

# Graph Theory

Dr. Irfan Yousuf

Department of Computer Science (New Campus)

UET, Lahore

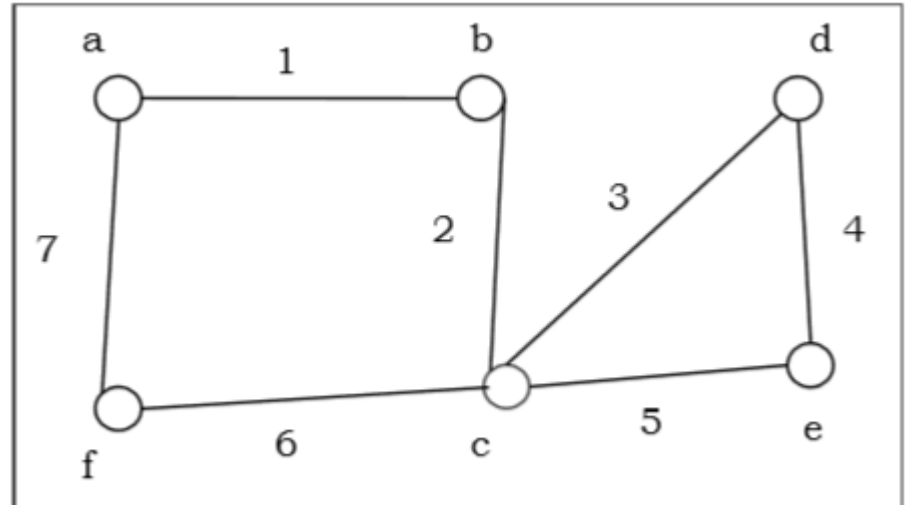
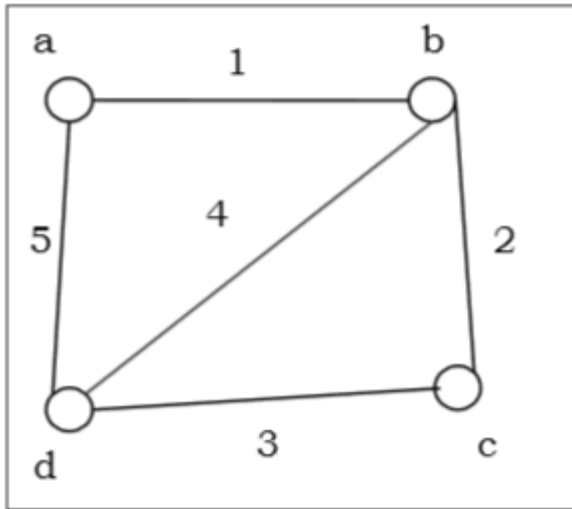
(Lecture # 17; March 22, 2023)

# Outline

- Eulerian and Hamiltonian Graphs

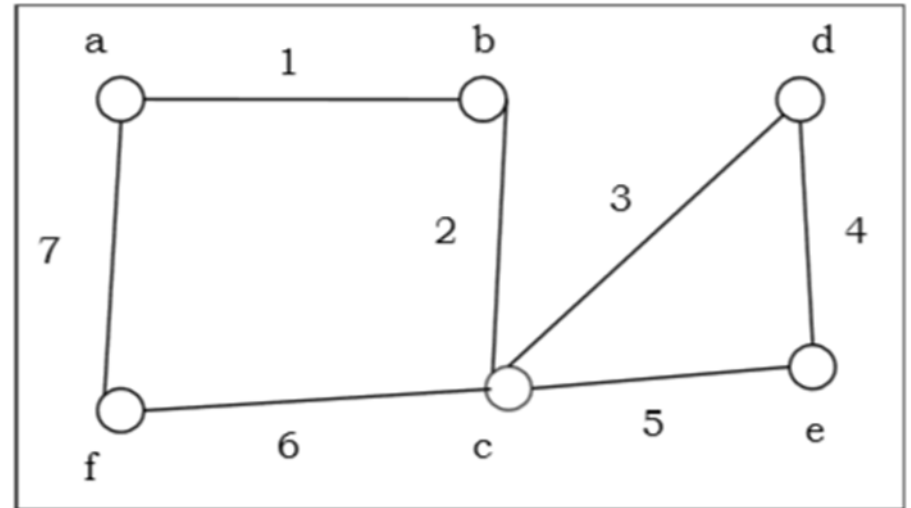
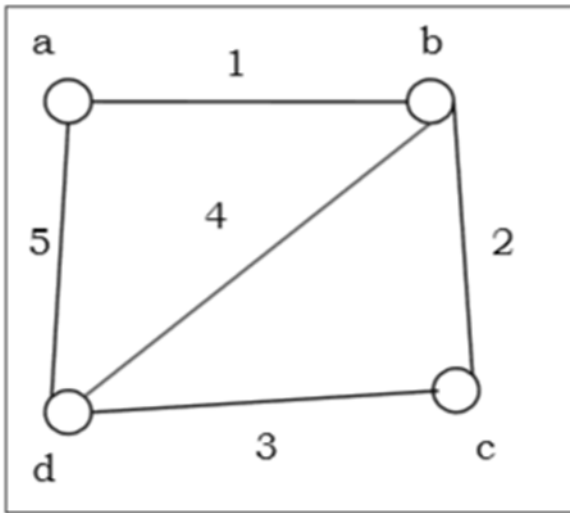
# Euler Graph

- A connected graph  $G$  is called an Euler graph, if there is a closed trail (or circuit) which includes every edge of the graph  $G$ .



# Euler Graph

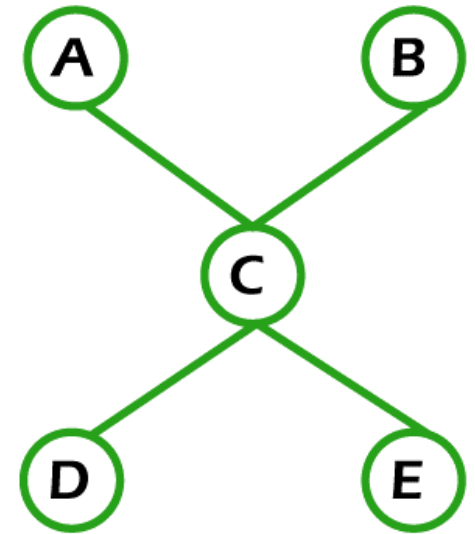
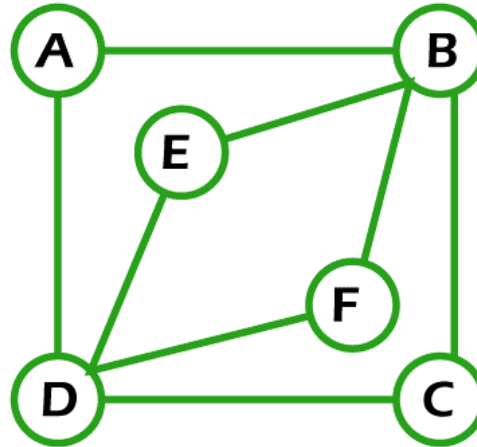
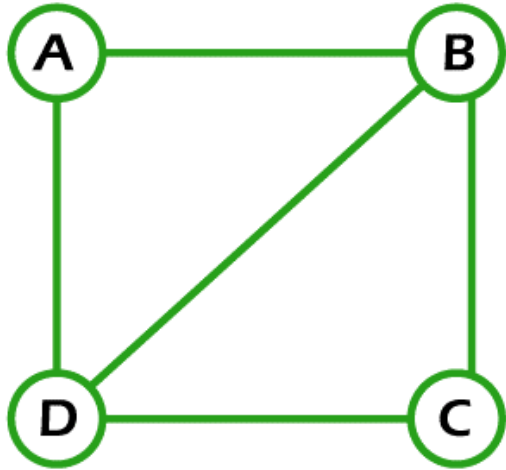
- Any connected graph is called as an Euler Graph if and only if all its vertices are of even degree.



# Euler Path

- We can also call the Euler path as Euler walk or Euler Trail.
- If there is a connected graph with a trail that has all the edges of the graph, then that type of trail will be known as the **Euler trail**.
- If there is a connected graph, which has a walk that passes through each and every edge of the graph only once, then that type of walk will be known as the **Euler walk**.
- An **Euler path** is a path that uses every edge of a graph exactly once.

# Euler Path

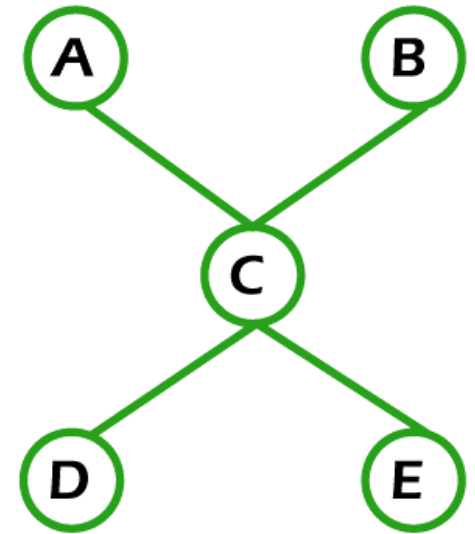
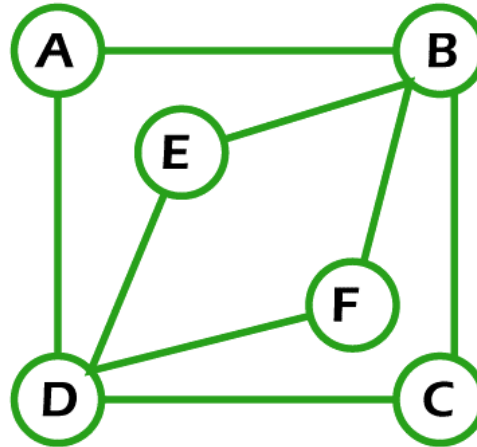
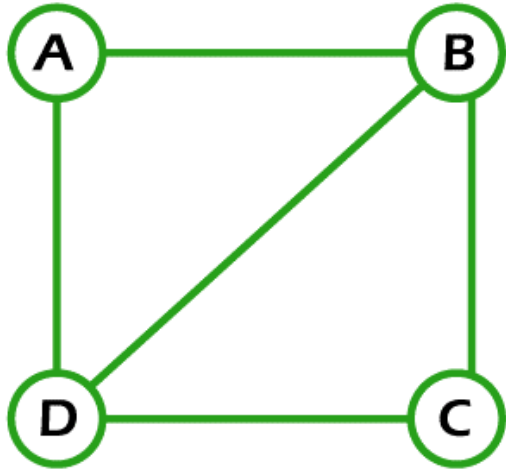


A connected graph has a Euler path if and only if the number of its vertices of odd degree is  $\leq 2$ .

# Euler Circuit

- Euler circuit is also known as Euler Cycle or Euler Tour.
- If there exists a Circuit in the connected graph that contains all the edges of the graph, then that circuit is called as an Euler circuit.
- If there exists a walk in the connected graph that starts and ends at the same vertex and visits every edge of the graph exactly once with or without repeating the vertices, then such a walk is called as an Euler circuit.
- An Euler trail that starts and ends at the same vertex is called as an Euler circuit.

# Euler Circuit

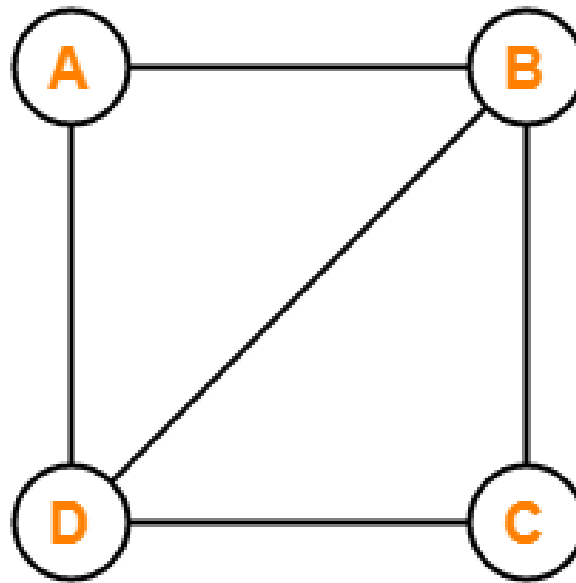


A graph will contain an Euler circuit if and only if all its vertices are of even degree.



# Semi-Euler Graph

- If a connected graph contains an Euler trail but does not contain an Euler circuit, then such a graph is called as a semi-Euler graph.



**Semi-Euler Graph**

# Euler Graph: Notes

- If the graph is connected and contains an Euler circuit, then it is an Euler graph.
- If all its vertices are of even degree, then graph contains an Euler circuit otherwise not.
- If the graph is connected and contains an Euler trail (not Euler Circuit), then graph is a semi-Euler graph otherwise not.
- If the number of vertices with odd degree are at most 2, then graph contains an Euler trail otherwise not.

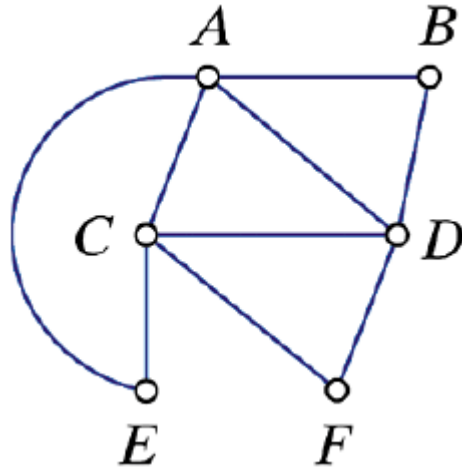
# How to find Euler Path or Circuit?

- Fleury's algorithm finds an Euler circuit or an Euler path in a connected graph.
- Fleury's algorithm is an elegant but inefficient algorithm.
- It proceeds by repeatedly removing edges from the graph in such way, that the graph remains Eulerian.

# Fleury's Algorithm

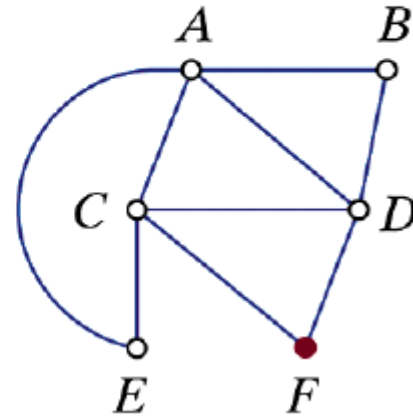
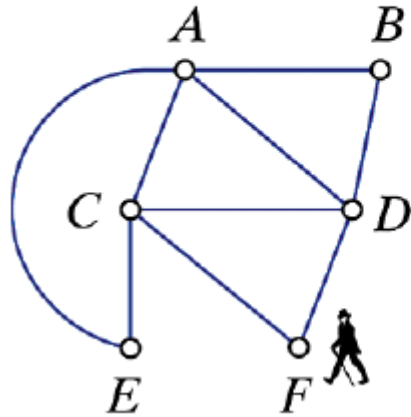
- **Preliminaries.** Make sure that the graph is connected and either (1) has no odd vertices (circuit) or (2) has just two odd vertices (path).
- **Start.** Choose a starting vertex. [In case (1) this can be any vertex; in case (2) it must be one of the two *odd* vertices.]
- **Intermediate steps.** At each step, if you have a choice, *don't choose a bridge of the yet-to-be-traveled part* of the graph. However, if you have only one choice, take it.
- **End.** When you can't travel any more, the circuit (path) is complete. [In case (1) you will be back at the starting vertex; in case (2) you will end at the other odd vertex.]

# Fleury's Algorithm



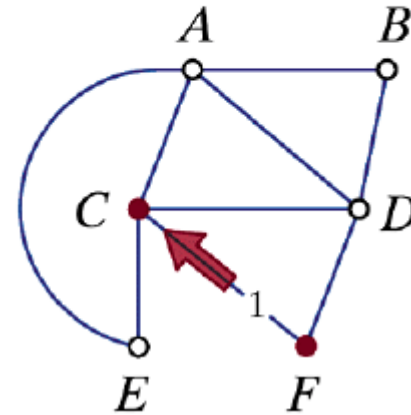
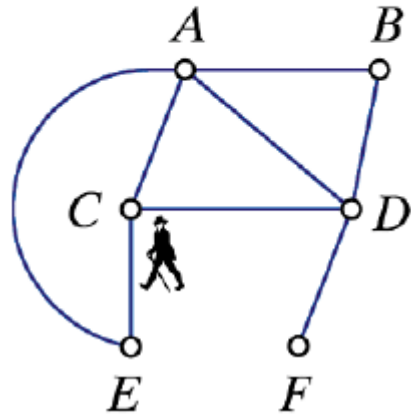
# Fleury's Algorithm

**Start:** We can pick any starting point we want. Let's say we start at  $F$ .



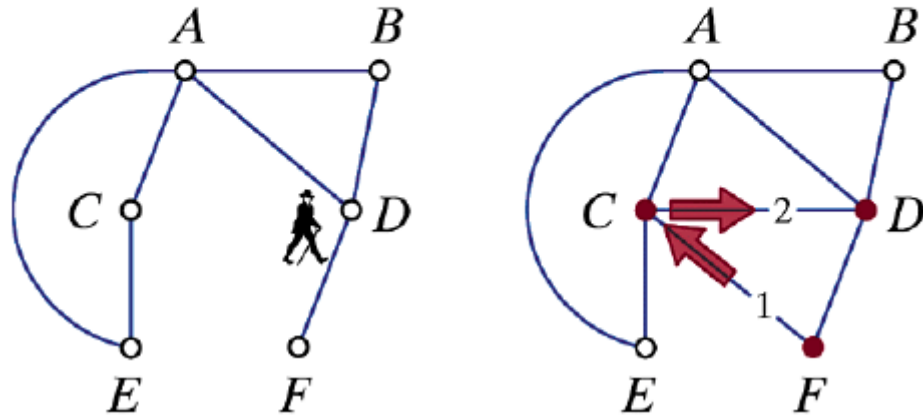
# Fleury's Algorithm

**Step 1:** Travel from  $F$  to  $C$ .  
(Could have also gone from  $F$  to  $D$ .)



# Fleury's Algorithm

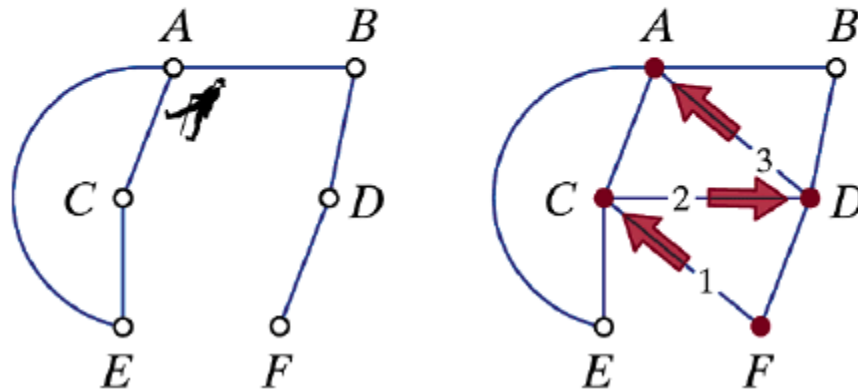
**Step 2:** Travel from  $C$  to  $D$ .  
(Could have also gone to  $A$  or to  $E$ .)





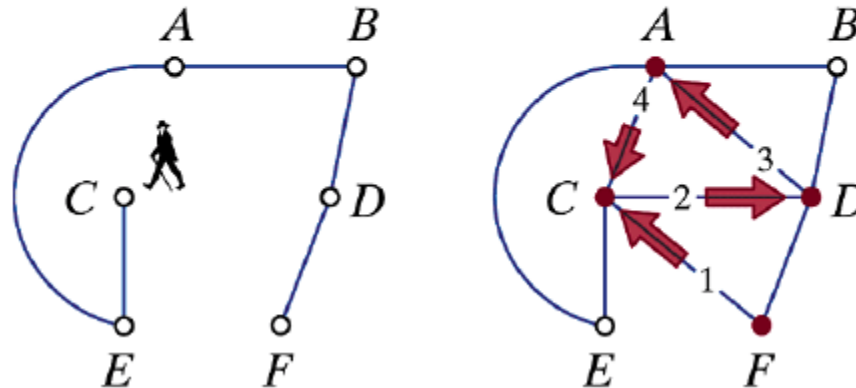
# Fleury's Algorithm

**Step 3:** Travel from  $D$  to  $A$ .  
(Could have also gone to  $B$  but not to  $F$ — $DF$  is a bridge!)



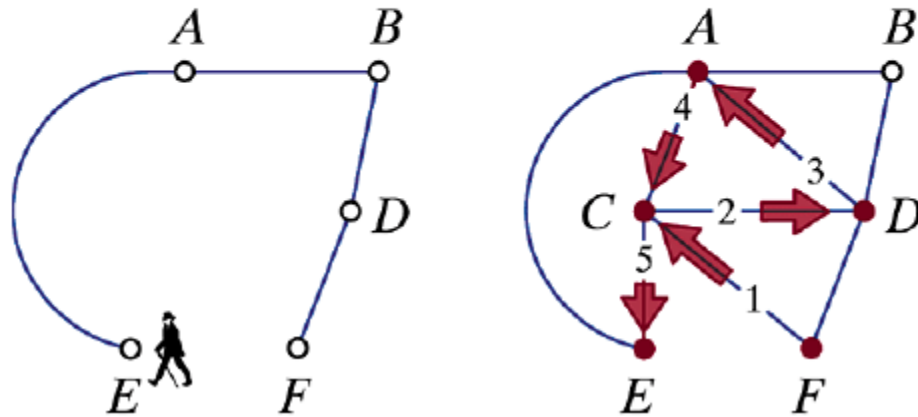
# Fleury's Algorithm

**Step 4:** Travel from  $A$  to  $C$ .  
(Could have also gone to  $E$  but not to  $B$ — $AB$  is a bridge!)



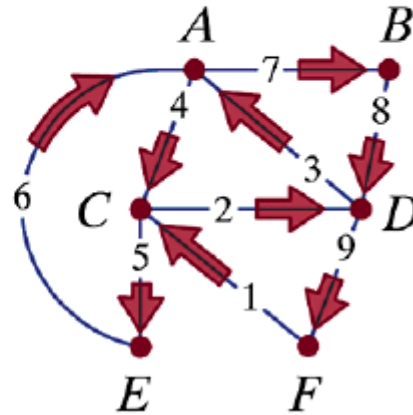
# Fleury's Algorithm

**Step 5:** Travel from  $C$  to  $E$ .  
(There is no choice!)



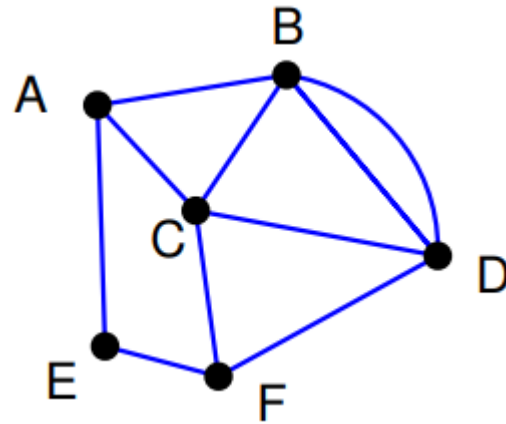
# Fleury's Algorithm

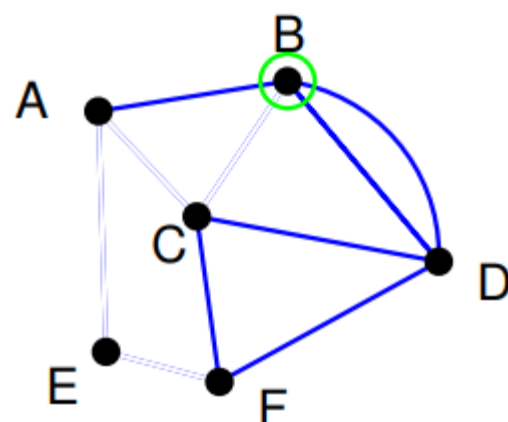
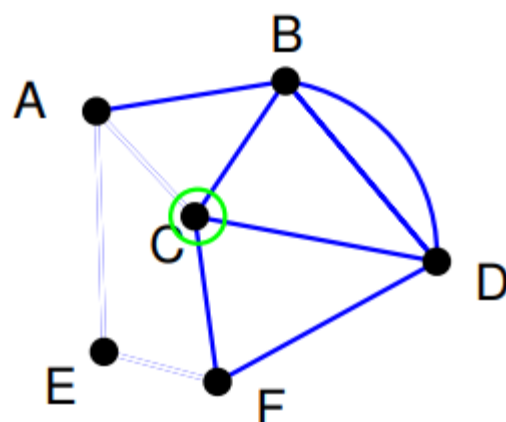
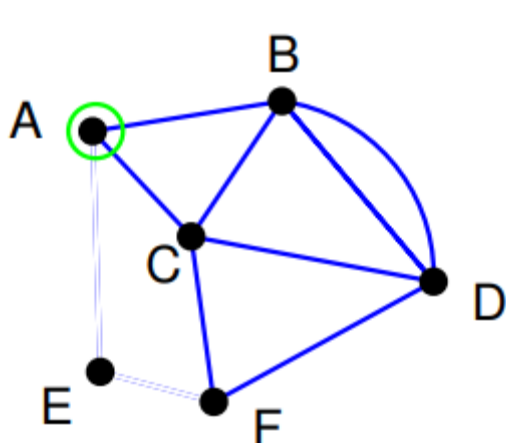
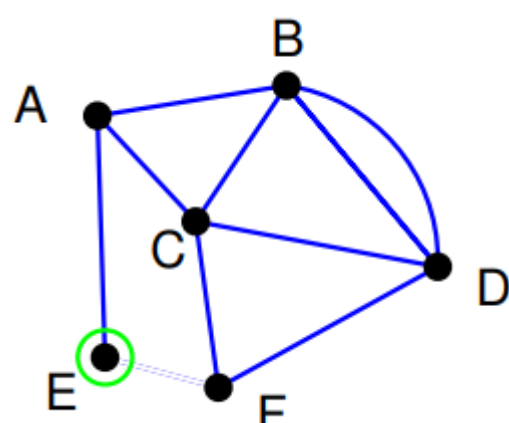
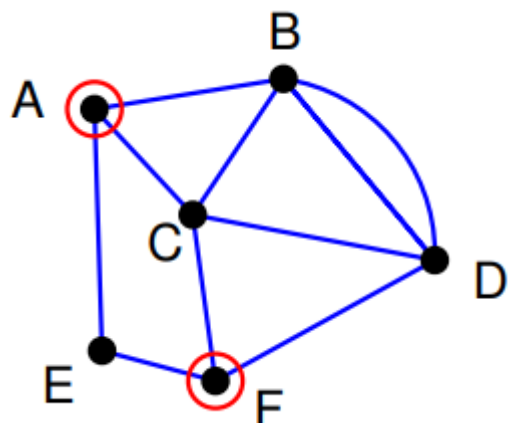
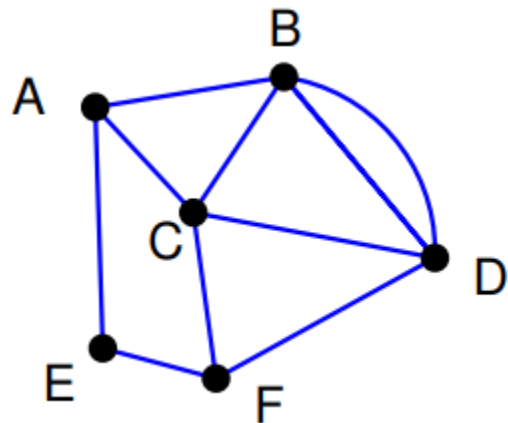
**Steps 6, 7, 8, and 9:** Only one way to go at each step.



# Fleury's Algorithm

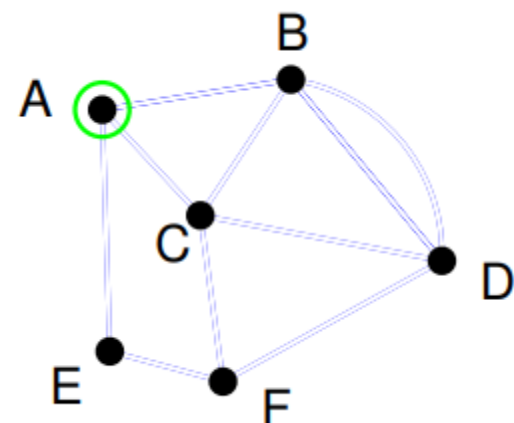
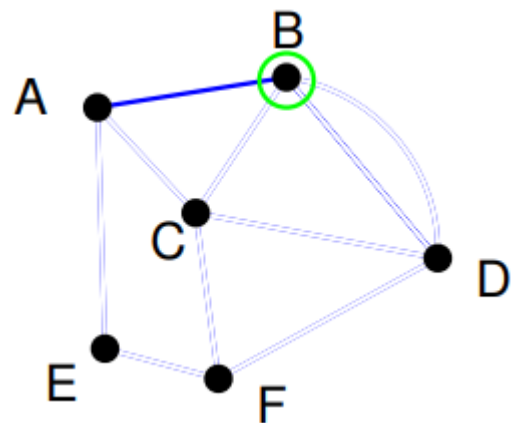
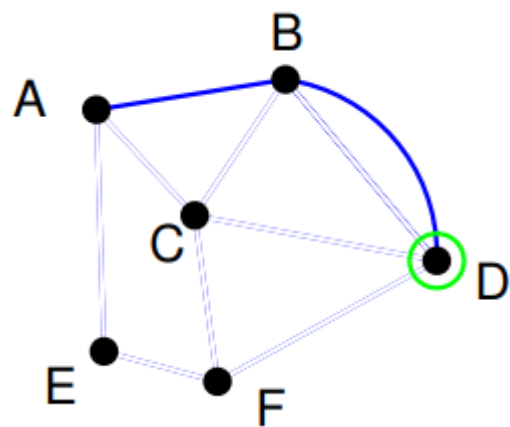
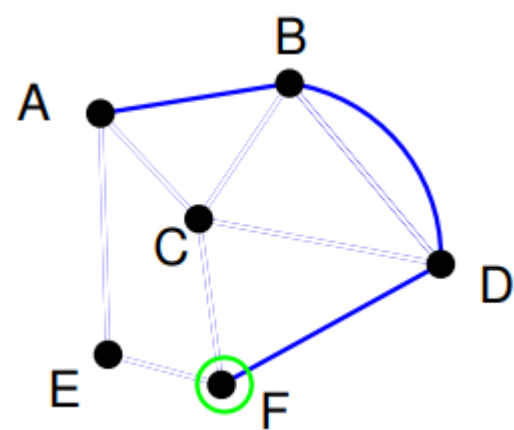
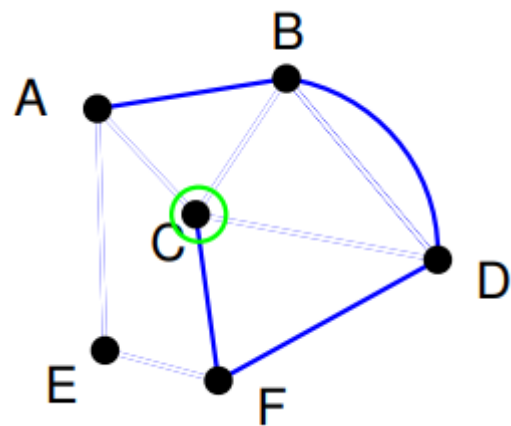
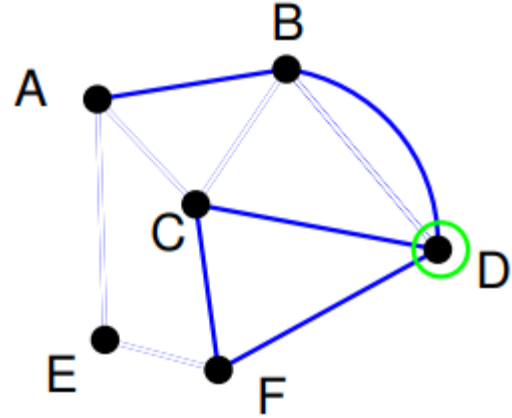
Find an Euler circuit in the graph.





Up until this point, the choices didn't matter.

But now, crossing the edge BA would be a mistake, because we would be stuck there.



**Euler Path: FEACBDFCAB**

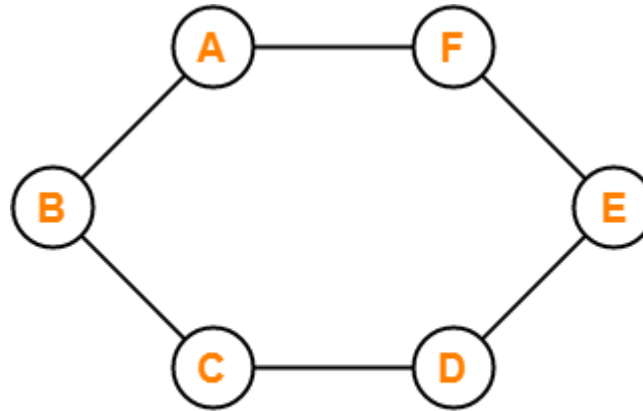
# Applications of Euler Graphs

- Euler paths and circuits can be used to solve many practical problems such as finding a path or circuit that traverses each
  - street in a neighborhood,
  - road in a transportation network,
  - connection in a utility grid,
  - link in a communications network.
- Other applications are found in the
  - layout of circuits,
  - network multicasting,
  - molecular biology, where Euler paths are used in the sequencing of DNA.



# Hamiltonian Graph

- A connected graph  $G$  is called Hamiltonian graph if there is a cycle which includes every vertex of  $G$  and the cycle is called **Hamiltonian cycle**.

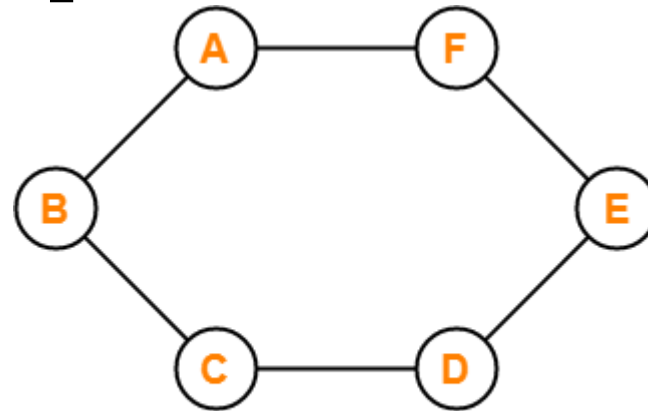


Example of Hamiltonian Graph

# Hamiltonian Graph

- **Dirac's Theorem:** A simple graph with  $n$  vertices ( $n \geq 3$ ) is Hamiltonian if every vertex has degree  $n/2$  or greater.
- **Ore's Theorem:** A simple graph with  $n$  vertices ( $n \geq 3$ ) is Hamiltonian if, for every pair of non-adjacent vertices, the sum of their degrees is  $n$  or greater.

# Hamiltonian Graph



Example of Hamiltonian Graph

$$\forall (\text{non-adjacent vertices pair } v, u) (\deg(v) + \deg(w) \geq n) \\ \Rightarrow \text{Graph is Hamiltonian}$$

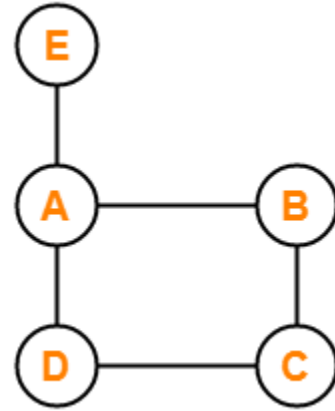
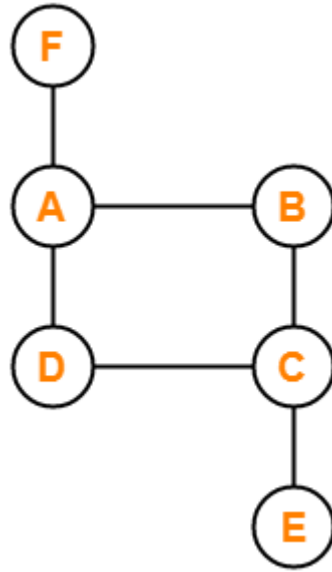
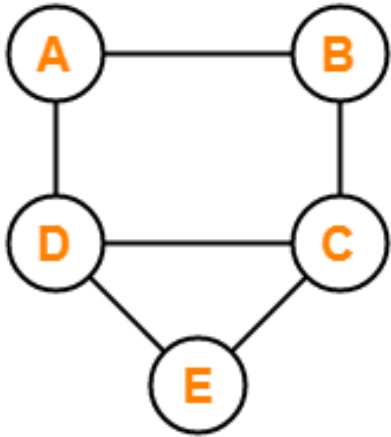
But this does not imply the reverse of it, that means,

$$\text{Graph is Hamiltonian} \nRightarrow \\ \forall (\text{non-adjacent vertices pair } v, u) (\deg(v) + \deg(w) \geq n)$$

# Hamiltonian Path

- If there exists a walk in the connected graph that visits every vertex of the graph exactly once without repeating the edges, then such a walk is called as a Hamiltonian path.
- If there exists a Path in the connected graph that contains all the vertices of the graph, then such a path is called as a Hamiltonian path.

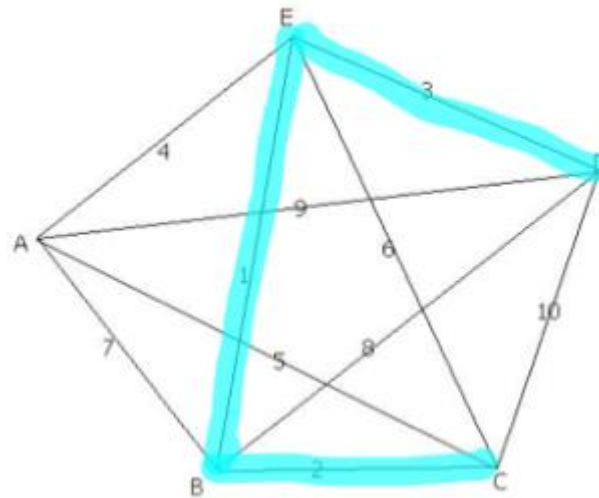
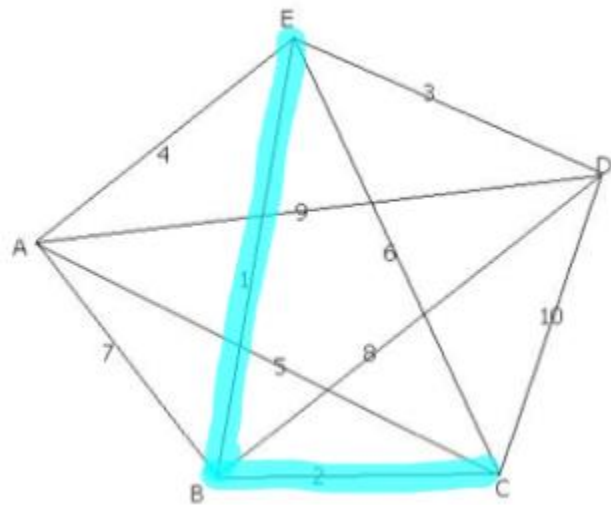
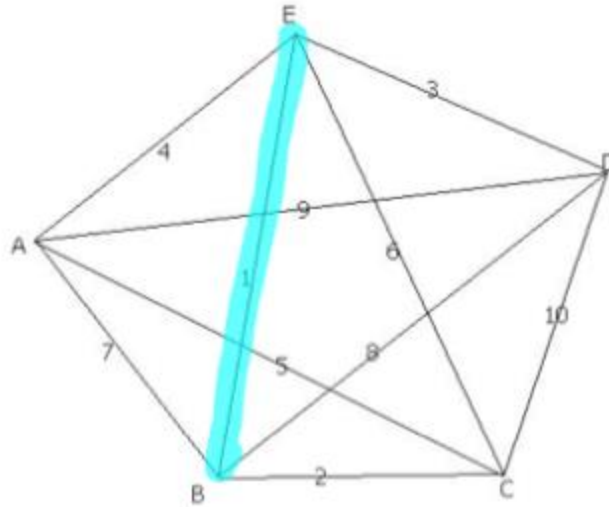
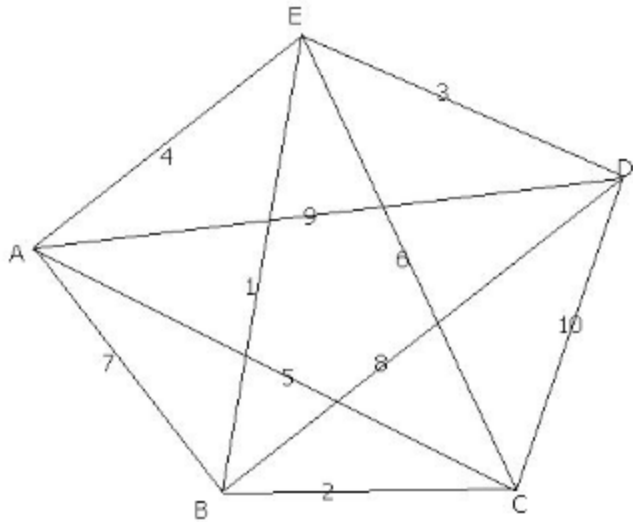
# Hamiltonian Path



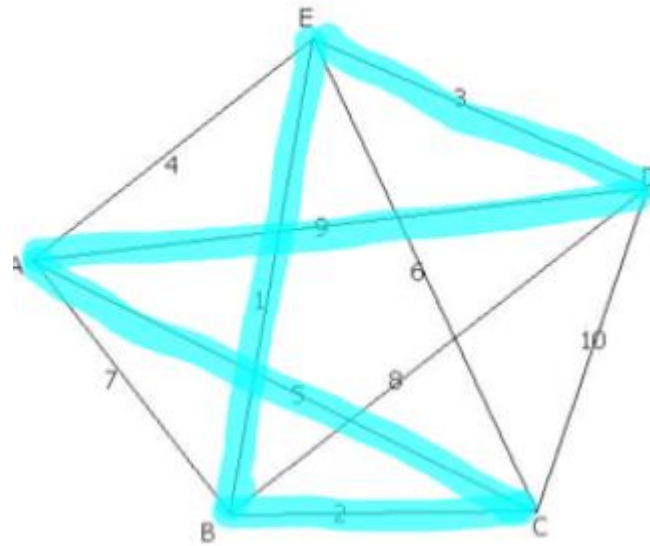
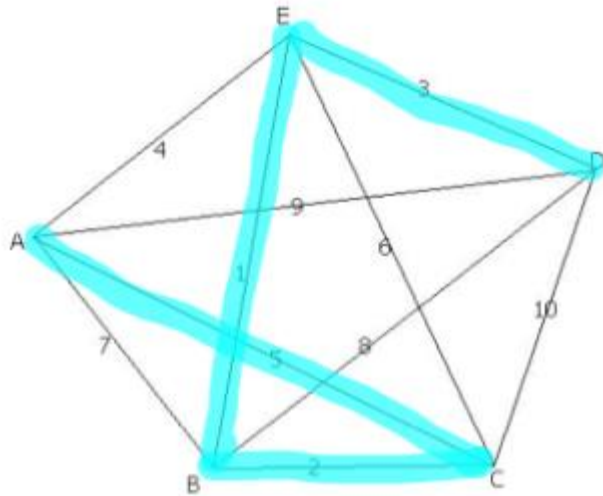
# How to find Hamiltonian Cycle / Circuit?

- Cheapest Link Algorithm
- Pick the link with the smallest weight first (if there is a tie, randomly pick one). Mark the corresponding edge in color.
- Pick the next cheapest link and mark the corresponding edge in color.
- Continue picking the cheapest link available. Mark the corresponding edge in color except when **a)** it closes a circuit or **b)** it results in three edges coming out of a single vertex.
- When there are no more vertices to link, close the red circuit.

# How to find Hamiltonian Cycle / Circuit?



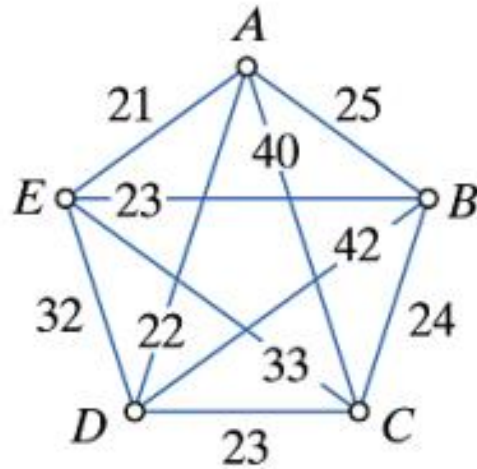
# How to find Hamiltonian Cycle / Circuit?



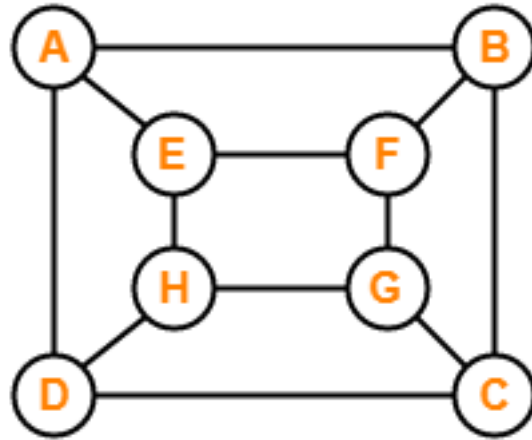
The solution is ACBEDA or ADEBCA with total weight of 20 miles.



# How to find Hamiltonian Cycle / Circuit?



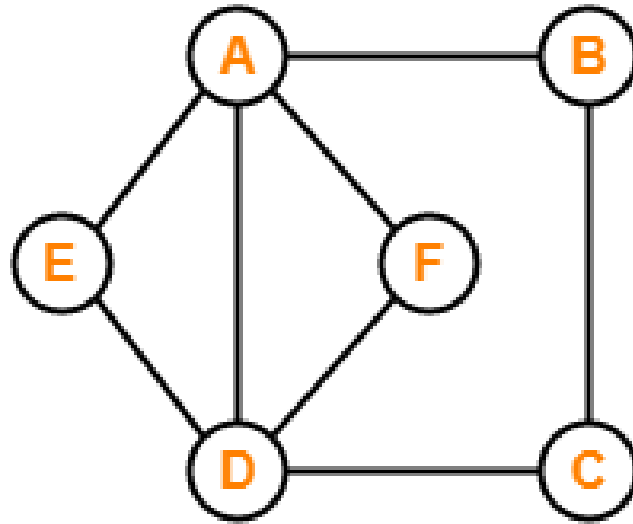
# Euler or Hamiltonian?



The graph contains both a Hamiltonian path (ABCDHGFE) and a Hamiltonian circuit (ABCDHGFEA).

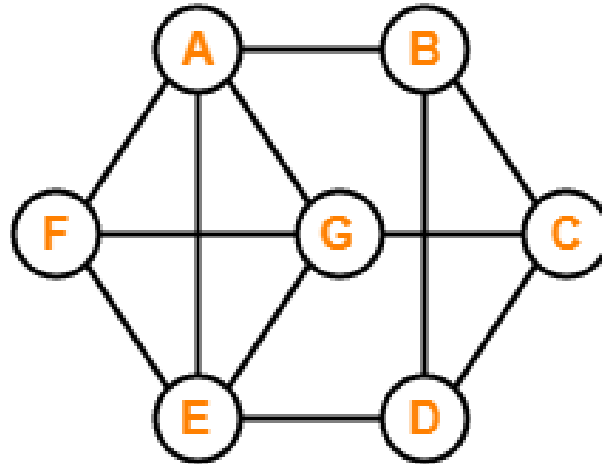
It is not an Euler graph.

# Euler or Hamiltonian?



- No Hamiltonian path nor it contains a Hamiltonian circuit.
- It is an Euler graph.

# Euler or Hamiltonian?



The graph contains both a Hamiltonian path (ABCDEFG) and a Hamiltonian circuit (ABCDEFGA).

It is not an Euler graph.

# Euler vs. Hamiltonian?

Property	Euler	Hamilton
Repeated visits to a given node allowed?	Yes	No
Repeated traversals of a given edge allowed?	No	No
Omitted nodes allowed?	No	No
Omitted edges allowed?	No	Yes

# Summary

- Euler Graphs
- Hamiltonian Graphs