## Tannakian reconstruction of Tambara modules

Or, why that "optics" formula?

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September 26, 2019

## 1 Definition of Tambara module

Let **M** be a monoidal category acting both on two arbitrary categories **C** and **D**. We write  $\underline{M}$  for the image of  $M \in \mathbf{M}$  both in  $[\mathbf{C}, \mathbf{C}]$  and  $[\mathbf{D}, \mathbf{D}]$ .

**Definition 1.** A Tambara module consists of a profunctor  $P: \mathbb{C}^{op} \times \mathbb{D} \to \mathbf{Sets}$  endowed with a family of morphisms  $\alpha_M: P(A, B) \to P(\underline{M}A, \underline{M}B)$  natural in both  $A \in \mathbb{C}$  and  $B \in \mathbb{D}$ , and dinatural in  $M \in \mathbb{M}$ ; which additionally makes the following diagrams commute.

Remark 2. The original definition of Tambara module [T<sup>+</sup>06] deals only with actions that arise from a monoidal product  $\otimes : \mathbf{C} \to [\mathbf{C}, \mathbf{C}]$ . We use the term Tambara module also for the more general concept, allowing for arbitrary monoidal actions.

We can extend Pastro and Street [PS08] construction of free Tambara module over a profunctor  $P \colon \mathbf{C}^{op} \times \mathbf{D} \to \mathbf{Sets}$  to the case of general monoidal actions. Tambara modules are equivalently algebras for a monad  $\Psi$  defined by

$$\Psi P(S,T) = \int^{M,X,Y} \mathbf{C}(S,\underline{M}X) \times \mathbf{D}(\underline{M}Y,T) \times P(X,Y).$$

We know how to contruct free Tambara modules. What is the free Tambara module over a representable functor hom((A, B), -)? We call it  $\mathbf{Optic}((A, B), -)$ , and it can be written as

$$\mathbf{Optic}((A,B),-) \cong \int^M \mathbf{C}(S,\underline{M}A) \times \mathbf{D}(\underline{M}B,T).$$

That is, the formula for optics is given by the free Tambara module on a representable functor.

## 2 Tannakian reconstruction

Milewski [Mil17], and then Boisseau and Gibbons [BG18], proved a unified profunctor representation theorem for optics, that is widely used in programming libraries such as

Kmett's lens [Kme18]. Milewski suggested to me that profunctor representation was surprisingly similar to Tannakian reconstruction; we will prove the theorem following the proof of Tannakian reconstruction (for, say, groups).

**Theorem 3.** Let  $\mathcal{U}_{(A,B)} \colon \mathcal{F} \to \mathbf{Sets}$  the functor that evaluates a Tambara module on the object (A,B). There exists an isomorphism

$$[\mathcal{L}, \mathbf{Sets}] (\mathcal{U}_{(A,B)}, \mathcal{U}_{(S,T)}) \cong \mathbf{Optic}((A,B), (S,T)),$$

natural on both (A, B) and (S, T).

*Proof.* The claim is that this theorem is precisely Tannakian reconstruction for Tambara modules. We first note that, by definition, the functor  $\mathcal{U}_{(A,B)}$  is represented by  $\mathbf{Optic}((A,B),-)$ , the free Tambara module over the hom-profunctor. In fact, for any Tambara module  $P \colon \mathbf{C}^{op} \times \mathbf{C} \to \mathbf{Sets}$ ,

$$\mathcal{U}_{(A,B)}P \cong \operatorname{Nat}(\operatorname{hom}((A,B),-),P) \cong \mathcal{T}_{\mathcal{C}}(\mathbf{Optic}((A,B),-),P).$$

Then, by Tannakian reconstruction,  $[\mathcal{T}, \mathbf{Sets}] (\mathcal{U}_{(A,B)}, \mathcal{U}_{(S,T)}) \cong \mathbf{Optic}((A,B), (S,T)).$ 

## References

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