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IT3090E - Databases

Chapter 5bis: Database design
Some complements

Made by the Database Teaching Team - SoICT

LEARNING POINTS

1. Functional Dependency
2. Armstrong's Axioms and secondary rules
3. Closure of a FD set, closure of a set of attributes
4. A minimal key
5. Equivalence of sets of FDs
6. Minimal Sets of FDs

LEARNING OBJECTIVES

Upon completion of this lesson, students will be able to:

1. Recall the concept of functional dependency, Armstrong's axioms and secondary rules
2. Identify closure of a FD set, closure of a set of attributes
3. Find a minimal key of a relation under a set of FDs
4. Identify the equivalence of sets of FDs and find the minimal cover of a set of FDs

1. Functional Dependency

1.1. Introduction

1.2. Definition

1. Functional Dependency

1.1. Introduction

- We have to deal with the problem of database design: anomalies, redundancies
- The single most important concept in relational schema design theory

1. Functional Dependency

1.2. Definition

- Suppose that $R = \{A_1, A_2, \ldots, A_n\}$, X and Y are non-empty subsets of R.
- A functional dependency (FD), denoted by $X \rightarrow Y$, specifies a constraint on the possible tuples that can form a relation state r of R. The constraint is that, for any two tuples t_1 and t_2 in r that have $t_1[X] = t_2[X]$, they must also have $t_1[Y] = t_2[Y]$.
  - X: the left-hand side of the FD
  - Y: the right-hand side of the FD

1. Functional Dependency

1.2. Definition

- This means that the values of the X component of a tuple uniquely (or functionally) determine the values of the Y component.
- A FD $X \rightarrow Y$ is trivial if $X \supseteq Y$
- If X is a candidate key of R, then $X \rightarrow R$

1. Functional Dependency

1.2. Definition

- Examples
  - $AB \rightarrow C$

A	B	C	D
a1	b1	c1	d1
a1	b1	c1	d2
a1	b2	c2	d1
a2	b1	c3	d1

subject

subject_id	name	credit	percentage_final_exam
IT3090	Databases	3	0.7
IT4843	Data integration	3	0.7
IT4868	Web mining	2	0.6
IT2000	Introduction to ICT	2	0.5
IT3020	Discrete Mathematics	2	0.7
IT3030	Computer Architectures	3	0.7

  - subject_id $\rightarrow$ name,
  - subject_id $\rightarrow$ credit,
  - subject_id $\rightarrow$ percentage_final_exam,
  - subject_id $\rightarrow$ {name, credit}

2. Armstrong's axioms

2. Armstrong's axioms

2.1. Armstrong's axioms

- $R = \{A_1, A_2, \dots, A_n\}$, X, Y, Z, W are subsets of R.
- XY denoted for $X \cup Y$

- Reflexivity
  - If $Y \subseteq X$ then $X \rightarrow Y$
- Augmentation
  - If $X \rightarrow Y$ then $XZ \rightarrow YZ$
- Transitivity
  - If $X \rightarrow Y$, $Y \rightarrow Z$ then $X \rightarrow Z$

2. Armstrong's axioms

2.2. Secondary rules

- Union
  - If $X \rightarrow Y$, $X \rightarrow Z$ then $X \rightarrow YZ$.
- Pseudo-transitivity
  - If $X \rightarrow Y$, $WY \rightarrow Z$ then $XW \rightarrow Z$.
- Decomposition
  - If $X \rightarrow Y$, $Z \subseteq Y$ then $X \rightarrow Z$

2. Armstrong's axioms

2.3. An example

- Given a set of FDs: $F = \{AB \to C, C \to A\}$
- Prove: $BC \to ABC$
  - From $C \to A$, we have $BC \to AB$ (Augmentation)
  - From $AB \to C$, we have $AB \to ABC$ (Augmentation)
  - And we can conclude $BC \to ABC$ (Transitivity)

3. Closure of a FD set, closure of a set of attributes

3.1. Closure of a FD set

3.2. Closure of a set of attributes

3.3. A problem

3. Closure of a FD set, closure of a set of attributes

3.1. Closure of a FD set

- Suppose that $F = \{A \rightarrow B, B \rightarrow C\}$ on R(A, B, C,...). We can infer many FDs such as:

  $A \rightarrow C$, $AC \rightarrow BC$,...

- Definition
  - Formally, the set of all dependencies that include F as well as all dependencies that can be inferred from F is called the closure of F, denoted by F^+.

- $F \models X \rightarrow Y$ to denote that the FD $X \rightarrow Y$ is inferred from the set of FDs F.

3. Closure of a FD set, closure of a set of attributes

3.2. Closure of a set of attributes

- Problem
  - We have F, and $X \rightarrow Y$, we have to check if $F \vDash X \rightarrow Y$ or not
- Should we calculate F^+? $\Rightarrow$ Closure of a set of attributes
- Definition
  - For each such set of attributes X, we determine the set X^+ of attributes that are functionally determined by X based on F; X^+ is called the closure of X under F.

3. Closure of a FD set, closure of a set of attributes

3.2. Closure of a set of attributes

- To find the closure of an attribute set X^+ under F

Input: A set F of FDs on a relation schema R, and a set of attributes X, which is a subset of R.

$X^0 := X$;
repeat
  for each functional dependency $Y \to Z$ in F do
    if $X^{i-1} \supseteq Y$ then $X^i := X^{i-1} \cup Z$;
    else $X^i := X^{i-1}$
until (X^i unchanged);
$X^+ := X^i$

3. Closure of a FD set, closure of a set of attributes

3.2. Closure of a set of attributes

- An example
  - Given R = {A, B, C, D, E, F} and F = {AB $\rightarrow$ C, BC $\rightarrow$ AD, D $\rightarrow$ E, CF $\rightarrow$ B}.

    Calculate $(AB)^+_F$

    - X^0 = AB
    - X^1 = ABC (from AB $\rightarrow$ C)
    - X^2 = ABCD (from BC $\rightarrow$ AD)
    - X^3 = ABCDE (from D $\rightarrow$ E)
    - X^4 = ABCDE
    - $(AB)^+_F$=ABCDE

3. Closure of a FD set, closure of a set of attributes

3.3. A Problem

- $X \rightarrow Y$ can be inferred from F if and only if $Y \subseteq X^+_F$
- $F \vDash X \rightarrow Y \Leftrightarrow Y \subseteq X^+_F$
- An example
  - Let R = {A, B, C, D, E}, F = {$A \rightarrow B$, $B \rightarrow CD$, $AB \rightarrow CE$}.

  Consider whether or not $F \vDash A \rightarrow C$

  - $(A)^+_F$ = ABCDE $\supseteq$ {C}

4. A minimal key

4.1. Definition

4.2. An algorithm to find a minimal key

4.3. An example

4. A minimal key

4.1. Definition

- Minimal key
  - Given $R = \{A_1, A_2, \ldots, A_n\}$, a set of FDs F
  - K is considered as a minimal key of R if:
    - $K \subseteq R$
    - $K \rightarrow R \in F^+$
    - Với $\forall K' \subset K$, thì $K' \rightarrow R \notin F^+$
  - $K^+ = R$ and $K \backslash \{A_i\} \rightarrow R \notin F^+$

4. A minimal key

4.2. An algorithm to find a minimal key

- To find a minimal key

Input: $R = \{A_1, A_2, \ldots, A_n\}$, a set of FDs F

- $Step^0$ $K^0 = R$
- $Step^i$ If $(K^{i-1} \setminus \{A_i\}) \rightarrow R$ then $K^i = K^{i-1} \setminus \{A_i\}$

  else $K^i = K^{i-1}$
- $Step^{n+1}$ $K = K^n$

4. A minimal key

4.3. An example

Given $R = \{A, B, C, D, E\}$, $F = \{AB \rightarrow C, AC \rightarrow B, BC \rightarrow DE\}$.

Find a minimal key

- Step^0: $K^0 = R = ABCDE$
- Step^1: Check if or not $(K^0 \setminus \{A\}) \rightarrow R$ (i.e, $BCDE \rightarrow R$).

  $(BCDE)^+ = BCDE \neq R$. Vậy $K^1 = K^0 = ABCDE$
- Step^2: Check if or not $(K^1 \setminus \{B\}) \rightarrow R$ (i.e, $ACDE \rightarrow R$).

  $(ACDE)^+ = ABCDE = R$. So, $K^2 = K^1 \setminus \{B\} = ACDE$
- Step^3: $K^3 = ACDE$
- Step^4: $K^4 = ACE$
- Step^5: $K^5 = AC$

We infer that AC is a minimal key

4. A minimal key

4.4. An other algorithm to find a minimal key

Input: $U = \{A_1, A_2, ..., A_n\}$, F

Output: a minimal key

- **Step 1:**

  VT = set of all attributes on the left-side of FD in F

  VP = set of all attributes on the right-side of FD in F

  $X = U \setminus VP$: set of attributes **that must be in K**

  $Y = VP \setminus VT$: set of attributes **that must NOT be in** K

  $Z = VP \cap VT$: set of attrbutes **that may be in K**

- **Step 2**: If $(X)^+ = U$ then X is the minimal key: K = X. End!
- **Step 3**: If $(X)^+ \neq U$ then

  - $K^0 = X \cup Z$
  - Repeat: check if we can remove any attribute from Z (similar to slide 22)
  - **$K = K^i$**

5. Equivalence of Sets of FDs

5.1. Definition

5.2. An example

5. Equivalence of Sets of FDs

5.1. Definition

- Definition:
  - A set of FDs F is said to cover another set of FDs G if every FD in G is also in F^+ (every dependency in G can be inferred from F).
- Check if F and G are equivalent:
  - Two sets of FDs F and G are equivalent if $F^+ = G^+$.
    - Therefore, equivalence means that every FD in G can be inferred from F, a nd every FD in F can be inferred from G;
    - That is, G is equivalent to F if both the conditions - G covers F and F covers G - hold.

5. Equivalence of Sets of FDs

5.2. An example

- Prove that $F = \{A \rightarrow C, AC \rightarrow D, E \rightarrow AD, E \rightarrow H\}$ and $G = \{A \rightarrow CD, E \rightarrow AH\}$ are equivalent
  - For each FD of F, prove that it is in G^+
    - $A \rightarrow C$: $(A)^+_G = ACD \supseteq C$, so $A \rightarrow C \in G^+$
    - $AC \rightarrow D$: $(AC)^+_G = ACD \supseteq D$, so $AC \rightarrow D \in G^+$
    - $E \rightarrow AD$: $(E)^+_G = EAHCD \supseteq AD$, so $E \rightarrow AD \in G^+$
    - $E \rightarrow H$: $(E)^+_G = EAHCD \supseteq H$, so $E \rightarrow H \in G^+$
    - $\Rightarrow F^+ \subseteq G^+$
  - For each FD of G, prove that it is in F^+ (the same)
    - $\Rightarrow G^+ \subseteq F^+$
  - $\Rightarrow F^+ = G^+$

6. A minimal cover of a set of FDs

6.1. Definition

6.2. An algorithm to find a minimal cover of a set of FDs

6.3. An example

6. A minimal cover of a set of FDs

6.1. Definition

- Minimal Sets of FDs
  - A set of FDs F to be minimal if it satisfies:
    - Every dependency in F has a single attribute for its right-hand side.
    - We cannot replace any dependency X → A in F with a dependency Y → A, where Y is a proper subset of X, and still have a set of dependencies that is equivalent to F.
    - We cannot remove any dependency from F and still have a set of dependencies that is equivalent to F.
  - A set of dependencies in a standard or canonical form and with no redundancies

6. A minimal cover of a set of FDs

6.2. An algorithm to find a minimal cover of a set of FDs

- Finding a Minimal Cover F for a Set of FDs G

Input: A set of FDs G.

1. Set F := G.

2. Replace each functional dependency $X \rightarrow \{A_1, A_2, \ldots, A_n\}$ in F by the n FDs $X \rightarrow A_1$, $X \rightarrow A_2$, ..., $X \rightarrow A_n$.

3. For each FD $X \rightarrow A$ in F

   for each attribute B that is an element of X

   if $\{\{F - \{X \rightarrow A\}\} \cup \{(X - \{B\}) \rightarrow A\}\}$ is equivalent to F

   then replace $X \rightarrow A$ with $(X - \{B\}) \rightarrow A$ in F.

4. For each remaining functional dependency $X \rightarrow A$ in F

if $\{F - \{X \rightarrow A\}\}$ is equivalent to F, then remove $X \rightarrow A$ from F.

6. A minimal cover of a set of FDs

6.3. An example

- $G = \{B \rightarrow A, D \rightarrow A, AB \rightarrow D\}$. We have to find the minimal cover of G.

  - All above dependencies are in canonical form

  - In step 2, we need to determine if $AB \rightarrow D$ has any redundant attribute on the left-hand side; that is, can it be replaced by $B \rightarrow D$ or $A \rightarrow D$?

  Since $B \rightarrow A$ then $AB \rightarrow D$ may be replaced by $B \rightarrow D$.

  We now have a set equivalent to original G, say G_1: $\{B \rightarrow A, D \rightarrow A, B \rightarrow D\}$.

  - In step 3, we look for a redundant FD in G_1. Using the transitive rule on $B \rightarrow D$ and $D \rightarrow A$, we conclude $B \rightarrow A$ is redundant.

  Therefore, the minimal cover of G is $\{B \rightarrow D, D \rightarrow A\}$

Remark

- Functional dependencies
- Armstrong axioms and their secondary rules
- Closure of a set of FDs
- Closure of a set of attributes under a set of FDs
- An algorithm to find a minimal key
- Equivalence of sets of FDs
- Finding a minimal set of a set of FDs

Summary

1. **Functional Dependency**
   - A FD X → Y: the values of the X component of a tuple uniquely (or functionally) determine the values of the Y component
2. **Armstrong 's Axioms and secondary rules**
   - Reflexivity, Augmentation, Transitivity
   - Union, Pseudo-transitivity, Decomposition
3. **Closure of a FD set, closure of a set of attributes**
   - All dependencies that can be inferred from F, include F, is called the closure of F, denoted by F^+
   - A set of attributes are functionally determined by X based on F
4. **A minimal key**
   - A minimal set of attributes can determine R
5. **Equivalence of sets of functional dependencies**
   - F is equivalent to G if every dependency in G can be inferred from F, and every dependency in F can be inferred from G
6. **A minimal cover of a set of FDs**
   - A set of dependencies in a standard or canonical form and with no redundancies

References

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- Elmasri and Navathe, Fundamentals of Database Systems, 6th edition, Addison-Wesley, 2 011.

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