

# Graph polynomials

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## 1 The chromatic polynomial

**Definition 1.** A *proper graph coloring* is an assignment of colors to vertices such that adjacent vertices are different colors.

**Question 1.** Given  $x$  colors, how many ways are there to color  $G$ ?

**Notation 1.**  $P_G(x)$  is defined to be the number of colorings of  $G$  using up to  $x$  colors.

For  $K_n$ , this is how it goes:

colors	1	2	3	...	$n-1$	$n$	$n+1$	...	$x$
# colorings	0	0	0	...	0	$n!$	$\binom{n+1}{k}n!$	...	$\binom{x}{n}n!$

For a tree  $T$ , we have:

colors	1	2	3	...	$x$
colorings	0	1	$3 \cdot 2^{n-1}$	...	$x(x-1)^{n-1}$

Let  $f(k)$  denote the number of different  $k$ -partitions of  $G$ . Then

$$P_G(x) = \sum_{k=1}^{|V|} f(k) \cdot x(x-1) \cdots (x-n+1)$$

“It’s a polynomial in  $x$ !”

## 2 Actually calculating that

How can we simplify a graph?

- Delete an edge to get  $G - e$ .
- Contract an edge to get  $G/e$ .

**Proposition 1.** The chromatic polynomial  $P_G$  satisfies

$$P_G = P_{G-e} - P_{G/e}.$$

## 3 The reliability polynomial

What is the probability that  $G$  remains connected after removing edges with probability  $1 - p$ ?

Call this  $R_G(p)$ . We have that

**Theorem 1.**

$$R_G(p) = pR_{G/e}(p) + (1-p)R_{G-e}(p)$$

In  $G/e$ ,  $e$  does not fail. In  $G - e$ ,  $e$  necessarily fails. This is essentially what the proof consists of.