

#### PART A – TABLE EXPLANATIONS

This project is to design a relational backend for an online service that allows people to communicate with others in their neighborhoods. For this project, the users can register, login, message other users who are their friends, add neighbors, reply to message in threads, apply to be in a block through a member only voting system. These are some of the functionalities. Below, I will explain the reasons behind how I establish a series of tables to manage a community-based platform where users can register, apply to specific blocks within neighborhoods, and interact through messages and relationships. I am using PostgreSQL, Flask in python to construct this application.

1. Users: This table is the center of my platform system. It stores user's unique SERIAL ID, email, password(hashed for security purpose), and their address.

```
CREATE TABLE Users (
3.
       user_id SERIAL PRIMARY KEY,
4.
       username VARCHAR(255) UNIQUE NOT NULL,
5.
       email VARCHAR(255) UNIQUE NOT NULL,
       password hash VARCHAR(255) NOT NULL,
6.
       address VARCHAR(255) NOT NULL,
7.
8.
       profile_text TEXT,
9.
       photo url VARCHAR(255)
10.);
```

Figure 1.1

The application will allow users to create user profile, update user information and delete user profiles. In the next stage of the project, I will have address in a separate table along with coordinates for more precise filtering of messages relating to a certain area of the neighborhood. Upon query 'SELECT \* FROM Users', you can see all users, address(to the right, not shown) and hashed password. Figure 1.2

13	13	wearaszdf	sdfgxg@sdrf.com	scrypt:32768:8:1\$03Rz220C1437qHW7\$e5d53ca27773fd1d909d86857380099133a2b6419e40816beb917eadfb3c5e13bcb3
14	14	first	first@first.com	scrypt: 32768: 8:1\$i9VFBstN2k4bkcST\$03d237be72a40944aba4b8232cb5cf15b6ca5b31ae4c6990653467472a957ad5b6e81e82b2b2b2b2b2b2b2b2b2b2b2b2b2b2b2b2b2b2b
15	15	test1	test1@test1.com	scrypt:32768:8:1\$FeA095QuVQkgt8FL\$868afc4d36023440f668a6a7ea6b678a99fbc7949cb347568b30b54b6c4664c08d405e
16	16	test2	test2.test2@example.com	scrypt:32768:8:1\$HSitle4XaiL9iKXx\$0d04e6bd86d63f480f17cfee770e8052d2864fa5f2eba80921240c9b7a2b5ea6b84a21d63
17	17	test3	test3.test3@example.com	scrypt:32768:8:1\$hqs904GecM25YQRc\$5a61b6cc5253c343923cd786b951708bb7f76a179514a173f11a5fee08426b41b575b
18	18	test4	test4.test4@example.com	scrypt:32768:8:1\$aJ5g9pNTMK5Vy7qm\$85b27022962fa4e69e0f7788c364a5ae249d2011ccf2b030fd76e7427769459a0c024
19	19	test5	test5.test5@example.com	scrypt:32768:8:1\$P5dF8bZiPSxzf2NU\$ad7bfc4859792c80919c907f1d9f8b8e33170cbbd5b41860846f9392efefe5471864096
20	20	test6	test6.test6@example.com	scrypt:32768:8:1\$w1oh94n6tWPpQViz\$28a4e5ef70bfcea59a9eba9fb27ac559f4df66ead4c3720ac8e83288d171030d6594a47

Figure 1.2

2. Neighborhoods: Neighborhoods are preexisting divisions with predetermined unique neighborhood name, description, and unique id. No users can change or add neighborhoods. In my project. Figure 2.2 shows a sample of my predefined neighborhoods.

```
CREATE TABLE Neighborhoods (
neighborhood_id SERIAL PRIMARY KEY,
name VARCHAR(200) UNIQUE NOT NULL,
description TEXT
);
```

Figure 2.1

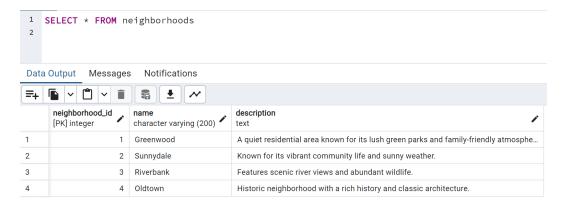


Figure 2.2

3. Blocks: Related to neighborhoods, this table stores details about specific blocks within neighborhoods, including a unique ID, the associated neighborhoodID, geographical center point, radius (for postGIS mapping purposes), and a textual description. For simplicity sake, I divided the community into four quadrants with two blocks in each quadrants. Figure 3.2 shows all predetermined blocks.

```
CREATE TABLE Blocks (
    block_id SERIAL PRIMARY KEY,
    neighborhood_id INT NOT NULL,
    name VARCHAR (200) NOT NULL,
    description VARCHAR (200),
    latitude FLOAT,
    longitude FLOAT,
    radius FLOAT,
    column location geography(Point, 4326),
    -- use postGIS extension for radius
    FOREIGN KEY (neighborhood_id) REFERENCES Neighborhoods(neighborhood_id)
```

Figure 3.1

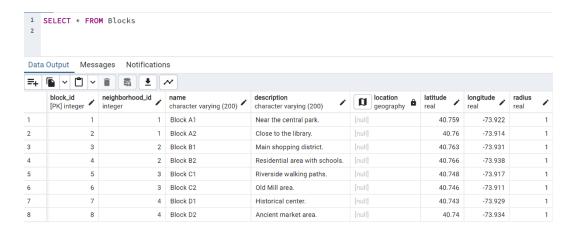


Figure 3.2

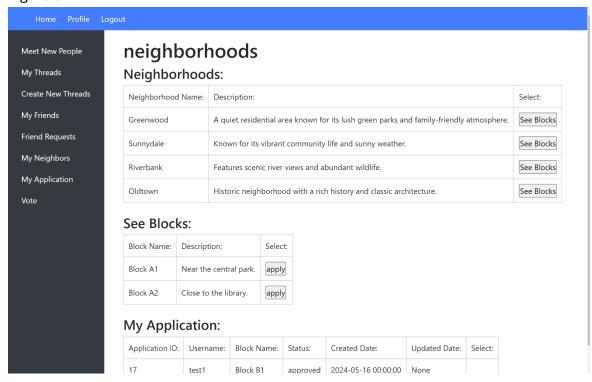


Figure 3.3

4. Block Memberships: Users should apply for block memberships. This table contains all the users\_id and their associated blocks. Furthermore, a column called status to indicate whether the user is an active member of the block or an in-active member of the block, in which case it would mean that the user has moved away from the block and no longer a member. When first applied, the status is default to pending.

```
CREATE TABLE Block_Memberships (
membership_id SERIAL PRIMARY KEY,
block_id INT NOT NULL,
```

```
user_id INT NOT NULL,
    status VARCHAR(20) DEFAULT 'pending' NOT NULL CHECK (status IN ('active',
'in-active')),
    join_date DATE NOT NULL,
    -- if membership canceled.
    update_date Date NOT NULL,
    FOREIGN KEY (block_id) REFERENCES Blocks(block_id),
    FOREIGN KEY (user_id) REFERENCES Users(user_id)
);
```

Figure 4.1

The block membership keeps a record of all user's status. The status is updated based on the approval of the application table via membership\_id which we will discuss next. There will be a delete membership option in which case, the record will update to inactive but the record is still kept for record keeping sake. Figure 4.2 shows a sample of all users with their membership status associated with a block. Once a user's membership status gets approved, the block name will appear on users' profile page along side their username, address, and email address etc.

1 2	SELECT * FROM Block_memberships								
Data	Data Output Messages Notifications								
	membership_id [PK] integer	block_id integer /	user_id integer	status character varying (20)	join_date /	update_date /			
1	1	1	1	active	2023-01-15	2023-01-15			
2	2	1	2	active	2023-02-20	2023-02-20			
3	3	2	3	active	2023-03-05	2023-04-01			
4	4	2	4	in-active	2023-03-05	2023-06-01			
5	5	3	5	active	2023-04-10	2023-04-10			
6	6	4	6	active	2023-05-15	2023-05-15			
7	7	4	7	active	2023-07-20	2023-07-20			
8	8	5	8	active	2023-08-25	2023-08-25			
9	9	6	9	in-active	2023-08-25	2023-09-30			
10	10	7	10	active	2023-08-25	2023-10-15			

Figure 4.2

5. Applications: Each applications record is uniquely identified by applicationID and it includes references to UserID for the applicant and blockID for the block they wish to join. The 'status' field indicates the current state of the application. It is default to 'pending' upon the creation of an instance. The status of the application can be changed based on the aggregated query of the Votes table. Figure 5.2 shows a sample of a newly created application that is pending.

```
application_id SERIAL PRIMARY KEY,
  block_id INT NOT NULL,
  applicant_id INT NOT NULL,
  status VARCHAR(20) DEFAULT 'pending' CHECK (status IN ('pending', 'approved',
  'rejected')),
  created_date TIMESTAMP DEFAULT CURRENT_TIMESTAMP,
  updated_date date,
  FOREIGN KEY (block_id) REFERENCES Blocks(block_id),
  FOREIGN KEY (applicant_id) REFERENCES Users(user_id)
);
```

Figure 5.1

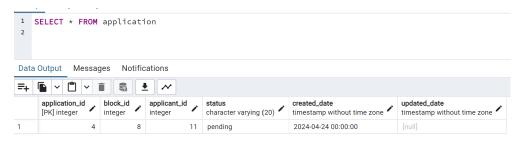


Figure 5.2

# My Application:

Application ID:	Username:	Block Name:	Status:	Created Date:	Updated Date:	Select:	
17	test1	Block A1	approved	2024-05-16 00:00:00	None		
Your application	n has been u	pdated to the n	ew block.				×

Figure 5.3 User may apply for a block, choose a different block while keeping the older application until the new one gets approved and the older application is dropped.

6. Votes: This table ties votes to a specific application a user makes to join a block. It will store the applicationID, voter\_id, status, created\_at, and updated\_at timestamps. Since there must be three votes who are already members of a block, the first three members are preexisting admin who are already in the block eligible to admit the first new member. From then on, all future voters would have to have the same block\_id as the associated application block\_id to vote. The voting system is simple. If a qualified user approves an application the vote\_count will increase by 1. There is a query that checks all votes to a particular application\_id and sums up the vote\_count. As soon as the vote count becomes greater or equal to 3, it triggers the approval of the application, changes the status of the application to approved, and the block\_memberships status becomes active.

```
CREATE TABLE Votes (
    application_id INT NOT NULL,
    voter_id INT NOT NULL,
    -- plus 1
    vote_count INT
    FOREIGN KEY (voter_id) REFERENCES Users(user_id),
    FOREIGN KEY (application_id) REFERENCES Application(application_id),
    -- each voter can vote once for each application
    PRIMARY KEY (application_id, voter_id),
    created_at TIMESTAMP,
    updated_at TIMESTAMP
);
```

Figure 6.1

```
cursor.execute("""SELECT SUM(CAST(vote_count AS INTEGER))
AS total votes
                                          FROM Votes
                                          WHERE application_id = %s;""",
(application id,))
                        current vote result = cursor.fetchone()
                        current vote = current vote result['total votes'] if
current vote result['total votes'] else 0
                        # Insert new vote
                        cursor.execute("""INSERT INTO Votes (application id,
voter_id, vote_count, created_at, updated_at)
                                          VALUES (%s, %s, %s, CURRENT TIMESTAMP,
CURRENT TIMESTAMP);"",
                                       (application id, session['user id'],
current_vote + 1))
                        # CHECK VOTE AGAIN
                        cursor.execute("""SELECT SUM(CAST(vote count AS INTEGER))
AS total votes
                                          FROM Votes
                                          WHERE application id = %s;""",
(application id,))
                        current vote result = cursor.fetchone()
                        new_vote = current_vote_result['total_votes'] if
current_vote_result['total_votes'] else 0
                        if new vote >= 3:
                            cursor.execute("""UPDATE Application
                                              SET status = 'approved'
```

Figure 6.2 current\_vote\_result checks the current tallie and stores it in a variable. Upon member approval, it increments the tallie by 1 and makes a comparison. If >=3, it triggers the update of application status and insertion of applicant into block\_membership.

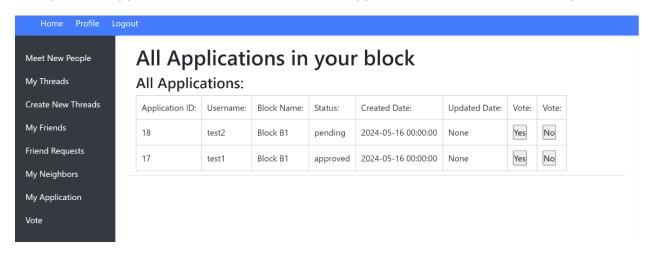


Figure 6.3 All votable applications is accessible to all qualified users.

7. Friendships: This table captures friendship relationships, which are mutual. If User A wants to be friends with User B, then the status is pending until User B accepts User A as friends. For efficiency, although the columns store user\_id1 and user\_id2 and it is always the case that user\_id1 is the user who makes the request. When storing, we store the smaller value as user\_id1 via a CHECK to ensure that user\_id is always less than user\_id2. As soon as user1 makes a friend request, it pops off the list of names that are either 'rejected', 'nill', or 'pending'. If user\_2 rejects friend request, user\_2 will appear back on user\_1's 'Meet new people' list.

```
CREATE TABLE Friendships (
    user_id1 INT NOT NULL,
    user_id2 INT NOT NULL,
    status VARCHAR(20) DEFAULT 'pending' CHECK (status IN ('pending', 'accepted',
    'rejected')),
    created_date TIMESTAMP DEFAULT CURRENT_TIMESTAMP,
    update_date TIMESTAMP DEFAULT CURRENT_TIMESTAMP,
```

```
-- when storng, store as smaller value as user_id1
   CHECK (user_id1 < user_id2), -- Ensure that user_id1 is always less than
user_id2
   FOREIGN KEY (user_id1) REFERENCES Users(user_id),
   FOREIGN KEY (user_id2) REFERENCES Users(user_id)
);</pre>
```

Figure 7.1

1	SELECT * FROM friendships								
Data	Data Output Messages Notifications								
=+									
	user_id1 integer	user_id2 integer	status character varying (20)	created_date timestamp without time zone	update_date timestamp without time zone				
1	1	11	accepted	2024-04-24 03:34:32.660079	2024-04-24 03:36:35.775566				
2	1	2	pending	2024-04-24 17:46:19.497431	2024-04-24 17:46:19.497431				
3	1	2	accepted	2024-04-24 17:46:19.497431	2024-04-24 17:46:19.497431				

Figure 7.2

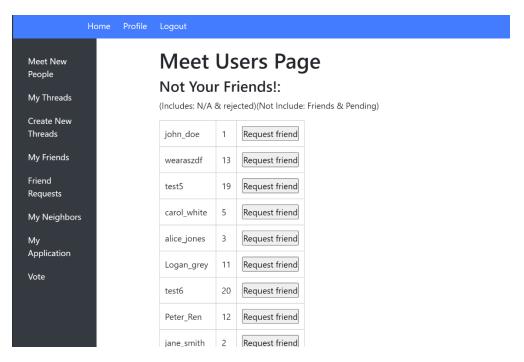
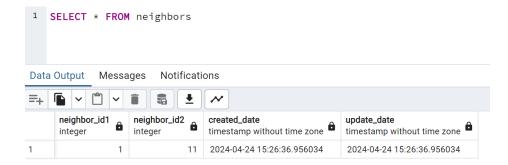


Figure 7.3

8. Neighbors: Similar to friends, this table captures neighbor relationships between two users, neighbor\_id1 and neighbor\_id2. However, the key difference is that neighbor is one directional. Neighbor 1 can assign neighbor 2 only. If the user of neighbor\_id2 wants to have a neighbor relationship with the user of neighbor\_id1, then user 2 would have to set

up a new record stating this intent. Here, it is ok to have neighbor 1 and neighbor 2 to have a duplicated relationship since it is recommended or allowed that both can request to be neighbor of each other. Mutual consent is not required.

```
CREATE TABLE Neighbors (
    neighbor_id1 INT NOT NULL,
    neighbor_id2 INT NOT NULL,
    created_date TIMESTAMP DEFAULT CURRENT_TIMESTAMP,
    update_date TIMESTAMP DEFAULT CURRENT_TIMESTAMP,
    FOREIGN KEY (neighbor_id1) REFERENCES Users(user_id),
    FOREIGN KEY (neighbor_id2) REFERENCES Users(user_id));
```



Figures 8.1 and 8.2

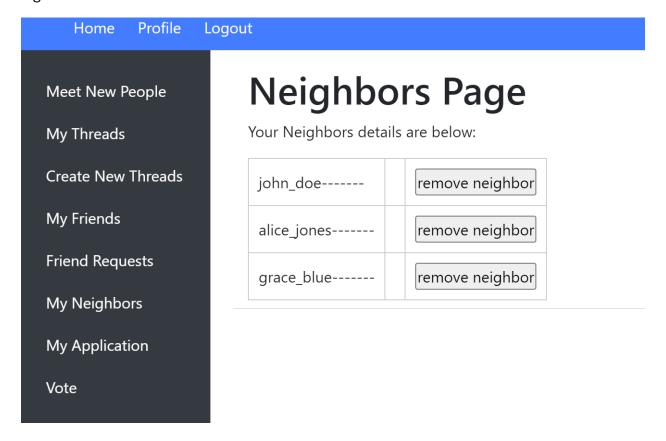


Figure 8.3, Neighbor list is instantly populated upon adding neighbor to your list. Once, user1 added user2, user2 can begin sending messages to user1.

9. Thread: Thread works a lot like emails. This table contains all the threads identified by unique threadID. There is no need for the id of the original message since all message is treated as equal and differentiated by each unique timestamp and messageIDs. There is a button for "New Thread" to create a new thread. In addition, sender needs to specify the recipient type of the message. Is it personal friends, neighbor, or block. This will determine the recipient\_id so the system know it is dealing with a user\_id, block\_id or a neighbor\_id; each will be a different function. Based om the reference ID, the system will query all the users who are in that block\_id or friend\_ID and display the message to them.

The Thread creation also returns the threadID, used for all message associated with this thread. I will demonstrate this later.

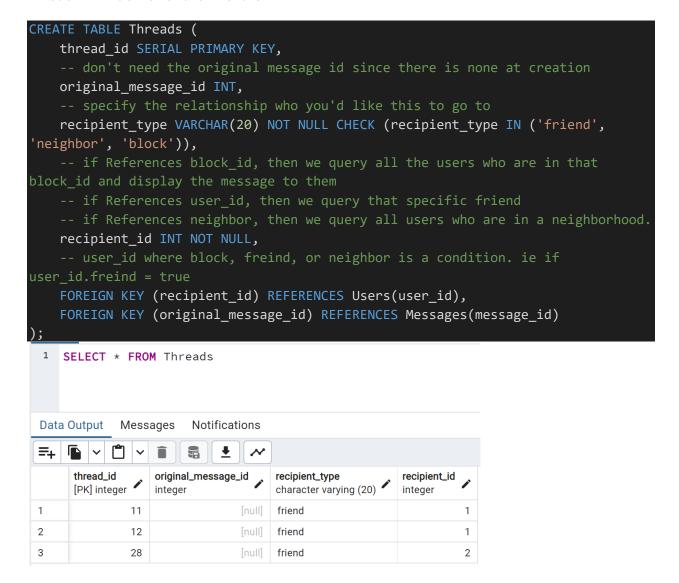


Figure 9.1 and 9.2

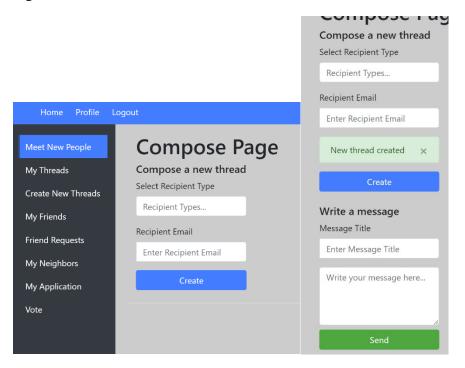


Figure 9.3, 9.4. Only friends can start threads. If check friend, a pop up textbox appears for user1 to write to user2 via the /newThread route.

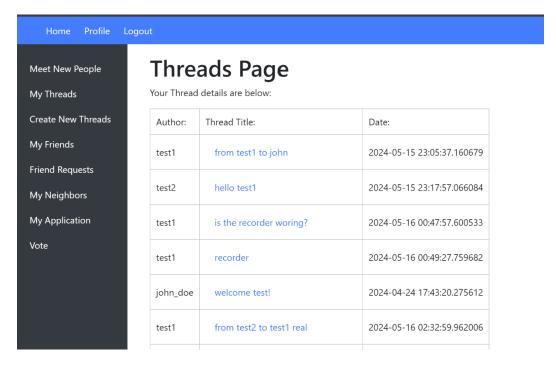


Figure 9.5 This page contains the title of each thread. User can click on the title of the thread to view all messages.

10.Messages. Manages messages posted by users. Each message is linked to a specific threadID and includes detailes like the author's ID, timestamp, title, body of the message, and coordinates related to the message content. The send\_to and reply\_to are IDs that tracks who are the target audience, ensuring only a certain group of people can view and reply to a particular message.

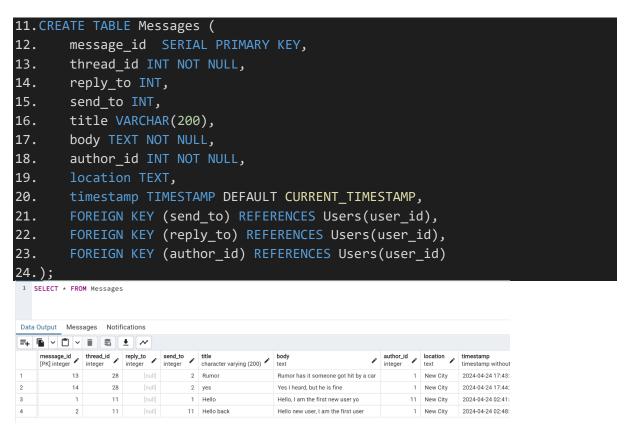
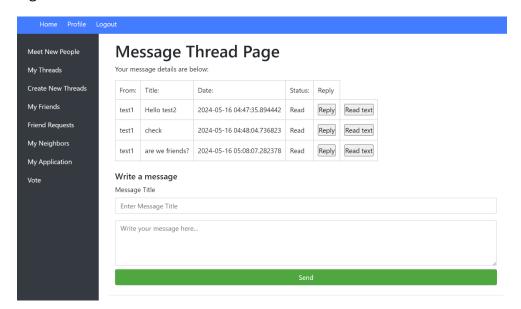


Figure 10.1 and 10.2



- Figure 10.3 user can view all messages in a thread. By clicking 'Read text' the body of the message will appear and the record\_read\_status will mark the status of the message read by session[user\_id]. User can also reply to a message, grabbing the message id from the click action and the thread\_id from the url.
- 11. Read\_status\_recorder: a table that stores information that tracks whether a message has been read or not. It tracks the message\_id and reader\_id so that upon loggin\_in, it remembers what th user has seen and only display the 'unread' message via its status.

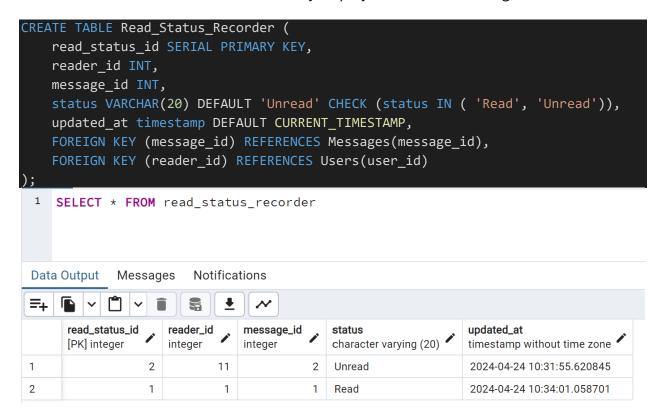


Figure 11.1 and 11.2

Figure 11.3 Upon clicking 'read text', an update query change the status to 'Read'.

#### PART B - SAMPLE QUERIRES

- 1) Create new user profile:
- 2) INSERT INTO Users (username, email, password\_hash, address, profile\_text, photo\_url) VALUES
- 3) ('Logan\_grey', 'logan.grey@example.com', 'hashpassword11', '104 Royal St,
   Beachside', 'Politician', 'http://example.com/photos/logangrey.jpg');
- 2) Update user information

```
UPDATE Users
SET email = 'newLogan.grey@example.com', address = '456 New Address St, New City'
```

- 3) Apply to block\_membership
  - 4) INSERT INTO Application (block\_id, applicant\_id) VALUES5) (8, 11);
  - 6) WHERE user\_id = 11; -- Assuming you know the user's ID is 11

    System will fetch block\_id from blocks where block name is "block D2". This will return block\_ID = 8. System will also fetch application\_id from Users where username is 'Logan Grey'. This will return 11.
- 4) Create new Thread

```
INSERT INTO Threads (recipient_type, recipient_id) Values ('friend', '1')
RETURNING thread_id;
```

Click 'create new thread', fetch 'type' by selecting from a drop down menu of ('friend', 'neighbor', or block). Fetch recipient\_id (user\_id) via email if type = friend or neighborhood\_id if type = neighborhood, or block\_id if type = block. The query will only query the user with the same recipient-type to show them the message. Finally, fetch RETURNING thread\_id from previously created thread, store in a variable for all future message related to this thread.

5) Create new message

```
INSERT INTO Messages (thread_id, send_to, title, body, author_id, location)
VALUES
(11, 1, 'Hello', 'Hello, I am the first new user', 11, 'New City');
```

Fetch recipient\_if from thread as send\_to. Fetch thread\_id previously stored as a variable, in this case, thread\_id = 11. And Author is the current user.

#### 6) reply to message

```
INSERT INTO Messages (thread_id, send_to, title, body, author_id, location)
VALUES
(11, 11, 'Hello back', 'Hello new user, I am the first user', 1, 'New City');
```

UserId = 1 is going to reply to userID 11 under the same thread. This page will only show messages with the same thread\_id.

## 7) Make friend request

User\_id1 makes friend request to User\_id2. If user\_id2 accepts, it updates this instance's status to 'accepted'. Again, we can take any user id but we store it as lest to greatest for easier filter later.

# 8) update friendship

```
UPDATE Friendships
SET status = 'accepted',
    update_date = CURRENT_TIMESTAMP
WHERE user_id1 = 1
    AND user_id2 = 11
```

9) become beighbors

```
INSERT INTO neighbors (neighbor_id1, neighbor_id2) VALUES
(1,11);
```

Neighbor relationships are not mutal, so there is no need for status update.

# 10) List all friends of user\_id11

```
SELECT
    CASE
        WHEN f.user_id1 = 11 THEN f.user_id2
        ELSE f.user_id1
    END AS friend_id,
    u.username,
    f.status,
    f.created_date,
    f.update_date
```

```
FROM Friendships f

JOIN Users u ON u.user_id = CASE

WHEN f.user_id1 = 11 THEN f.user_id2

ELSE f.user_id1

END

WHERE (11 = f.user_id1 OR 11 = f.user_id2)

AND f.status = 'accepted';

-- return just the user_ids of friends

SELECT

CASE

WHEN f.user_id1 = 11 THEN f.user_id2

ELSE f.user_id1

END AS friend_id

FROM Friendships f

WHERE (11 = f.user_id1 OR 11 = f.user_id2)

AND f.status = 'accepted';
```

Figure 10.1.1

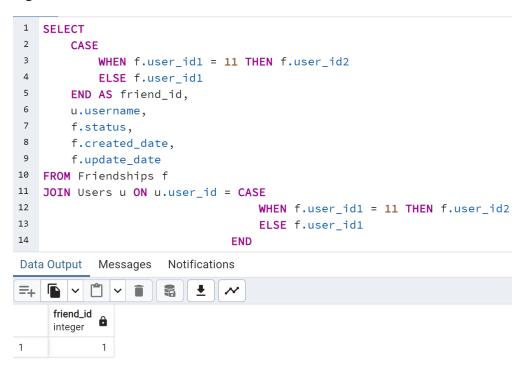


Figure 10.1.2

Because I have user\_id1 and user\_id2, I need to check all permutations of user 11 being either userid\_1 or userId\_2. The their status is accepted then return the user ids of just the friends. The query is seen in figure 10.1.1 and the resulting 1 friend can be seen in figure 10.1.2.

11)insert and update new record of read\_status\_recorder

```
INSERT INTO Read_status_recorder (reader_id, message_id) Values
(1, 1),
(11, 2);
-- Update
UPDATE Read_status_recorder
SET status = 'Read',
    updated_at = CURRENT_TIMESTAMP
WHERE reader_id = 1
AND message_id = 1;
```

Result can be seen in figure 11.2

12) Show only unread messages to user\_id=11



Figure 12.1.1 Result of unread messages to user\_id=11

13) Show only unread messages to and from userID = 11 and FROM and TO userID=11's friends

```
SELECT m.message_id
 JOIN read_status_recorder r ON m.message_id = r.message_id
 JOIN (
         SELECT
             CASE
                 WHEN f.user_id1 = 11 THEN f.user_id2
                 ELSE f.user id1
             END AS friend_id
         FROM Friendships f
         WHERE 11 IN (f.user_id1, f.user_id2)
             AND f.status = 'accepted'
 ) friends ON (m.send_to = friends.friend_id AND m.author_id = 11)
 OR (m.author_id = friends.friend_id AND m.send_to = 11)
 Where r.status = 'Unread';
ata Output Messages Notifications
F 🖺 🗸 📋 🗸 🗂 😤 💆 🥢
  message_id
[PK] integer
```

First select only user11's friends. Then join on the condition of the appropriate sender and receiver IDs where the status is 'Unread'.

#### 14) Insert Neighbors

```
INSERT INTO Neighbors (neighbor_id1, neighbor_id2, created_date, update_date)
VALUES (1, 11, CURRENT_TIMESTAMP, CURRENT_TIMESTAMP);
```

To ensure strangers cannot send emails to strangers. Neighbor messages an be viewd and received based on invitation only. More specifically, only user1 can set up and receive messages and only user2 can send messages. If user2 wants to receive messages, user2 needs to setup need relationship with user1 since the relationship is one directional. Users can get each others information in-person for their first meeting.

## 15)SELECT all userId=11's neighbors

```
16)SELECT u.user_id, n.neighbor_id2 FROM Users u
17)JOIN neighbors n ON u.user_id = n.neighbor_id1
18)WHERE u.user_id =1
```

16) Fetch all unread messages across all available types (ie all friend, all neighbors, all block) containing sentence like "Hello"

```
SELECT m.message_id
FROM Messages m
JOIN read_status_recorder r ON m.message_id = r.message_id
LEFT JOIN (
    SELECT
    CASE
```

```
WHEN f.user_id1 = 1 THEN f.user_id2
                         ELSE f.user_id1
                END AS friend id
        FROM Friendships f
        WHERE 1 IN (f.user_id1, f.user_id2)
                AND f.status = 'accepted'
) friends ON (m.send to = friends.friend id AND m.author id = 1)
        OR (m.author_id = friends.friend_id AND m.send_to = 1)
LEFT JOIN (
        SELECT DISTINCT bm.user_id
        FROM Block Memberships bm
        JOIN Block Memberships bm2 ON bm.block id = bm2.block id
        WHERE bm2.user id = 1
) blocks ON (m.send to = blocks.user id AND m.author id = 1)
        OR (m.author_id = blocks.user_id AND m.send_to = 1)
LEFT JOIN (
        SELECT u.user_id, n.neighbor_id1, n.neighbor_id2
        FROM Users u
        JOIN neighbors n ON u.user id = n.neighbor id1
        WHERE u.user id = 1
) neighbors ON (m.send to = neighbors.neighbor id1 AND m.author id = 1)
        OR (m.author_id = neighbors.neighbor_id2 AND m.send_to = 1)
WHERE r.status = 'Unread'
        AND LOWER(body) LIKE LOWER('%Hello %')
        AND (friends.friend_id IS NOT NULL OR blocks.user_id IS NOT NULL OR
neighbors.user_id IS NOT NULL);
     1 SELECT m.message_id
     2 FROM Messages m
        JOIN read_status_recorder r ON m.message_id = r.message_id
            SELECT
                CASE
                     WHEN f.user_id1 = 1 THEN f.user_id2
                     ELSE f.user_id1
                END AS friend_id
           FROM Friendships f
           WHERE 1 IN (f.user_id1, f.user_id2)
AND f.status = 'accepted'
    12
    ) friends ON (m.send_to = friends.friend_id AND m.author_id = 1)
            OR (m.author_id = friends.friend_id AND m.send_to = 1)
          SELECT DISTINCT bm.user_id
    17
             FROM Block_Memberships bm
            JOIN Block_Memberships bm2 ON bm.block_id = bm2.block_id
             WHERE bm2.user_id = 1
    20 ) blocks ON (m.send_to = blocks.user_id AND m.author_id = 1)
             OR (m.author_id = blocks.user_id AND m.send_to = 1)
22 LEFT JOIN (
    Data Output Messages Notifications

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          message_id
[PK] integer
```

This query needs to satisfy either one of the three sub queries, being friend = accepted or block matchs users.block or user\_id matches neighbor\_id. Furthermore, the message as to

be unread and the sentence is changed to lower case for all possible entry of the target words. The three tables are join using LEFT JOIN, allowing for a NULL result if the relationship condition is not met.

#### 17) Location filter

```
CREATE EXTENSION postgis;
CREATE TABLE blocks_area (
    blocks area id SERIAL PRIMARY KEY,
    geom GEOMETRY(Point, 4326)
);
-- Insert latitude and longitude values into the location column
UPDATE Blocks
SET location = ST_SetSRID(ST_MakePoint(longitude, latitude), 4326);
--FILTER DISTNACES
SELECT *
FROM Messages
WHERE ST_DWithin(
    location::geography, -- Cast the location column to geography type
    ST_SetSRID(ST_MakePoint(-73.9857, 40.7484), 4326)::geography, -- Cast the
point to geography type
    200 -- Distance in meters
```

Added colums for latitude and longitude to the blocks table. The data type for these columns are 'geography(Point, 4326) which specifies the spatial reference system used by PostGIS. For the final project, I would like to use google map to test exact longitude and latitude values using the PostGIS functions. '200' is the distance threshold in meters from the user's home.

The setup handles the users interactions within the platform. Improvements should be made regarding the scalability of certain tables like read\_status\_recorder as it is not useful to keep a table of all the read and unread status in long-term storage. Coordinates need to be further tested with real world values and make the corresponding adjustment to the block and message tables where more precise location is needed.