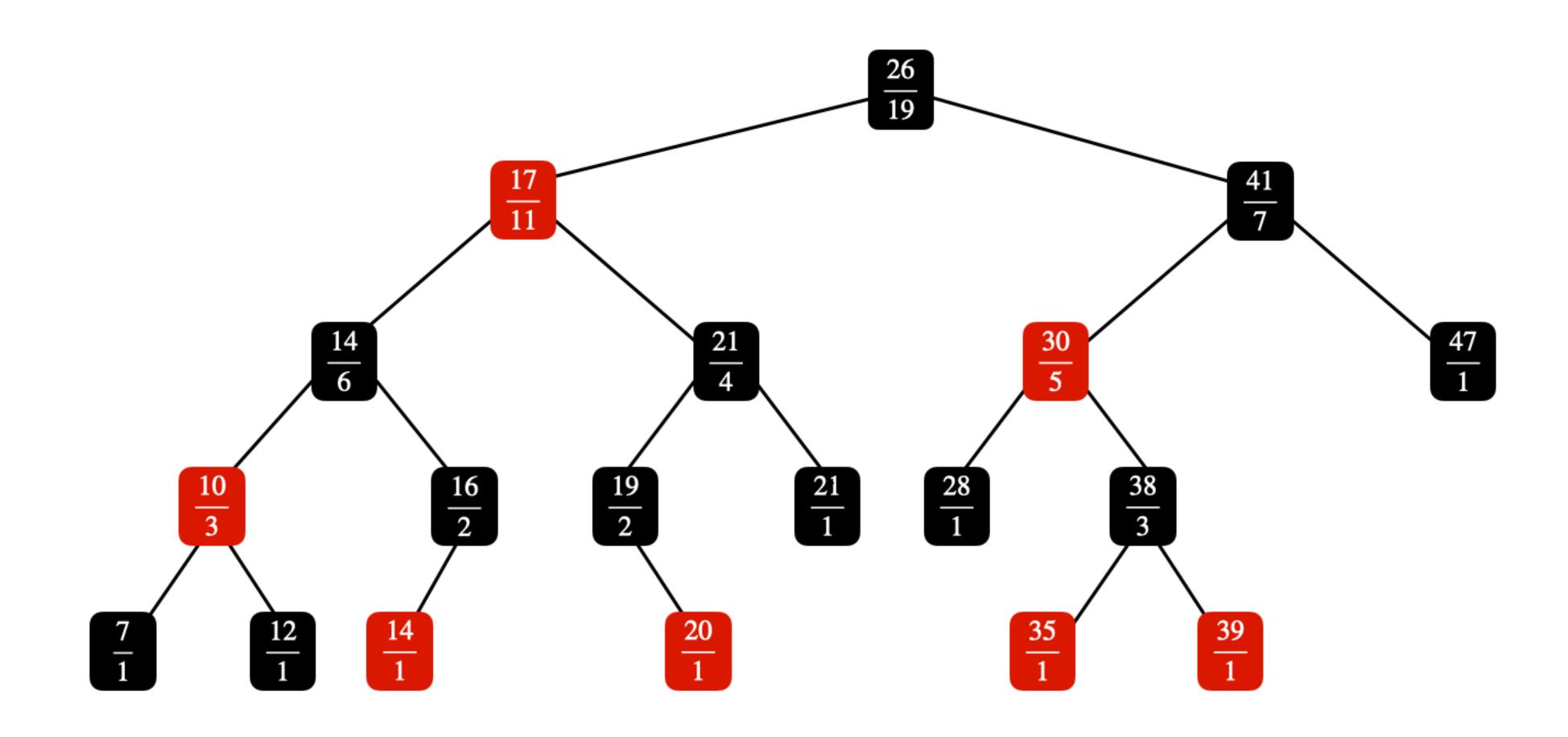
#### Lecture 11

Augmenting Data Structures

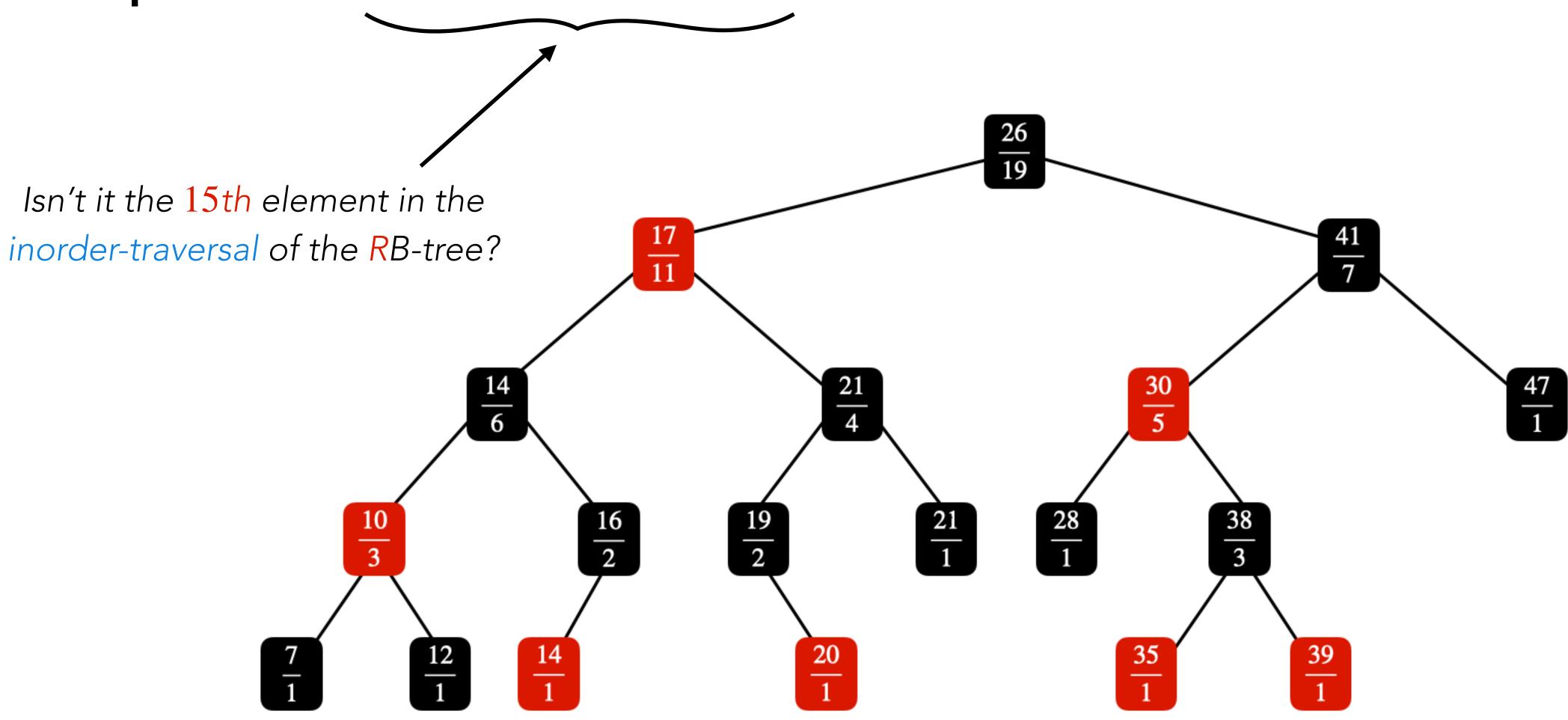
Recall that rank of an element is its position in the sorted order (w.r.t. keys) of the set.

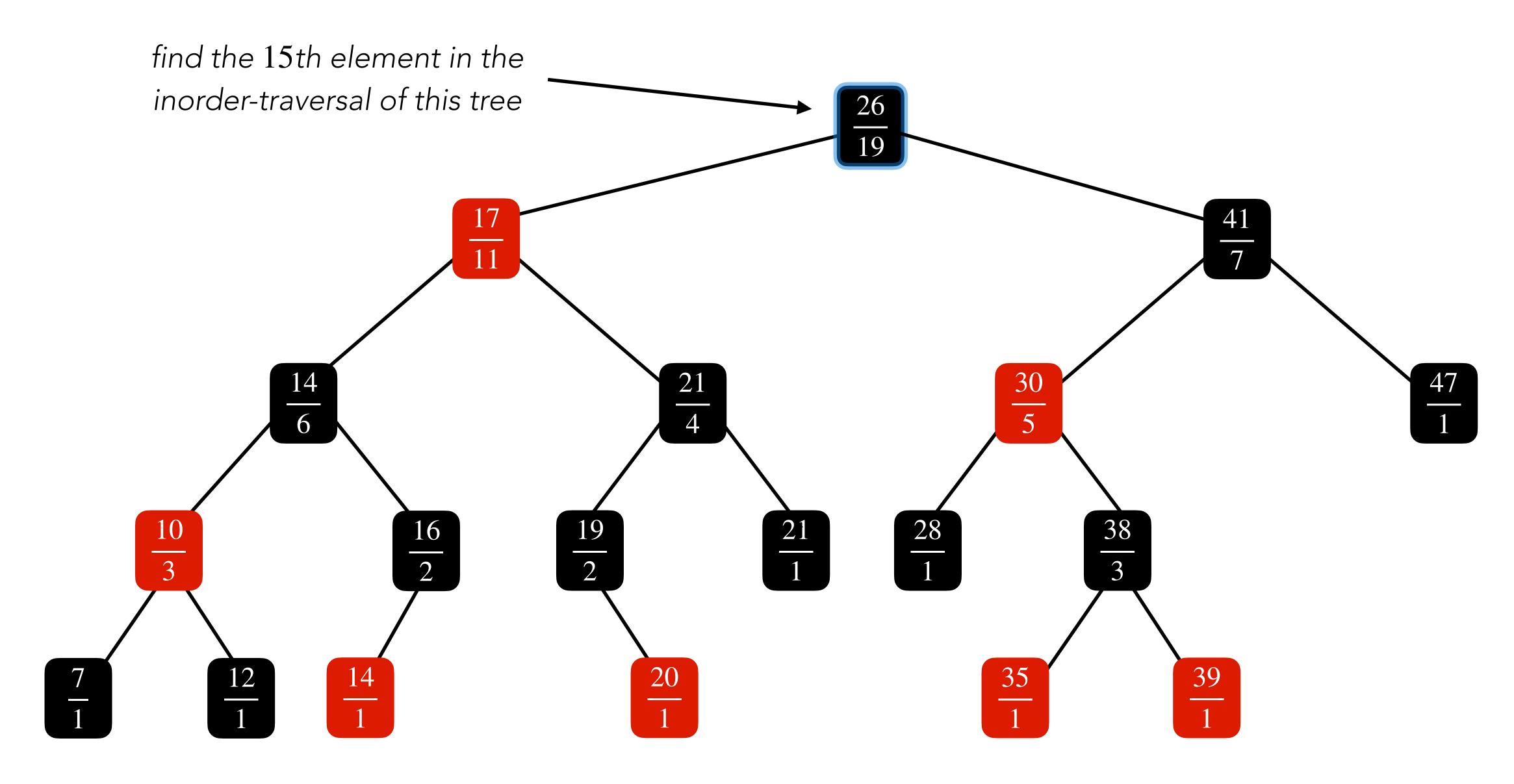
**Example:** Find an element with 15th rank in the below set or RB-tree.

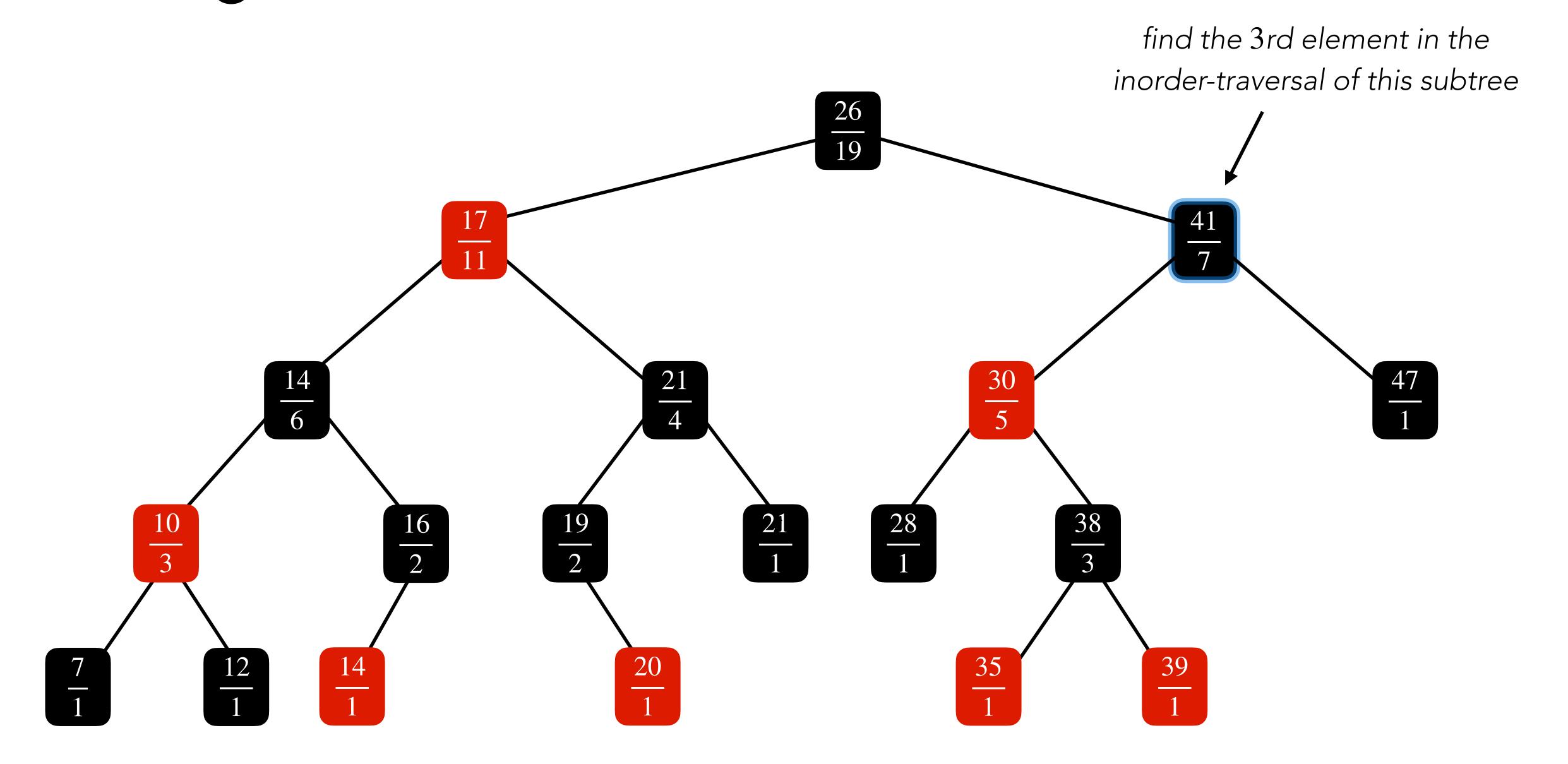
**Example:** Find an element with 15th rank in the below set or RB-tree.

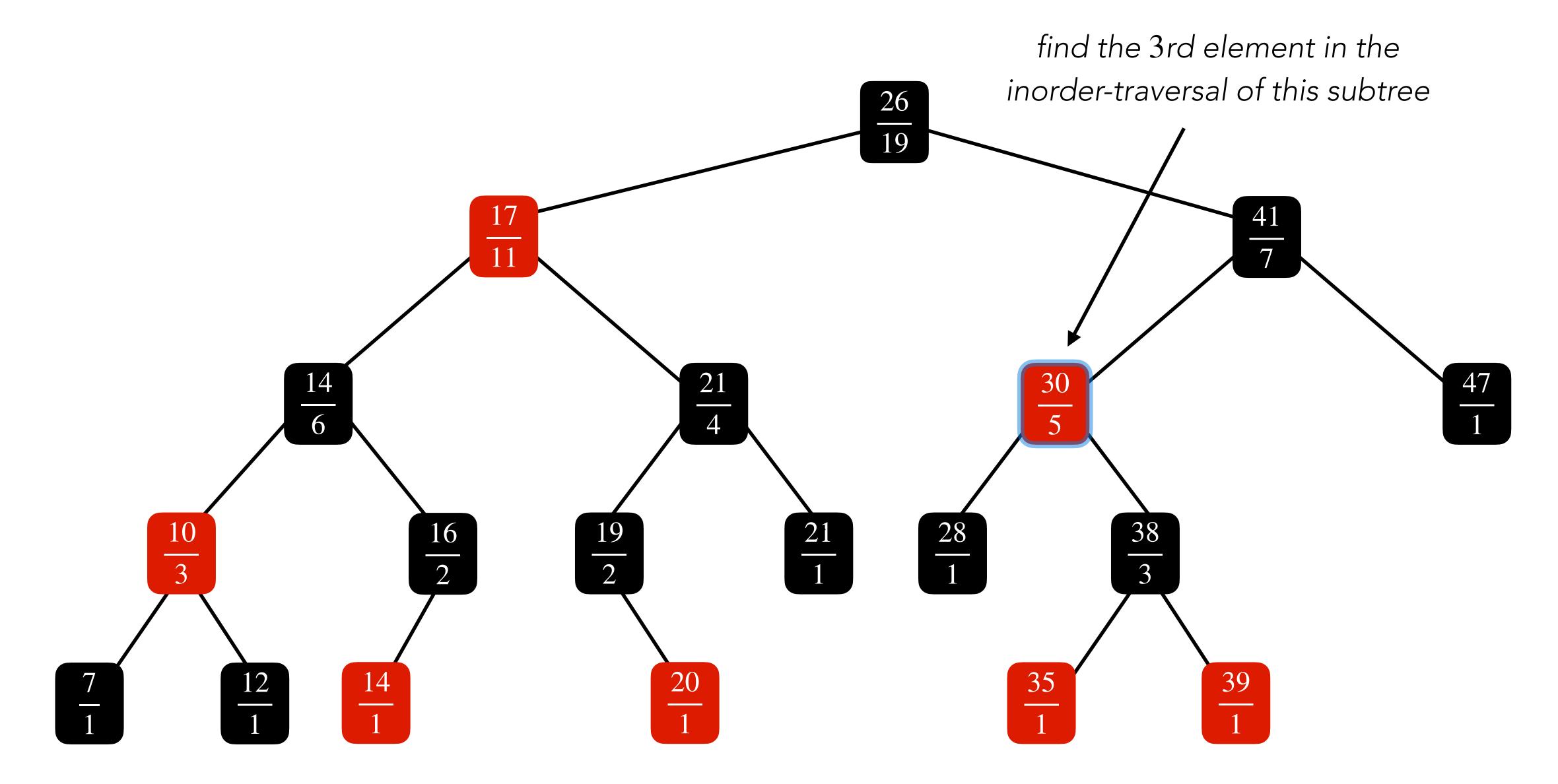


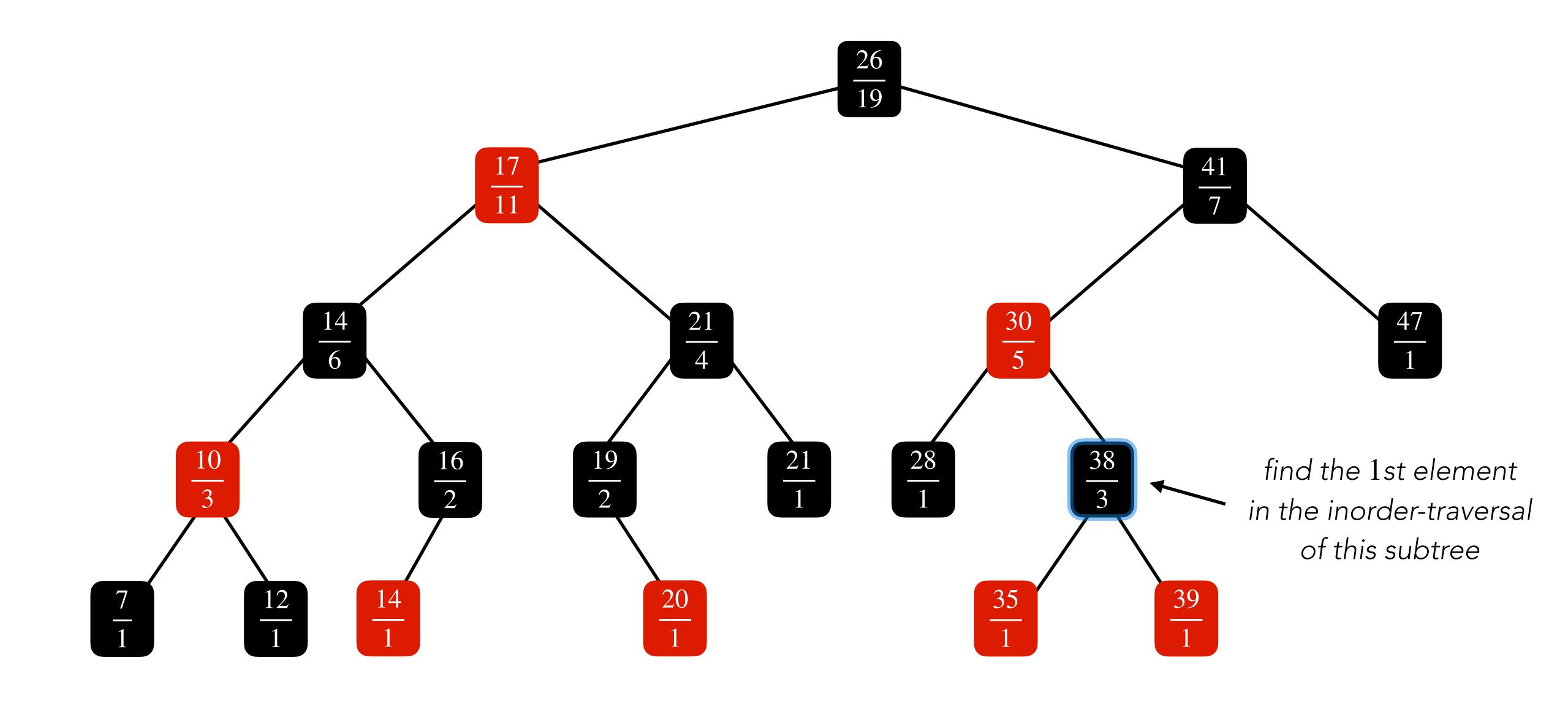
**Example:** Find an element with 15th rank in the below set or RB-tree.

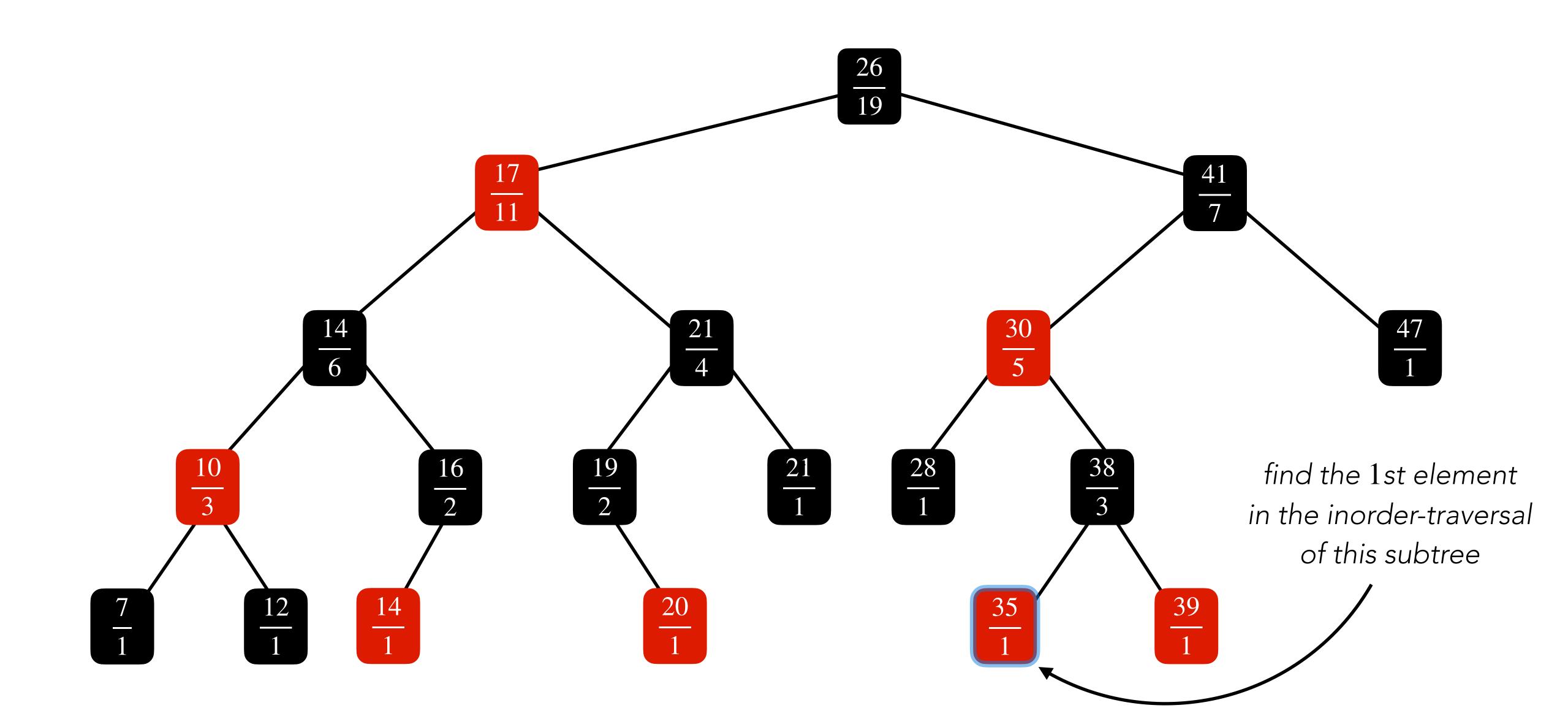


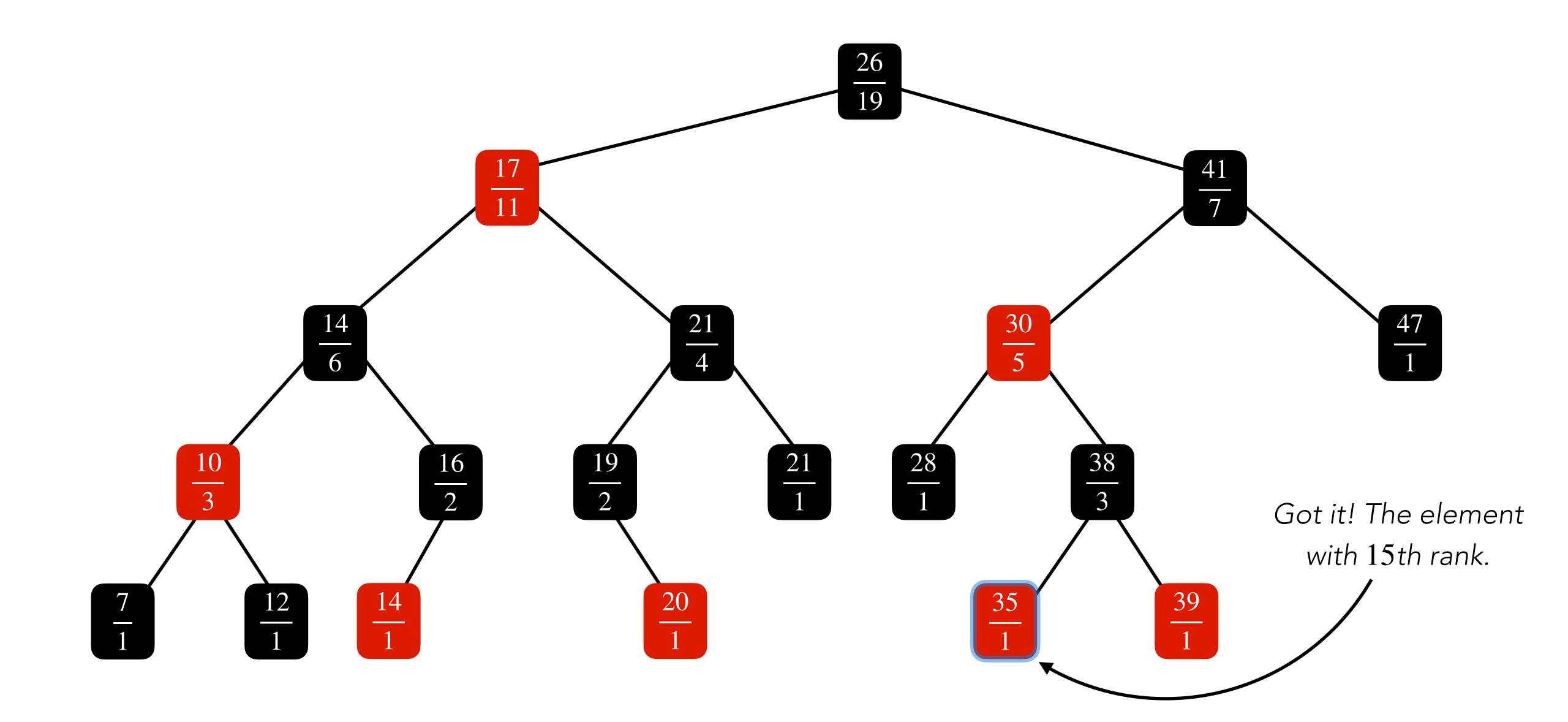








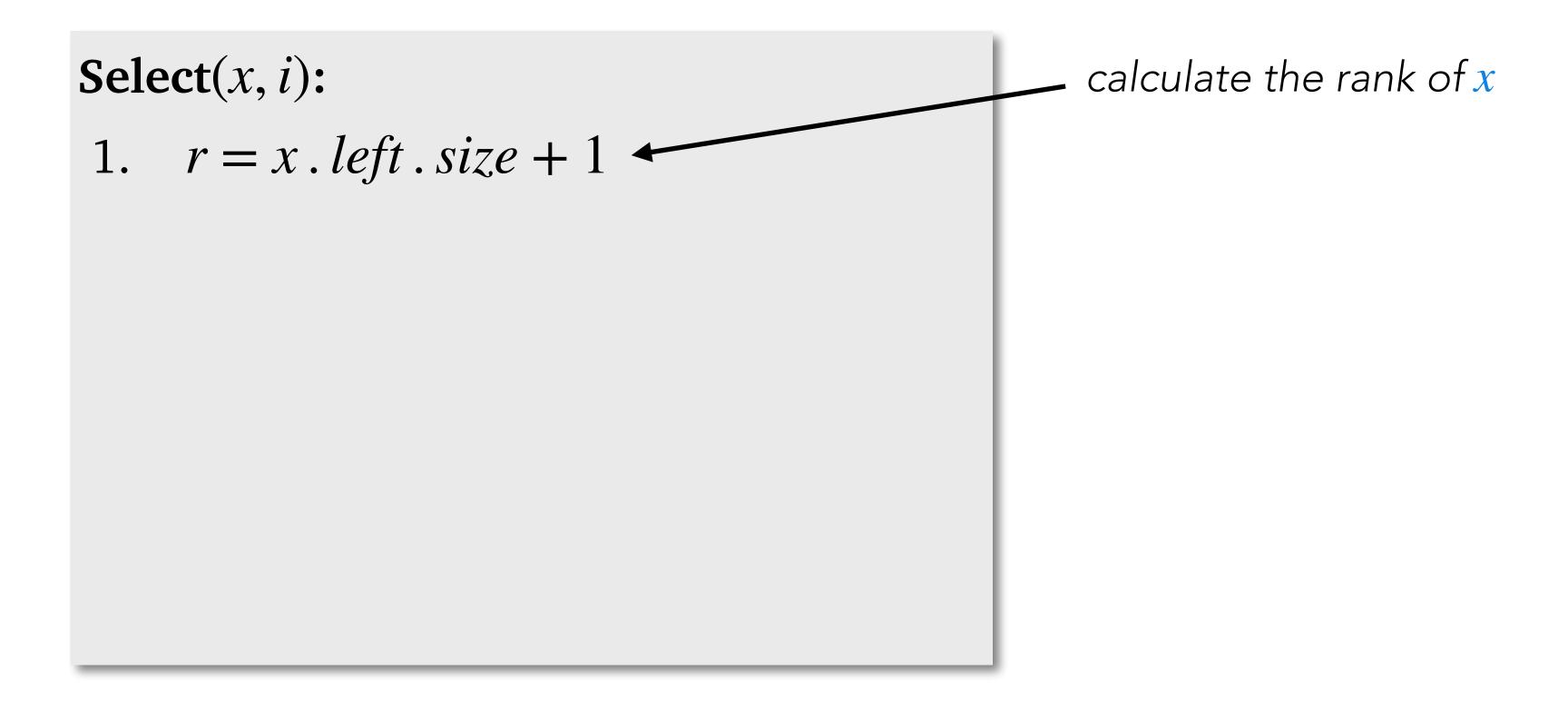


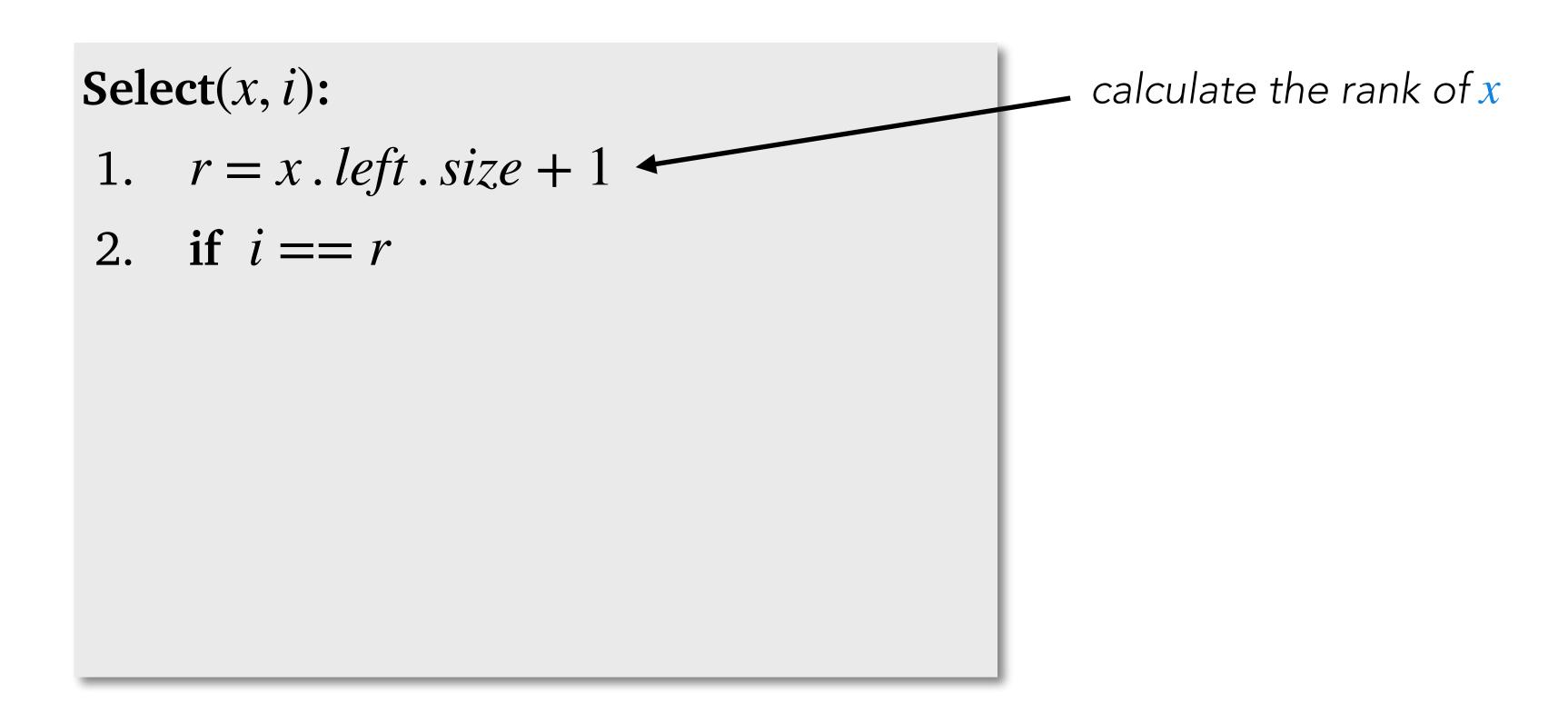


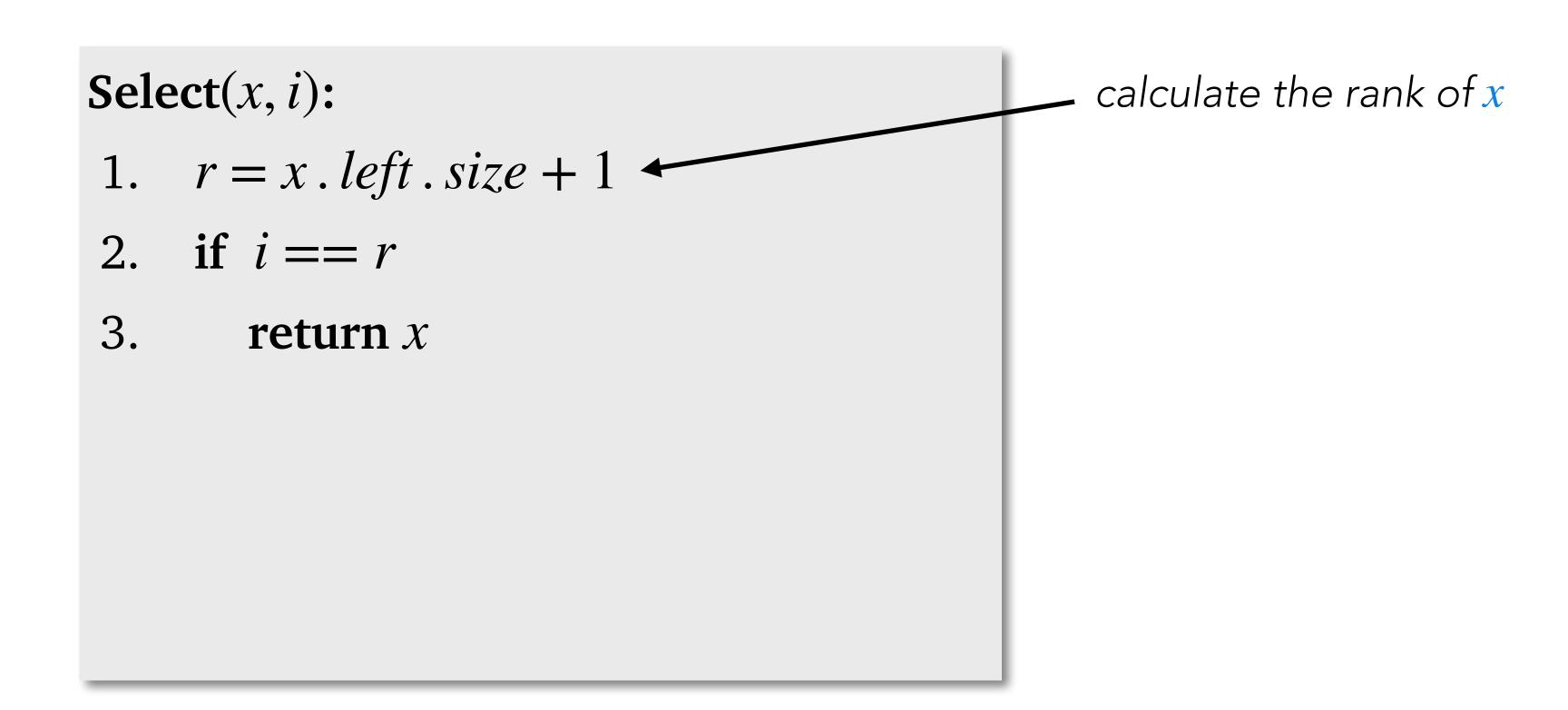
```
Select(x, i):
```

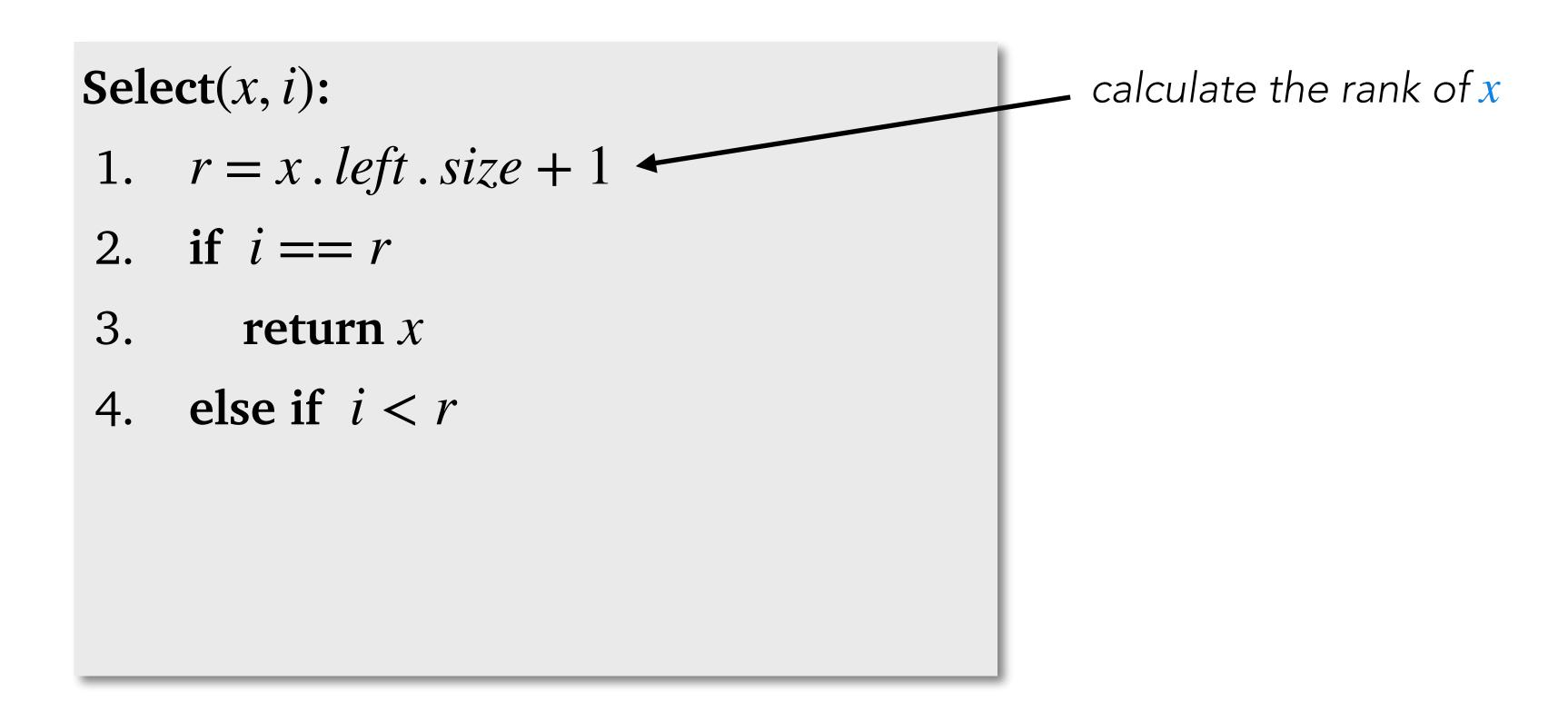
```
Select(x, i):

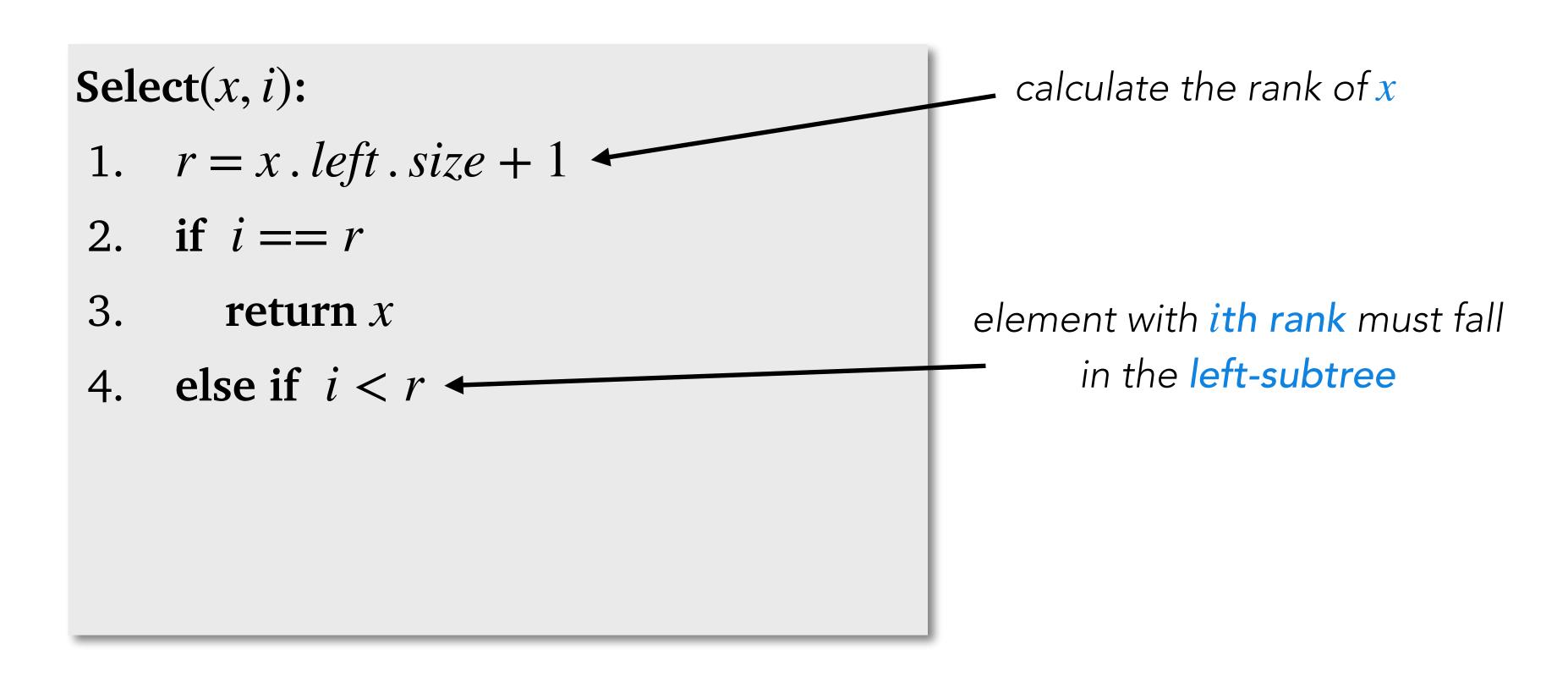
1. r = x \cdot left \cdot size + 1
```

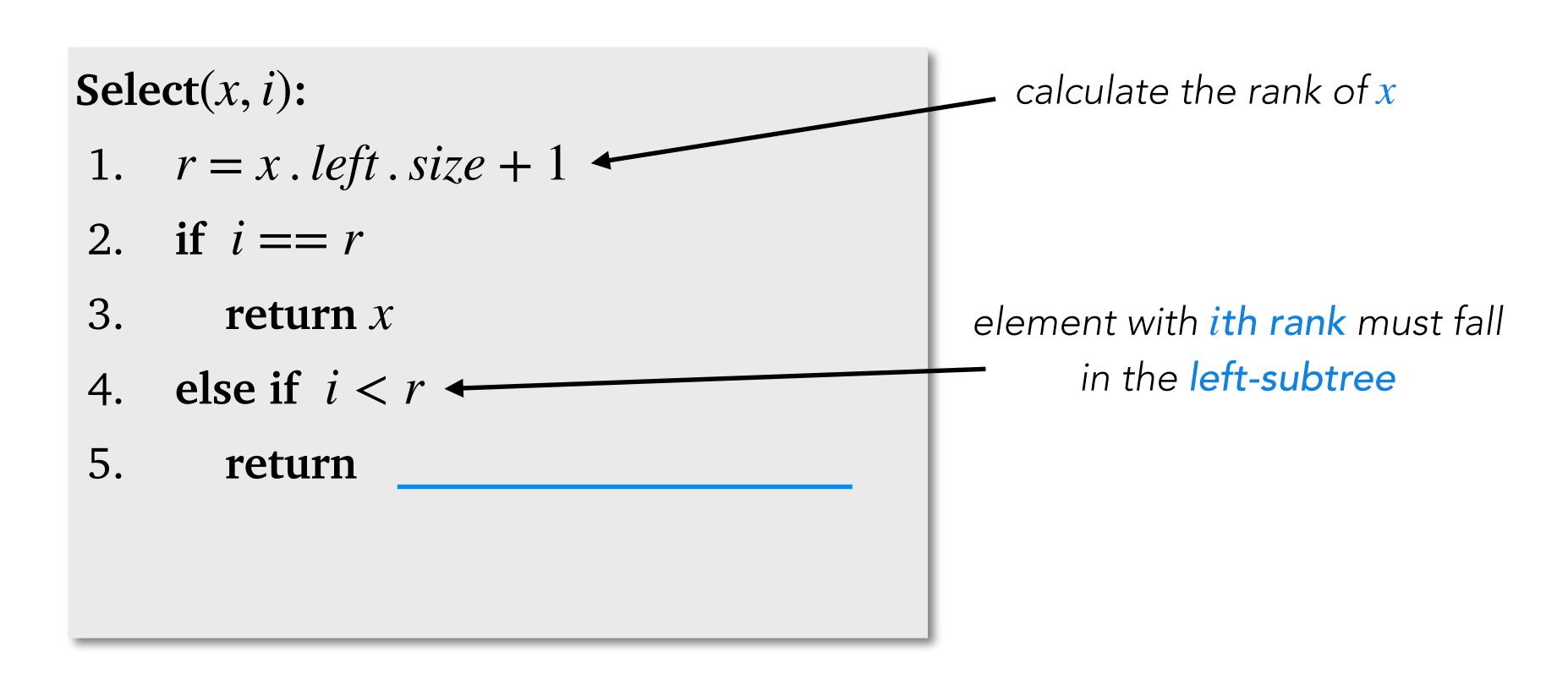


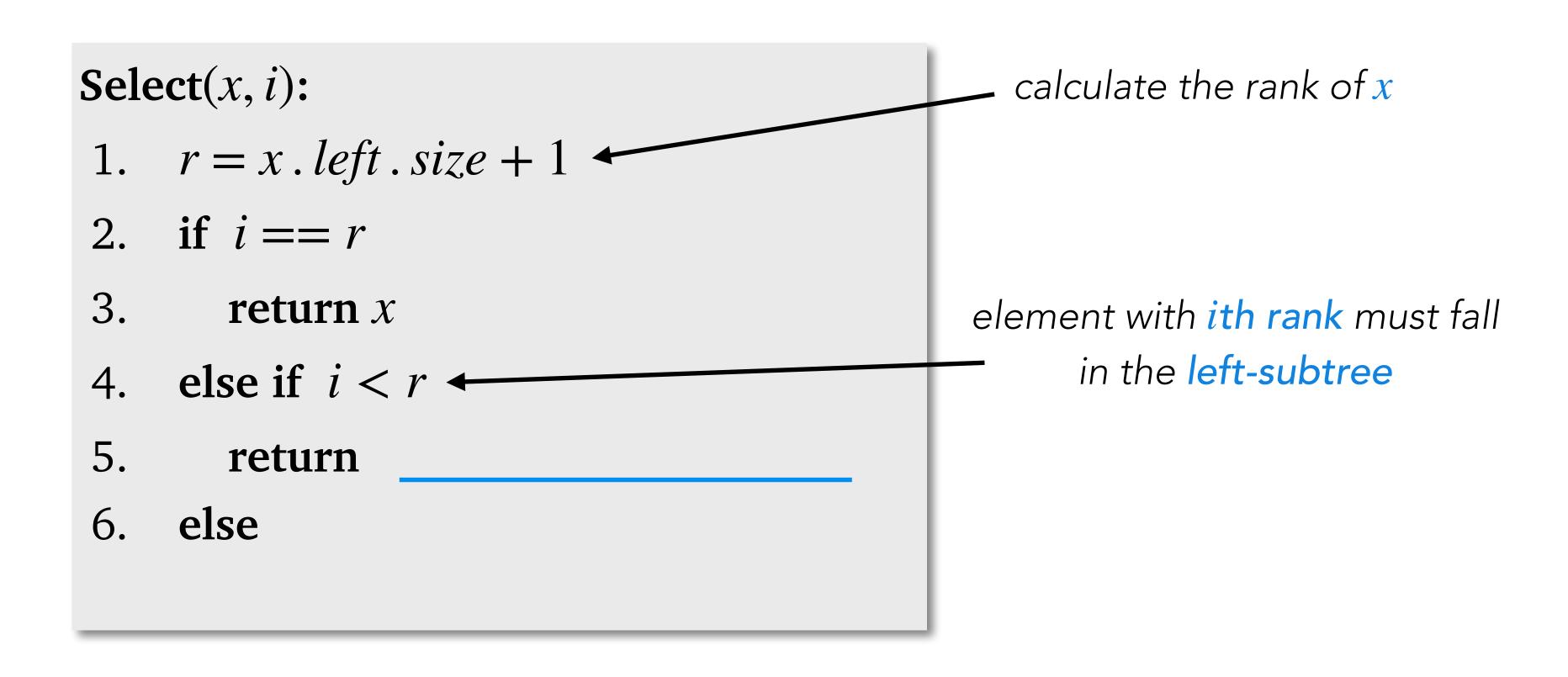


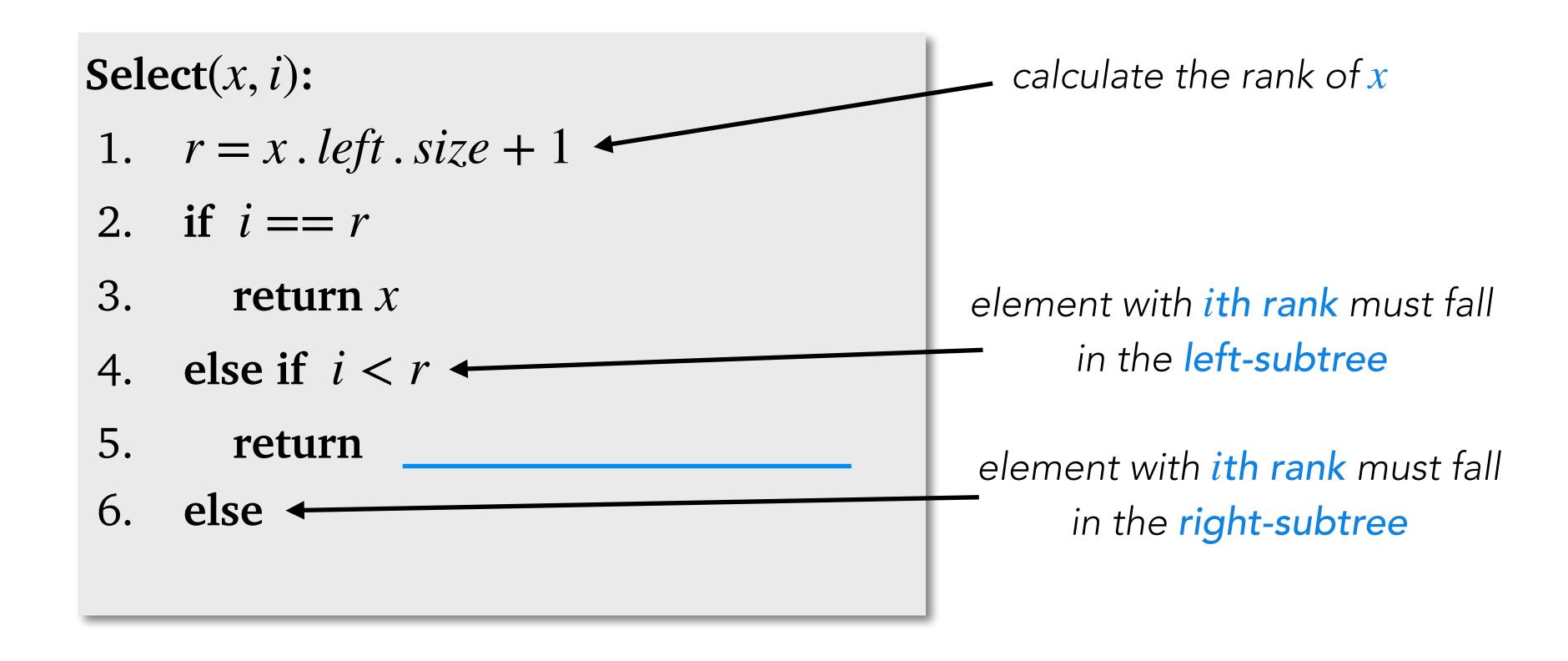


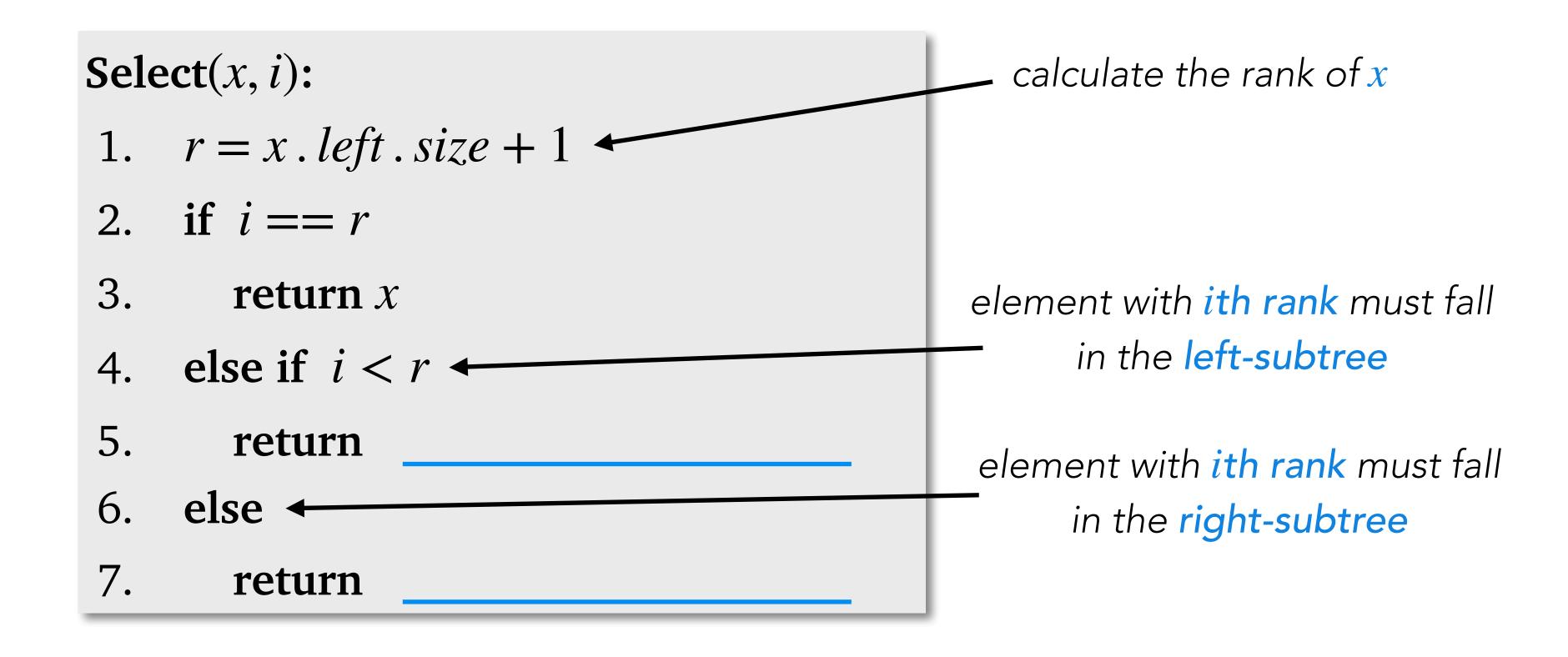


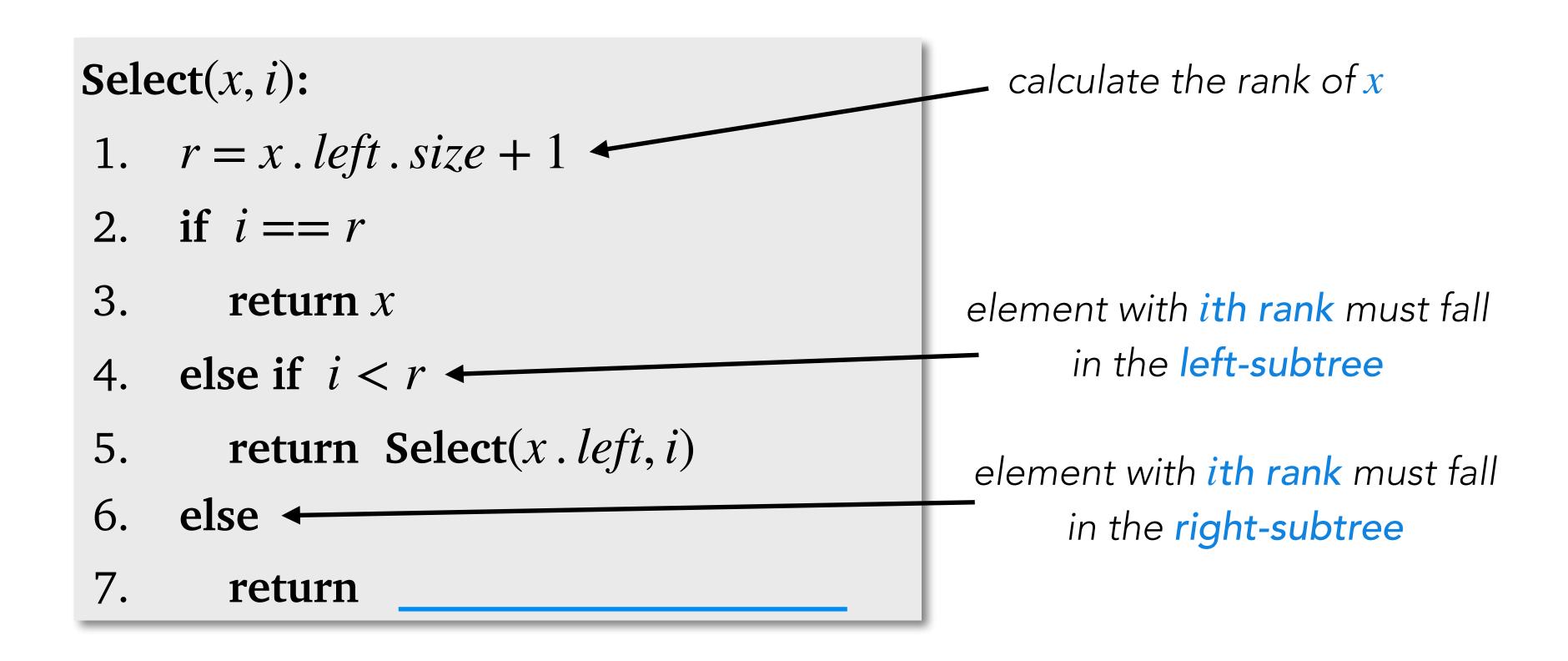


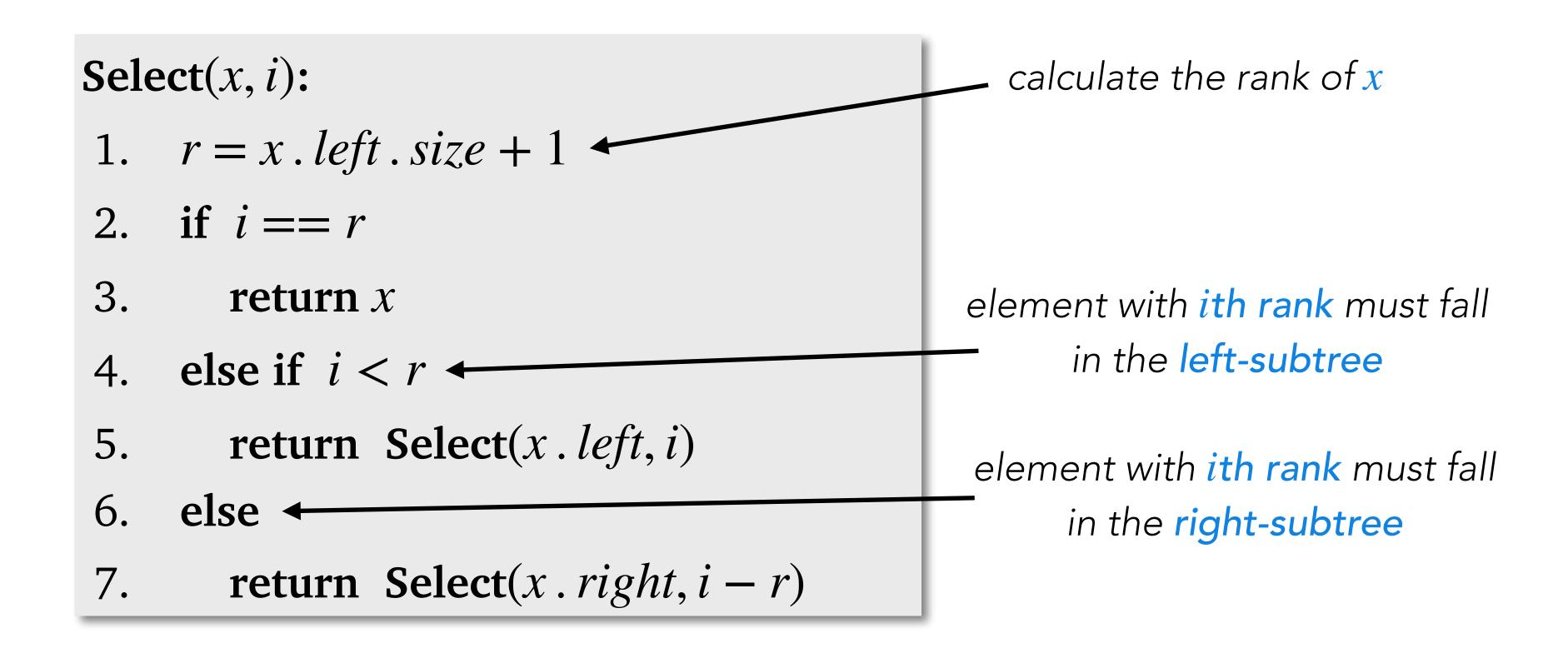




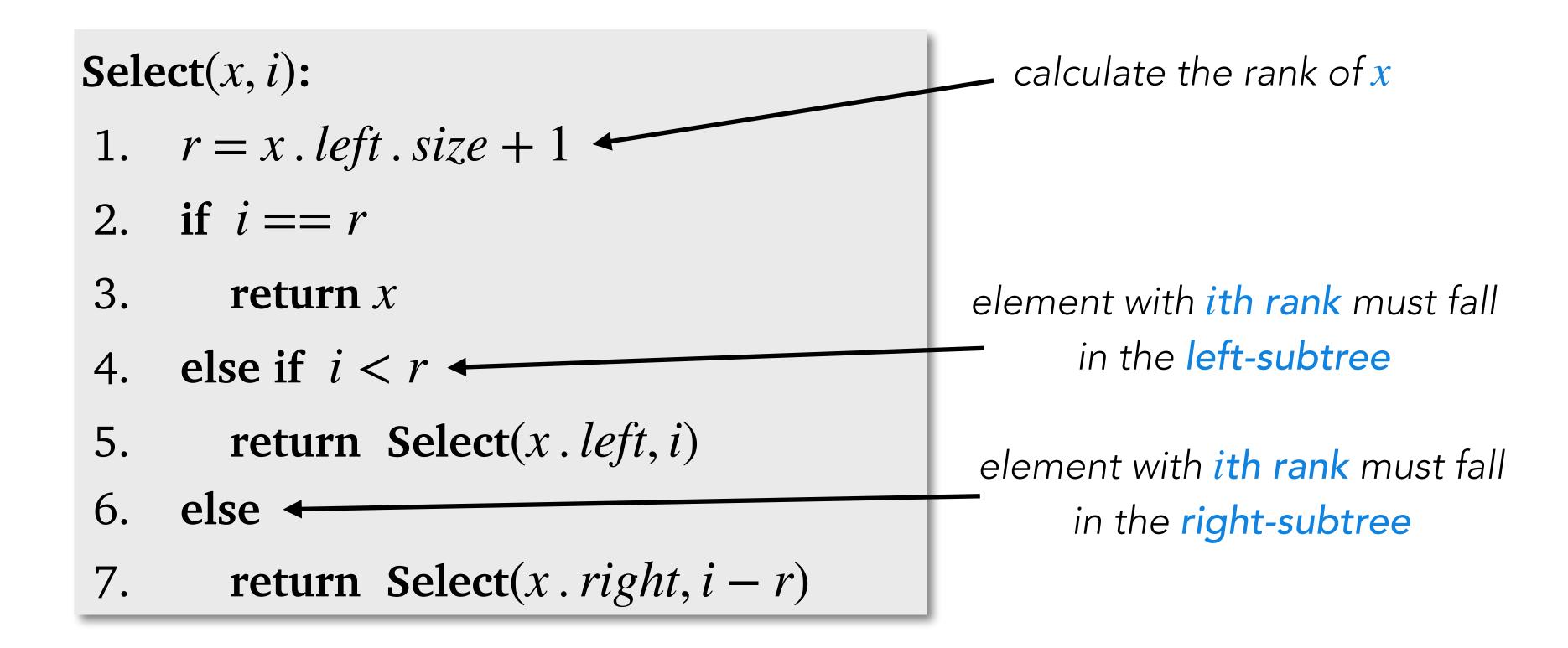




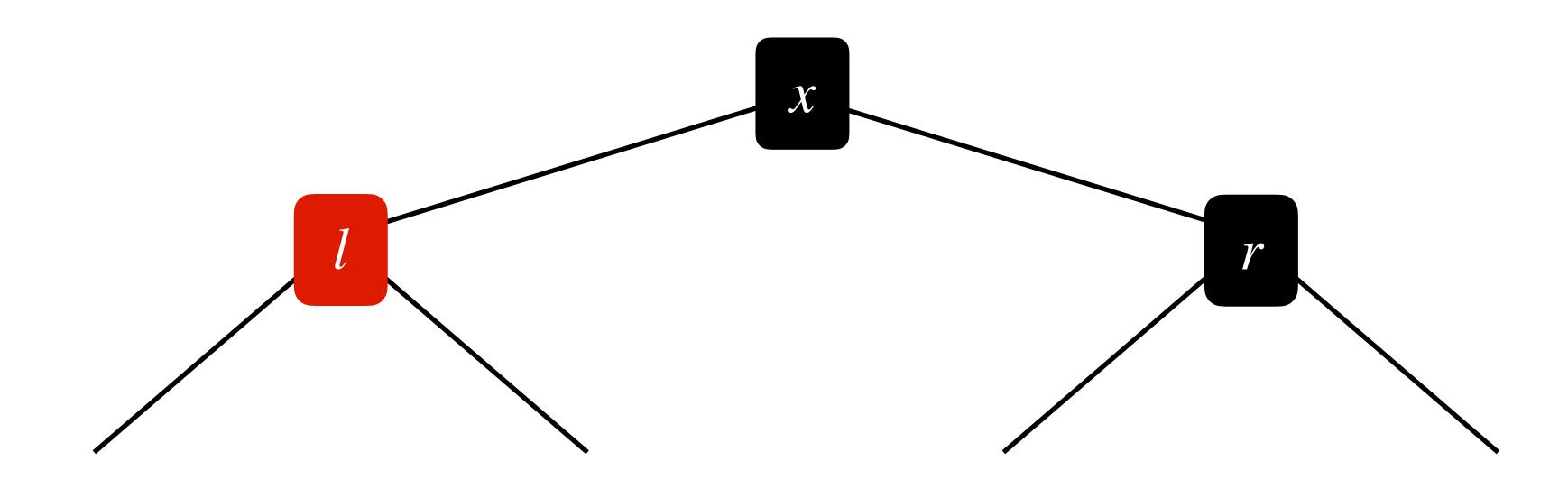


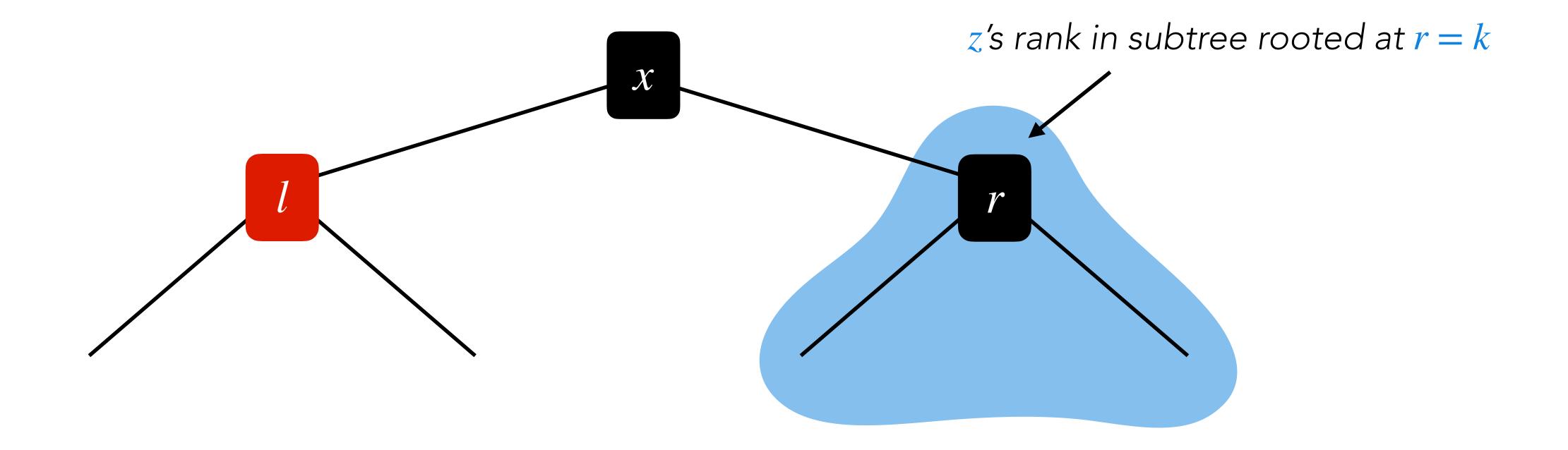


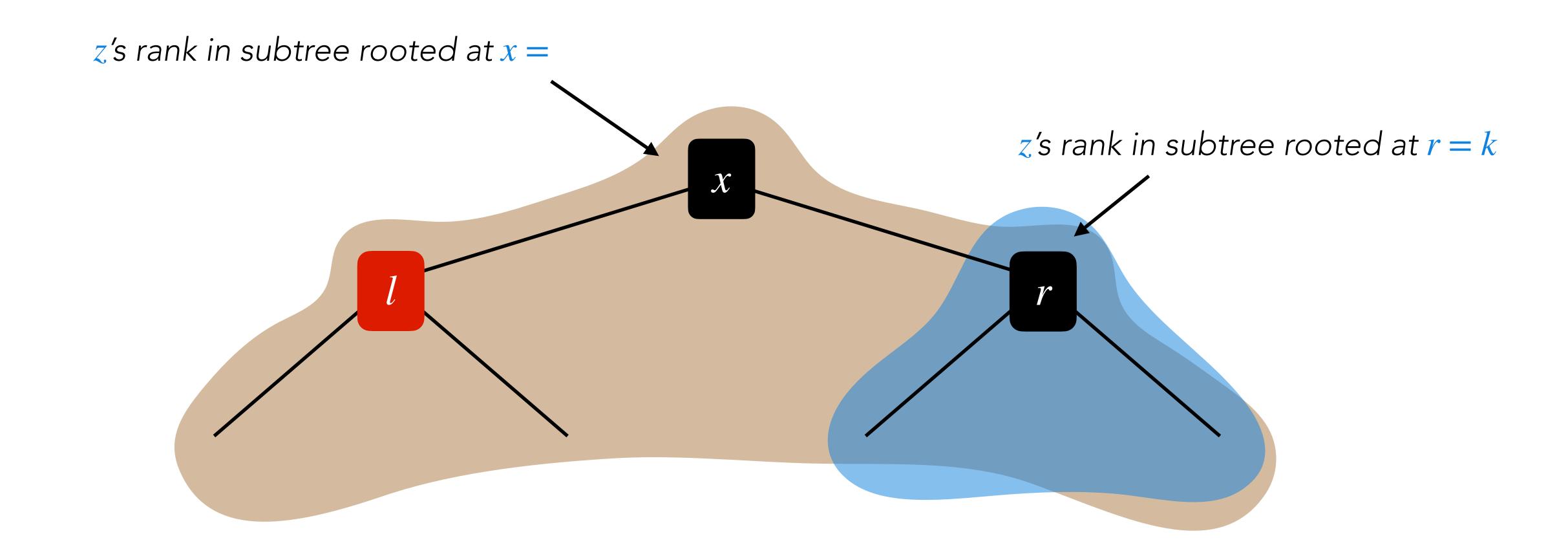
To find the element with *i*th rank in T call **Select**(T. root, i).



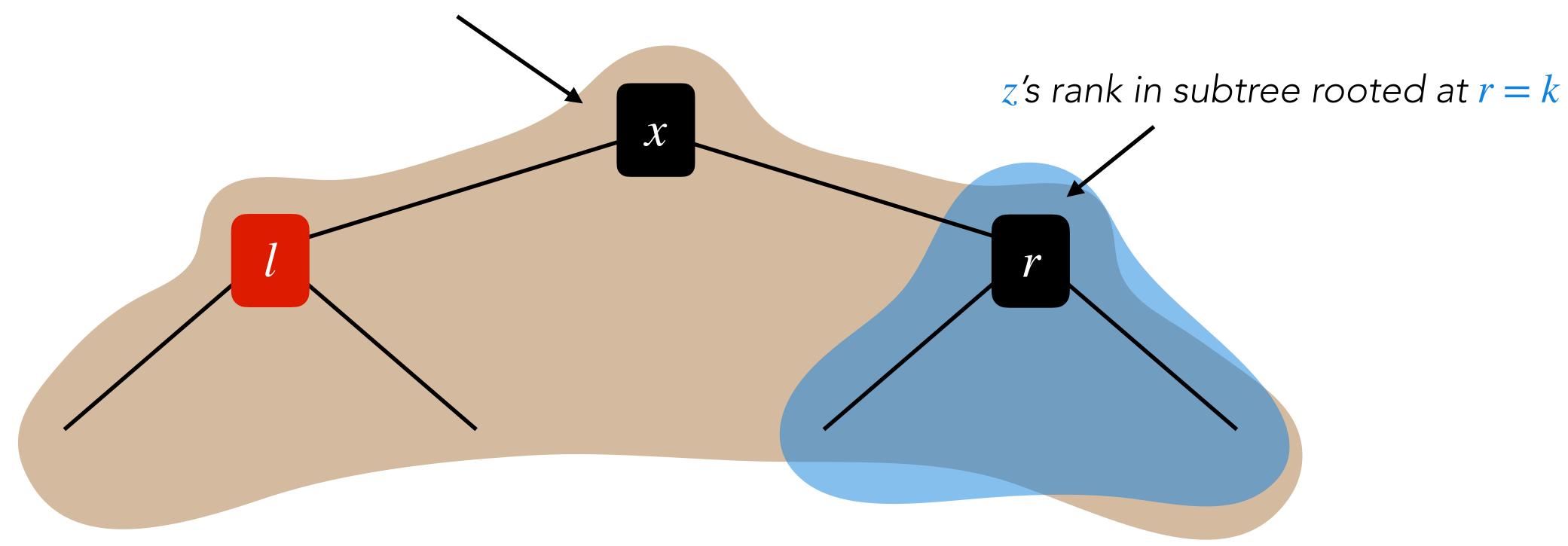
Time Complexity:  $O(h) = O(\log n)$  as with every recursive call algorithm goes one level down.

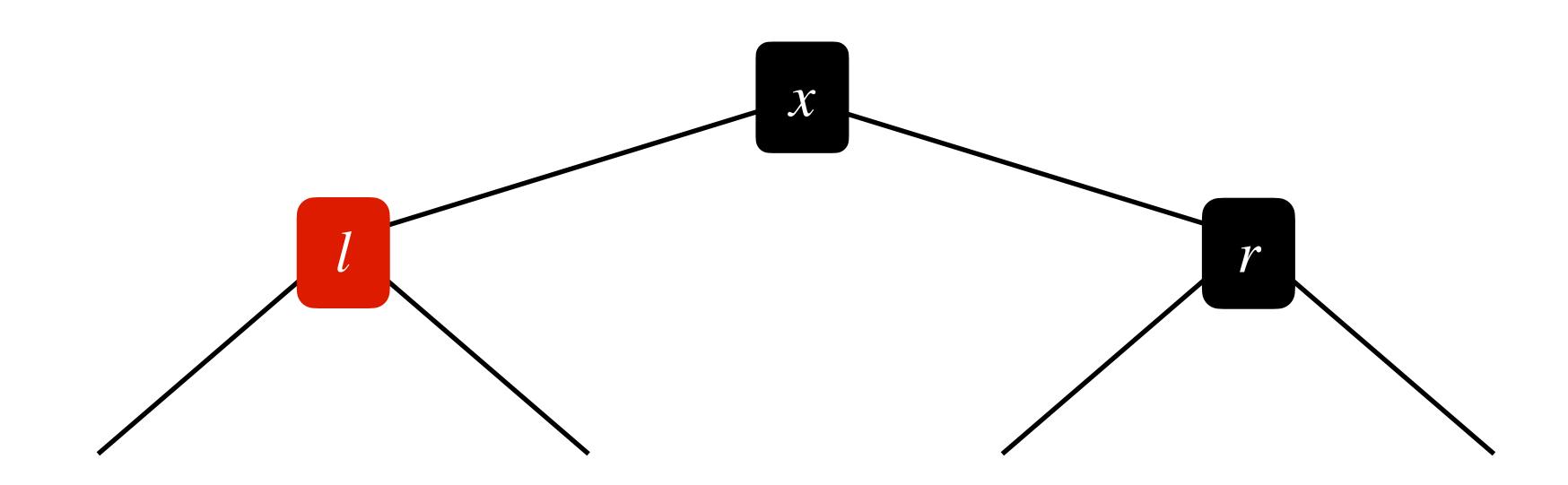


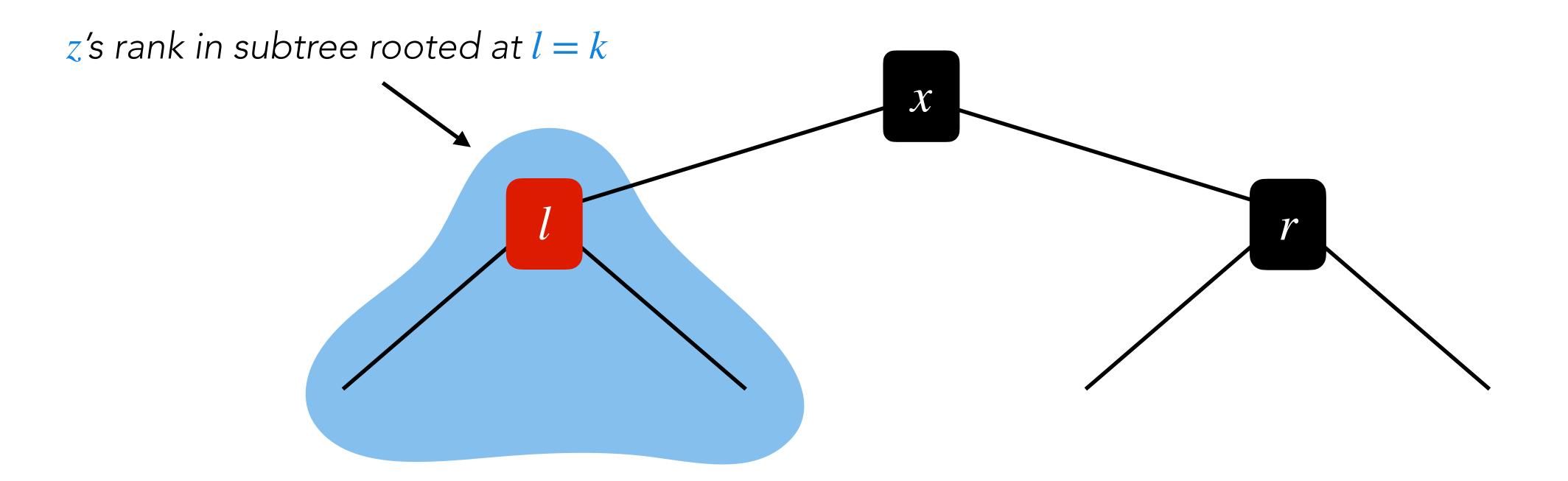


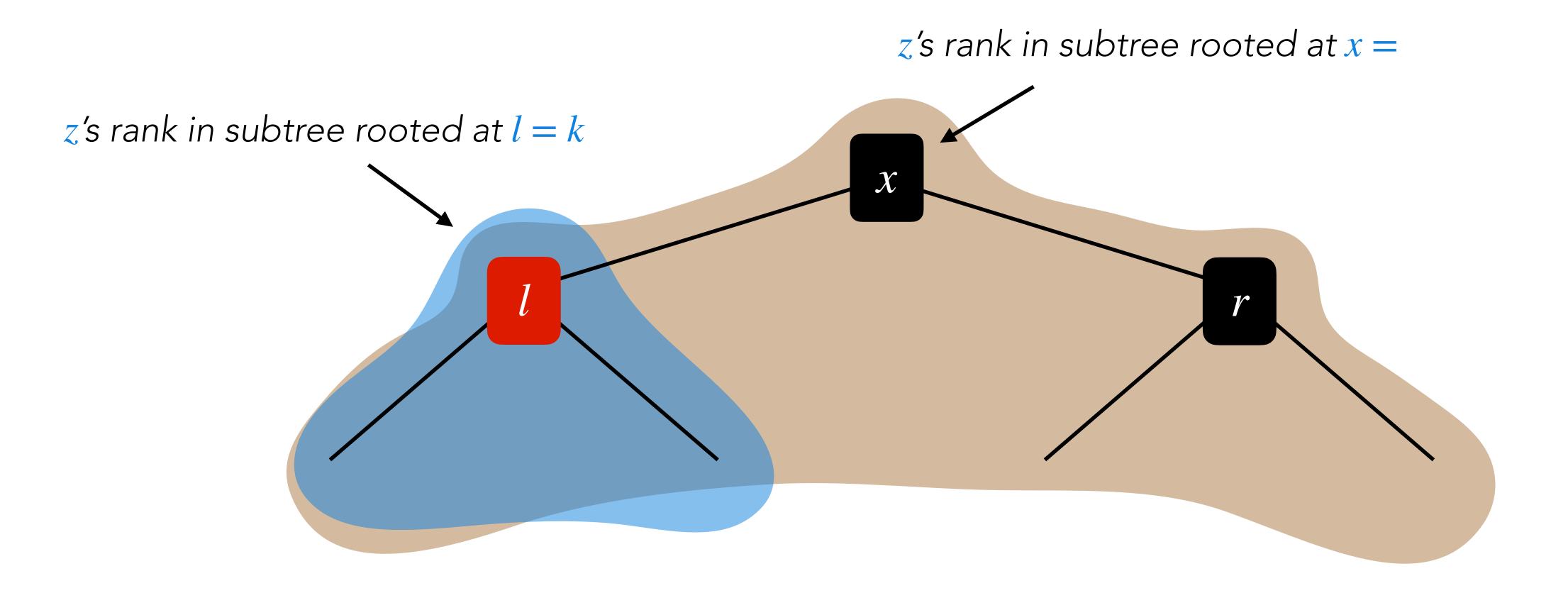


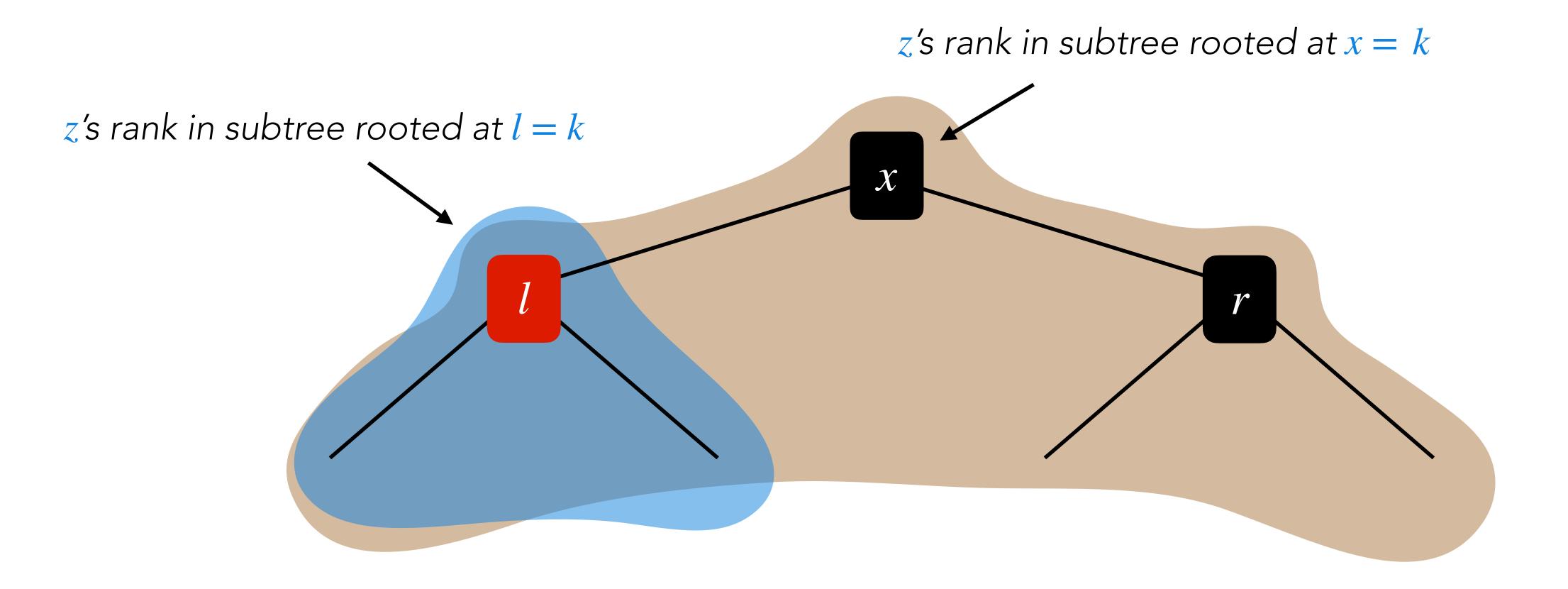
z's rank in subtree rooted at x = k + l. size + 1





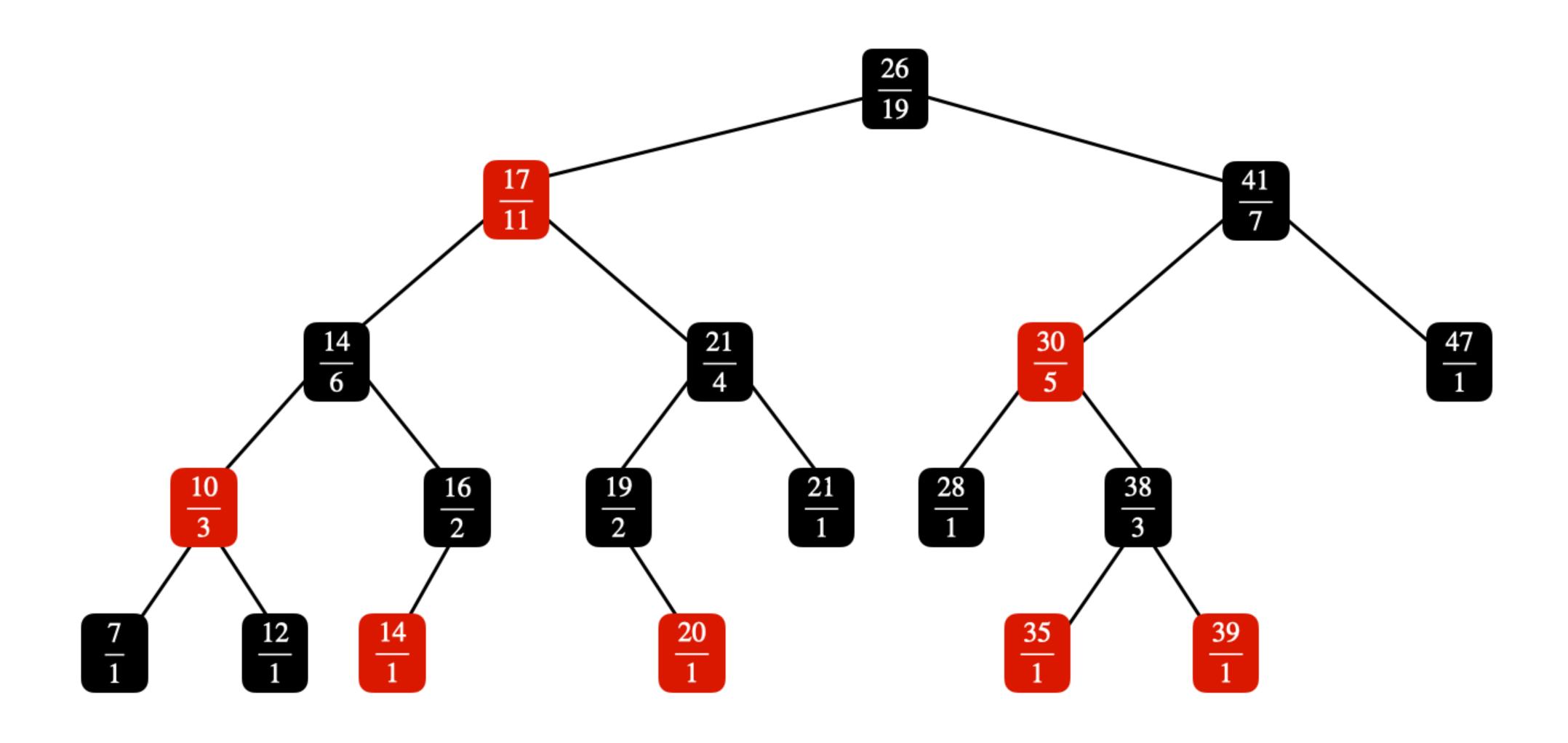




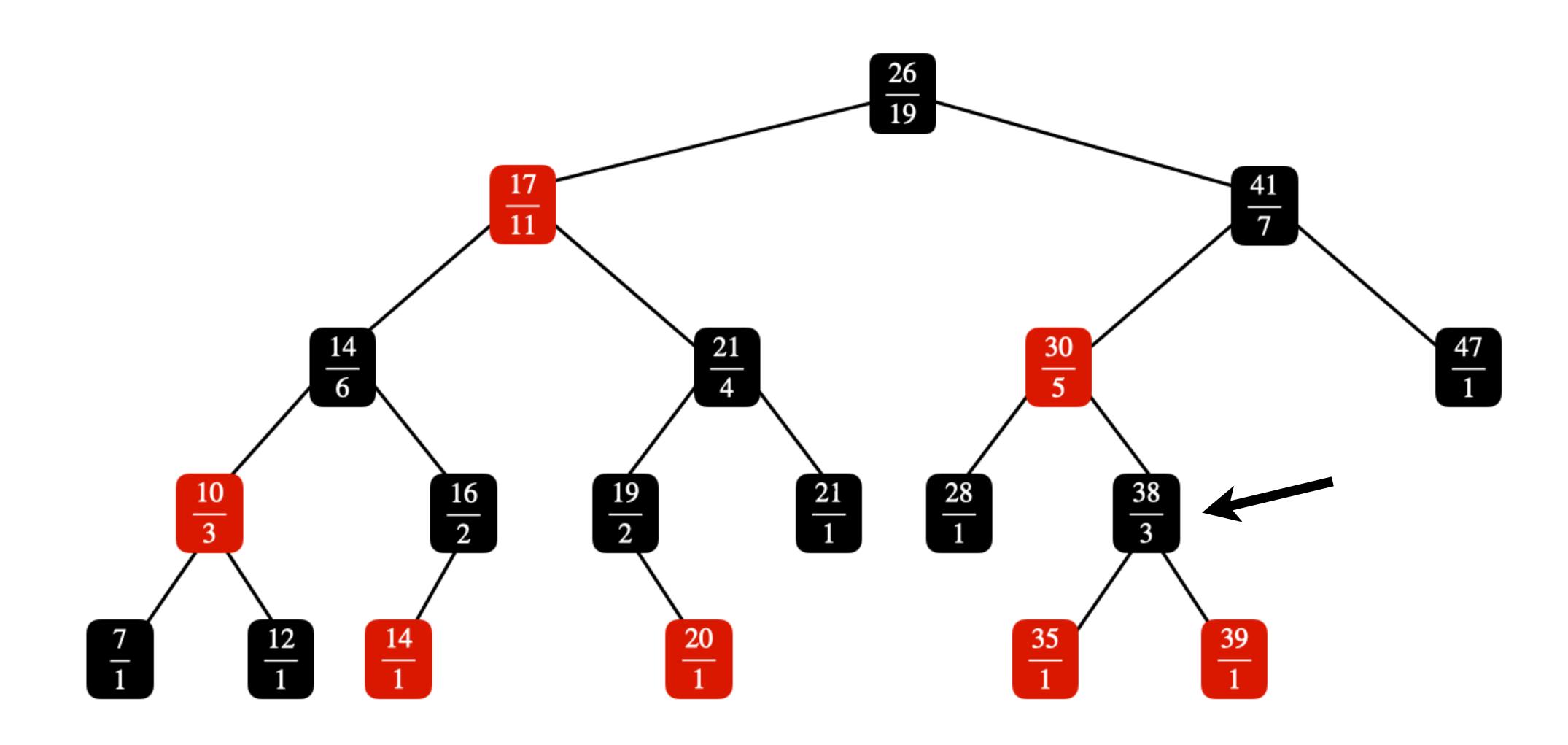


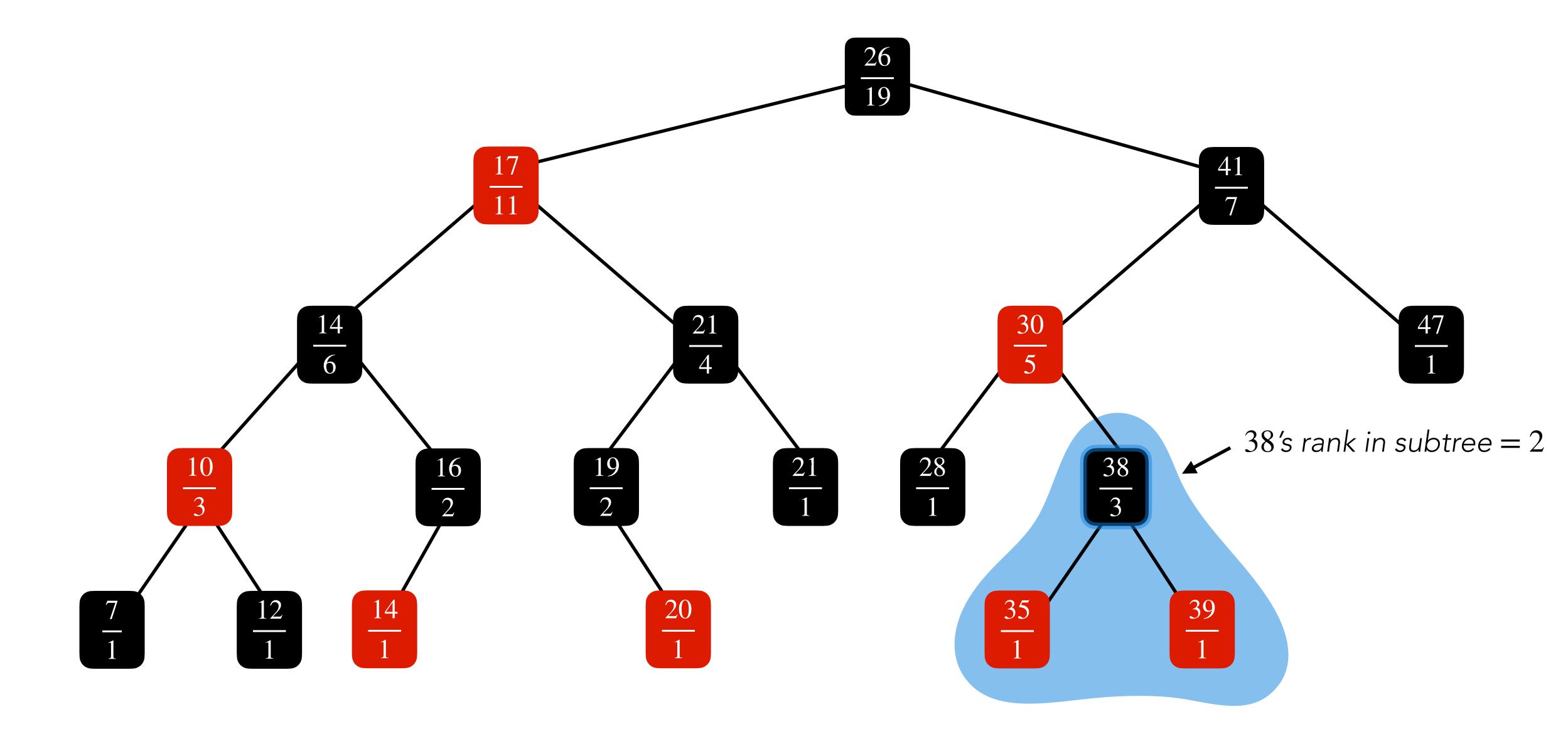
**Example:** Find the rank of 38 in the below set or RB-tree.

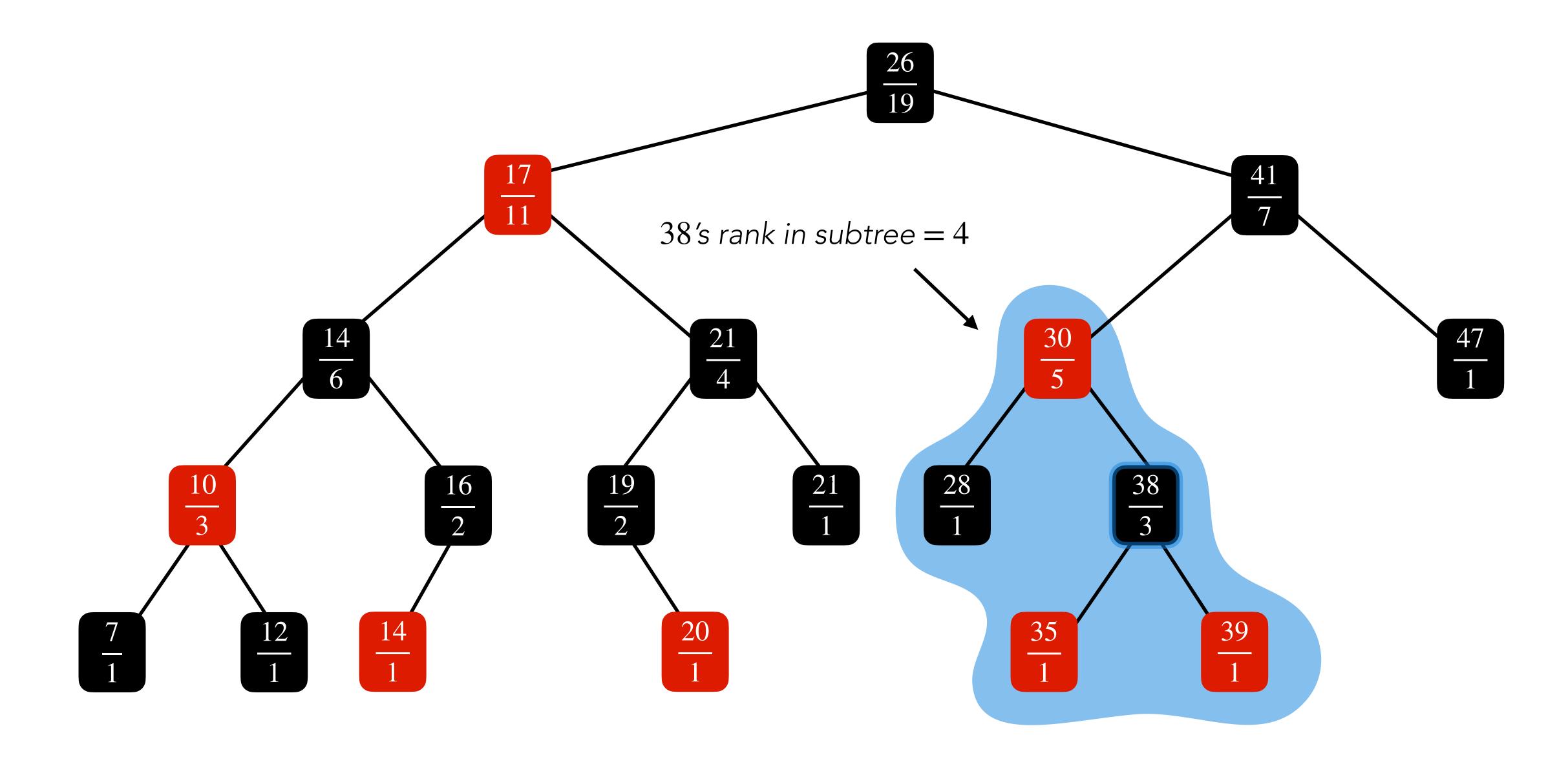
**Example:** Find the rank of 38 in the below set or RB-tree.

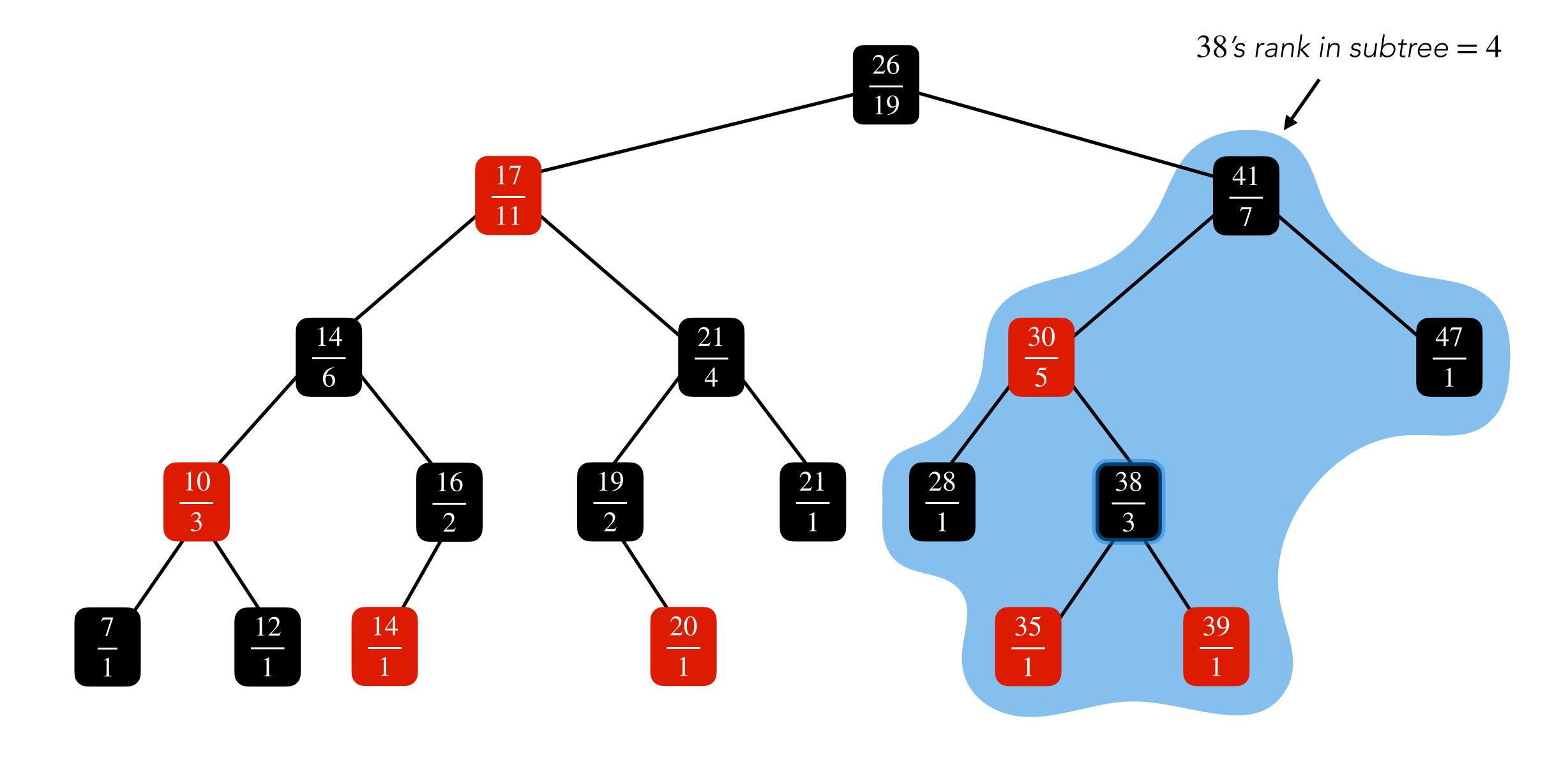


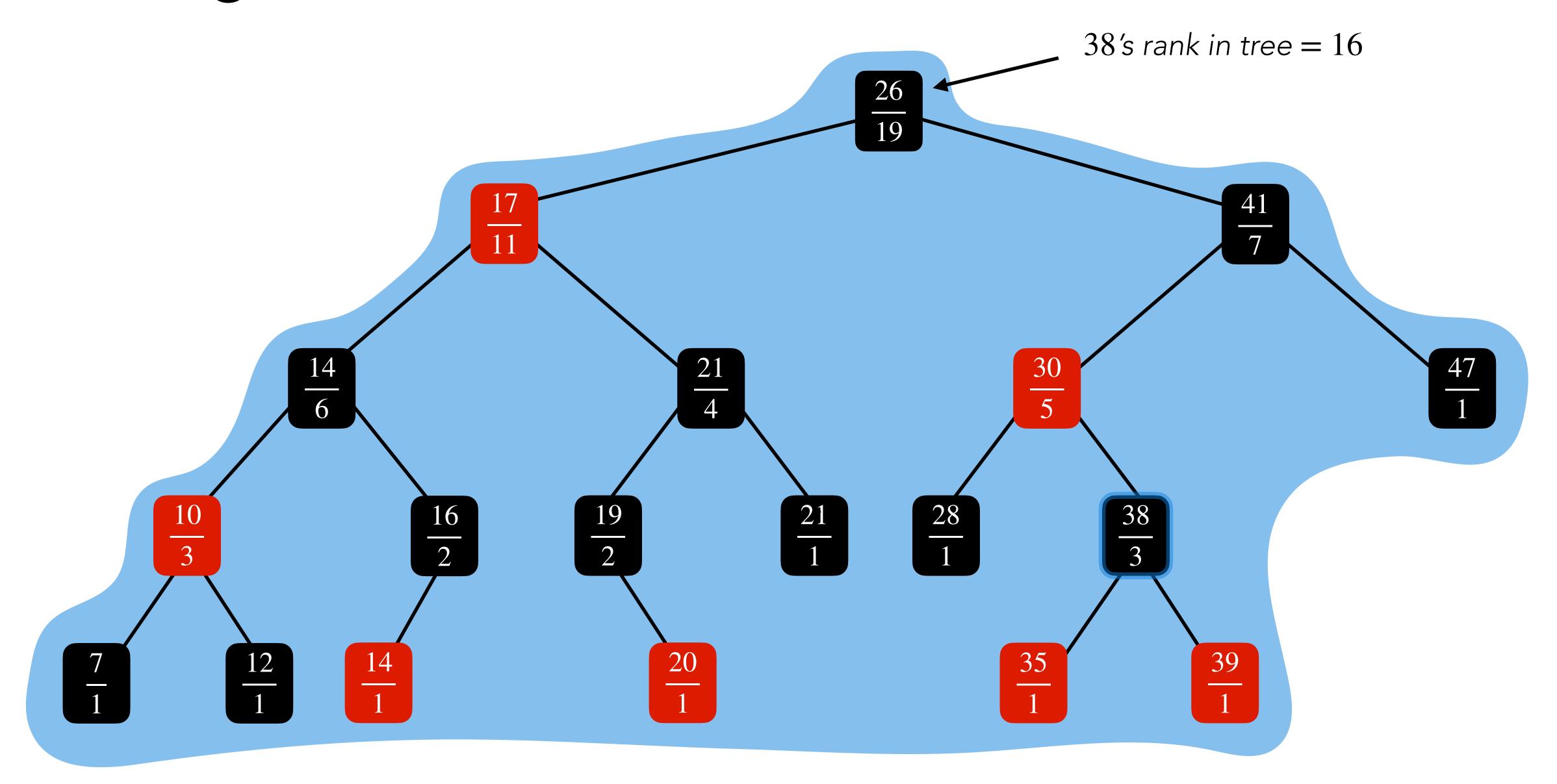
**Example:** Find the rank of 38 in the below set or RB-tree.







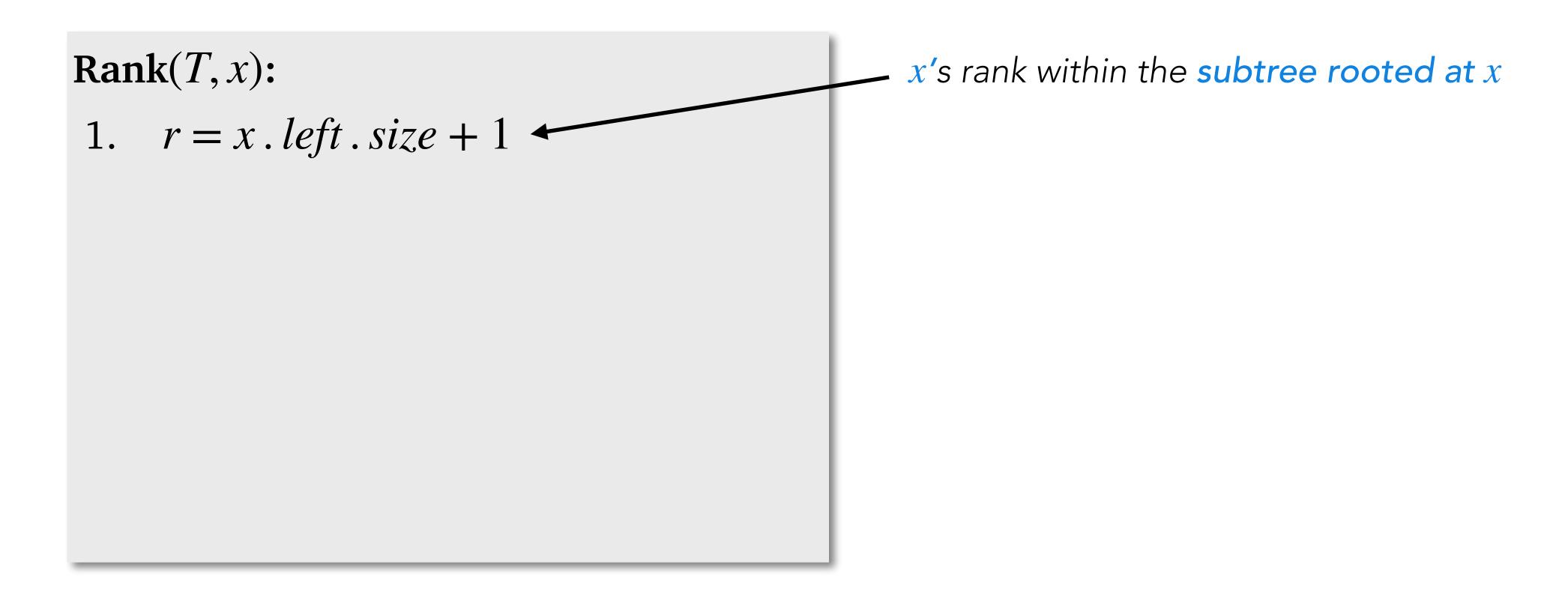


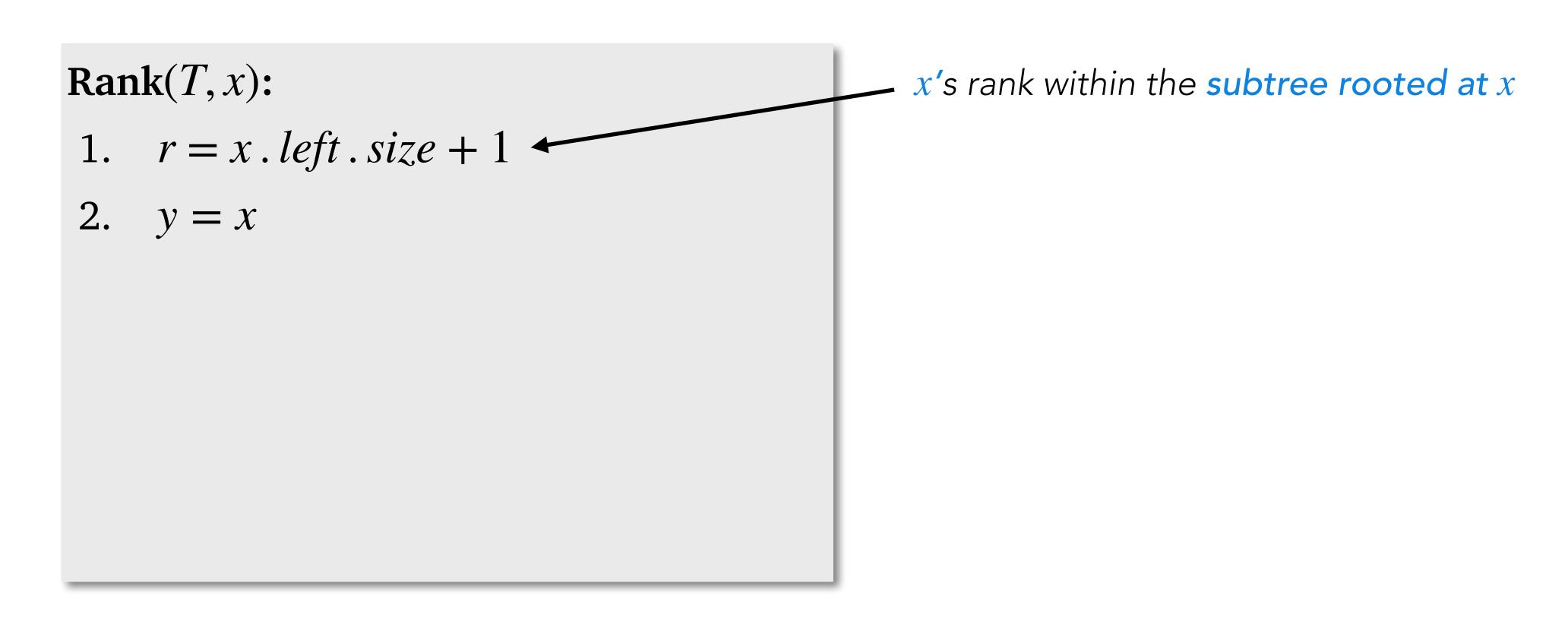


```
Rank(T, x):
```

```
Rank(T, x):

1. r = x \cdot left \cdot size + 1
```





To find the rank of an element x in T call Rank(T, x).

#### Rank(T, x):

- 1.  $r = x \cdot left \cdot size + 1$
- 2. y = x
- 3. while  $y \neq T$ . root

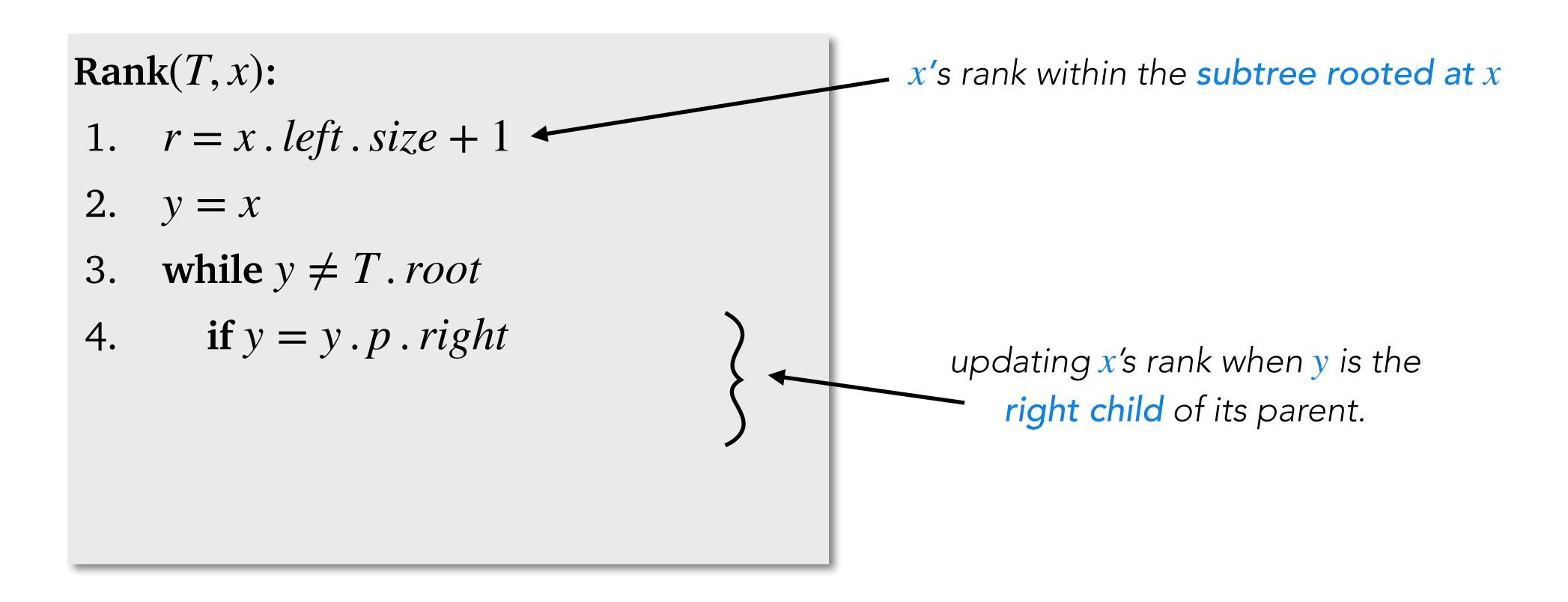
x's rank within the subtree rooted at x

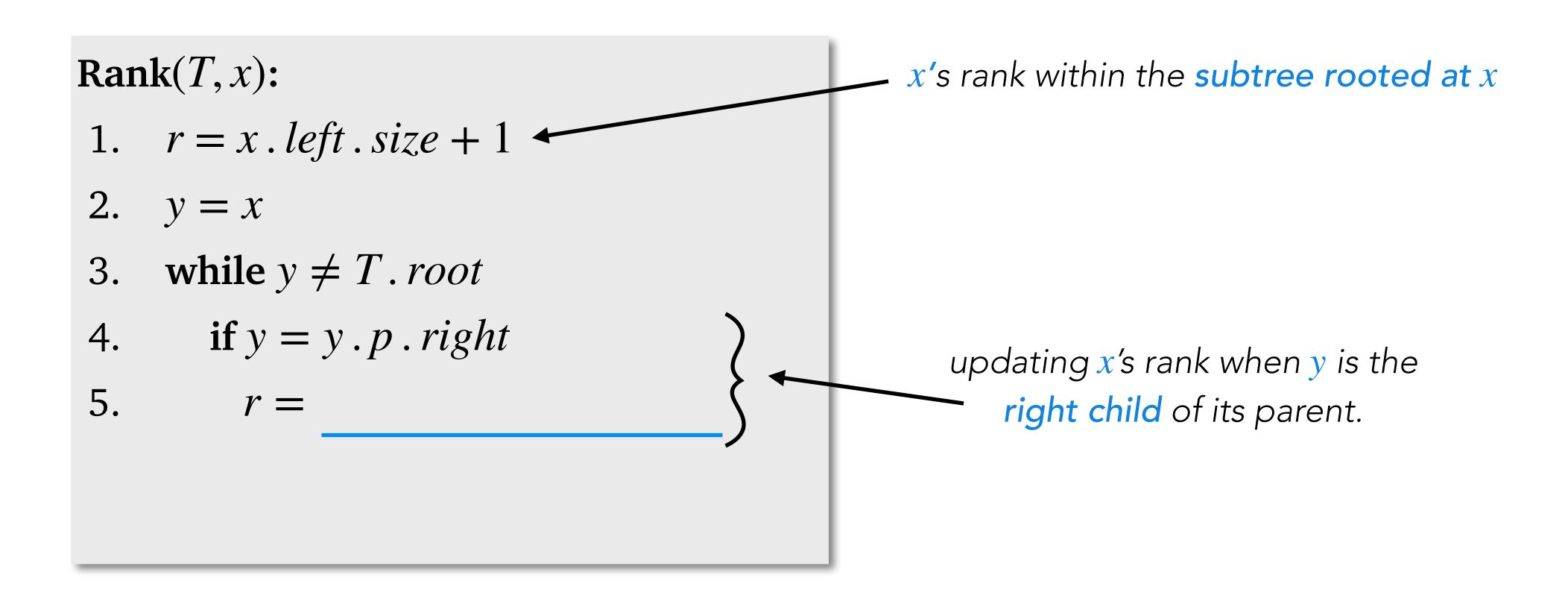
To find the rank of an element x in T call Rank(T, x).

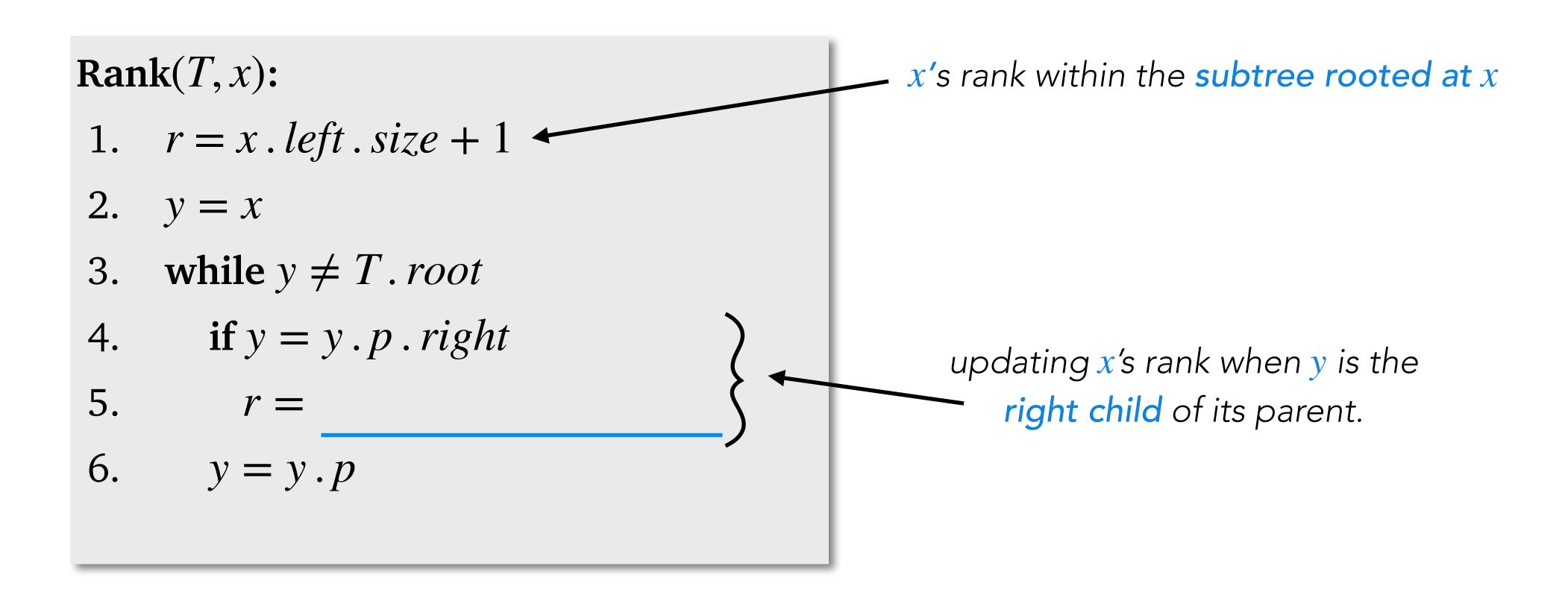
#### Rank(T, x):

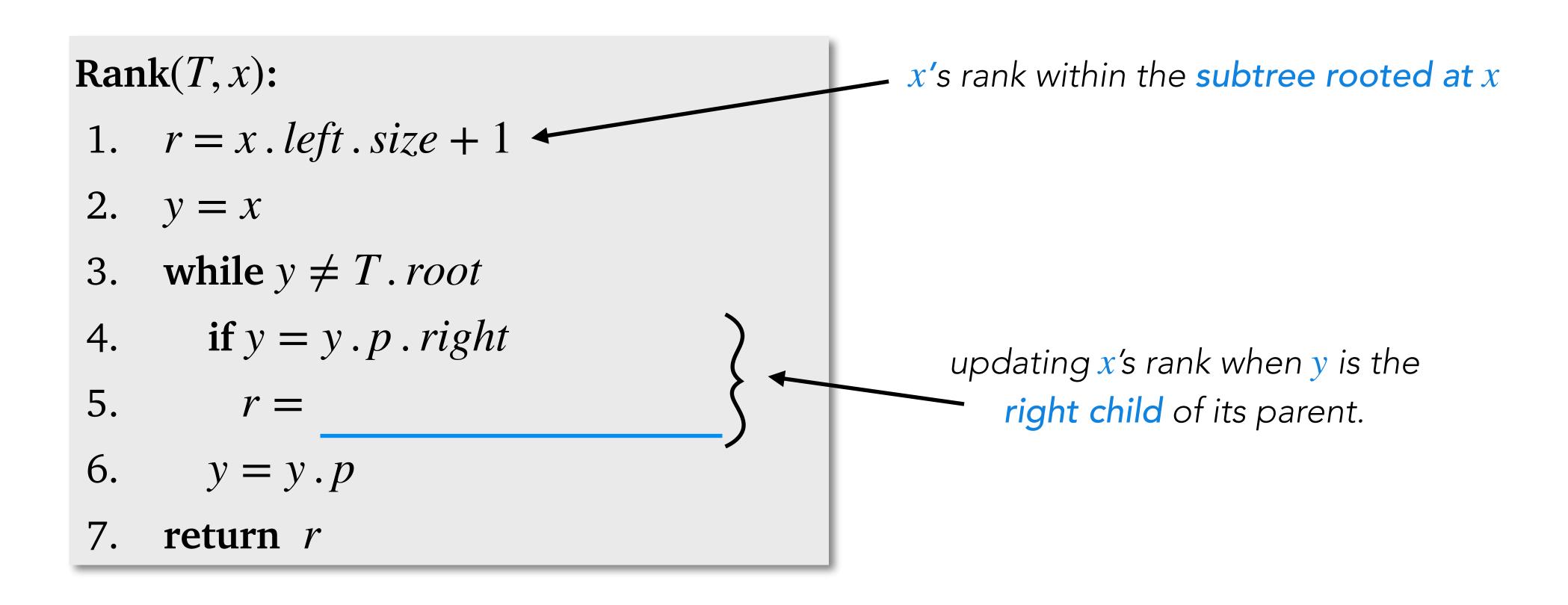
- 1.  $r = x \cdot left \cdot size + 1$
- 2. y = x
- 3. while  $y \neq T$ . root
- 4. if  $y = y \cdot p \cdot right$

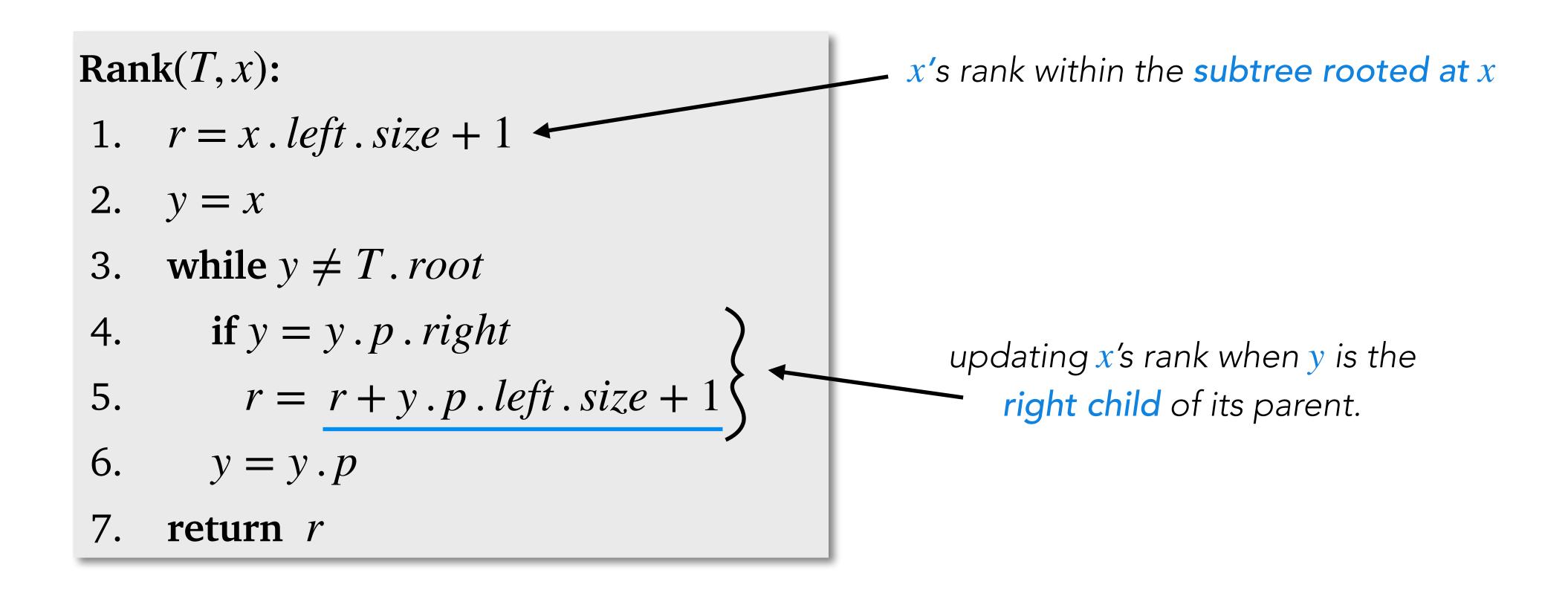
x's rank within the subtree rooted at x



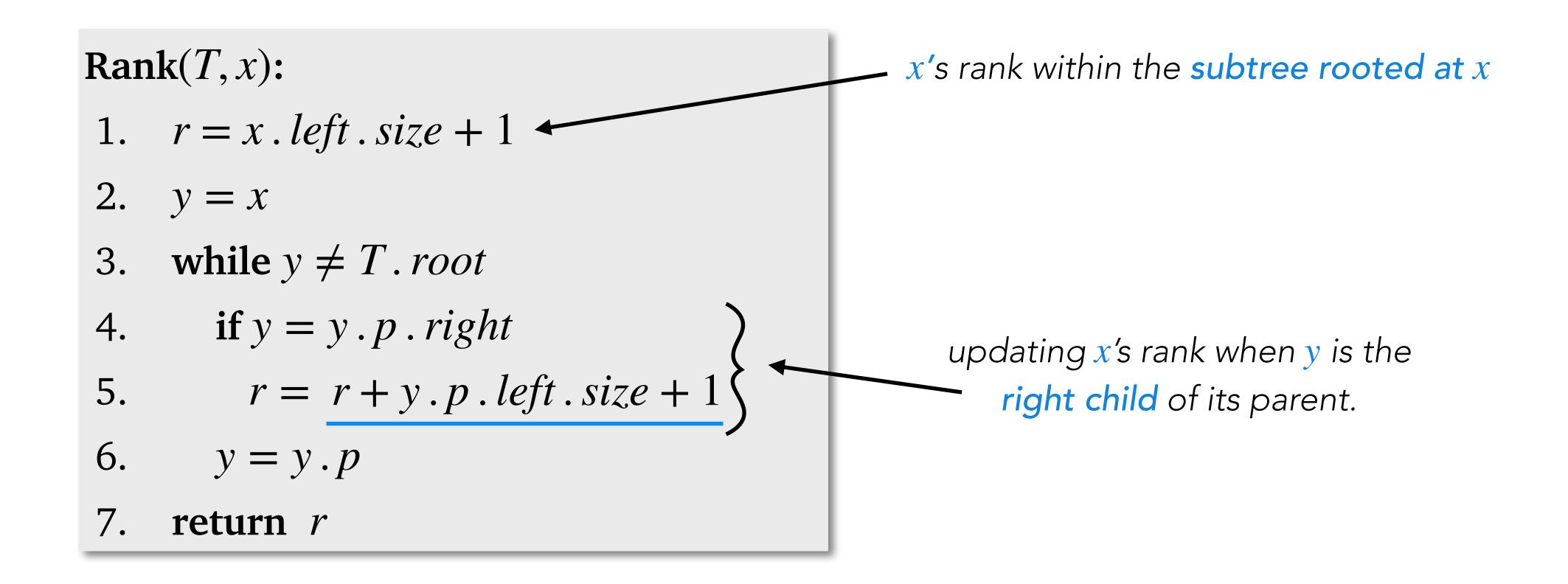




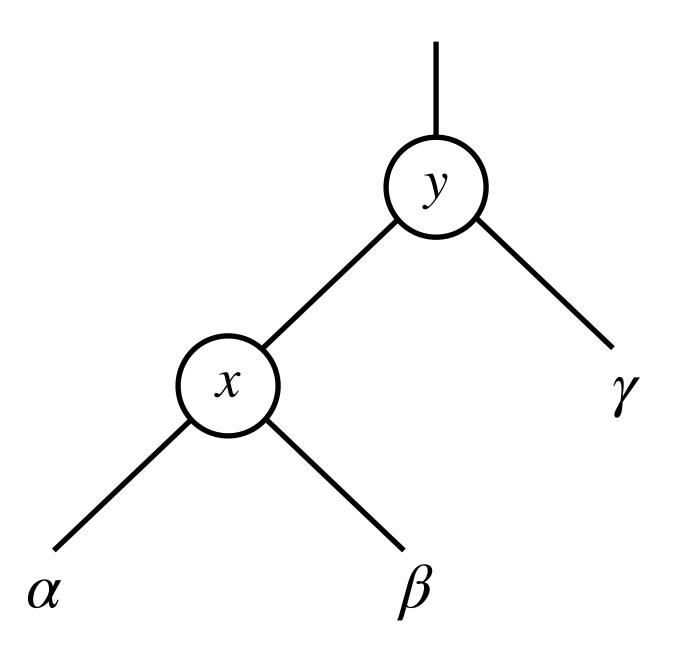


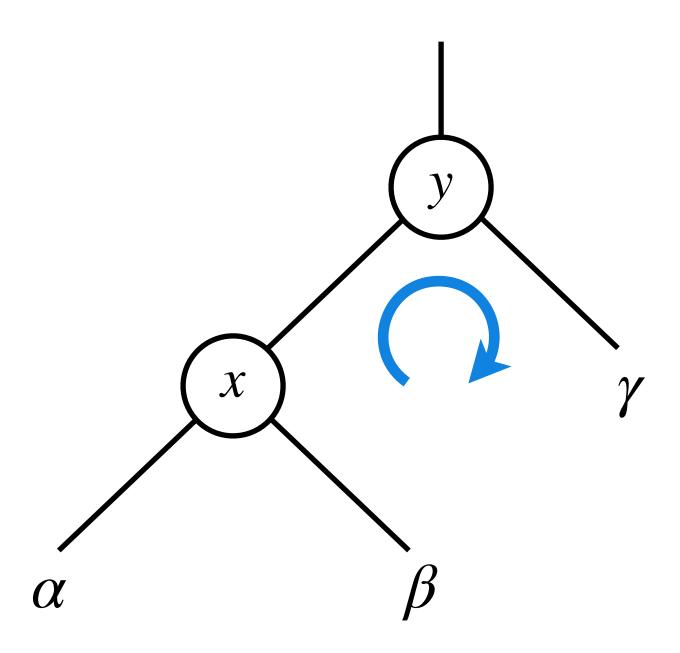


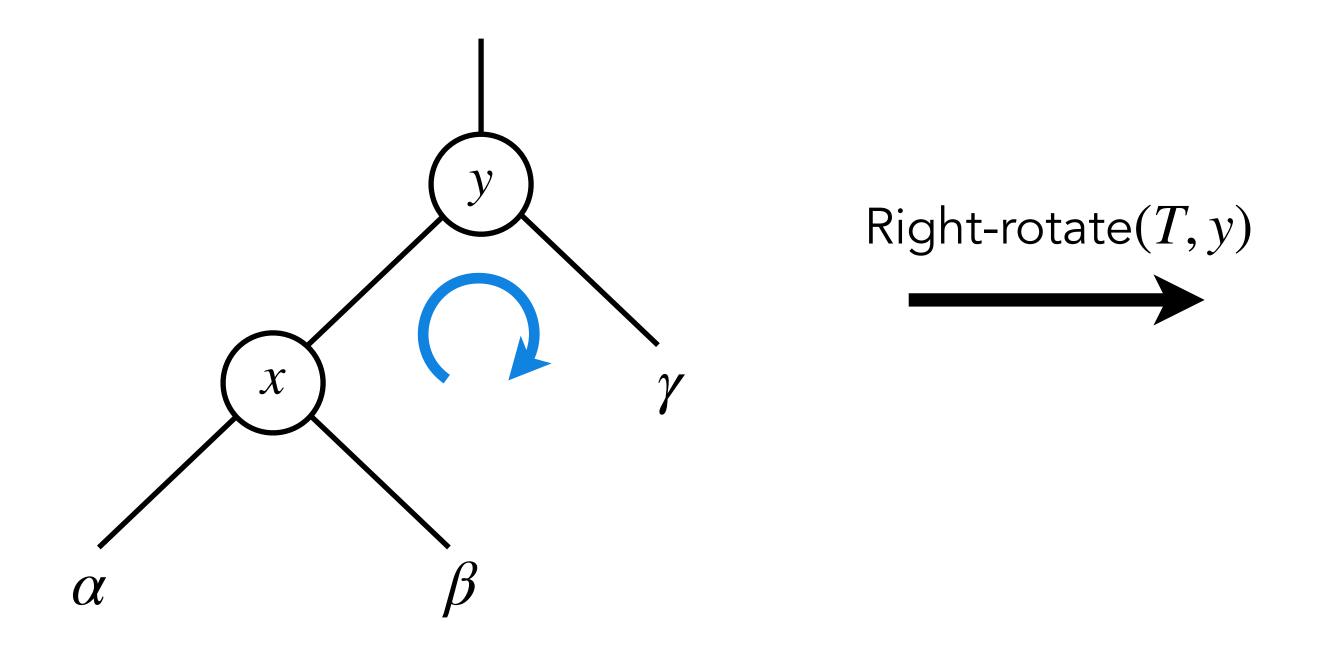
To find the rank of an element x in T call Rank(T, x).

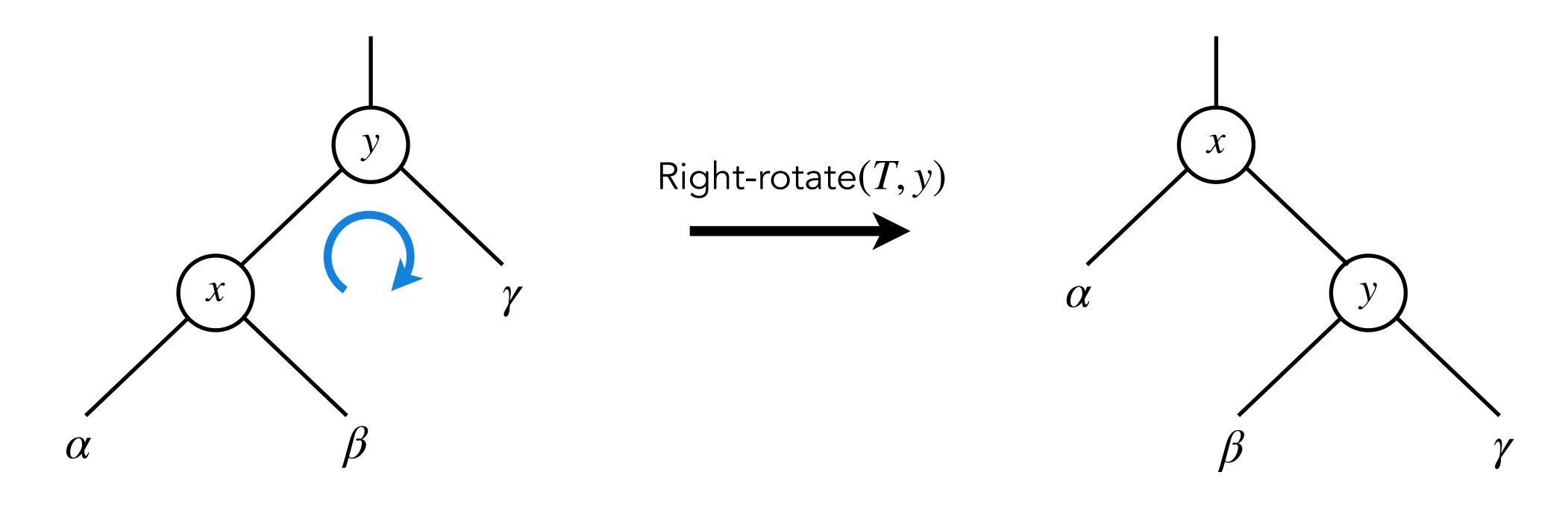


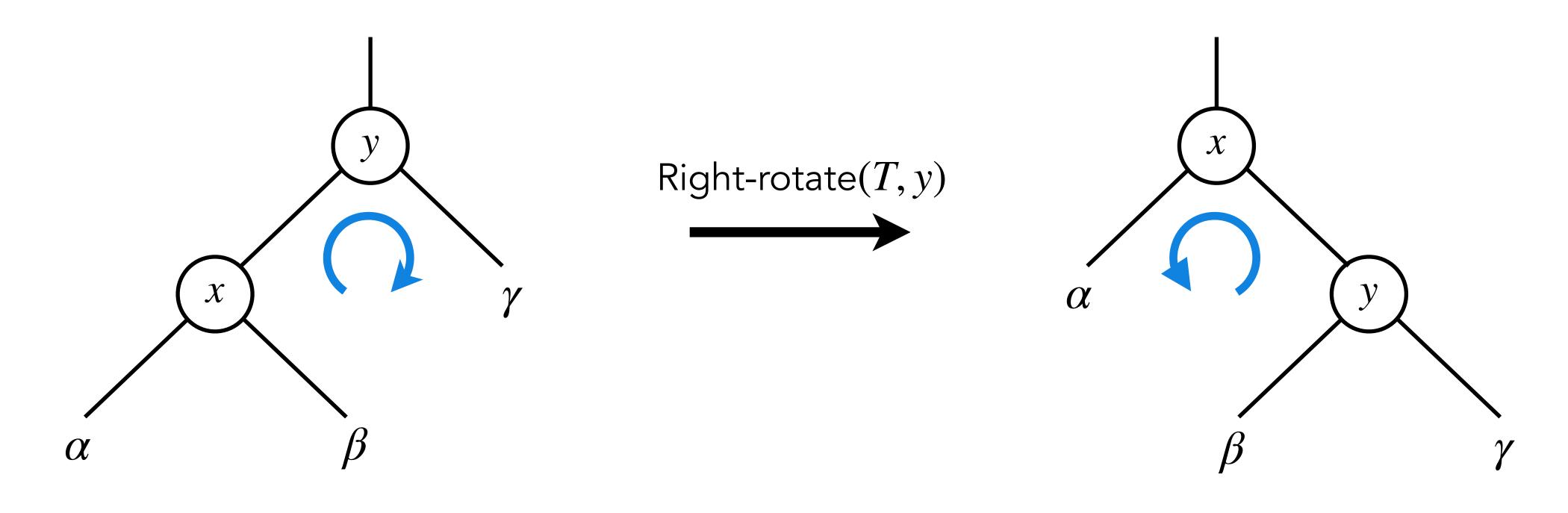
Time Complexity:  $O(h) = O(\log n)$  as with every iteration of while loop, y moves closer to root.

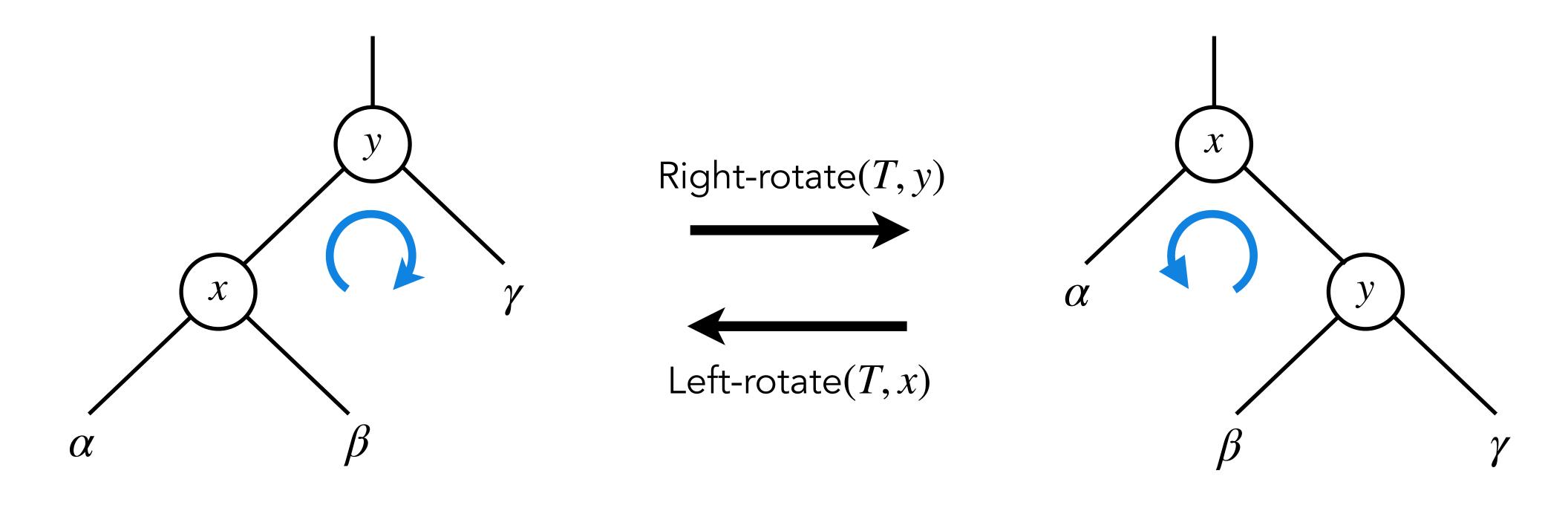




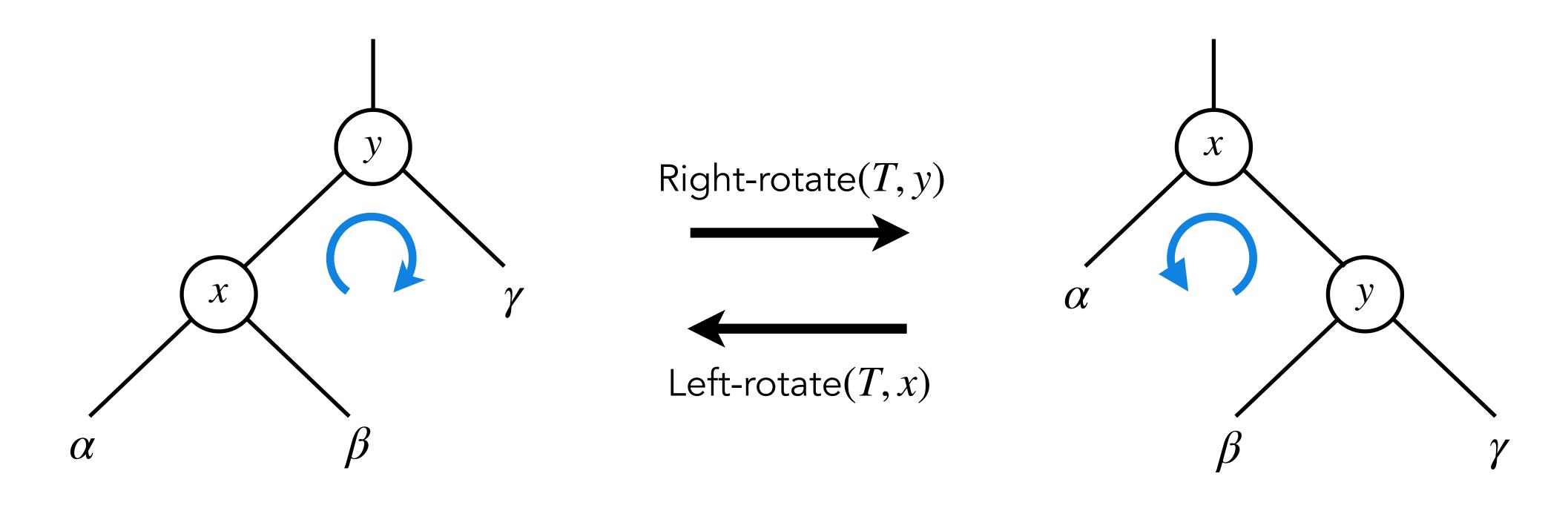




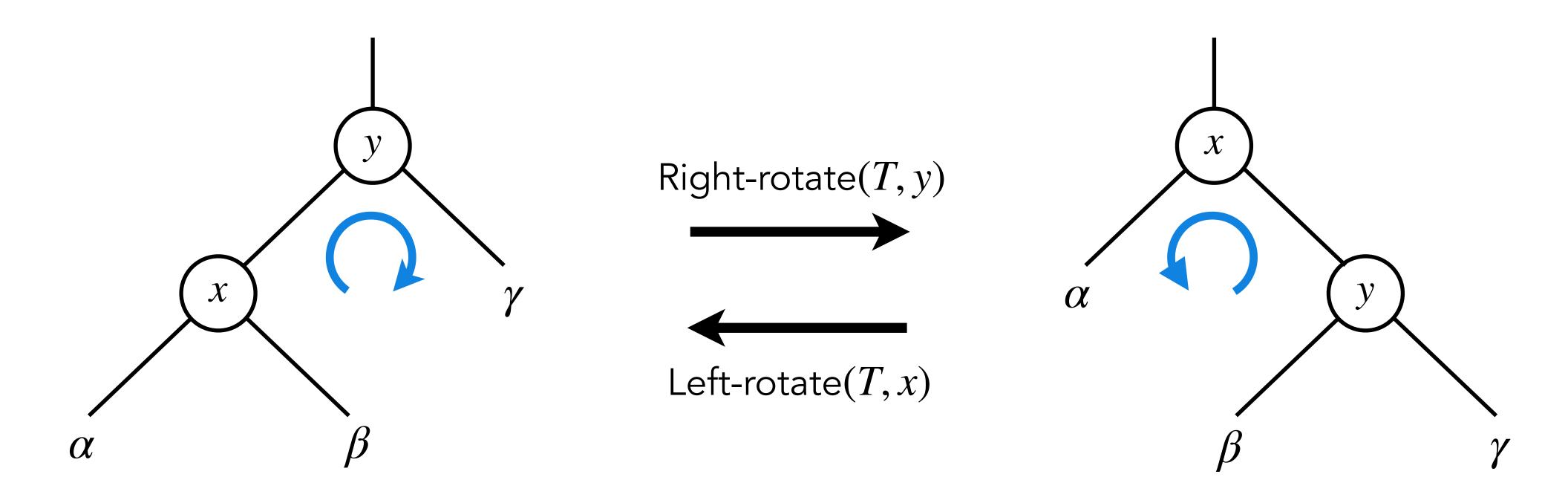




Maintaining subtree sizes during rotations is easy.

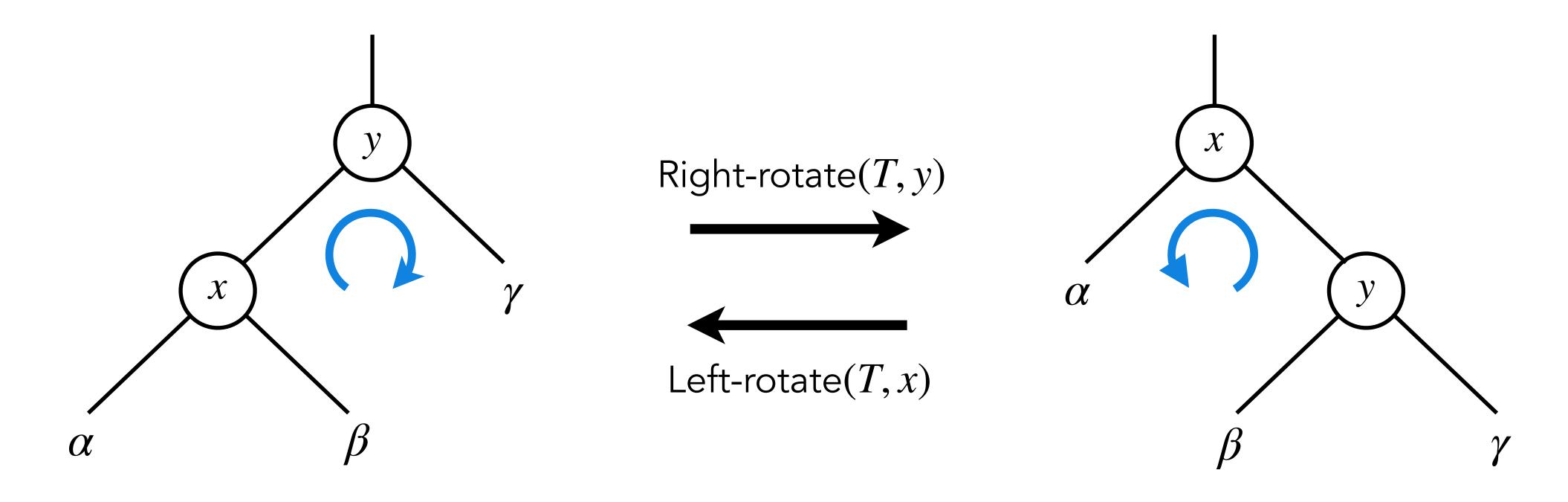


Maintaining subtree sizes during rotations is easy.



1. 
$$y$$
 .  $size =$ 

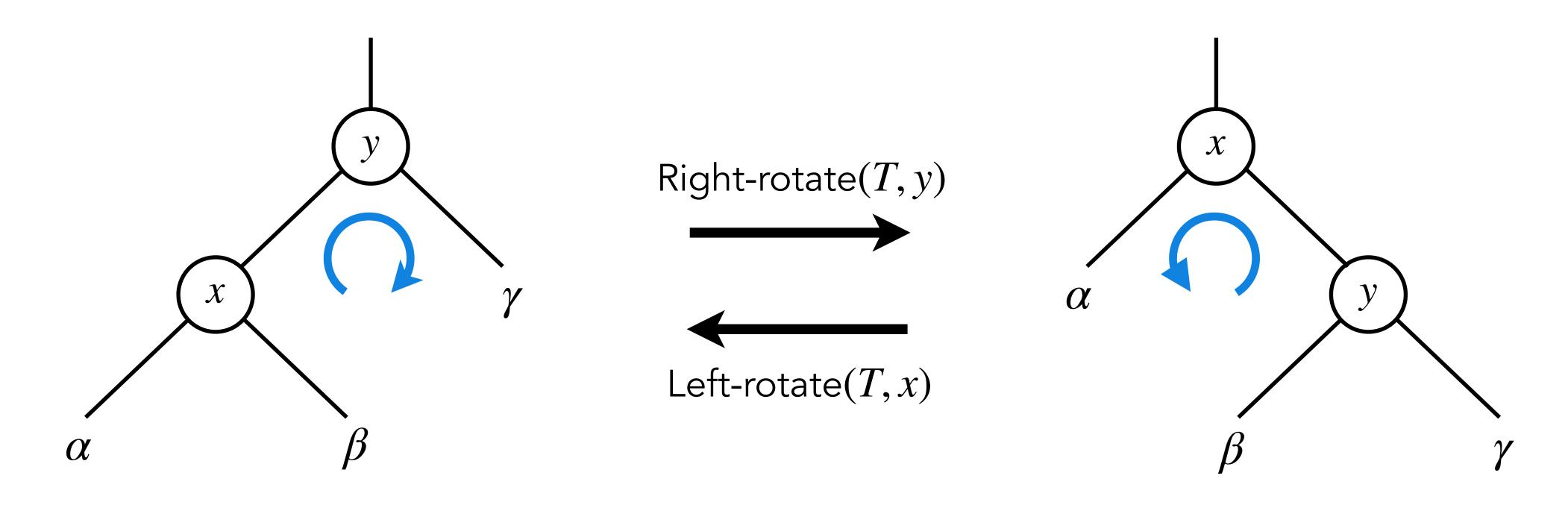
Maintaining subtree sizes during rotations is easy.



$$1. y. size =$$

$$2. x.size =$$

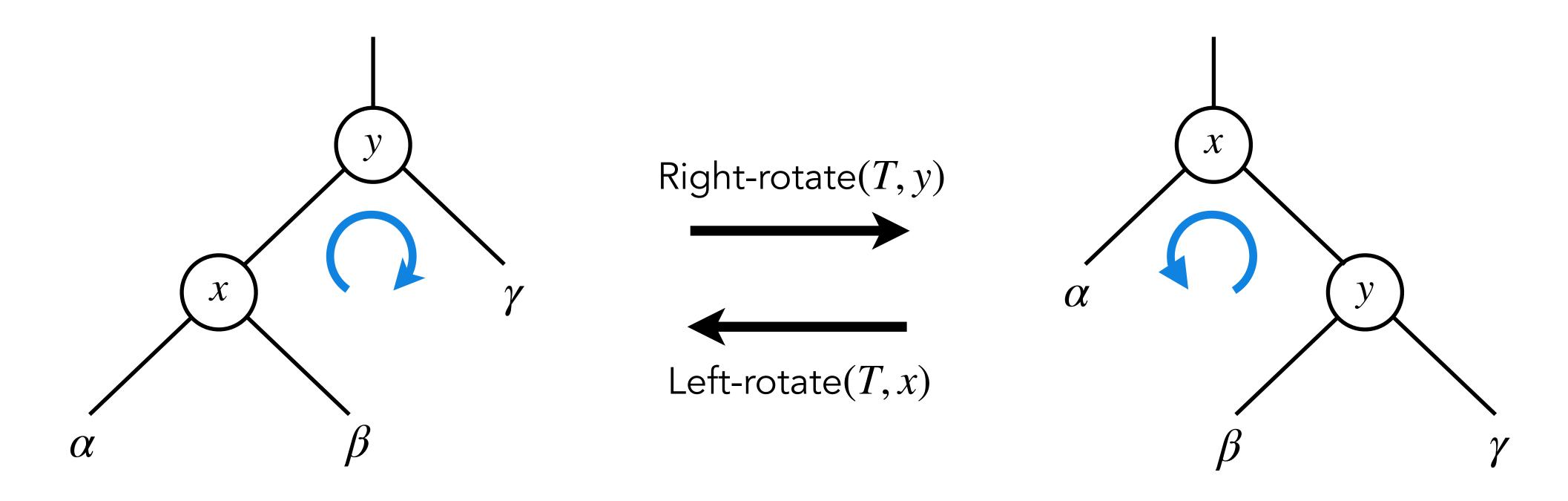
Maintaining subtree sizes during rotations is easy.



1. 
$$y.size = x.size$$

$$2. x.size =$$

Maintaining subtree sizes during rotations is easy.



- 1. y.size = x.size
- 2. x.size = x.left.size + x.right.size + 1