KENNESAW STATE UNIVERSITY

ICPC Cheat Sheet

Dues Ex Machina

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Graph Theory

Dijsktra's Shortest Path Algorithm

```
// Number of vertices in the graph
#define V 9
// A utility function to find the vertex with minimum distance value, from
// the set of vertices not yet included in shortest path tree
int minDistance(int dist[], bool sptSet[])
   // Initialize min value
   int min = INT_MAX, min_index;
   for (int v = 0; v < V; v++)
     if (sptSet[v] == false && dist[v] <= min)</pre>
         min = dist[v], min_index = v;
   return min_index;
// A utility function to print the constructed distance array
int printSolution(int dist[], int n)
  printf("Vertex Distance from Source\n");
   for (int i = 0; i < V; i++)
      printf("%d \t\t %d\n", i, dist[i]);
}
// Funtion that implements Dijkstra's single source shortest path algorithm
// for a graph represented using adjacency matrix representation
void dijkstra(int graph[V][V], int src)
     int dist[V];
                      // The output array. dist[i] will hold the shortest
                      // distance from src to i
     bool sptSet[V]; // sptSet[i] will true if vertex i is included in
shortest
                     // path tree or shortest distance from src to i is
finalized
     // Initialize all distances as INFINITE and stpSet[] as false
     for (int i = 0; i < V; i++)
        dist[i] = INT_MAX, sptSet[i] = false;
     // Distance of source vertex from itself is always 0
     dist[src] = 0;
     \ensuremath{//} Find shortest path for all vertices
     for (int count = 0; count < V-1; count++)</pre>
```

```
// Pick the minimum distance vertex from the set of vertices not
       // yet processed. u is always equal to src in first iteration.
       int u = minDistance(dist, sptSet);
       // Mark the picked vertex as processed
       sptSet[u] = true;
       // Update dist value of the adjacent vertices of the picked vertex.
       for (int v = 0; v < V; v++)
         // Update dist[v] only if is not in sptSet, there is an edge from
         // u to v, and total weight of path from src to v through u is
         // smaller than current value of dist[v]
         if (!sptSet[v] && graph[u][v] && dist[u] != INT_MAX
                                       && dist[u]+graph[u][v] < dist[v])
            dist[v] = dist[u] + graph[u][v];
     }
     // print the constructed distance array
    printSolution(dist, V);
}
Dijsktra's Algorithm for Adjacency List Representation
// A structure to represent a node in adjacency list
struct AdjListNode
    int dest;
    int weight;
    struct AdjListNode* next;
};
// A structure to represent an adjacency liat
struct AdjList
    struct AdjListNode *head; // pointer to head node of list
};
// A structure to represent a graph. A graph is an array of adjacency lists.
// Size of array will be V (number of vertices in graph)
struct Graph
{
    int V;
    struct AdjList* array;
};
// A utility function to create a new adjacency list node
struct AdjListNode* newAdjListNode(int dest, int weight)
    struct AdjListNode* newNode =
            (struct AdjListNode*) malloc(sizeof(struct AdjListNode));
   newNode->dest = dest;
    newNode->weight = weight;
```

```
newNode->next = NULL;
    return newNode;
}
// A utility function that creates a graph of V vertices
struct Graph* createGraph(int V)
    struct Graph* graph = (struct Graph*) malloc(sizeof(struct Graph));
    graph->V = V;
    // Create an array of adjacency lists. Size of array will be V
    graph->array = (struct AdjList*) malloc(V * sizeof(struct AdjList));
    // Initialize each adjacency list as empty by making head as NULL
    for (int i = 0; i < V; ++i)
        graph->array[i].head = NULL;
   return graph;
}
// Adds an edge to an undirected graph
void addEdge(struct Graph* graph, int src, int dest, int weight)
    // Add an edge from src to dest. A new node is added to the adjacency
    // list of src. The node is added at the begining
    struct AdjListNode* newNode = newAdjListNode(dest, weight);
    newNode->next = graph->array[src].head;
    graph->array[src].head = newNode;
    // Since graph is undirected, add an edge from dest to src also
   newNode = newAdjListNode(src, weight);
    newNode->next = graph->array[dest].head;
    graph->array[dest].head = newNode;
}
// Structure to represent a min heap node
struct MinHeapNode
{
    int v;
    int dist;
};
// Structure to represent a min heap
struct MinHeap
{
                  // Number of heap nodes present currently
    int size;
    int capacity; // Capacity of min heap
                 // This is needed for decreaseKey()
    int *pos;
    struct MinHeapNode **array;
};
// A utility function to create a new Min Heap Node
struct MinHeapNode* newMinHeapNode(int v, int dist)
{
```

```
struct MinHeapNode* minHeapNode =
           (struct MinHeapNode*) malloc(sizeof(struct MinHeapNode));
    minHeapNode->v = v;
    minHeapNode->dist = dist;
   return minHeapNode;
}
// A utility function to create a Min Heap
struct MinHeap* createMinHeap(int capacity)
    struct MinHeap* minHeap =
         (struct MinHeap*) malloc(sizeof(struct MinHeap));
   minHeap->pos = (int *)malloc(capacity * sizeof(int));
   minHeap->size = 0;
   minHeap->capacity = capacity;
   minHeap->array =
         (struct MinHeapNode**) malloc(capacity * sizeof(struct
MinHeapNode*));
   return minHeap;
// A utility function to swap two nodes of min heap. Needed for min heapify
void swapMinHeapNode(struct MinHeapNode** a, struct MinHeapNode** b)
    struct MinHeapNode* t = *a;
    *a = *b;
    *b = t;
}
// A standard function to heapify at given idx
// This function also updates position of nodes when they are swapped.
// Position is needed for decreaseKey()
void minHeapify(struct MinHeap* minHeap, int idx)
    int smallest, left, right;
    smallest = idx;
    left = 2 * idx + 1;
    right = 2 * idx + 2;
    if (left < minHeap->size &&
        minHeap->array[left]->dist < minHeap->array[smallest]->dist )
      smallest = left;
    if (right < minHeap->size &&
        minHeap->array[right]->dist < minHeap->array[smallest]->dist )
      smallest = right;
    if (smallest != idx)
        // The nodes to be swapped in min heap
        MinHeapNode *smallestNode = minHeap->array[smallest];
        MinHeapNode *idxNode = minHeap->array[idx];
        // Swap positions
        minHeap->pos[smallestNode->v] = idx;
```

```
minHeap->pos[idxNode->v] = smallest;
        // Swap nodes
        swapMinHeapNode(&minHeap->array[smallest], &minHeap->array[idx]);
        minHeapify(minHeap, smallest);
    }
}
// A utility function to check if the given minHeap is ampty or not
int isEmpty(struct MinHeap* minHeap)
    return minHeap->size == 0;
// Standard function to extract minimum node from heap
struct MinHeapNode* extractMin(struct MinHeap* minHeap)
{
    if (isEmpty(minHeap))
        return NULL;
    // Store the root node
    struct MinHeapNode* root = minHeap->array[0];
    // Replace root node with last node
    struct MinHeapNode* lastNode = minHeap->array[minHeap->size - 1];
    minHeap->array[0] = lastNode;
    // Update position of last node
    minHeap->pos[root->v] = minHeap->size-1;
    minHeap->pos[lastNode->v] = 0;
    // Reduce heap size and heapify root
    --minHeap->size;
   minHeapify(minHeap, 0);
   return root;
}
// Function to decreasy dist value of a given vertex v. This function
// uses pos[] of min heap to get the current index of node in min heap
void decreaseKey(struct MinHeap* minHeap, int v, int dist)
    // Get the index of v in heap array
    int i = minHeap->pos[v];
    // Get the node and update its dist value
    minHeap->array[i]->dist = dist;
    // Travel up while the complete tree is not hepified.
    // This is a O(Logn) loop
   while (i && minHeap->array[i]->dist < minHeap->array[(i - 1) / 2]->dist)
        // Swap this node with its parent
```

```
minHeap->pos[minHeap->array[i]->v] = (i-1)/2;
        minHeap \rightarrow pos[minHeap \rightarrow array[(i-1)/2] \rightarrow v] = i;
        swapMinHeapNode(&minHeap->array[i], &minHeap->array[(i - 1) / 2]);
        // move to parent index
        i = (i - 1) / 2;
    }
}
// A utility function to check if a given vertex
// 'v' is in min heap or not
bool isInMinHeap(struct MinHeap *minHeap, int v)
   if (minHeap->pos[v] < minHeap->size)
     return true;
   return false;
// A utility function used to print the solution
void printArr(int dist[], int n)
    printf("Vertex Distance from Source\n");
    for (int i = 0; i < n; ++i)
        printf("%d \t\t %d\n", i, dist[i]);
}
// The main function that calulates distances of shortest paths from src to
// vertices. It is a O(ELogV) function
void dijkstra(struct Graph* graph, int src)
    int V = graph->V;// Get the number of vertices in graph
    int dist[V];
                   // dist values used to pick minimum weight edge in cut
    // minHeap represents set E
    struct MinHeap* minHeap = createMinHeap(V);
    // Initialize min heap with all vertices. dist value of all vertices
    for (int v = 0; v < V; ++v)
    {
        dist[v] = INT_MAX;
        minHeap->array[v] = newMinHeapNode(v, dist[v]);
        minHeap -> pos[v] = v;
    // Make dist value of src vertex as 0 so that it is extracted first
    minHeap->array[src] = newMinHeapNode(src, dist[src]);
    minHeap->pos[src]
                        = src;
    dist[src] = 0;
    decreaseKey(minHeap, src, dist[src]);
    // Initially size of min heap is equal to V
    minHeap->size = V;
    // In the followin loop, min heap contains all nodes
```

```
// whose shortest distance is not yet finalized.
    while (!isEmpty(minHeap))
        // Extract the vertex with minimum distance value
        struct MinHeapNode* minHeapNode = extractMin(minHeap);
        int u = minHeapNode->v; // Store the extracted vertex number
        // Traverse through all adjacent vertices of u (the extracted
        // vertex) and update their distance values
        struct AdjListNode* pCrawl = graph->array[u].head;
        while (pCrawl != NULL)
            int v = pCrawl->dest;
            // If shortest distance to v is not finalized yet, and distance
to v
            // through u is less than its previously calculated distance
            if (isInMinHeap(minHeap, v) && dist[u] != INT_MAX &&
                                           pCrawl->weight + dist[u] < dist[v])</pre>
            {
                dist[v] = dist[u] + pCrawl->weight;
                // update distance value in min heap also
                decreaseKey(minHeap, v, dist[v]);
            pCrawl = pCrawl->next;
        }
    }
    // print the calculated shortest distances
   printArr(dist, V);
```

Prim's Minimum Spanning Tree for Adjacency List Representationi

```
// A structure to represent a node in adjacency list
struct AdjListNode
{
   int dest;
   int weight;
   struct AdjListNode* next;
};

// A structure to represent an adjacency liat
struct AdjList
{
   struct AdjListNode *head; // pointer to head node of list
};

// A structure to represent a graph. A graph is an array of adjacency lists.
// Size of array will be V (number of vertices in graph)
struct Graph
{
   int V;
```

```
struct AdjList* array;
};
// A utility function to create a new adjacency list node
struct AdjListNode* newAdjListNode(int dest, int weight)
    struct AdjListNode* newNode =
            (struct AdjListNode*) malloc(sizeof(struct AdjListNode));
   newNode->dest = dest;
   newNode->weight = weight;
   newNode->next = NULL;
   return newNode;
}
// A utility function that creates a graph of V vertices
struct Graph* createGraph(int V)
    struct Graph* graph = (struct Graph*) malloc(sizeof(struct Graph));
    graph->V = V;
    // Create an array of adjacency lists. Size of array will be V
    graph->array = (struct AdjList*) malloc(V * sizeof(struct AdjList));
     // Initialize each adjacency list as empty by making head as NULL
    for (int i = 0; i < V; ++i)
        graph->array[i].head = NULL;
   return graph;
}
// Adds an edge to an undirected graph
void addEdge(struct Graph* graph, int src, int dest, int weight)
    // Add an edge from src to dest. A new node is added to the adjacency
    // list of src. The node is added at the begining
    struct AdjListNode* newNode = newAdjListNode(dest, weight);
    newNode->next = graph->array[src].head;
    graph->array[src].head = newNode;
    // Since graph is undirected, add an edge from dest to src also
   newNode = newAdjListNode(src, weight);
   newNode->next = graph->array[dest].head;
   graph->array[dest].head = newNode;
// Structure to represent a min heap node
struct MinHeapNode
    int v;
    int key;
};
// Structure to represent a min heap
struct MinHeap
```

```
{
                 // Number of heap nodes present currently
    int capacity; // Capacity of min heap
    int *pos;
                 // This is needed for decreaseKey()
    struct MinHeapNode **array;
};
// A utility function to create a new Min Heap Node
struct MinHeapNode* newMinHeapNode(int v, int key)
    struct MinHeapNode* minHeapNode =
           (struct MinHeapNode*) malloc(sizeof(struct MinHeapNode));
   minHeapNode->v = v;
   minHeapNode->key = key;
   return minHeapNode;
}
// A utilit function to create a Min Heap
struct MinHeap* createMinHeap(int capacity)
{
    struct MinHeap* minHeap =
         (struct MinHeap*) malloc(sizeof(struct MinHeap));
   minHeap->pos = (int *)malloc(capacity * sizeof(int));
    minHeap->size = 0;
   minHeap->capacity = capacity;
   minHeap->array =
         (struct MinHeapNode**) malloc(capacity * sizeof(struct
MinHeapNode*));
   return minHeap;
// A utility function to swap two nodes of min heap. Needed for min heapify
void swapMinHeapNode(struct MinHeapNode** a, struct MinHeapNode** b)
    struct MinHeapNode* t = *a;
    *a = *b;
    *b = t;
}
// A standard function to heapify at given idx
// This function also updates position of nodes when they are swapped.
// Position is needed for decreaseKev()
void minHeapify(struct MinHeap* minHeap, int idx)
    int smallest, left, right;
    smallest = idx;
    left = 2 * idx + 1;
    right = 2 * idx + 2;
    if (left < minHeap->size &&
        minHeap->array[left]->key < minHeap->array[smallest]->key )
      smallest = left;
    if (right < minHeap->size &&
        minHeap->array[right]->key < minHeap->array[smallest]->key )
```

```
smallest = right;
    if (smallest != idx)
        // The nodes to be swapped in min heap
        MinHeapNode *smallestNode = minHeap->array[smallest];
        MinHeapNode *idxNode = minHeap->array[idx];
        // Swap positions
        minHeap->pos[smallestNode->v] = idx;
        minHeap->pos[idxNode->v] = smallest;
        // Swap nodes
        swapMinHeapNode(&minHeap->array[smallest], &minHeap->array[idx]);
        minHeapify(minHeap, smallest);
}
// A utility function to check if the given minHeap is ampty or not
int isEmpty(struct MinHeap* minHeap)
   return minHeap->size == 0;
}
// Standard function to extract minimum node from heap
struct MinHeapNode* extractMin(struct MinHeap* minHeap)
    if (isEmpty(minHeap))
        return NULL;
    // Store the root node
    struct MinHeapNode* root = minHeap->array[0];
    // Replace root node with last node
    struct MinHeapNode* lastNode = minHeap->array[minHeap->size - 1];
    minHeap->array[0] = lastNode;
    // Update position of last node
    minHeap->pos[root->v] = minHeap->size-1;
    minHeap->pos[lastNode->v] = 0;
    // Reduce heap size and heapify root
    --minHeap->size;
   minHeapify(minHeap, 0);
   return root;
}
// Function to decreasy key value of a given vertex v. This function
// uses pos[] of min heap to get the current index of node in min heap
void decreaseKey(struct MinHeap* minHeap, int v, int key)
    // Get the index of v in heap array
```

```
int i = minHeap->pos[v];
    // Get the node and update its key value
    minHeap->array[i]->key = key;
    // Travel up while the complete tree is not hepified.
    // This is a O(Logn) loop
    while (i && minHeap->array[i]->key < minHeap->array[(i - 1) / 2]->key)
        // Swap this node with its parent
        minHeap->pos[minHeap->array[i]->v] = (i-1)/2;
        minHeap \rightarrow pos[minHeap \rightarrow array[(i-1)/2] \rightarrow v] = i;
        swapMinHeapNode(&minHeap->array[i], &minHeap->array[(i - 1) / 2]);
        // move to parent index
        i = (i - 1) / 2;
    }
}
// A utility function to check if a given vertex
// 'v' is in min heap or not
bool isInMinHeap(struct MinHeap *minHeap, int v)
   if (minHeap->pos[v] < minHeap->size)
     return true;
  return false;
}
// A utility function used to print the constructed MST
void printArr(int arr[], int n)
{
    for (int i = 1; i < n; ++i)
       printf("%d - %d\n", arr[i], i);
// The main function that constructs Minimum Spanning Tree (MST)
// using Prim's algorithm
void PrimMST(struct Graph* graph)
    int V = graph->V;// Get the number of vertices in graph
    int key[V];
                    // Key values used to pick minimum weight edge in cut
    // minHeap represents set E
    struct MinHeap* minHeap = createMinHeap(V);
    // Initialize min heap with all vertices. Key value of
    // all vertices (except 0th vertex) is initially infinite
    for (int v = 1; v < V; ++v)
        parent[v] = -1;
        key[v] = INT_MAX;
        minHeap->array[v] = newMinHeapNode(v, key[v]);
        minHeap->pos[v] = v;
    }
```

```
// is extracted first
   key[0] = 0;
   minHeap->array[0] = newMinHeapNode(0, key[0]);
    minHeap->pos[0]
                     = 0;
    // Initially size of min heap is equal to V
   minHeap->size = V;
    // In the followin loop, min heap contains all nodes
    // not yet added to MST.
   while (!isEmpty(minHeap))
        // Extract the vertex with minimum key value
        struct MinHeapNode* minHeapNode = extractMin(minHeap);
        int u = minHeapNode->v; // Store the extracted vertex number
        // Traverse through all adjacent vertices of u (the extracted
        // vertex) and update their key values
        struct AdjListNode* pCrawl = graph->array[u].head;
        while (pCrawl != NULL)
            int v = pCrawl->dest;
            // If v is not yet included in MST and weight of u-v is
            // less than key value of v, then update key value and
            if (isInMinHeap(minHeap, v) && pCrawl->weight < key[v])</pre>
                key[v] = pCrawl->weight;
                parent[v] = u;
                decreaseKey(minHeap, v, key[v]);
            pCrawl = pCrawl->next;
    }
    // print edges of MST
   printArr(parent, V);
}
Prim's Minimum Spanning Tree Algorithm
// Number of vertices in the graph
#define V 5
// A utility function to find the vertex with minimum key value, from
// the set of vertices not yet included in MST
int minKey(int key[], bool mstSet[])
   // Initialize min value
   int min = INT_MAX, min_index;
```

// Make key value of 0th vertex as 0 so that it

```
for (int v = 0; v < V; v++)
     if (mstSet[v] == false && key[v] < min)</pre>
         min = key[v], min_index = v;
  return min_index;
}
// A utility function to print the constructed MST stored in parent[]
int printMST(int parent[], int n, int graph[V][V])
  printf("Edge Weight\n");
   for (int i = 1; i < V; i++)
      printf("%d - %d
                      %d \n", parent[i], i, graph[i][parent[i]]);
// Function to construct and print MST for a graph represented using
adiacency
// matrix representation
void primMST(int graph[V][V])
     int parent[V]; // Array to store constructed MST
     int key[V]; // Key values used to pick minimum weight edge in cut
     bool mstSet[V]; // To represent set of vertices not yet included in MST
     // Initialize all keys as INFINITE
     for (int i = 0; i < V; i++)
        key[i] = INT_MAX, mstSet[i] = false;
     // Always include first 1st vertex in MST.
     key[0] = 0;
                 // Make key 0 so that this vertex is picked as first
vertex
     parent[0] = -1; // First node is always root of MST
     // The MST will have V vertices
     for (int count = 0; count < V-1; count++)</pre>
        // Pick thd minimum key vertex from the set of vertices
        // not yet included in MST
        int u = minKey(key, mstSet);
        // Add the picked vertex to the MST Set
        mstSet[u] = true;
        // Update key value and parent index of the adjacent vertices of
        // the picked vertex. Consider only those vertices which are not yet
        // included in MST
        for (int v = 0; v < V; v++)
           // graph[u][v] is non zero only for adjacent vertices of m
           // mstSet[v] is false for vertices not yet included in MST
           // Update the key only if graph[u][v] is smaller than key[v]
          if (graph[u][v] && mstSet[v] == false && graph[u][v] < key[v])</pre>
             parent[v] = u, key[v] = graph[u][v];
     }
```

```
// print the constructed MST
printMST(parent, V, graph);
}
```

Kruskal's Minimum Spanning Tree Algorithm

```
// a structure to represent a weighted edge in graph
struct Edge
    int src, dest, weight;
};
// a structure to represent a connected, undirected and weighted graph
struct Graph
    // V-> Number of vertices, E-> Number of edges
    int V, E;
    // graph is represented as an array of edges. Since the graph is
    // undirected, the edge from src to dest is also edge from dest
    // to src. Both are counted as 1 edge here.
    struct Edge* edge;
};
// Creates a graph with V vertices and E edges
struct Graph* createGraph(int V, int E)
    struct Graph* graph = (struct Graph*) malloc( sizeof(struct Graph) );
    graph->V = V;
    graph->E = E;
    graph->edge = (struct Edge*) malloc( graph->E * sizeof( struct Edge ) );
    return graph;
}
// A structure to represent a subset for union-find
struct subset
    int parent;
    int rank;
};
// A utility function to find set of an element i
// (uses path compression technique)
int find(struct subset subsets[], int i)
    // find root and make root as parent of i (path compression)
    if (subsets[i].parent != i)
        subsets[i].parent = find(subsets, subsets[i].parent);
    return subsets[i].parent;
```

```
}
// A function that does union of two sets of x and y
// (uses union by rank)
void Union(struct subset subsets[], int x, int y)
    int xroot = find(subsets, x);
    int yroot = find(subsets, y);
    // Attach smaller rank tree under root of high rank tree
    // (Union by Rank)
    if (subsets[xroot].rank < subsets[yroot].rank)</pre>
        subsets[xroot].parent = yroot;
    else if (subsets[xroot].rank > subsets[yroot].rank)
        subsets[yroot].parent = xroot;
    // If ranks are same, then make one as root and increment
    // its rank by one
    else
    {
        subsets[yroot].parent = xroot;
        subsets[xroot].rank++;
    }
}
// Compare two edges according to their weights.
// Used in qsort() for sorting an array of edges
int myComp(const void* a, const void* b)
    struct Edge* a1 = (struct Edge*)a;
    struct Edge* b1 = (struct Edge*)b;
    return al->weight > bl->weight;
}
// The main function to construct MST using Kruskal's algorithm
void KruskalMST(struct Graph* graph)
    int V = graph->V;
    struct Edge result[V]; // This will store the resultant MST
    int e = 0; // An index variable, used for result[]
    int i = 0; // An index variable, used for sorted edges
    // Step 1: Sort all the edges in non-decreasing order of their weight
    // If we are not allowed to change the given graph, we can create a copy
of
    // array of edges
    qsort(graph->edge, graph->E, sizeof(graph->edge[0]), myComp);
    // Allocate memory for creating V ssubsets
    struct subset *subsets =
        (struct subset*) malloc( V * sizeof(struct subset) );
    // Create V subsets with single elements
    for (int v = 0; v < V; ++v)
```

```
subsets[v].parent = v;
    subsets[v].rank = 0;
// Number of edges to be taken is equal to V-1
while (e < V - 1)
    // Step 2: Pick the smallest edge. And increment the index
    // for next iteration
    struct Edge next_edge = graph->edge[i++];
    int x = find(subsets, next_edge.src);
    int y = find(subsets, next_edge.dest);
    // If including this edge does't cause cycle, include it
    // in result and increment the index of result for next edge
    if(x != y)
        result[e++] = next_edge;
        Union(subsets, x, y);
    // Else discard the next_edge
}
// print the contents of result[] to display the built MST
printf("Following are the edges in the constructed MST\n");
for (i = 0; i < e; ++i)
    printf("%d -- %d == %d\n", result[i].src, result[i].dest,
                                               result[i].weight);
return;
```

Dynamic Programming

Longest Increasing Subsequence

```
/* lis() returns the length of the longest increasing subsequence in
    arr[] of size n */
int lis( int arr[], int n )
{
    int *lis, i, j, max = 0;
    lis = (int*) malloc( sizeof( int ) * n );

    /* Initialize LIS values for all indexes */
    for ( i = 0; i < n; i++ )
        lis[i] = 1;

    /* Compute optimized LIS values in bottom up manner */
    for ( i = 1; i < n; i++ )</pre>
```

```
for ( j = 0; j < i; j++ )
        if ( arr[i] > arr[j] && lis[i] < lis[j] + 1)
        lis[i] = lis[j] + 1;

/* Pick maximum of all LIS values */
for ( i = 0; i < n; i++ )
    if ( max < lis[i] )
        max = lis[i];

/* Free memory to avoid memory leak */
free( lis );

return max;
}</pre>
```

Longest Common Subsequence

```
int lcs( char *X, char *Y, int m, int n )
{
   if (m == 0 || n == 0)
     return 0;
   if (X[m-1] == Y[n-1])
     return 1 + lcs(X, Y, m-1, n-1);
   else
     return max(lcs(X, Y, m, n-1), lcs(X, Y, m-1, n));
}
```

Minimum Cost Path

```
int minCost(int cost[R][C], int m, int n)
{
    int i, j;

    // Instead of following line, we can use int tc[m+1][n+1] or
    // dynamically allocate memoery to save space. The following line is
    // used to keep te program simple and make it working on all compilers.
    int tc[R][C];

    tc[0][0] = cost[0][0];

/* Initialize first column of total cost(tc) array */
    for (i = 1; i <= m; i++)
        tc[i][0] = tc[i-1][0] + cost[i][0];

/* Initialize first row of tc array */
    for (j = 1; j <= n; j++)
        tc[0][j] = tc[0][j-1] + cost[0][j];

/* Construct rest of the tc array */</pre>
```

Matrix Chain Multiplication

```
// Matrix Ai has dimension p[i-1] x p[i] for i = 1..n
int MatrixChainOrder(int p[], int n)
    /* For simplicity of the program, one extra row and one extra column are
       allocated in m[][]. Oth row and Oth column of m[][] are not used */
    int m[n][n];
    int i, j, k, L, q;
    /* m[i,j] = Minimum number of scalar multiplications needed to compute
       the matrix A[i]A[i+1]...A[j] = A[i...j] where dimention of A[i] is
       p[i-1] x p[i] */
    // cost is zero when multiplying one matrix.
    for (i = 1; i < n; i++)
        m[i][i] = 0;
    // L is chain length.
    for (L=2; L<n; L++)
    {
        for (i=1; i<=n-L+1; i++)
            j = i+L-1;
            m[i][j] = INT_MAX;
            for (k=i; k<=j-1; k++)
                // q = cost/scalar multiplications
                q = m[i][k] + m[k+1][j] + p[i-1]*p[k]*p[j];
                if (q < m[i][j])
                    m[i][j] = q;
   return m[1][n-1];
```

Knapsack Problem

```
// Returns the maximum value that can be put in a knapsack of capacity W
int knapSack(int W, int wt[], int val[], int n)
{
   int i, w;
   int K[n+1][W+1];

   // Build table K[][] in bottom up manner
   for (i = 0; i <= n; i++)
   {
      for (w = 0; w <= W; w++)
      {
        if (i==0 || w==0)
            K[i][w] = 0;
        else if (wt[i-1] <= w)
            K[i][w] = max(val[i-1] + K[i-1][w-wt[i-1]], K[i-1][w]);
      else
            K[i][w] = K[i-1][w];
   }
}
return K[n][W];
}</pre>
```

Longest Palindromic Subsequence

```
// Returns the length of the longest palindromic subsequence in seq
int lps(char *str)
{
  int n = strlen(str);
  int i, j, cl;
  int L[n][n]; // Create a table to store results of subproblems

// Strings of length 1 are palindrome of lentgh 1
  for (i = 0; i < n; i++)
    L[i][i] = 1;</pre>
```

```
// Build the table. Note that the lower diagonal values of table are
    // useless and not filled in the process. The values are filled in a
    // manner similar to Matrix Chain Multiplication DP solution (See
    // http://www.geeksforgeeks.org/archives/15553). cl is length of
    // substring
    for (cl=2; cl<=n; cl++)
    {
        for (i=0; i<n-cl+1; i++)
            j = i+cl-1;
            if (str[i] == str[j] && cl == 2)
               L[i][j] = 2;
            else if (str[i] == str[j])
               L[i][j] = L[i+1][j-1] + 2;
            else
               L[i][j] = max(L[i][j-1], L[i+1][j]);
        }
    }
   return L[0][n-1];
}
```

Maximum Sum Increasing Subsequence

```
/* maxSumIS() returns the maximum sum of increasing subsequence in arr[] of
    size n */
int maxSumIS( int arr[], int n )
{
    int *msis, i, j, max = 0;
    msis = (int*) malloc( sizeof( int ) * n );

    /* Initialize msis values for all indexes */
```

```
for ( i = 0; i < n; i++ )
    msis[i] = arr[i];

/* Compute maximum sum values in bottom up manner */
for ( i = 1; i < n; i++ )
    for ( j = 0; j < i; j++ )
        if ( arr[i] > arr[j] && msis[i] < msis[j] + arr[i])
        msis[i] = msis[j] + arr[i];

/* Pick maximum of all msis values */
for ( i = 0; i < n; i++ )
    if ( max < msis[i] )
        max = msis[i];

/* Free memory to avoid memory leak */
free( msis );

return max;
}</pre>
```

Ugly Numbers

```
/* Function to get the nth ugly number*/
unsigned getNthUglyNo(unsigned n)
   unsigned *ugly =
             (unsigned *)(malloc(sizeof(unsigned)*n));
    unsigned i2 = 0, i3 = 0, i5 = 0;
    unsigned i;
    unsigned next_multiple_of_2 = 2;
    unsigned next_multiple_of_3 = 3;
    unsigned next_multiple_of_5 = 5;
    unsigned next_ugly_no = 1;
    *(ugly+0) = 1;
    for(i=1; i<n; i++)
       next_ugly_no = min(next_multiple_of_2,
                           next_multiple_of_3,
                           next_multiple_of_5);
       *(ugly+i) = next_ugly_no;
       if(next_ugly_no == next_multiple_of_2)
           i2 = i2+1;
           next_multiple_of_2 = *(ugly+i2)*2;
       if(next_ugly_no == next_multiple_of_3)
           i3 = i3+1;
           next_multiple_of_3 = *(ugly+i3)*3;
       if(next_ugly_no == next_multiple_of_5)
```

Bellman-Ford Algorithm

```
struct Edge
    int src, dest, weight;
};
// a structure to represent a connected, directed and weighted graph
struct Graph
    // V-> Number of vertices, E-> Number of edges
    int V, E;
    // graph is represented as an array of edges.
    struct Edge* edge;
};
// Creates a graph with V vertices and E edges
struct Graph* createGraph(int V, int E)
    struct Graph* graph = (struct Graph*) malloc( sizeof(struct Graph) );
   graph->V = V;
   graph->E = E;
   graph->edge = (struct Edge*) malloc( graph->E * sizeof( struct Edge ) );
   return graph;
}
// A utility function used to print the solution
void printArr(int dist[], int n)
   printf("Vertex Distance from Source\n");
    for (int i = 0; i < n; ++i)
        printf("%d \t\t %d\n", i, dist[i]);
// The main function that finds shortest distances from src to all other
// vertices using Bellman-Ford algorithm. The function also detects negative
// weight cycle
void BellmanFord(struct Graph* graph, int src)
    int V = graph->V;
    int E = graph->E;
    int dist[V];
```

```
// Step 1: Initialize distances from src to all other vertices as
INFINITE
    for (int i = 0; i < V; i++)
        dist[i] = INT_MAX;
    dist[src] = 0;
    // Step 2: Relax all edges |V| - 1 times. A simple shortest path from src
    // to any other vertex can have at-most |V| - 1 edges
    for (int i = 1; i \le V-1; i++)
    {
        for (int j = 0; j < E; j++)
            int u = graph->edge[j].src;
            int v = graph->edge[j].dest;
            int weight = graph->edge[j].weight;
            if(dist[u] + weight < dist[v])
                dist[v] = dist[u] + weight;
    }
    // Step 3: check for negative-weight cycles. The above step guarantees
    // shortest distances if graph doesn't contain negative weight cycle.
    // If we get a shorter path, then there is a cycle.
    for (int i = 0; i < E; i++)
        int u = graph->edge[i].src;
        int v = graph->edge[i].dest;
        int weight = graph->edge[i].weight;
        if (dist[u] + weight < dist[v])</pre>
            printf("Graph contains negative weight cycle");
    printArr(dist, V);
    return;
}
```

Floyd-Warshall Algorithm

```
// Number of vertices in the graph
#define V 4

/* Define Infinite as a large enough value. This value will be used
  for vertices not connected to each other */
#define INF 99999

// A function to print the solution matrix
void printSolution(int dist[][V]);

// Solves the all-pairs shortest path problem using Floyd Warshall algorithm
void floydWarshell (int graph[][V])
```

```
{
    /* dist[][] will be the output matrix that will finally have the shortest
      distances between every pair of vertices */
    int dist[V][V], i, j, k;
    /* Initialize the solution matrix same as input graph matrix. Or
       we can say the initial values of shortest distances are based
       on shortest paths considering no intermediate vertex. */
    for (i = 0; i < V; i++)
        for (j = 0; j < V; j++)
            dist[i][j] = graph[i][j];
    /* Add all vertices one by one to the set of intermediate vertices.
      ---> Before start of a iteration, we have shortest distances between
all
      pairs of vertices such that the shortest distances consider only the
      vertices in set \{0, 1, 2, ... k-1\} as intermediate vertices.
      ----> After the end of a iteration, vertex no. k is added to the set of
      intermediate vertices and the set becomes \{0, 1, 2, ... k\} */
    for (k = 0; k < V; k++)
        // Pick all vertices as source one by one
        for (i = 0; i < V; i++)
            // Pick all vertices as destination for the
            // above picked source
            for (j = 0; j < V; j++)
                // If vertex k is on the shortest path from
                // i to j, then update the value of dist[i][j]
                if (dist[i][k] + dist[k][j] < dist[i][j])
                    dist[i][j] = dist[i][k] + dist[k][j];
            }
        }
    }
    // Print the shortest distance matrix
   printSolution(dist);
}
```

Binary Search Tree

```
/* A Dynamic Programming based function that calculates minimum cost of
   a Binary Search Tree. */
int optimalSearchTree(int keys[], int freq[], int n)
{
    /* Create an auxiliary 2D matrix to store results of subproblems */
    int cost[n][n];

    /* cost[i][j] = Optimal cost of binary search tree that can be
        formed from keys[i] to keys[j].
        cost[0][n-1] will store the resultant cost */
```

```
// For a single key, cost is equal to frequency of the key
    for (int i = 0; i < n; i++)
        cost[i][i] = freq[i];
    // Now we need to consider chains of length 2, 3, \dots .
    // L is chain length.
    for (int L=2; L<=n; L++)</pre>
        // i is row number in cost[][]
        for (int i=0; i<=n-L+1; i++)
            // Get column number j from row number i and chain length L
            int j = i+L-1;
            cost[i][j] = INT_MAX;
            // Try making all keys in interval keys[i..j] as root
            for (int r=i; r<=j; r++)
               // c = cost when keys[r] becomes root of this subtree
               int c = ((r > i)? cost[i][r-1]:0) +
                        ((r < j)? cost[r+1][j]:0) +
                       sum(freq, i, j);
               if (c < cost[i][j])</pre>
                  cost[i][j] = c;
   return cost[0][n-1];
}
```