



Identifying Minimal Set of FDS or
Standard Form of FDs or
Irreducible set of FDs or
Canonical Cover

Standard Form of FDs or Canonical Cover

A canonical cover of F is a minimal set of functional dependencies equivalent to F, having no redundant dependencies or redundant parts of dependencies.

- It is denoted by F_c
- A set of FDs F is minimal if it satisfy the following conditions:
 - F and F_c must be equivalent (i.e., F logically implies all dependencies in F_c and vice-versa)
 - Every FD of F_c is simple (i.e., in the RHS there will be single attribute in every FD.



- Fc is left reduced (set of FDs, where no extraneous left attribute present).
- Fc is non-redundant.

NOTE: Let F be set of FDs over relation R and let A \rightarrow B, a FD in F, is said to be redundant, if and only if it will be equivalent to $F = F - \{A \rightarrow B\}$

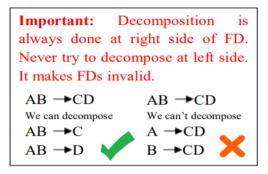


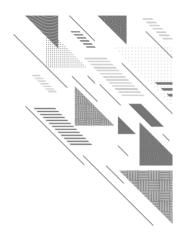
Step-I: Decompose all the FDs such that R.H.S contain only single attribute.



$$F = \{A \rightarrow BC\}$$

$$F_c = \{A \rightarrow B, A \rightarrow C\}$$





Step-II: Find the redundant FDs and delete them from FD set.

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F = \{A \rightarrow B, B \rightarrow C, A \rightarrow C\}
A^{+} \text{ (take } A^{+} \text{ without considering } A \rightarrow B) = \{A, C\},
It \text{ means } A \rightarrow B \text{ is must, we can't get B from any other FD.}
B^{+} = \{B\}, \text{ it means } B \rightarrow C \text{ is also must.}
A^{+} = \{A, B, C\}, \text{ thus, } A \rightarrow C \text{ is redundant.}
F = \{A \rightarrow B, B \rightarrow C\}
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Note: Once the non essential FD is identified, then do not include that non essential FD while computing the closure of attributes further. Exclude that FD immediately.

NOTE: You and directly remove all the Trivial FDs, if exists.



Step-III: Find redundant/extraneous attribute on L.H.S and delete them.

- If we are able to remove an attribute from a functional dependency without changing the closure of the set of functional dependencies, that attribute is called as Extraneous Attribute. It is possible when there are more than one attribute on left hand side.

$$F = \{AB \rightarrow C\}$$

This FD has more than one attribute on the Left side, it may contain extra attribute at left side.

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Q.1 F = $\{A \rightarrow C, AC \rightarrow D, E \rightarrow AD, E \rightarrow H\}$

Find the canonical cover for the given set of FDs over F.

Step-I: Decompose all the FDs to get single attribute on RHS.

$$A \rightarrow C$$

$$AC \rightarrow D$$

$$\mathsf{E} \to \mathsf{A}$$

$$\mathsf{E}\to\mathsf{D}$$

$$E \rightarrow H$$

Step-II: Remove redundant FDs by checking the effect of their non-existence one by one.

$$A \rightarrow C$$
 A⁺ = {A}. thus, A \rightarrow C is must.

$$AC \rightarrow D$$
 AC⁺ = {A, C}, thus, AC \rightarrow D is must

$$E \rightarrow A$$
 $E^+ = \{E, D, H\}$, thus $E \rightarrow A$ is must.

$$E \rightarrow D$$
 $E^+ = \{E, A, H, C, D\}$, thus $E \rightarrow D$ is redundant.

$$E \rightarrow H$$
 $E^+ = \{E, A, C, D\}$, thus $E \rightarrow H$ is must.

$$A \rightarrow C$$

$$AC \rightarrow D$$

$$E \rightarrow A$$

$$E \rightarrow H$$

Step-III: Remove extraneous attribute in FDs having more than one attribute on LHS.

$$A \rightarrow C$$

$$AC \rightarrow D$$

$$E \rightarrow A$$

$$E \rightarrow H$$

In AC -> D, C is extraneous since A+ = AC

$$F_C = \{ A \rightarrow C, A \rightarrow D, E \rightarrow A, E \rightarrow H \}$$

$$F_C = \{ A \rightarrow CD, E \rightarrow AH \}$$



Q.2 F = {A \rightarrow B, C \rightarrow B, D \rightarrow ABC, AC \rightarrow D}.

Find the canonical cover for the given set of FDs over F.

Step-I: $A \rightarrow B$

 $C \rightarrow B$

 $D \rightarrow A$

 $\mathsf{D}\to\mathsf{B}$

 $D \rightarrow C$

 $AC \rightarrow D$

Step-II: $A \rightarrow B$

 $C \rightarrow B$

 $D \rightarrow A$

 $D \rightarrow B$

 $D \rightarrow C$

 $AC \rightarrow D$

Step-III: $AC \rightarrow D$

 $A^+ = AB$

 $C^+ = CB$

Neither A, nor C is extraneous attribute.

 $F_C: \{A \rightarrow B, C \rightarrow B, D \rightarrow AC, AC \rightarrow D\}$



Given a relational Schema R(A, B, C, D) and set of Function Dependency FD = { B \rightarrow A, AD \rightarrow BC, C \rightarrow ABD }.

Find the canonical cover?

Step-I: Split FDs

- 1. $B \rightarrow A$
- 2. AD \rightarrow B (using decomposition inference rule on AD \rightarrow BC)
- 3. AD \rightarrow C (using decomposition inference rule on AD \rightarrow BC)
- 4. $C \rightarrow A$ (using decomposition inference rule on $C \rightarrow ABD$)
- 5. $C \rightarrow B$ (using decomposition inference rule on $C \rightarrow ABD$)
- 6. $C \rightarrow D$ (using decomposition inference rule on $C \rightarrow ABD$)

Step-II: Remove redundant FDs

 $B \rightarrow A$, $AD \rightarrow C$, $C \rightarrow B$, $C \rightarrow D$

Step-III: Remove extraneous attributes

 $B \rightarrow A$, $AD \rightarrow C$, $C \rightarrow BD$

Q. 4 Given a relational Schema R(W, X, Y, Z) and set of Function Dependency FD = $\{W \rightarrow X, Y \rightarrow X, Z \rightarrow WXY, WY\}$

 \rightarrow Z }. Find the canonical cover?

Step-I: Split FDs

- 1. $W \rightarrow X$
- 2. $Y \rightarrow X$
- 3. $Z \rightarrow W$ (using decomposition inference rule on $Z \rightarrow WXY$)
- 4. $Z \rightarrow X$ (using decomposition inference rule on $Z \rightarrow WXY$)
- 5. $Z \rightarrow Y$ (using decomposition inference rule on $Z \rightarrow WXY$)
- 6. WY → Z

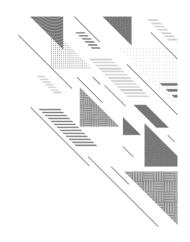
Step-II: Remove redundant FDs

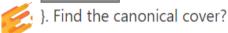
Hence resultant FD = { W \rightarrow X, Y \rightarrow X, Z \rightarrow W, Z \rightarrow Y, WY \rightarrow Z }

Step-III: Remove extraneous attributes

Hence resultant FD = { W \rightarrow X, Y \rightarrow X, Z \rightarrow W, Z \rightarrow Y, WY \rightarrow Z } and we can rewrite as:

FD = { W \rightarrow X, Y \rightarrow X, Z \rightarrow WY, WY \rightarrow Z } is Canonical Cover of FD = { W \rightarrow X, Y \rightarrow X, Z \rightarrow WXY, WY \rightarrow Z }.





Step-I: Split FDs

- 1. V → W
- 2. $VW \rightarrow X$
- 3. $Y \rightarrow V$ (using decomposition inference rule on $Y \rightarrow VXZ$)
- 4. $Y \rightarrow X$ (using decomposition inference rule on $Y \rightarrow VXZ$)
- 5. $Y \rightarrow Z$ (using decomposition inference rule on $Y \rightarrow VXZ$)

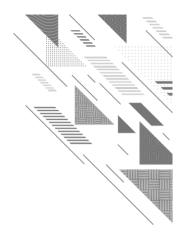
Step-II: Remove redundant FDs

Hence resultant FD = { $V \rightarrow W$, $VW \rightarrow X$, $Y \rightarrow V$, $Y \rightarrow Z$ }.

Step-III: Remove extraneous attributes

Hence resultant FD = { $V \rightarrow W, V \rightarrow X, Y \rightarrow V, Y \rightarrow Z$ } and we can rewrite as

 $FD = \{ \ V \rightarrow WX, \ Y \rightarrow VZ \ \} \ is \ Canonical \ Cover \ of \ FD = \{ \ V \rightarrow W, \ VW \rightarrow X, \ Y \rightarrow VXZ \ \}.$



Suppose a relational schema R(w x y z), and set of functional dependency as following



$$F: \{ x \rightarrow w, \\ wz \rightarrow xy, \\ y \rightarrow wxz \}$$

Find the canonical cover F_c (Minimal set of functional dependency).

Step-I: Decompose FDs

$$x \rightarrow w$$
,
 $wz \rightarrow x$,
 $wz \rightarrow y$,
 $y \rightarrow w$,
 $y \rightarrow x$,
 $y \rightarrow z$

Step-II: Remove redundant FDs

$$x \rightarrow w$$

$$wz \rightarrow y$$

$$y \rightarrow x$$

$$y \rightarrow z$$

Step-III: Remove extraneous attributes

$$F_{c} : \{ x \rightarrow w, \\ wz \rightarrow y, \\ y \rightarrow xz \}$$