1 - Who is on board?

```
% Introduce all candidates.
   candidate(granger).
2
   candidate(malfoy).
   candidate(potter).
   candidate(weasley).
   \% All persons in the board are candidates.
   \verb| allCandidates(board(C, T, S))| := candidate(C), candidate(T), candidate(S).
   % Two candidates are distinct.
10
   distinct(C1, C2) :- candidate(C1), candidate(C2), C1 \= C2.
11
   \% All persons in the board are distinct.
13
   allDistinct(board(C, T, S)) := distinct(C, T), distinct(C, S), distinct(T, S).
14
15
   % Being a member of the board.
16
   partOfBoard(P, board(P, _, _)).
17
   partOfBoard(P, board(_, P, _)).
18
   partOfBoard(P, board(_, _, P)).
19
   % Being not a member of the board.
21
   notPartOfBoard(P, board(C, T, S)) :- distinct(P, C), distinct(P, T), distinct(P, S).
23
   % Either Potter or Malfoy.
24
   potterOrMalfoy(B) :- partOfBoard(potter, B), notPartOfBoard(malfoy, B).
   potterOrMalfoy(B) :- partOfBoard(malfoy, B), notPartOfBoard(potter, B).
potterOrMalfoy(B) :- notPartOfBoard(malfoy, B), notPartOfBoard(potter, B).
26
27
   % Malfoy only if Granger is chairperson.
29
   malfoyOnlyIfGranger(board(granger, T, S)) :- partOfBoard(malfoy, board(granger, T, S)).
30
                                                :- notPartOfBoard(malfoy, B).
31
   malfoyOnlyIfGranger(B)
32
   % Weasley only if Potter is part of board.
33
   weaslyOnlyIfPotter(B) :- partOfBoard(weasley, B), partOfBoard(potter, B).
34
35
   weaslyOnlyIfPotter(B) :- notPartOfBoard(weasley, B).
   \mbox{\ensuremath{\mbox{\%}}} Potter only if Granger not secretary.
37
   \verb|potterWithoutGranger(board(C, T, granger))| := \verb|notPartOfBoard(potter, board(C, T, granger))|.
38
   potterWithoutGranger(board(_, _, S))
                                                :- distinct(S, granger).
39
40
   % Granger only if Weasley not chairperson.
41
   42
   grangerWithoutWeasley(board(C, _, _))
43
   % Solve the problem.
45
46
   solve(B) :- allCandidates(B),
                allDistinct(B),
47
                potterOrMalfoy(B),
48
49
                malfoyOnlyIfGranger(B),
                weaslyOnlyIfPotter(B),
50
51
                potterWithoutGranger(B),
                grangerWithoutWeasley(B).
```

The solutions are as follows.

```
?- solve(B).
B = board(potter, granger, weasley);
B = board(granger, potter, weasley);
B = board(granger, weasley, potter);
false.
```

2 - Predicates

First we recapitulate the family example.

```
mother(david,
                      linda).
   mother(elizabeth, mary).
2
   mother(james,
                      jennifer).
   mother(jennifer,
                      patricia).
4
   mother (michael,
                     mary).
   mother(robert,
                      patricia).
   husband(jennifer, david).
   husband(linda,
                      william).
9
   husband(mary,
10
                      robert).
   husband(patricia, john).
11
12
   male(david).
13
   male(james).
14
   male(john).
15
16
   male(michael).
17
   male(robert).
   male(william).
18
19
   female(elizabeth).
20
   female(jennifer).
21
22
   female(linda).
   female(mary).
23
   female(patricia).
24
25
   father(C, F) := mother(C, M), husband(M, F).
26
27
   brother (P, B) := male(B), mother(P, M), mother(B, M), B = P.
28
29
   sister(P, S) := female(S), mother(P, M), mother(S, M), S = P.
```

Now we introduce a helper predicate for the parents of a person.

```
parent(C, P) :- mother(C, P).
parent(C, P) :- father(C, P).
```

With this helper predicate we can define son very easily.

```
son(S, P) :- male(S), parent(S, P).
```

The predicate cousin can be defined as follows.

```
cousin(C, P) :- parent(C, X), brother(X, B), parent(P, B).
cousin(C, P) :- parent(C, X), sister(X, S), parent(P, S).
```

The predicate brother_in_law can be defined as follows.

```
brother_in_law(BL, P) :- husband(P, H), brother(H, BL).
brother_in_law(BL, P) :- husband(W, P), brother(W, BL).
```

Finally, the predicate ancestor.

```
ancestor(A, P) :- parent(P, A).
ancestor(A, P) :- parent(P, X), ancestor(A, X).
```

First we recapitulate the append predicate.

```
append([], Ys, Ys).
append([X|Xs], Ys, [X|Zs]) :- append(Xs, Ys, Zs).
```

Now we can implement the predicates lookup and member2.

```
lookup(K, KVs, V) :- append(_, [(K, V)|_], KVs).
member2(X, Xs) :- append(_, [X|Ys], Xs), append(_, [X|_], Ys).
```

The predicate reverse can be implemented in a naive and an efficient way.

```
reverse([], []).
reverse([X|Xs], Ys) :- reverse(Xs, Zs), append(Zs, [X], Ys).

reverse2(Xs, Ys) :- reverseAcc(Xs, [], Ys).

reverseAcc([], Ys, Ys).
reverseAcc([X|Xs], Ys, Zs) :- reverseAcc(Xs, [X|Ys], Zs).
```

The predicate sublist can be implemented as follows.

```
sublist(Xs, Ys) :- append(Zs, _, Ys), append(_, Xs, Zs).
```

3 - Binary Numbers

We begin with addition.

```
add(o,
                             Υ,
      add(X,
                                            X).
                            ο,
      add(pos(X), pos(Y), pos(Z)) :- addP(X, Y, Z).
      addP(i,
                                    o(i)).
                          i,
                           o(Y), i(Y)).
      addP(i,
      addP(i,
                           i(Y), o(Z)) :- addP(i, Y, Z).
      addP(o(X), i,
                                     i(X)).
      addP(o(X), o(Y), o(Z)) :- addP(X, Y, Z).
      addP(o(X), i(Y), i(Z)) :- addP(X, Y, Z).
addP(i(X), i, o(Z)) :- addP(i, X, Z).
addP(i(X), o(Y), i(Z)) :- addP(X, Y, Z).
10
11
12
      \mathtt{addP}(\mathtt{i}(\mathtt{X})\,,\,\,\mathtt{i}(\mathtt{Y})\,,\,\,\mathtt{o}(\mathtt{Z}))\,\,:-\,\,\mathtt{addP}(\mathtt{i}\,,\,\,\mathtt{X}\,,\,\,\mathtt{A})\,,\,\,\mathtt{addP}(\mathtt{A}\,,\,\,\mathtt{Y}\,,\,\,\mathtt{Z})\,.
```

We define subtraction with addition.

```
sub(X, Y, Z) :- add(Y, Z, X).
```

We test addition.

```
% 0 + 2 = S
?- add(o, pos(o(i)), S).
S = pos(o(i));
false.
% 1 + 2 = S
?- add(pos(i), pos(o(i)), S).
S = pos(i(i));
false.
% X + Y = 2
?- add(X, Y, pos(o(i))).
X = 0,
Y = pos(o(i));
                   % 0 + 2 = 2
X = pos(o(i)),
                   % 2 + 0 = 2
Y = o;
X = Y, Y = pos(i); % 1 + 1 = 2
^CAction (h for help) ? abort
% Execution Aborted
```

In the last query we noticed that the search did not terminate. This is due to the fact that the Prolog system tries increasingly large numbers for X and Y, the sum of which is of course not 2. We will fix this later.

We test subtraction.

```
% 4 - 2 = D
?- sub(pos(o(o(i))), pos(o(i)), D).
D = pos(o(i));
false.
% 0 - Y = 1
?- sub(o, Y, pos(i)).
false.
% 2 - X = Y
?- sub(pos(o(i)), X, Y).
X = 0,
Y = pos(o(i));
                   % 2 - 0 = 2
X = pos(o(i)),
Y = o;
                   % 2 - 2 = 0
X = Y, Y = pos(i); % 2 - 1 = 1
^CAction (h for help) ? abort
% Execution Aborted
```

Here the same problem occurs.

To solve the problem of guessing too large numbers, we first define a predicate less(X, Y), which is fulfilled if X is smaller than Y.

```
less(o, pos(_)).
less(pos(X), pos(Y)) :- lessP(X, Y).

lessP(i, o(_)).
lessP(i, i(_)).
lessP(o(X), o(Y)) :- lessP(X, Y).
lessP(o(X), i(X)).
lessP(o(X), i(Y)) :- lessP(X, Y).
lessP(i(X), o(Y)) :- lessP(X, Y).
lessP(i(X), o(Y)) :- lessP(X, Y).
lessP(i(X), i(Y)) :- lessP(X, Y).
```

With this predicate we can specify that the sum of two numbers > 0 is greater than each summand.

```
add(o,
                       Y).
   add(X,
                       X).
               ο,
2
   add(pos(X), pos(Y), pos(Z)) := lessP(X, Z), lessP(Y, Z), addP(X, Y, Z).
   addP(i,
                   o(i)).
              o(Y), i(Y)).
   addP(i,
   addP(i,
              i(Y), o(Z)) :- addP(i, Y, Z).
   addP(o(X), i,
                   i(X)).
   addP(o(X), o(Y), o(Z)) :- addP(X, Y, Z).
   10
11
   addP(i(X), o(Y), i(Z)) :- addP(X, Y, Z).
   {\tt addP(i(X),\ i(Y),\ o(Z))\ :-\ addP(i,\ X,\ A),\ addP(A,\ Y,\ Z).}
13
```

Subtraction stays the same.

```
sub(X, Y, Z) :- add(Y, Z, X).
```

Now we can also solve those requests where the system did not terminate before.

4 - Boolean Operations

```
% (x /\ y) \/ z
% (x /\ y) \/ ((y /\ z) /\ z)
% (x /\ (¬y) /\ z) \/ ((z /\ y) \/ z)
and(true, true, true).
and(true, false, false).
and(false, _, false).
or(true, _, true).
or(false, Y, Y).
not(true, false).
not(false, true).
ex1(X, Y, Z, Res) := and(X, Y, XaY), or(XaY, Z, Res).
ex2(X, Y, Z, Res) :- and(X, Y, XaY),
                     and(Y, Z, YaZ),
                     and(YaZ, Z, YaZaZ),
                     or(XaY, YaZaZ, Res).
ex3(X, Y, Z, Res) :- not(Y, NotY),
                     and(X, NotY, XaNotY),
                     and(XaNotY, Z, XaNotYaZ),
                     and(Z, Y, ZaY),
                     or(ZaY, Z, ZaYoZ),
                     or(XaNotYaZ, ZaYoZ, Res).
\% Which results do you get for the values `X = true`, `Y = false` and `Z = true`?
% ?- ex1(true, false, true, Res).
% Res = true.
% ?- ex2(true, false, true, Res).
% Res = false.
% ?- ex3(true, false, true, Res).
% Res = true;
% false.
```

```
% Which assignments yield `true` as result?
%
% ?- ex1(X, Y, Z, true).
% X = Y, Y = true;
% X = Z, Y = false, Z = true;
% X = false, Z = true.
% ?- ex2(X, Y, Z, true).
% X = Y, Y = Z, Z = true;
% X = Y, Y = true, Z = false;
% X = false, Y = Z, Z = true;
% false.
% ?- ex3(X, Y, Z, true).
% X = Y, Y = Z, Z = true;
% X = false, Y = Z, Z = true;
% X = Z, Y = false, Z = true;
% X = Y, Y = false, Z = true;
% false.
% Is the third equation dependent on `x` or `z`?
% ?- ex3(X, Y, Z, Res).
% X = Y, Y = Z, Z = Res, Res = true;
% X = Y, Y = true, Z = Res, Res = false;
% X = false, Y = Z, Z = Res, Res = true;
% X = Z, Y = true, Z = Res, Res = false;
% X = Z, Y = false, Z = Res, Res = true;
% X = true, Y = Z, Z = Res, Res = false;
% X = Y, Y = false, Z = Res, Res = true;
% X = Y, Y = Z, Z = Res, Res = false.
% Here it becomes clear that the result ('Res') is ultimately always
\% dependent on the allocation of the variable `Z`!
\mbox{\ensuremath{\mbox{\%}}} This fact is clarified by the following queries.
% ?- ex3(X, Y, true, false).
% false.
%
% ?- ex3(X, Y, false, true).
% false.
%
\mbox{\ensuremath{\mbox{\%}}} The queries are either false
% ?- ex3(X, Y, true, true).
% X = Y, Y = true;
% X = false, Y = true;
% X = true, Y = false;
% X = Y, Y = false.
% ?- ex3(X, Y, false, false).
% X = Y, Y = true;
% X = false, Y = true;
% X = true, Y = false;
% X = Y, Y = false.
```

```
\% or hold for every value of `X` und `Y`.
\mbox{\ensuremath{\mbox{\%}}} The equation is also independent of 'X'.
% ?- ex3(true, Y, Z, false).
% Y = true, Z = false;
% Y = Z, Z = false.
% ?- ex3(false, Y, Z, false).
% Y = true, Z = false;
% Y = Z, Z = false.
\% No matter how 'X' is instantiated, the result is 'false' if 'Y' and 'Z'
\mbox{\ensuremath{\mbox{\%}}} are instantiated accordingly.
%
\% ?- ex3(false, Y, Z, true).
% Y = Z, Z = true;
% Y = false, Z = true.
% ?- ex3(true, Y, Z, true).
% Y = Z, Z = true;
% Y = false, Z = true.
%
\% The same applies for the result `true`.
```