

SSN LAB ASSIGNMENT: SYMMETRICAL ENCRYPTION

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1.DES

```
1.  setKey(0101010101010101)
    encryptDES(416c690000000000)
    IP:  L0=07000205, R0=00060600
    Rnd1  f(R0=00060600, SK1=00 00 00 00 00 00 00 00 ) = c8b8cbfc
    Rnd2  f(R1=cfb8c9f9, SK2=00 00 00 00 00 00 00 00 ) = ac9062f5
    Rnd3  f(R2=ac9664f5, SK3=00 00 00 00 00 00 00 00 ) = ef63339a
    Rnd4  f(R3=20dbfa63, SK4=00 00 00 00 00 00 00 00 ) = e4c0c3ed
    Rnd5  f(R4=4856a718, SK5=00 00 00 00 00 00 00 00 ) = b688d6c5
    Rnd6  f(R5=96532ca6, SK6=00 00 00 00 00 00 00 00 ) = 078e4a35
    Rnd7  f(R6=4fd8ed2d, SK7=00 00 00 00 00 00 00 00 ) = 8731e059
    Rnd8  f(R7=1162ccff, SK8=00 00 00 00 00 00 00 00 ) = ec5804d6
    Rnd9  f(R8=a380e9fb, SK9=00 00 00 00 00 00 00 00 ) = c5907157
    Rnd10 f(R9=d4f2bda8, SK10=00 00 00 00 00 00 00 00 ) = 498041a2
    Rnd11 f(R10=ea00a859, SK11=00 00 00 00 00 00 00 00 ) = eb0493e5
    Rnd12 f(R11=3ff62e4d, SK12=00 00 00 00 00 00 00 00 ) = 8459728c
    Rnd13 f(R12=6e59dad5, SK13=00 00 00 00 00 00 00 00 ) = 08057433
    Rnd14 f(R13=37f35a7e, SK14=00 00 00 00 00 00 00 00 ) = 45050caf
    Rnd15 f(R14=2b5cd67a, SK15=00 00 00 00 00 00 00 00 ) = 1c0ad75b
    Rnd16 f(R15=2bf98d25, SK16=00 00 00 00 00 00 00 00 ) = 4aaf3da9
    FP:  L=ff950aac, R=31f6753d
    returns ff950aac31f6753d
```

Addition

At start we have 64 bits of input (416c690000000000 hex = "Ali" ASCII) and 64 bit key (0101010101010101).

We convert our text to binary:

[illegible]

Then we convert our hex key to binary:

```
0b1000000001000000001000000001000000001000000001000000001 =
0x0101010101010101
```

Next step will be IP of our input text:

```
! [sbox] (/home/mrzizik/screenshots/sbox.png)
```

I wrote a small script for doing this in python:

```
1 ip = (58, 50, 42, 34, 26, 18, 10, 2, 60, 52, 44, 36, 28, 20, 12, 4, 62,
54, 46, 38, 30, 22, 14, 6, 64, 56, 48, 40, 32, 24, 16, 8, 57, 49, 41, 33,
25, 17, 9, 1, 59, 51, 43, 35, 27, 19, 11, 3, 61, 53, 45, 37, 29, 21, 13, 5,
63, 55, 47, 39, 31, 23, 15, 7)
2 input = [0, 1, 0, 0, 0, 0, 0, 0, 1, 0, 1, 1, 0, 1, 1, 0, 0, 0, 1, 1, 0, 1,
0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0,
0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0]
3 output = []
4 for i in xrange(64):
5     output.append(input[ip[i]-1])
6 print output
```

After that we have got such permuted list

```
000001110000000000000001000000101000000000000001100000011000000000
```

After that we do IP and divide key to two parts 28 bits each:

```
! [keypermut](/home/mrzizik/screenshots/keypermut.png)
```

[illegible][illegible]

Then we shift our key, but in our situation (all are zeros) shifting wouldn't change anything.

For some rounds we make shift 2, not one.

Then we make second permutation of key, that's also useless for us.

We have round 1 key. In that manner we generate keys for all rounds.

Then we divide our input to 2 parts and go to first round of encryption

```
![cipher](/home/mrzizik/screenshots/cipher.png)
```

```
Rnd8 f(R7=1162ccff, SK8=00 00 00 00 00 00 00 00 ) = ec5804d6
```

2. Firstly we make internal permutation of key, while we permute it by special box and make it 56 bit instead of 64.

57	49	41	33	25	17	09	01
58	50	42	34	26	18	10	02
59	51	43	35	27	19	11	03
60	52	44	36	63	55	47	39
31	23	15	07	62	54	46	37
30	22	14	06	61	53	45	37
29	21	13	05	28	20	12	04

then we spit it in two parts each 28 bit.

Then we shift both of them left (1,2,9,16 rounds - 1 bit, else - 2 bits) and send to compressing P box

14	17	11	24	01	05	03	28
15	06	21	10	23	19	12	04
26	08	16	07	27	20	13	02
41	52	31	37	47	55	30	40
51	45	33	48	44	49	39	56
34	53	46	42	50	36	29	32

that makes it 48 bit from 56 for each round of encryption.

Wrote with explanations in 1 answer

3. It's because the first step of key generation, when we make 56bit key from 64bit one.

In our example we have in binary:

```
0000 0001 0000 0001
0000 0001 0000 0001
0000 0001 0000 0001
0000 0001 0000 0001
```

And the first step cuts out 8, 16, 24, 32, 40, 48, 56, 64, and left us with all zero's in key.

2.AES

1. Diffusion is reached by 16 rounds of SubBytes, ShiftRows and Mix Columns with applying round key. Also the last permutation grants more diffusion
2. Confusion elements are rcon and sbox of key scheduling, as all key scheduling process. It makes one bit modification in input key become huge change in key, that we are going to use and in output ciphertext.

Many rounds of crypting when output of round becomes input of next, and diffusion elements also grants more confusion.

3.Bonus: RC4

After bruteforce attack found that keys are "adwtg" and "495706". There is book of Darwin encrypted in them.

```
1 The Project Gutenberg EBook of On the Origin of Species, by Charles Darwin
2
3 This eBook is for the use of anyone anywhere at no cost and with
4 almost no restrictions whatsoever. You may copy it, give it away or
5 re-use it under the terms of the Project Gutenberg License included
6 with this eBook or online at www.gutenberg.org
7
8
9 Title: On the Origin of Species
10      1st Edition
11
12 Author: Charles Darwin
13
14 Release Date: Release Date: March, 1998 [EBook #1228]
15 Posting Date: November 23, 2009
16
17
```

1. (a) How did you identify the encrypted files ?

Firstly i just bruteforced all of them by lowercase alphabet, and after crackin one, tried other with digits, but then i opened them in python and saw that 2 encrypted files have different pattern of content

```

E" IovC b us et n1
</S
% * W1 \ , 7 6 W l = Fc2 cg j * b / c ) ( E )
> 3zI4          v < Bh vA
, M

E ) } n ] L A s
: : e ; d / w ! M P w M ? P F 4 " 2 > J v 1 ` n d X R 6 y Q f 覓
8 | r v R Q ) o T # x t ^ D
5 k { *
1 v 4 h 3 U k * ^ : ! 6 > X < J = D N

K & g 0 Y e I c 3 ( X [ o : 2 M G $
/ J 0 I | 0 0 ` R 0 0 h ) 0 0 0 P ] Ü /
o Q 9 V d 6 } ^ 6 -
7 7 C T V V + ^ = ,
8 F M } " v ' T 9 ~ # Q , { 0
9
e M @ m v ! a I z L ,
K t A "

4 , + [ f ; " a x H A ? ~ 2 a
+ G u B t I , * , e K 8 s ] b " l n 8 [ D ; ] C &
T T ~ J t
T : 2 z t } / Z d v [ ' ( o | x m g l Y / c ! * $ G
X (

```

But those with trash inside are like this:

```

cPPdB^YR[AsEEVYZPKStqXW^[V|YAQQ~A^_\W_wdHPZ]UBZLwXPA[ ]FpQCD^V839:e[^K\v_^
XQFR_CCPPACTX^XZI^]RTWMGYVE]X@_\[ZJ@P]SBP@X<9VTXVGD]XG\GDCZTL\VZCD_YAJ[UGVE
m_DZYLW_AJQAWXER\MQFRNZK9:CVMF\YEBVQ\FE[RA\F]BX^M\UcEW_\WDtBLPWUCTt\ZQ^BVQ
[ZXEUVS5?N]DYCP\JUs\XSVF^][Q[\QE@0BSEEVYZPKS^AP5?
4>=;g^LY\~]L]\CZPQ[ [V`G]VPQC<9DL|PYEZ XV839:pFCPZKr[VJY\GuRE0\W9:
<9e]Y\UCTsYA\cV[1]wV[ [yQCP_
i=;cXKAPZWwVLP~^ERUW\F
<9:2yXZWDRP[q^V_^K]4>=;>=cerelvre{~kif{vtl~adt}u}g~us|xsvze{rzkw}wsg`tp~}
f=5?4>=;>=hGVPERVSW@cDVyFJWXTA:2839:
<9:283{~g}vfyvzyzcavtqpj=;>=wg`xtgjppjqbgrcqzwwqycvabtwjtzqczyaqqbgemr~xuuxj
u}vt:2839:sJ[ ]XF\T@|TKCY_ux<9:2s\X^Dws`XTeWLXXtRWVVS YRR[u]^_RRY[uEPj[SXVCQ
PJ=;>=y@M\_Cx^~_DAYYY{VaRKPFXYVDqLFY_Tptcu]T^XU@nZ@UWTew@WPe[R5?n[B]W839:
<9{w}{~

```

2. (b) What is the effective key strength for each of the keys?

Effective key strength for key containing 6 digits is 10^6

Effective key strength for lower case character key with length 5 is 26^5

3. (c)

```

def worker(base):
    #read 64KB from the file
    data = open(FILE_NAME, 'rb').read(2**16)
    #generate all the strings of KEY_LENGTH length and check them
    #We know prior that the key starts with a. Remove the next two
lines for generic behavior
    if string.ascii_lowercase in ALPHABET:
        base = tuple(['a']) + base
        startms = time.time()*1000.0
        coll = len(ALPHABET)**(KEY_LENGTH-len(base))
        print coll
        for i in itertools.product(ALPHABET, repeat=KEY_LENGTH-
len(base)):
            check(''.join(base + i), data)
            endms = time.time()*1000.0
            print (1000/((endms-startms)/coll))

```

2771 attempts

4. (d)

I set up CPU_COUNT to 3 and changed serial method to

```
worker(parallel())
```

It gave me 2771*3 attempts in sec.

5. (e)

$(100^6)/2771 = 360880548$ seconds or something like 11 years in one thread