gain insight into register indirect addressing mode.

AVR ADVANCED ASSEMBLY LANGUAGE PROGRAMMING 6.3 REGISTER INDIRECT ADDRESSING MODE

Register indirect with displacement

Suppose we want to read a byte that is a few bytes higher than where the Z register points to. To do so we can increment the Z register so that it points to the desired location and then read it. But there is an easier way; we can use the register indirect with displacement. In this addressing mode a fixed number is added to the Z register. For example, if we want to read from the location that is 5 bytes after the location to which Z points, we can write the following instruction:

The general format of the instruction is as follows:

```
LDD Rd, Z+q ;load from Z+q into Rd
```

where q is a number between 0 to 63, and Rd is any of the general purpose registers.

AVR ADVANCED ASSEMBLY LANGUAGE PROGRAMMING 6.3 REGISTER INDIRECT ADDRESSING MODE

To store a byte of data in a data memory location using the register indirect with displacement addressing mode we can use STD (Store with Displacement). The instruction is as follows:

For example, the following instruction writes the contents of R20 into the location that is five bytes away from where Z points to:

```
STD Z+5,R20 ;store R20 into location Z+5
```

AVR ADVANCED ASSEMBLY LANGUAGE PROGRAMMING 6.3 REGISTER INDIRECT ADDRESSING MODE

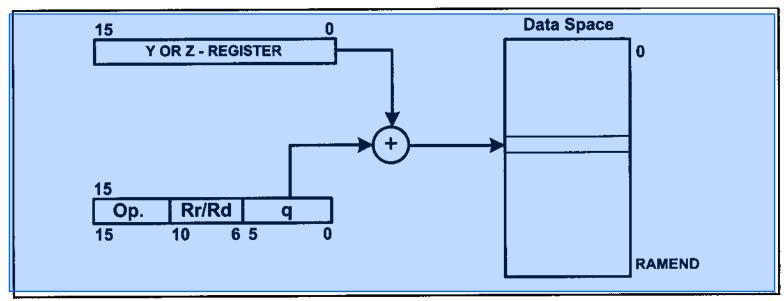


Figure 6-10. Register Indirect with Displacement

AVR ADVANCED ASSEMBLY LANGUAGE PROGRAMMING 6.3 REGISTER INDIRECT ADDRESSING MODE

Example 6-11

Write a function that adds the contents of three continuous locations of data space and stores the result in the first location. The Z register should point to the first location before the function is called.

Solution:

```
.INCLUDE "M32DEF.INC"
     LDI R16, HIGH (RAMEND) ; initialize the stack pointer
     OUT SPH, R16
     LDI R16, LOW (RAMEND)
     OUT SPL, R16
                            ;initialize the Z register
     LDI ZL,0x00
     LDI ZH.2
     CALL ADD3LOC
                           :call add3loc
                            ;loop forever
HERE: JMP
           HERE
ADD3LOC:
                       ;R21 = 0
     LDI R21,0
     LD R20,Z ;R20 = contents of location Z
LDD R16,Z+1 ;R16 = contents of location Z+1
     ADD R20, R16
                          ;R20 = R20 + R16
                          ;branch if carry cleared
;increment R21 if carry occurred
     BRCC L1
     INC R21
                           ;R16 = contents of location Z+2
     LDD R16,Z+2
L1:
                            ;R20 = R20 + R16
     ADD R20,R16
                           ;branch if carry cleared
     BRCC L2
                          ;increment R21
     INC R21
                           ;store R20 into location Z
     ST Z,R20
     STD Z+1,R21
                           store R21 into location Z+1
     RET
```

.DB (define byte) and fixed data in program ROM

The .DB data directive is widely used to allocate ROM program (code) memory in byte-sized chunks. When .DB is used to define fixed data, the numbers can be in decimal, binary, hex, or ASCII formats. The .DB directive is widely used to define ASCII strings.

Assume that we have burned the following fixed data into the program ROM of an AVR chip. Give the contents of each ROM location starting at \$500. See Appendix F for the hex values of the ASCII characters.

```
:MY DATA IN FLASH ROM
     .ORG $500
DATA1: .DB 1,8,5,3
                ; DECIMAL (1C in hex)
DATA2: . DB 28
DATA3: .DB 0b00110101
                       ;BINARY (35 in hex)
DATA4: . DB 0x39
                       ; HEX
     .ORG 0x510
DATA4: .DB 'Y' ; single ASCII char
DATA5: .DB '2','0','0','5'; ASCII numbers
     .ORG $516
DATA6: .DB "Hello ALI" ; ASCII string
```

Example 6-12

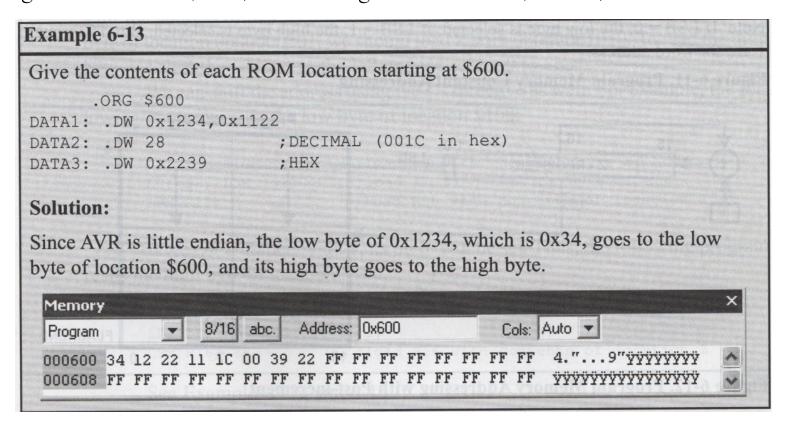
.DB (define byte) and fixed data in program ROM

In Example 6-12 notice that each location of program memory is 2 bytes, whereas the . DB directive allocates byte-sized chunks. If we allocate a few bytes of data using the . DB directive, the first byte goes to the low byte of ROM location; the second byte goes to the high byte of ROM location; the third byte goes to the low byte of the next location of program ROM; and so on.

In the cases in which we allocate an odd number of ROM locations using . DB, the assembler will automatically make the number of allocated locations even by placing a zero into the high byte of the last location.

In Example 6-12 notice also that we must use single quotes (') for a single character and double quotes (") for a string.

AVR assembly also allows the use of .DW in place of . DB to define values greater than 255 (0xFF) but not larger than 65,535 (0xFFFF).



Reading table elements in the AVR

Now, we need to have a register pointer to point to the data to be fetched from the code space. The Z register can be used for this purpose. For this reason we can call it register indirect flash addressing mode.

In AVR terminology, there are two register indirect flash addressing modes:

- program memory constant addressing and
- program memory addressing with post-increment.

In the program memory constant addressing mode, the content of Z does not change when the instruction is executed, which is why is called constant addressing; whereas in the program memory addressing with post-increment, the content of Z increments after each execution.

The Instruction

"LPM Rn,Z"

uses program memory constant addressing mode, while

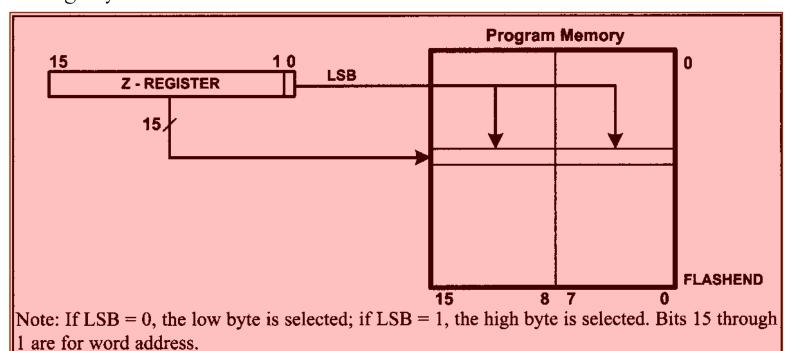
"LPM Rn,Z+"

uses program memory addressing with post-increment.

There is a group of AVR instructions designed for table processing. Table 6.6 shows the instructions for table reading in the AVR.

Table 6-6: AVR Table Read Instructions					
Instruction	Function	Description			
LPM Rn,Z	Load from Program Memory	After read, Z stays the same			
LPM Rn,Z+	Load from Program Memory with post-inc	c. Reads and increments Z			
Note: The by	e of data is read into the Rn register from c	ode space pointed to by Z.			

As you know, in the AVR, each location of the program memory is 2 bytes. So, we should mention if we want to read the low byte or the high byte. The least significant bit (LSB) of the Z register indicates whether the low byte or the high byte should be read.



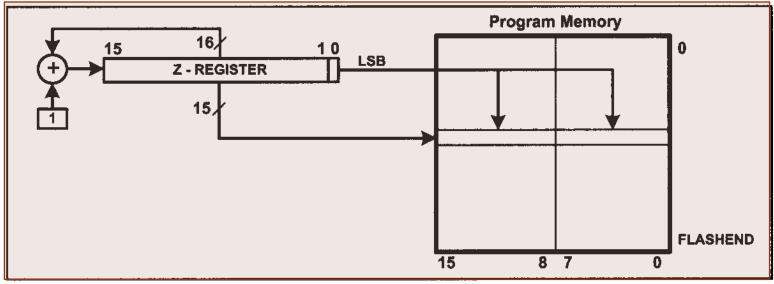


Figure 6-12. Program Memory Addressing with Post-increment

Figure 6-13b shows the value that should be loaded into the Z register in order to address each byte of the program memory.

For example, to address the low byte of location \$0002, we should load the Z register with \$0005, as shown below:

```
LDI ZH, 0 \times 00; load ZH with 0 \times 00 (the high byte of addr.)
LDI ZL, 0x05
                   ;load ZL with 0x05 (the low byte of addr.)
LPM R16, Z
                    :load R16 with contents of location Z
```

Low	High	Address		Low	High	Address
0000 0000 0000 0000	0000 0000 0000 0001	0000 0000 0000 0000		\$0000	\$0001	\$0000
0000 0000 0000 0010	0000 0000 0000 0011	0000 0000 0000 0001		\$0002	\$0003	\$0001
0000 0000 0000 0100	0000 0000 0000 0101	0000 0000 0000 0010	П	\$0004	\$0005	\$0002
0000 0000 0000 0110	0000 0000 0000 0111	0000 0000 0000 0011	П	\$0006	\$0007	\$0003
0000 0000 0000 1000	0000 0000 0000 1001	0000 0000 0000 0100		\$0008	\$0009	\$0004
0000 0000 0000 1010	0000 0000 0000 1011	0000 0000 0000 0101		\$000A	\$000B	\$0005
1111 1111 1111 1100	1111 1111 1111 1101	0111 1111 1111 1110		\$FFFC	\$FFFD	\$7FFE
1111 1111 1111 1110	1111 1111 1111 1111	0111 1111 1111 1111		\$FFFE	\$FFFF	\$7FFF
Figure 6-13a. Values of Z (in Binary) Figure 6-13b. Values of Z						

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As you see in Figure 6-13a, to read the low byte of each location we should shift the address of that location one bit to the left. For instance, to access the low byte of location 0b00000101, we should load Z with 0b000001010.

To read the high byte, we shift the address to the left and we set bit 0 to one.

For example, the following program reads the low byte of location \$100:

```
LDI ZH, HIGH($100<<1) ; load ZH with the high byte of addr.
LDI ZL, LOW ($100<<1) ; load ZL with the low byte of addr.
          ;load R16 with contents of location Z
LPM R16, Z
```

If we OR a number with 1, its bit 0 will be set. Thus, the following program reads the high byte of location \$100.

```
LDI ZH, HIGH(($100<<1/)|1)
LDI ZL, LOW (($100<<1) 1
            ;load R16 with contents of location Z
LPM R16, Z
```

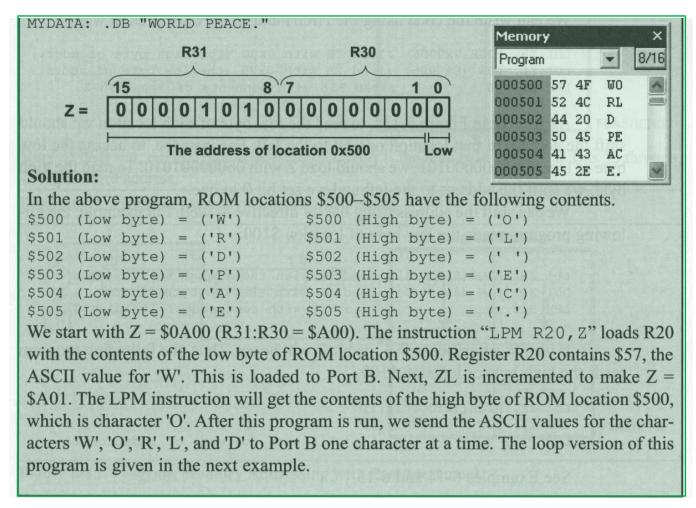
AVR ADVANCED ASSEMBLY LANGUAGE PROGRAMMING

6.4 LOOK-UP TABLE AND TABLE F

Example 6-14

In this program, assume that the phrase "WORLD PEACE" is burned into ROM location starting at \$500, and the program is burned into ROM locations starting at 0. Analyze how program works and state where "WORLD PEASE." is stored after this program is run.

1	. INCLUD	E "M32D	EF.INC"
2		.ORG	\$0000
3		LDI	R20,0xFF
4		OUT	DDRB,R20
5		LDI	ZL, LOW (MYDATA<<1)
6		LDI	ZH, HIGH (MYDATA<<1)
7		LPM	R20,Z
8		OUT	PORTB,R20
9		INC	ZL
10		LPM	R20,Z
11		OUT	PORTB,R20
12		INC	ZL
13		LPM	R20,Z
14		OUT	PORTB,R20
15		INC	ZL
16		LPM	R20,Z
17		OUT	PORTB,R20
18		INC	ZL
19		LPM	R20,Z
20		OUT	PORTB,R20
21		INC	ZL
22		LPM	R20,Z
23		OUT	PORTB,R20
24		INC	ZL
25	HERE:	RJMP	HERE
26			
27		.ORG	\$500
28	MYDATA:	.DB	"WORLD PEACE."
			11/27/202



Example 6-15

Assuming that program ROM space starting at \$500 contains "WORLD PEACE.". Write a program to send characters to Port B one byte at a time.

```
;This method uses a counter
     .include "m32def.inc"
              .ORG
                       $0000
                       R16,11
              LDI
                       R20, OxFF
              LDI
              OUT
                       DDRB, R20
                       ZL, LOW (MYDATA << 1)
              LDI
                       ZH, HIGH (MYDATA<<1)
              LDI
     L1:
              LPM
                       R20, Z
10
                       PORTB, R20
              OUT
11
              INC
                       ZL
12
                       R16
              DEC
13
                       L1
              BRNE
14
     HERE:
              RJMP
                       HERE
15
16
     ; data is burned into code (program)
17
     ;space starting at $500
18
                       $500
              .ORG
19
     MYDATA: .DB
                       "WORLD PEASE."
```

```
:This method uses a null char for
     ; end of string
     .include "m32def.inc"
            .ORG $0000
            LDI R20,0xFF
            OUT DDRB, R20
            LDI ZL, LOW (MYDATA << 1)
            LDI ZH, HIGH (MYDATA << 1)
 8
     L1:
            LPM R20,Z
10
            CPI R20,0
11
            BREQ HERE
12
            OUT PORTB, R20
13
            INC
                    ZL
14
            RJMP L1
15
     HERE: RJMP
                 HERE
16
17
     ; data is burned into code (program)
18
     ; space starting at $500
19
            .ORG $500
20
     MYDATA: .DB "WORLD PEASE.", 0
```

Auto-increment option for Z

Using the "INC ZL" instruction to increment the pointer can cause a problem when an address such as \$5FF is incremented. The carry will not propagate into ZH. The AVR gives us the option of LPM Rn, Z+ (load program memory with post-increment) as shown in Table 6-6.

Example 6-16

```
;burn into ROM starting at 0
    .include "m32def.inc"
                    $0000
            .ORG
                   R20, OxFF
            LDI
            OUT DDRB, R20
            LDI ZL, LOW (MYDATA << 1)
            LDI
                   ZH, HIGH (MYDATA<<1)
    L1:
            LPM
                   R20, Z+
                   R20,0
            CPI
10
            BREQ
                    HERE
            OUT PORTB, R20
12
            RJMP
                    T.1
13
    HERE: RJMP
14
15
     ;burned into code (program) space
16
     ; starting at $500
17
                    $500
            .ORG
    MYDATA: .DB "WORLD PEASE.", 0
```

Example 6-17

Assume that ROM space starting at \$100 contains the message "The Promise of World Peace". Write a program to bring this message into the CPU one byte at a time and place the byte in Ram locations starting at \$140.

```
. EQU
        RAM BUF = 0x140
.ORG
        $0000
.INCLUDE "M32DEF.INC"
        LDI
                R20,0xFF
                DDRB, R20
        OUT
        LDI
                ZH, HIGH (MYDATA<<1)
        LDI
                ZL, LOW (MYDATA<<1)
        LDI
                XH, HIGH (RAM BUF)
        LDI
                XL, LOW (RAM BUF)
                R20, Z+
        LPM
        CPI
                R20,0
        BREQ
                HERE
                X+,R20
        ST
        RJMP
                L1
HERE:
        RJMP
                HERE
        .ORG
                0x100
                "The Promise of World Peace", 0
MYDATA:
        .DB
```

Look-up table

The look-up table is a widely used concept in microcontroller programming. It allows access to elements of a frequently used table with minimum operations. As an example, assume that for a certain application we need $4+x^2$ values in the range of 0 to 9. We can use a look-up table instead of calculating the values, which takes some time. In the AVR, to get the table element we add the index to the address of the look-up table. This is shown in Examples 6-18 through 6-20.

.ORG

\$0000

Example 6-18

Assume that the lower three bits of Port C are connect to three switches. Write a program to send the following ASCII characters to Port D based on the BEGIN: status of the switches.

000	.0,
001	·1·
010	'2'
011	'3'
100	'4'
101	'5'
110	·6'
111	·7·

```
.INCLUDE "M32DEF.INC"
        LDI
                R16,0x0
        OUT
                DDRC,R16
                R16, OxFF
        LDI
        OUT
                DDRD, R16
                ZH, HIGH (ASCII TABLE << 1)
        LDI
                R16, PINC
        IN
                R16,0b00000111
        ANDI
                ZL, LOW (ASCII TABLE << 1)
        LDI
                ZL,R16
        ADD
                R17, Z
        LPM
                PORTD, R17
        OUT
        RJMP
                BEGIN
;lookup table for ASCII numbers 0-7
.ORG
        0x20
ASCII TABLE:
        .DB '0','1','2','3','4','5','6','7'
```

Example 6-19

Write a program to get the x value from Port B and sends x^2 to Port C. Assume that PB3-PB0 has x value of 0-9. use a lookup table instead of a multiply instruction.

```
$0000
.ORG
.INCLUDE "M32DEF.INC"
        LDI
                R16,0x0
                DDRB, R16
        OUT
                R16,0xFF
        LDI
        OUT
                DDRC, R16
                ZH, HIGH (XSQR TABLE << 1)
        LDI
                ZL, LOW (XSQR TABLE << 1)
        LDI
                R16, PINB
        IN
                R16,0x0F
        ANDI
                ZL, R16
        ADD
                R18,Z
        LPM
                PORTC, R18
        OUT
        RJMP
                L1
;lookup table for square of numbers 0-9
.ORG
        0x10
XSQR TABLE:
        .DB 0,1,4,9,16,25,36,49,64,81
```

Example 6-20

Write a program to get the x value from Port B and send x^2+2x+3 to Port C. Assume PB3-PB0 has the x value of 0-9. Use a look-up table instead of a multiply instruction.

```
$0000
.ORG
.INCLUDE "M32DEF.INC"
        LDI
                 R16,0x0
                 DDRB, R16
        OUT
                 R16,0xFF
        LDI
                 DDRC,R16
        OUT
                 ZH, HIGH (TABLE << 1)
        LDI
                 ZL, LOW (TABLE << 1)
        LDI
                 R16, PINB
        IN
                 R16,0x0F
        ANDI
        ADD
                 ZL, R16
                 R18,Z
        LPM
                 PORTC, R18
        OUT
        RJMP
                 L1
.ORG
        0x10
TABLE:
         .DB 3,11,18,27,38,51,66,83,102
```

Accessing a look-up table in RAM

The look-up table elements can also be in RAM instead of ROM. Sometimes we need to bring in the elements of the look-up table from RAM because the elements are dynamic and can change. In the AVR, we can do that using the pointers.

Writing table elements in AVR

In AVR we also have the SPM instruction, which allows us to write (store) data into program memory.

In this section, we provide more programming examples of bit manipulation using the bit-addressable and byte-addressable options of the AVR family.

In Table 6-7, some of the bit-oriented instructions are given. Notice that the bit-oriented instructions use only one addressing mode, the direct addressing mode.

Table 6-7: Single-Bit (Bit-Oriented) Instructions for AVR

Instruction	Function
SBI A,b	Set Bit b in I/O register
CBI A,b	Clear Bit b in I/O register
SBIC A,b	Skip next instruction if Bit b in I/O register is Cleared
SBIS A,b	Skip next instruction if Bit b in I/O register is Set
BST Rr,b	Bit store from register Rr to T
BLD Rd,b	Bit load from T to Rd
SBRC Rr,b	Skip next instruction if Bit b in Register is Cleared
SBRS Rr,b	Skip next instruction if Bit b in Register is Set
BRBS s,k	Branch if Bit s in status register is Set
BRBC s,k	Branch if Bit s in status register is Cleared

Note: A can be any location of the I/O register.

Manipulating the bits of general purpose registers Setting the bits

The SBR (Set Bits in Register) instruction sets the specified bits in the general purpose register. It has the following format:

Operation:

(i) $Rd \leftarrow Rd \vee K$

Syntax:

Operands:

Program Counter:

(i) SBR Rd,K

 $16 \le d \le 31, 0 \le K \le 255$

PC ← PC + 1

16-bit Opcode:

0110	KKKK	dddd	KKKK

Status Register (SREG) and Boolean Formula:

	Т					_	
-	-	-	⇔	0	⇔	⇔	-

For example, in the following program the SBR instruction sets the bits 2, 5, and 6 regardless of their previous values.

```
LDI R17,0b01011001 ;R17 = 0x59
SBR R17,0b01100100 ;set bits 2, 5, and 6 in register R17
```

When execution of the above instructions is finished, R17 contains 0x7D. Notice that the SBR instruction is a byte-oriented instruction as it manipulates the whole byte at one time.

Clearing the bits

The CBR (Clear Bits in Register) instruction clears the specified bits in the general purpose register. It has the following format:

Operation:

(i) Rd ← Rd • (\$FF - K)

Syntax: Operands:

Program Counter:

(i) CBR Rd,K

 $16 \le d \le 31, 0 \le K \le 255$

 $PC \leftarrow PC + 1$

16-bit Opcode: (see ANDI with K complemented)

Status Register (SREG) and Boolean Formula:

I	Т	Н	s	V	N	Z	С
-	_	_	\Leftrightarrow	0	⇔	⇔	_

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For example, in the following program the CBR instruction clears the bits 2, 5, and 6 regardless of their previous values.

```
LDI R17,0b01011001 ;R17 = 0x59 CBR R17,0b01100100 ;clear bits 2, 5, and 6 in register R17 After the execution of the above instructions, R17 contains 0x19.
```

Copying a bit

As we saw in Chapter 2, one of the bits in the SREG (status register) is named T (temporary), which is used when we want to copy a bit of data from one GPR to another GPR.

BST – Bit Store from Bit in Register to T Flag in SREG

Description:

Stores bit b from Rd to the T Flag in SREG (Status Register).

Operation:

(i) $T \leftarrow Rd(b)$

Syntax:

Operands:

Program Counter:

(i)

BST Rd,b

 $0 \le d \le 31, 0 \le b \le 7$

PC ← PC + 1

16-bit Opcode:

1111	101d	dddd	ddd0
------	------	------	------

BLD – Bit Load from the T Flag in SREG to a Bit in Register

Description:

Copies the T Flag in the SREG (Status Register) to bit b in register Rd.

Operation:

(i) Rd(b) ← T

Syntax:

Operands:

Program Counter:

(i)

BLD Rd,b

 $0 \le d \le 31, 0 \le b \le 7$

PC ← PC + 1

16 bit Opcode:

1111	100d	dddd	ddd0

For example, the following program copies bit 3 from R17 to bit 5 in register R19:

```
BST R17,3 ;store bit 3 from R17 to the T flag
BLD R19,5 ;copy the T flag to bit 5 in R19
```

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Example 6-21

A switch is connected to pin PB4. Write a program to get the status of the switch and save it in D0 of internal RAM location 0x200.

Solution:

```
.EQU MYREG = 0x200 ; set aside loc 0x200
            DDRB,0 ;make PBO an input
      CBI
            R17, PINB ; R17 = PINB
      IN
     BST R17,4 ;T = PINB.4

LDI R16,0x00 ;R16 = 0

BLD R16,0 ;R16.0 = T
            MYREG,R16 ; copy R16 to location $200
      STS
HERE: JMP
            HERE
```

SBRS – Skip if Bit in Register is Set

Description:

This instruction tests a single bit in a register and skips the next instruction if the bit is set.

Operation:

(i) If Rr(b) = 1 then $PC \leftarrow PC + 2$ (or 3) else $PC \leftarrow PC + 1$

Syntax:

Operands:

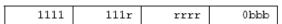
Program Counter:

(i) SBRS Rr,b

 $0 \le r \le 31, 0 \le b \le 7$

PC ← PC + 1, Condition false - no skip PC ← PC + 2, Skip a one word instruction PC ← PC + 3, Skip a two word instruction

16-bit Opcode:



For example, in the following program the "LDI R20,0x55" instruction will not be executed since bit 3 of R17 is set.

```
LDI R17,0b0001010
SBRS R17,3 ;skip next instruction if Bit 3 in R17 is set
LDI R20,0x55
LDI R30,0x33
```

SBRC – Skip if Bit in Register is Cleared

Description:

This instruction tests a single bit in a register and skips the next instruction if the bit is cleared.

Operation:

If Rr(b) = 0 then $PC \leftarrow PC + 2$ (or 3) else $PC \leftarrow PC + 1$ (i)

Syntax:

Operands:

Program Counter:

SBRC Rr.b (i)

 $0 \le r \le 31, 0 \le b \le 7$

PC ← PC + 1, Condition false - no skip PC ← PC + 2, Skip a one word instruction PC ← PC + 3, Skip a two word instruction

16-bit Opcode:

1111	110r	rrrr	ddd0

For example, in the following program the "LDI R20, 0x55" instruction will not be executed since bit 2 of R16 is cleared.

R16,0b0001010 LDI

SBRC R16,2 ; skip next instruction if Bit 2 in R16 is cleared R20,0x55 LDI

LDI

Example 6-22

A switch is connected to pin PC7. Using the SBRS instruction, write a program to check the status of the switch and perform the following:

```
.INCLUDE "M32DEF.INC"
                      ; include a file according to the IC you use
           DDRC,7
                      ; make PC7 an input
     CBI
     LDI
          R16,0xFF
     OUT
           DDRD, R16 ; make Port D an output port
        R20, PINC ; R20 = PINC
AGAIN: IN
                      ;skip next line if Bit PC7 is set
     SBRS R20,7
                     :it must be LOW
     RJMP
           OVER
          R16, 'Y' ;R16 = 'Y' ASCII letter Y
     LDI
          PORTD, R16
                     ;issue R16 to PD
     OUT
          AGAIN
                     :we could use JMP instead
     RJMP
OVER: LDI R16, 'N' ;R16 = 'N' ASCII letter N
                      ;issue R16 to PORTD
     OUT PORTD, R16
     RJMP
           AGAIN
                      ;we can use JMP too
```

Manipulating the bits of I/O registers

As discussed before, we can set and clear the lower 32 I/O registers (addresses 0 to 31) using the SBI (Set bit in 1/O register) and CBI (Clear bit in I/O register) instructions. For example, the following two instructions set the PORTA. 1 and clear the PORTB.4, respectively:

```
SBI PORTA,1 ;set Bit 1 in PORTA
CBI PORTB,4 ;clear Bit 4 in PORTB
```

Example 6-23

Write a program to toggle PB2 a total of 200 times.

Solution:

```
LDI R16,200 ;load the count into R16
SBI DDRB,2 ;DDRB.1 = 1, make RB1 an output
AGAIN:SBI PORTB,2 ;set bit PB2 (toggle PB2)
CBI PORTB,2 ;clear bit PB2 (toggle PB2)
DEC R16 ;decrement R16
BRNE AGAIN ;continue until counter is zero
```

Status register bit-addressability

Checking a flag bit

BRBS - Branch if Bit in SREG is Set

Description:

Conditional relative branch. Tests a single bit in SREG and branches relatively to PC if the bit is set. This instruction branches relatively to PC in either direction (PC - $63 \le$ destination \le PC + 64). The parameter k is the offset from PC and is represented in two's complement form.

Operation:

(i) If SREG(s) = 1 then $PC \leftarrow PC + k + 1$, else $PC \leftarrow PC + 1$

Syntax:

Operands:

(i) BRBS s,k

 $0 \le s \le 7$, $-64 \le k \le +63$

Program Counter:

 $PC \leftarrow PC + k + 1$

PC ← PC + 1, if condition is false

16-bit Opcode:

1111	00kk	kkkk	ksss

Status Register (SREG) and Boolean Formula:

ı	Т	Н	S	V	N	Z	С
_	_	_	_	_	_	_	_

BRBC - Branch if Bit in SREG is Cleared

Description:

Conditional relative branch. Tests a single bit in SREG and branches relatively to PC if the bit is cleared. This instruction branches relatively to PC in either direction (PC - $63 \le$ destination \le PC + 64). The parameter k is the offset from PC and is represented in two's complement form.

Operation:

(i) If SREG(s) = 0 then $PC \leftarrow PC + k + 1$, else $PC \leftarrow PC + 1$

Syntax:

Operands:

Program Counter:

(i) BRBC s,k

 $0 \le s \le 7$, $-64 \le k \le +63$

 $PC \leftarrow PC + k + 1$

PC ← PC + 1, if condition is false

16-bit Opcode:

Status Register (SREG) and Boolean Formula:

			S				
_	_	_	_	_	_	_	_

For example, in the following program the LDI instruction is not executed when the carry flag is set:

```
BRBS 0,L1 ;branch if status flag bit 0 is set LDI R20,3
```

We can write the same program using the "BRCS LI" instruction as follows:

```
BRCS L1 ;branch if carry flag is set LDI R20,3
```

Manipulating a bit

BSET - Bit Set in SREG

Description:

Sets a single Flag or bit in SREG.

Operation:

(i) SREG(s) ← 1

Syntax:

Operands:

Program Counter:

(i) BSET s

 $0 \le s \le 7$

 $PC \leftarrow PC + 1$

16-bit Opcode:

1001	0100	0sss	1000

Status Register (SREG) and Boolean Formula:

-	Т			-		_	
\Leftrightarrow	⇔	⇔	⇔	\Leftrightarrow \Leftrightarrow		⇔	⇔

BCLR - Bit Clear in SREG

Description:

Clears a single Flag in SREG.

Operation:

(i) SREG(s) ← 0

Syntax:

Operands:

Program Counter

(i) BCLR s

 $0 \le s \le 7$

 $PC \leftarrow PC + 1$

16-bit Opcode:

1001	0100	1sss	1000
			l

Status Register (SREG) and Boolean Formula:

-	Т		-	 _	
\Leftrightarrow	⇔	\$ → ⇔ ⇔		\$ \$	⇔

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For example, the following instruction sets the carry flag.

```
BSET 0 ;set bit 0 (carry flag)
```

As another example, the instruction "BSET 2" sets the N (Negative) flag.

For example, the following instruction clears the carry flag.

```
BCLR 0 ;clear bit 0 (carry flag)
```

As another example, the instruction "BCLR 1" clears the Z (Zero) flag.

A more convenient way is to use the CLZ instruction, as shown in Table 6-9.

Table 6-9: Manipulating the Flags of the Status Register

Instructio	n Action		Instruction	Action	
SEC	Set Carry	C = 1	CLC	Clear Carry	C = 0
SEZ	Set Zero	Z = 1	CLZ	Clear Zero	Z = 0
SEN	Set Negative	N = 1	CLN	Clear Negative	N = 0
SEV	Set overflow	V = 1	CLV	Clear overflow	V = 0
SES	Set Sign	S = 1	CLS	Clear Sign	S = 0
SEH	Set Half carry	H = 1	CLH	Clear Half carry	H = 0
SET	Set Temporary	T = 1	CLT	Clear Temporary	T = 0
SEI	Set Interrupt	I = 1	CLI	Clear Interrupt	I = 0

Internal RAM bit-addressability The internal RAM is not bit-addressable.

So, in order to manipulate a bit of the internal RAM location, you should bring it into the general purpose register and then manipulate it, as shown in Examples .

Example 6-25

Write a program to see if the internal RAM location \$195 contains an even value. If so, send it to Port B. If not, make it even and then send it to Port B.

Solution 1:

```
MYREG = 0 \times 195
                              :set aside loc 0x195
.EOU
      LDI
            R16,0xFF
      OUT
            DDRB,R16
                               ;make Port B an output port
AGAIN:LDS
            R16, MYREG
            R16,0
                              ;bit test DO, skip if set
      SBRS
                              ;it must be LOW
      RJMP
            OVER
                              ; clear bit D0 = 0
            R16,0b00000001
      CBR
OVER: OUT PORTB, R16
                              ; copy it to Port B
                              ; we can use RJMP too
      JMP
            AGAIN
```

Solution 2:

```
.EQU MYREG = 0 \times 195
                              :set aside loc 0x195
      LDI
            R16,0xFF
                              ; make Port B an output port
      OUT
            DDRB, R16
            R16, MYREG
AGAIN:LDS
      CBR
            R16,0b00000001
                              ; clear bit D0 = 0
           PORTB, R16
                              ; copy it to Port B
OVER: OUT
            AGAIN
                              :we can use RJMP too
      JMP
```

Example 6-27

Write a program to see if the internal RAM location \$137 contains an even value. If so, write 0x55 into location \$200. If not, write 0x63 into location \$200.

Solution:

```
.EOU MYREG = 0 \times 137
                       :set aside location 0x137
.EOU RESULT= 0x200
     LDS
           R16, MYREG
                       ; skip if clear Bit D0 of R16 register is clr
     SBRC R16,0
          OVER
                       :it is odd
     RJMP
     LDI
          R16,0x55
     STS
           RESULT, R16
     RJMP HERE
OVER: LDI
          R16,0x63
          RESULT, R16
     STS
HERE: RJMP HERE
```

6.6 ACCESSING EEPROM IN AVR

The data in SRAM will be lost if the power is disconnected. However, we need a place to save our data to protect them against power failure. EEPROM memory can save stored data even when the power is cut off. In this section we will show how to write to EEPROM memory and how to access it.

Table 6-10: Size of EEPROM Memory in ATmega Family

Chip	Bytes	Chip	Bytes	Chip	Bytes
ATmega8	512	ATmega16	512	ATmega32	1024
ATmega64	2048	ATmega128	4096	ATmega256RZ	4096
ATmega640	4096	ATmega1280	4096	ATmega2560	4096

EEPROM registers

There are three I/O registers that are directly related to EEPROM. These are **EECR** (EEPROM Control Register), **EEDR** (EEPROM Data Register), and **EEARH-EEARL** (EEPROM Address Register High-Low). Each of these registers is discussed in detail in this section.

EEPROM Data Register (EEDR)

To write data to EEPROM, you have to write it to the EEDR register and then transfer it to EEPROM. Also, if you want to read from EEPROM you have to read from EEDR. In other words, EEDR is a bridge between EEPROM and CPU.

EEPROM Address Register (EEARH and EEARL)

The EEARH:EEARL registers together make a 16-bit register to address each location in EEPROM memory space. When you want to read from or write to EEPROM, you should load the EEPROM location address in EEARs. As you see in Figure 6-15, only 10 bits of the EEAR registers are used in ATmega32. Because ATmega32 has 1024-byte EEPROM locations, we need 10 bits to address each location in EEPROM space

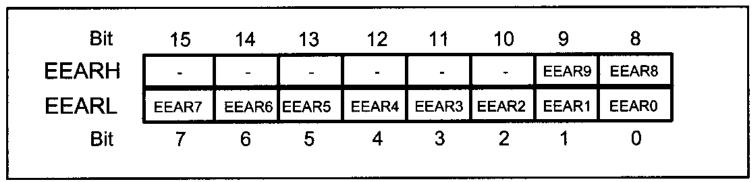


Figure 6-15. EEPROM Address Registers

EEPROM Control Register (EECR)

The EECR register is used to select the kind of operation to perform on. The operation can be start, read, and write. In Figure 6-16 you see the bits of the EECR register. The bits are as follows:

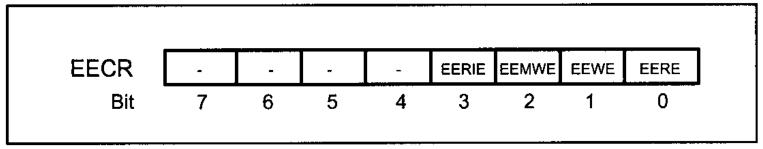


Figure 6-16. EEPROM Control Registers

EEPROM Read Enable (EERE): Setting this bit to one will cause a read operation if EEWE is zero. When a read operation starts, one byte of EEPROM will be read into the EEPROM Data Register (EEDR). The EEAR register specifies the address of the desired byte.

EEPROM Write Enable (EEWE) and EEPROM Master Write Enable (EEMWE): When EEMWE is set, setting EEWE within four clock cycles will start a write operation. If EEMWE is zero, setting EEWE to one will have no effect.

When you set EEMWE to one, the hardware clears the bit to zero after four clock cycles. This prevents unwanted write operations on EEPROM contents. Notice that you cannot start read or write operations before the last write operation is finished. You can check for this by polling the EEWE bit. If EEWE is zero it means that EEPROM is ready to start a new read or write operation.

EEPROM Ready Interrupt Enable (EERIE): In Chapter 10 you will learn about interrupts in AVR.

As you see in Figure 6- 16, bits 4 to 7 of EECR are unused at the present time and are reserved.

Programming the AVR to write on EEPROM

To write on EEPROM the following steps should be followed. Notice that steps 2 and 3 are optional, and the order of the steps is not important. Also note that you cannot do anything between step 4 and step 5 because the hardware clears the EEMWE bit to zero after four clock cycles.

- 1. Wait until EEWE becomes zero.
- 2. Write new EEPROM address to EEAR (optional).
- 3. Write new EEPROM data to EEDR (optional).
- 4. Set the EEMWE bit to one (in EECR register).
- 5. Within four clock cycles after setting EEMWE, set EEWE to one.

See Example 6-28 to see how we write a byte on EEPROM.

Example 6-28

.INCLUDE "M16DEF.INC"

EEDR, R16

EECR, EEWE

Write an AVR program to store 'G' into location 0x005F of EEPROM.

Solution:

OUT SBI

SBI

```
:wait for last write to finish
WAIT:
SBIC EECR, EEWE
                :check EEWE to see if last write is finished
RJMP WAIT
                :wait more
     R18,0 ;load high byte of address to R18
LDI
    R17,0x5F ; load low byte of address to R17
LDI
    EEARH, R18 ; load high byte of address to EEARH
OUT
    EEARL, R17 ; load low byte of address to EEARL
OUT
    R16, 'G' ; load 'G' to R16
LDI
```

Run and simulate the code on AVR Studio to see how the content of the EEPROM changes after the last line of code. Enter four NOP instructions before the last line, change the 'G' to 'H', and run the code again. Explain why the code doesn't store 'H' at location 0x005F of EEPROM.

; load R16 to EEPROM Data Register

EECR, EEMWE ; set Master Write Enable to one

;set Write Enable to one

Programming the AVR to read from EEPROM

To read from EEPROM the following steps should be taken. Note that step 2 is optional.

- 1. Wait until EEWE becomes zero.
- 2. Write new EEPROM address to EEAR (optional).
- 3. Set the EERE bit to one.
- 4. Read EEPROM data from EEDR.

See Example 6-29 to see how we read a byte from EEPROM.

Example 6-29

Write an AVR program to read the content of location 0x005F of EEPROM into PORTB.

Solution:

```
.INCLUDE "M16DEF.INC"
           R16.0xFF
     LDI
     OUT
           DDRB, R16
WAIT:
                       :wait for last write to finish
     SBIC EECR, EEWE
                       ; check EEWE to see if last write is finished
     RJMP
          WAIT
                       :wait more
           R18,0
     LDI
                       ; load high byte of address to R18
                       ; load low byte of address to R17
     LDI R17.0x5F
     OUT
         EEARH, R18 ; load high byte of address to EEARH
     OUT
           EEARL, R17
                       ; load low byte of address to EEARL
     SBI
           EECR, EERE
                       ;set Read Enable to one
           R16, EEDR ; load EEPROM Data Register to R16
     IN
           PORTB, R16
                       ;out R16 to PORTB
     OUT
```

Initializing EEPROM

In Section 6-4, you saw how to allocate program memory using the .DB directive. We can also allocate and initialize the EEPROM using the .DB directive. If we write .ESEG before a definition, the variable will be located in the EEPROM, whereas .CSEG before a definition causes the variable to be allocated in the code (program) memory. By default the variables are located in the program memory. For example, the following code allocates locations \$10 and \$11 of EEPROM for DATAl and DATA2, and initializes them with \$95 and \$19, respectively:

.ESEG .ORG \$10 DATA1: .DB \$95 DATA2: .DB \$19

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6.6 ACCESSING EEPROM IN AVR

The following code allocates DATA1 and DATA3 in program memory and DATA2 in EEPROM:

```
; by default it is located in code memory
DATA1:
                 $10
           .DB
     .ESEG
                 $20 ;it is located in EEPROM
DATA2:
           .DB
DATA3:
           .DB
                 $35 ;it is located in EEPROM
      .CSEG
DATA4:
                 $45 ;it is located in code memory
           .DB
```

Example 6-30

Write a program that counts how many times a system has been powered up.

Solution:

```
.INCLUDE "M32DEF.INC"
      LDI
            R20, HIGH (RAMEND)
     OUT
           SPH,R20
           R20, LOW (RAMEND)
      LDI
           SPL,R20
                              ;initialize stack pointer
      OUT
            XH, HIGH (COUNTER)
      LDI
      LDI
           XL, LOW (COUNTER)
                              ;X points to COUNTER
      CALL LOAD FROM EEPROM
                              ; load R20 with value of COUNTER
      INC
            R20
                              :increment R20
      CALL STORE IN EEPROM
                              ;store R20 in EEPROM
HERE: RJMP HERE
```

```
:----Load R20 with contents of location X of EEPROM
LOAD FROM EEPROM:
     SBIC EECR. EEWE
     RJMP LOAD FROM EEPROM ; wait while EEPROM is busy
     OUT EEARH, XH
     OUT EEARL, XL ; EEAR = X
     SBI EECR, EERE ; set Read Enable to one
     IN R20, EEDR
                           ;load EEPROM Data Register to r20
     RET
;----Store R20 into location X of EEPROM
STORE IN EEPROM:
     SBIC EECR, EEWE
     RJMP STORE IN EEPROM ; wait while EEPROM is busy
     OUT EEARH, XH
                           ; EEAR = X
     OUT EEARL, XL
     OUT EEDR, R20
     SBI EECR, EEMWE ; set Master Write Enable to one
     SBI EECR, EEWE
                           ;write EEDR into EEPROM
     RET
; -----EEPROM
.ESEG
.ORG 0
COUNTER: .DB 0
COUNTER is initialized with $0. Then, it is incremented on each power-up.
```

Checksum byte in EEPROM

The checksum will detect any corruption of the contents of ROM. One cause of ROM corruption is current surge, either when the system is turned on, or during operation. To ensure data integrity in ROM, the checksum process uses what is called a *checksum byte*. *The checksum byte is an extra byte* that is tagged to the end of a series of bytes of data. To calculate the checksum byte of a series of bytes of data, the following steps can be taken:

- 1. Add the bytes together and drop the carries.
- 2. Take the 2's complement of the total sum, and that is the checksum byte, which becomes the last byte of the series.

To perform a checksum operation, add all the bytes, including the checksum byte. The result must be zero. If it is not zero, one or more bytes of data have been changed (corrupted).

Checksum program

The checksum generation and testing program is given in subroutine form. Five subroutines perform the following operations:

- 1. Retrieve the data from EEPROM.
- 2. Test the checksum byte for any data error.
- 3. Initialize variables if the checksum byte is corrupted.
- 4. Calculate the checksum byte.
- 5. Store the data in EEPROM.

Each of these subroutines can be used in other applications.

Example 6-31 shows how to manually calculate the checksum for a list of values.

Example 6-31

Assume that we have 4 bytes of hexadecimal data: \$25, \$62, \$3F, and \$52.

- (a) Find the checksum byte.
- (b) Perform the checksum operation to ensure data integrity.
- (c) If the second byte, \$62, has been changed to \$22, show how the checksum method detects the error.

Solution:

(a) Find the checksum byte.

\$25

+ \$62

+ \$3F

+ \$52

\$118

(Dropping the carry of 1, we have \$18. Its 2's complement is \$E8. Therefore, the checksum byte is \$E8.)

(b) Perform the checksum operation to ensure data integrity.

\$25

+ \$62

+ \$3F

+ \$52

+ \$E8

\$200 (Dropping the carries, we see 00, indicating that the data is not corrupted.)

(c) If the second byte, \$62, has been changed to \$22, show how the checksum method detects the error.

\$25

+ \$22

+ \$3F

+ \$52

+ \$E8

\$1C0 (Dropping the carry, we get \$C0, which is not 00. This means that the data is corrupted.)

BCD to ASCII conversion program

Many RTCs (real-time clocks) provide time and date in BCD format. To display the BCD data on an LCD or a PC screen, we need to convert it to ASCII. Program 6-2

- (a) transfers packed BCD data from program ROM to data RAM,
- (b) converts packed BCD to ASCII, and
- (c) sends the ASCII to port B for display.

Binary (hex) to ASCII conversion program

Many ADC (analog-to-digital converter) chips provide output data in binary (hex). To display the data on an LCD or PC screen, we need to convert it to ASCII. The code for the binary-to-ASCII conversion is shown in Program 6-3.

Notice that the subroutine gets a byte of 8-bit binary (hex) data from Port B and converts it to decimal digits, and the second subroutine converts the decimal digits to ASCII digits and saves them.

We are saving the low digit in the lower address location and the high digit in higher address location.

This is referred to as the little-endian convention (i.e., low byte to low location, and high byte to high location).

What is a macro and how is it used?

Macros allow the programmer to write the task (code to perform a specific job) once only, and to invoke it whenever it is needed.

Macro definition

Every macro definition must have three parts, as follows:

.MACRO name

•••••

. ENDMACRO

What goes between the .MACRO and .ENDMACRO directives is called the *body* of the macro. The name must be unique and must follow Assembly language naming conventions. A macro can take up to 10 parameters. The parameters can be referred to as @0 to @9 in the body of the macro.

For example, moving immediate data into I/O register data RAM is a widely used service, but there is no instruction for that. We can use a macro to do the job as shown in the following code:

.MACRO	LOADIC			
LDI	R20,@1			
OUT	@0,R20			
.ENDMACRO				

The following are three examples of how to use the above macro:

```
    LOADIO PORTA, 0x20 ;send value 0x20 to PORTA
    .EQU VAL_1 = 0xFF
        LOADIO DDRC, VAL_1
    LOADIO SPL, 0x55 ;send value $55 to SPL
```

```
AVR A[; Program 6-4: toggling Port B using macros ] \MMING
       .INCLUDE "M32DEF.INC"
       .MACRO LOADIO
            LDI R20,01
Now exai
            OUT @0, R20
       .ENDMACRO
         -----time delay macro
       .MACRO DELAY
            LDI @0,@1
       BACK:
            NOP
            NOP
            NOP
            NOP
            DEC
                @ ()
            BRNE BACK
       .ENDMACRO
                   ----- starts
             .ORG 0
            LOADIO DDRB, 0xFF ; make PORTB output
            LOADIO PORTB, 0x55 ; PORTB = 0x55
       L1:
            DELAY R18,0x70 ;delay
            LOADIO PORTB, 0xAA ; PB = 0xAA
mashhoun@i
            DELAY R18,0x70
                                 ;delay
                                                      11/27/2023
```

RJMP

L1

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.INCLUDE directive

The .INCLUDE directive allows a programmer to write macros and save them in a file, and later bring them into any program file. For example, assume that the following widely used macros were written and then saved under the filename "MYMACRO1.MAC".

Assuming that the LOADIO and DELAY macros are saved on a disk under the filename "MYMACRO.MAC", the .INCLUDE directive can be used to bring this file into any ".asm" file and then the program can call upon any of the macros as many times as needed. When a file includes all macros, the macros are listed at the beginning of the ".lst" file and, as they are expanded, will be part of the program.

To understand this, see Program 6-5.

```
;Program 6-5: toggling Port B using macros
.INCLUDE "M32DEF.INC"
.INCLUDE "MYMACRO1.MAC" ;get macros from macro file
;------program starts
.ORG 0
LOADIO DDRB,0xFF
L1: LOADIO PORTB,0x55
DELAY R18,0x70
LOADIO PORTB,0xAA
DELAY R18,0x70
RJMP L1
```

. LISTMAC directive

Using the .LISTMAC directive we can **turn on the display of the bodies of macros in the list file. For example, examine** the following
code:

The assembler provides the following code in the .lst file:

						.MACRO	LOADIO
.INCLU	JDE "M3	2DEF.INC"				LDI	R20,@1
.MACRO		LOADIO				TUO	@0,R20
	LDI	R20,@1				.ENDMACRO	
	OUT	@0,R20					
.ENDMA	ACRO		000000	e240			
			000001	bb4b		LOADIO	PORTA,0x20
	LOADIO	PORTA, 0x20	000002	e543			
	LOADIO	DDRA,0x53	000003	bb4a		LOADIO	DDRA, 0×53
HERE:	JMP	HERE	000004	940c	0004	HERE:JMP	HERE

If we add the .LISTMAC directive to the above code:

```
.MACRO
                                                                    LOADIO
.MACRO
              LOADIO
                                                                    R20,01
                                                              LDI
                                                                    @0,R20
       LDI
              R20,01
                                                              OUT
              @0,R20
                                                         .ENDMACRO
       OUT
                                                         .LISTMAC
.ENDMACRO
.LISTMAC
                                                      +LDI R20 , 0x20
                                     000000 e240
       LOADIO
                     PORTA, 0x20
                                                      +OUT PORTA , R20
                                     000001 bb4b
                                                                           PORTA, 0x20
                      DDRA, 0x53
                                                              LOADIO
       LOADIO
HERE: JMP
              HERE
                                     000002 e543
                                                      +LDI R20 , 0x53
                                     000003 bb4a
                                                      +OUT DDRA , R20
                                                                           DDRA, 0 \times 53
                                                              LOADIO
                                     000004 940c 0004 HERE:JMP
                                                                    HERE
```

The + indicates that the code is from the macro.