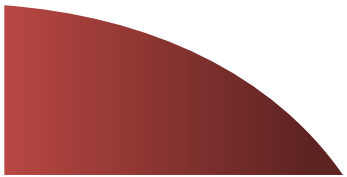


Closure of attribute sets





What is a closure of attribute sets?

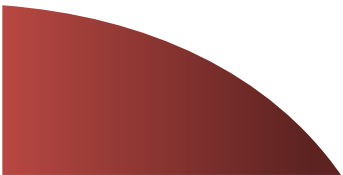
- Given a set of attributes α , the closure of α under F is the **set of attributes that are functionally determined by α under F** .
- It is denoted by α^+ .
- **Use of Closure of Attributes:**
 - To identify the additional FDs.
 - To identify the keys
 - To identify the equivalence of FD
 - To identify the standard form, canonical form or irreducible set of FD.

Algorithm

→ Algorithm to compute α^+ , the closure of α under F

→ Steps

1. $\text{result} = \alpha$
2. *while* (changes to result) *do*
 - for each $\beta \rightarrow \gamma$ in F *do*
 - *begin*
 - if $\beta \subseteq \text{result}$ then $\text{result} = \text{result} \cup \gamma$
 - else $\text{result} = \text{result}$
 - *end*





Closure of attribute sets [Example]

- Consider the relation schema $R = (A, B, C, G, H, I)$.

For this relation, a set of functional dependencies F can be given as

$$F = \{A \rightarrow B, A \rightarrow C, CG \rightarrow H, CG \rightarrow I, B \rightarrow H\}$$

Find out the closure of A^+ , B^+ , C^+ , $(AG)^+$.

$$A^+ = \{ABCH\}$$

$$B^+ = \{BH\}$$

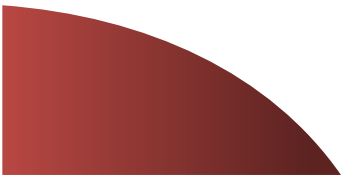
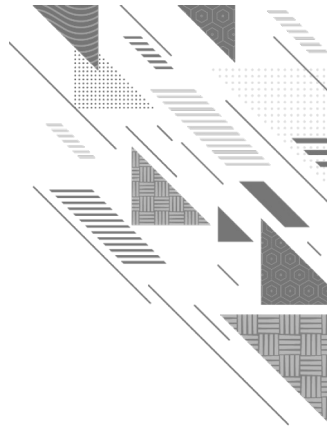
$$C^+ = \{C\}$$

→ Step 1.

$$\text{result} = a \Rightarrow \text{result} = AG$$

$A \rightarrow B$	$A \subseteq AG$	result = ABG
$A \rightarrow C$	$A \subseteq ABG$	result = ABCG
$CG \rightarrow H$	$CG \subseteq ABCG$	result = ABCGH
$CG \rightarrow I$	$CG \subseteq ABCGH$	result = ABCGHI
$B \rightarrow H$	$B \subseteq ABCGHI$	result = ABCGHI

$$AG^+ = \{ABCGHI\}$$





Closure of attribute sets [Exercise]

- Consider the relation schema $R = (A, B, C, D, E)$.

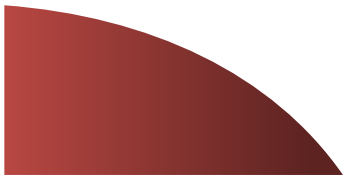
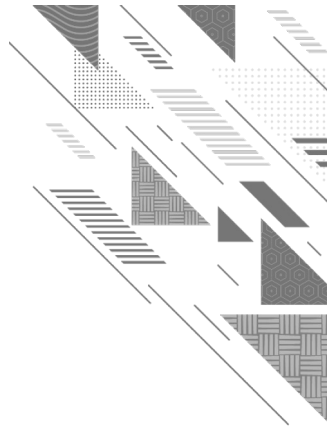
For this relation, a set of functional dependencies F can be given as

$$F = \{A \rightarrow B, B \rightarrow D, C \rightarrow DE, CD \rightarrow AB\}$$

Find out the closure of A^+ , B^+ , C^+ , D^+ , E^+ , $(CD)^+$, $(AD)^+$.

$$CK = \{C\}$$

A^+	$\{A, B, D\}$
B^+	$\{B, D\}$
C^+	$\{C, D, E, A, B\}$
D^+	$\{D\}$
E^+	$\{E\}$
$(CD)^+$	$\{A, B, C, D, E\}$
$(AD)^+$	$\{A, D, B\}$





Closure of attribute sets [Exercise]

Given functional dependencies (FDs) for relational schema $R = (A, B, C, D, E)$:

$F = \{A \rightarrow BC, CD \rightarrow E, B \rightarrow D, E \rightarrow A\}$

Find Closure for A

Find Closure for CD

Find Closure for B

Find Closure for BC

Find Closure for E

Answer

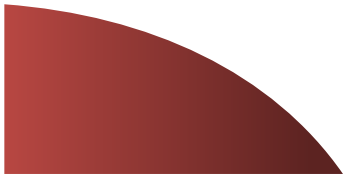
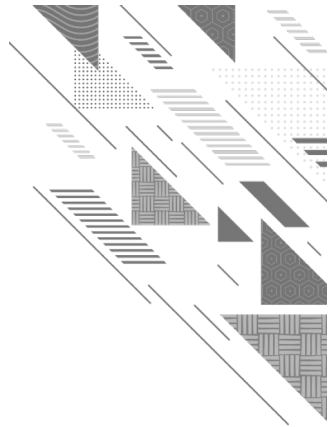
$A^+ = ABCDE$

$CD^+ = ABCDE$

$B^+ = BD$

$BC^+ = ABCDE$

$E^+ = ABCDE$





Exercise:

Q1.) Given $R(A, B, C, D, E)$ and $F = \{A \rightarrow BC, CD \rightarrow E, B \rightarrow D, E \rightarrow A\}$, **then, $A^+ = ?$**

$$A^+ = \{A, B, C, D, E\}$$

Q2.) Given $R(A, B, C, D, E, F)$ and $F = \{AB \rightarrow C, BC \rightarrow AD, D \rightarrow E, CF \rightarrow B\}$, **then, $AB \rightarrow F$ is a member of FD Set ?**

