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B20CS014

Question- Define Three agents: A, B, and a Trusted Third Party T. A and T share a symmetric key ANT. B and T share a symmetric key BAT. A wants to establish a symmetric session key AKB shared with B. Protocol development using SPAN+AVISPA:

- 1 Specifying protocol and properties
- 2 Debugging specification using animation: Find the blocking transition, monitor the

variables

- 3 Attack discovery, strengthening the protocol
- 4 Tuning and optimizing the protocol

We have to define a simple protocol with SPAN+AVISPA, in which there are three agents A, B, and trusted third-party T. In which A and T are going to share a symmetric key K_AT, B, and T are going to share a symmetric key K_BT and the objective of this protocol is to establish a symmetric session key K_AB shared by A and B. Now to develop the protocol we will test three protocols and every time we will try to optimize that protocol and jump to a new protocol that will be an optimized version of previous one.

1. keyexchangeprotocol_1

So first we will go to the key exchange protocol in the RESSI_tutorial part of the SPAN VirtualBox and will open keyexchangeprotocol_1.

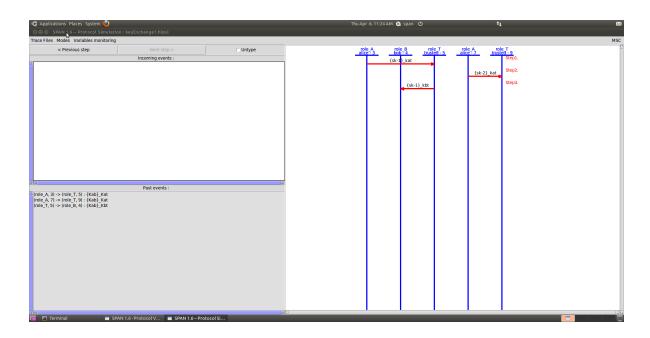
But when we execute it using ATSE it shows that the protocol is unsafe, below result is attached along with hlpsl code.

Specifying protocol and properties

1. A -> T: {Kab}_Kat

2. T -> B: {Kab}_Kbt

Protocol simulation:



 Debugging specification using animation: Find the blocking transition, monitor the variables

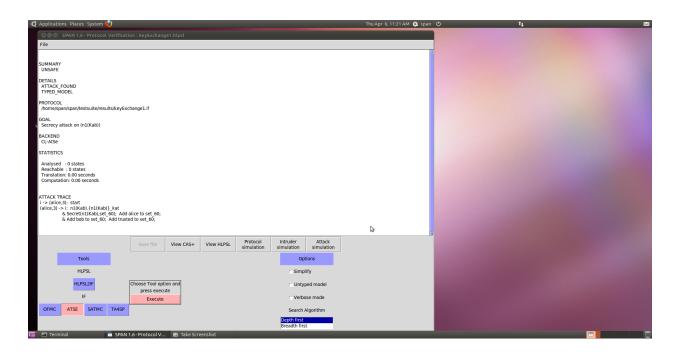
A is going to sends a key to T, he generates a key K_AB and is going to cipher it with that share symmetric key K_AT and then T receive it, he can decrypt because he has the key K_AT and then cipher it with K_AB and sends it to B.

So A has a symmetric key K_AT, T has two symmetric keys K_AT, K_BT, and B has of course one symmetric key that is K_BT.

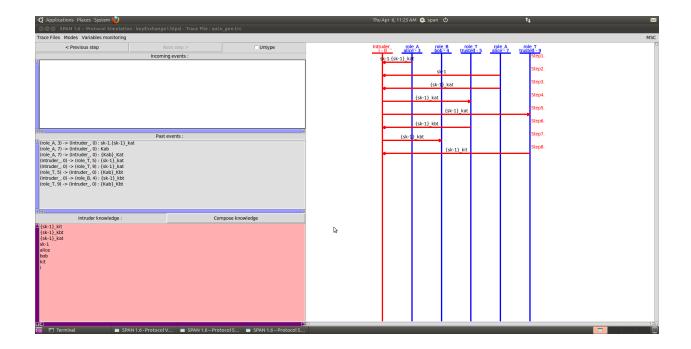
```
/\ secret(Kab', sec_1, {A, B, T})
/\ SND(Kab') %% Unsafe protocol but claimed SAFE!,
                                %% Because of the bugs in the spec.
end role
role role_T(T:agent,A:agent,B:agent,Kat,Kbt:symmetric_key,SND,RCV:channel(dy))
played_by T
def=
local
State:nat,
                 Kab:symmetric_key
init
State := 0
transition
1. State=0 /\ RCV({Kab'}_Kat) =|> State':=1 /\ SND({Kab'}_Kbt)
end role
role role_B(B:agent, A:agent, T:agent, Kbt:symmetric_key, SND, RCV:channel(dy))
played_by B
def=
local
State:nat,
                Kab:symmetric_key
init
State := 0
transition
1. State=0 /\ RCV({Kab'}_Kbt) =|> State':=1
end role
role session(A:agent, B:agent, T:agent, Kat, Kbt:symmetric_key)
def=
SND3, RCV3, SND2, RCV2, SND1, RCV1: channel(dy)
composition
                 role_A(A, B, T, Kat, SND1, RCV1) / 
role_B(B,A,T,Kbt,SND2,RCV2) /\
                role_T(T, A, B, Kat, Kbt, SND3, RCV3)
end role
role environment()
def=
const
kat,kbt,kit:symmetric_key,
                alice, bob, trusted: agent,
                sec_1, auth_1:protocol_id
intruder_knowledge = {alice,bob,kit}
composition
session(alice, bob, trusted, kat, kbt) /\
session(alice,i,trusted,kat,kit)
end role
```

• Attack discovery, strengthening the protocol

this protocol is unsafe we can see below.



below is the intruder simulation:



Tuning and optimizing the protocol

Since there are three bugs in the code of keyexchangeprotocol_1 the protocol is UNSAFE. and that's why we will check keyexchangeprotocol_5 which is the optimized version of keyexchangeprotocol_1.

2. keyexchangeprotocol_5

In this protocol, there are identities and norms everywhere. it's very big and complex protocol. In this protocol there is secrecy and mutual authentication between A and <u>B.Now</u> we will see all results and check all steps are valid for protocol development or not.

Specifying protocol and properties

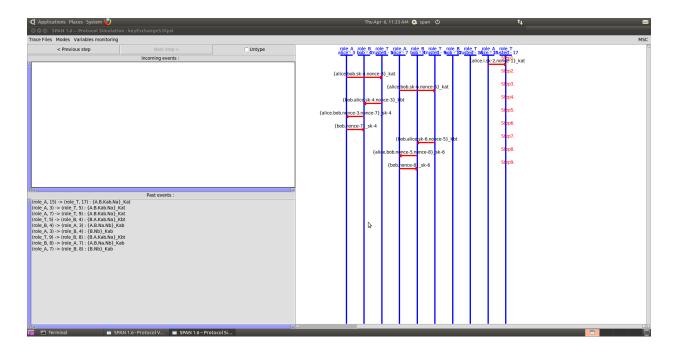
1. A -> T : {A,B,Kab,Na}_Kat

2. T -> B : {B,A,Kab,Na}_Kbt

3. $B -> A : \{A,B,Na,Nb\}_{Kab}$

4. A -> B : {A,B,Nb}_Kab

Protocol simulation



in above protocol simulation we can see that A - T using key K_AT. B-T using key K_BT and A - B using key K_AB.

 Debugging specification using animation: Find the blocking transition, monitor the variables

there is one session. but we can put more than one session in parallel. A is communicating with T using K_AT. B is communicating with T using the symmetric key K_BT. and after that A and B are talking with the key K_AB generated by trusted third party T.

So A has a symmetric key K_AT, T has two symmetric keys K_AT, K_BT, and B has of course one symmetric key that is K_BT.

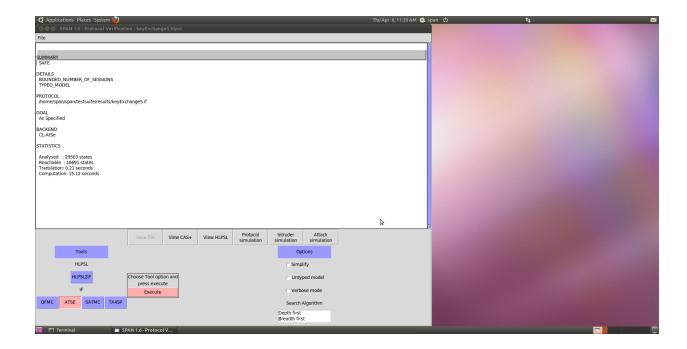
```
%% Key exchange protocol, secured for secrecy, mutual authentication
%%%of A and B (but unoptimized)
%% 1. A -> T : {A,B,Kab,Na}_Kat
%% 2. T -> B : {B,A,Kab,Na}_Kbt
%% 3. B -> A : {A,B,Na,Nb}_Kab
%% 4. A -> B : {A,B,Na,Nb}_Kab
role role_A(A:agent,B:agent,T:agent,Kat:symmetric_key,SND,RCV:channel(dy))
played_by A
```

```
def=
local
State:nat,
                Na, Nb:text,
                Kab:symmetric_key
init
State := 0
transition
1. State=0 /\ RCV(start) =|>
                   State':=1 /\ Na':=new() /\ Kab':=new() /\ SND({A.B.Kab'.Na'}_Kat)
/\ secret(Kab¹, sec_1, {A, B, T})
2. State=1 /\ RCV({A.B.Na.Nb'}_Kab) =|> State':=2 /\ SND({B.Nb'}_Kab)
                   %% A checks that he receives the same nonce
                   %% that he sent at step 1.
                   /\ request(A,B,auth_1,Na)
                   %% A hopes that Nb will permit to authenticate him
                   /\ witness(B, A, auth_2, Nb')
end role
role role_T(T:agent,A:agent,B:agent,Kat,Kbt:symmetric_key,SND,RCV:channel(dy))
played_by T
def=
State:nat,Na:text,Kab:symmetric_key
init
State := 0
transition
1. State=0 /\ RCV({A.B.Kab'.Na'}_Kat) =|>
                   State':=1 /\ SND({B.A.Kab'.Na'}_Kbt)
end role
role role_B(B:agent, A:agent, T:agent, Kbt:symmetric_key, SND, RCV:channel(dy))
played_by B
def=
local
State:nat,Na,Nb:text,Kab:symmetric_key
init
State := 0
transition
1. State=0 /\ RCV({B.A.Kab'.Na'}_Kbt) =|>
                   State':=1 /\ Nb':= new() /\ SND({A.B.Na'.Nb'}_Kab')
                   %% B hopes that Na will permit to authenticate him
                   /\ witness(B,A,auth_1,Na')
                2. State=1 /\ RCV({B.Nb}_Kab) =|> State':=2
                %% B checks that he receives the same nonce
```

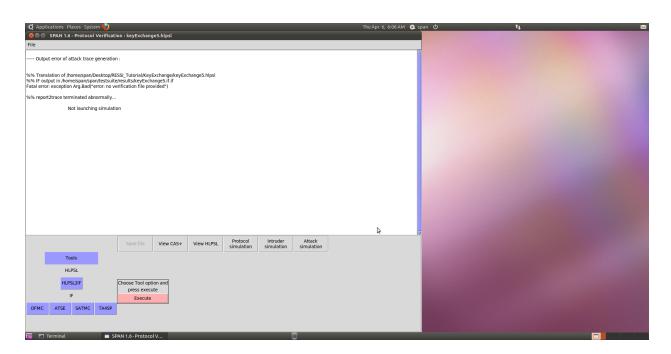
```
%% that he sent at step 1.
               /\ request(A, B, auth_2, Nb)
end role
role session(A:agent,B:agent,T:agent,Kat,Kbt:symmetric_key)
def=
local
SND3, RCV3, SND2, RCV2, SND1, RCV1: channel(dy)
composition
               role_A(A,B,T,Kat,SND1,RCV1) /\
role_B(B,A,T,Kbt,SND2,RCV2) /\
               role_T(T, A, B, Kat, Kbt, SND3, RCV3)
end role
role environment()
def=
const
kat,kbt,kit:symmetric_key,
%% we add a symmetric key: kit shared between the intruder and T
               alice, bob, trusted:agent,
               sec_1, auth_1, auth_2:protocol_id
composition
                    %% We run the regular session
session(alice, bob, trusted, kat, kbt)
                      %% in parallel with another regular session
                /\ session(alice, bob, trusted, kat, kbt)
%
                      %% and a session between the intruder (with key kit) and bob
                /\ session(i,bob,trusted,kit,kbt)
                      %% and a session between alice and the intruder (with key kit)
                /\ session(alice,i,trusted,kat,kit)
end role
goal
secrecy_of sec_1
       authentication_on auth_1
       authentication_on auth_2
end goal
environment()
```

Attack discovery, strengthening the protocol

this protocol is safe that is optimized version of protol 1 but it is not optimized.



we can say that it is safe as there is no attack simulation



• Tuning and optimizing the protocol

keyexchangeprotocol_5 is SAFE but not optimized. and that's why we will check keyexchangeprotocol_6 which is the optimized version of keyexchangeprotocol_5 and 1.

3. keyexchangeprotocol_6

this protocol is simplified, safe and optimized over all key exchange protocol.

• Specifying protocol and properties

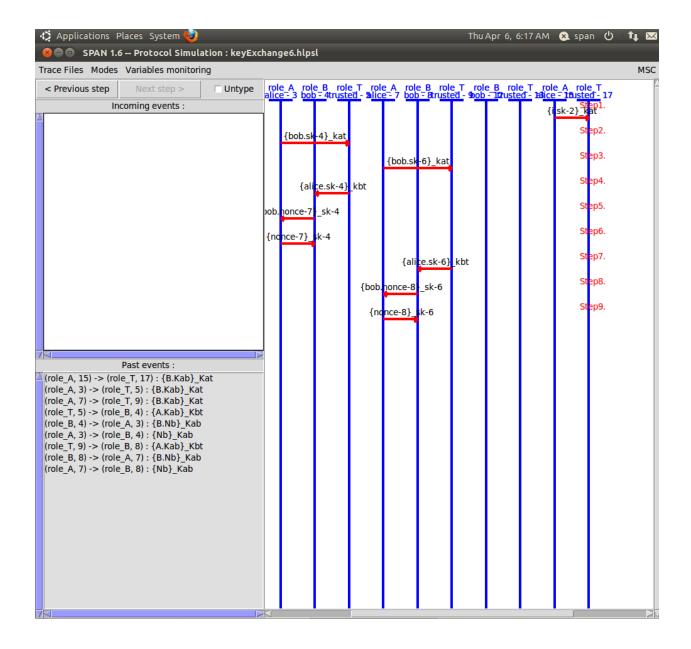
1. A -> T : {B,Kab}_Kat

2. **T -> B**: **{A,Kab}_Kbt**

3. **B** -> **A** : **{B,Nb}**_**Kab**

4. A -> B : {Nb}_Kab

Protocol simulation



in above protocol simulation we can see that A - T using key K_AT. B-T using key K_BT and A - B using key K_AB.

 Debugging specification using animation: Find the blocking transition, monitor the variables

there is one session. but we can put more than one session in parallel. A is communicating with T using K_AT. B is communicating with T using the symmetric key

K_BT. and after that A and B are talking with the key K_AB generated by trusted third party T.

So A has a symmetric key K_AT, T has two symmetric keys K_AT, K_BT, and B has of course one symmetric key that is K_BT.

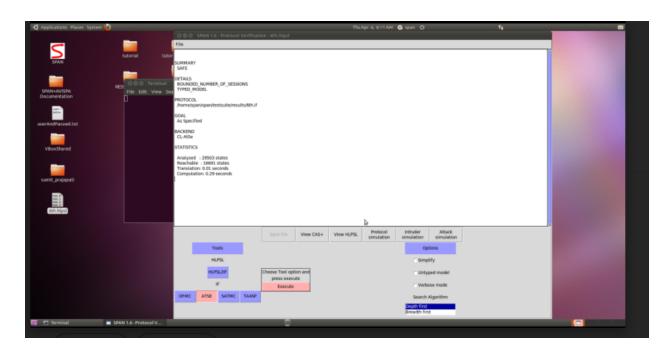
```
%%% Key exchange protocol, secured for secrecy, mutual authentication of
%%A and B (optimized)
%%% 1. A -> T : {B,Kab}_Kat
%%% 2. T -> B : {A,Kab}_Kbt
%%% 3. B -> A : {B,Nb}_Kab
%%% 4. A -> B : {Nb}_Kab
role role_A(A:agent,B:agent,T:agent,Kat:symmetric_key,SND,RCV:channel(dy))
played_by A
def=
local
State:nat,
               Na, Nb:text,
               Kab:symmetric_key
init
State := 0
transition
1. State=0 /\ RCV(start) =|>
                  /\ secret(Kab', sec_1, {A, B, T})
2. State=1 \land RCV({B.Nb'}_Kab) =|> State':=2 \land SND({Nb'}_Kab)
                  %% A checks that B uses the same key
                  %% that he sent at step 1.
                  /\ request(A, B, auth_1, Kab)
                  %% A hopes that Nb will permit to authenticate him
                  /\ witness(A, B, auth_2, Nb')
end role
role role_T(T:agent,A:agent,B:agent,Kat,Kbt:symmetric_key,SND,RCV:channel(dy))
played_by T
def=
State:nat, Na:text, Kab:symmetric_key
init
State := 0
transition
1. State=0 /\ RCV({B.Kab'}_Kat) =|>
                  State':=1 /\ SND({A.Kab'}_Kbt)
end role
```

```
role role_B(B:agent,A:agent,T:agent,Kbt:symmetric_key,SND,RCV:channel(dy))
played_by B
def=
local
State:nat,Na,Nb:text,Kab:symmetric_key
State := 0
transition
1. State=0 /\ RCV({A.Kab'}_Kbt) =|>
                   State':=1 /\ Nb':= new() /\ SND({B.Nb'}_Kab')
                   %% B hopes that Kab will permit to authenticate him
                   /\ witness(B, A, auth_1, Kab')
                2. State=1 /\ RCV({Nb}_Kab) =|> State':=2
                %% B checks that he receives the same nonce
                %% that he sent at step 1.
                /\ request(B, A, auth_2, Nb)
end role
role session(A:agent, B:agent, T:agent, Kat, Kbt:symmetric_key)
def=
local
SND3, RCV3, SND2, RCV2, SND1, RCV1: channel(dy)
composition
                role_A(A,B,T,Kat,SND1,RCV1) /\
role_B(B,A,T,Kbt,SND2,RCV2) /\
                role_T(T,A,B,Kat,Kbt,SND3,RCV3)
end role
role environment()
def=
const
kat,kbt,kit:symmetric_key,
                                 %% we add a symmetric key: kit
shared between the intruder and T
                alice, bob, trusted:agent,
                sec_1, auth_1, auth_2:protocol_id
intruder_knowledge = {alice,bob,kit}
                                        %% ... and we give it to the intruder
composition
                      %% We run the regular session
session(alice, bob, trusted, kat, kbt)
                      %% in parallel with another regular session
                  /\ session(alice, bob, trusted, kat, kbt)
                      %% and a session between the intruder (with key kit) and bob
                  /\ session(i, bob, trusted, kit, kbt)
                      %% and a session between alice and the intruder (with key kit)
                  /\ session(alice,i,trusted,kat,kit)
end role
```

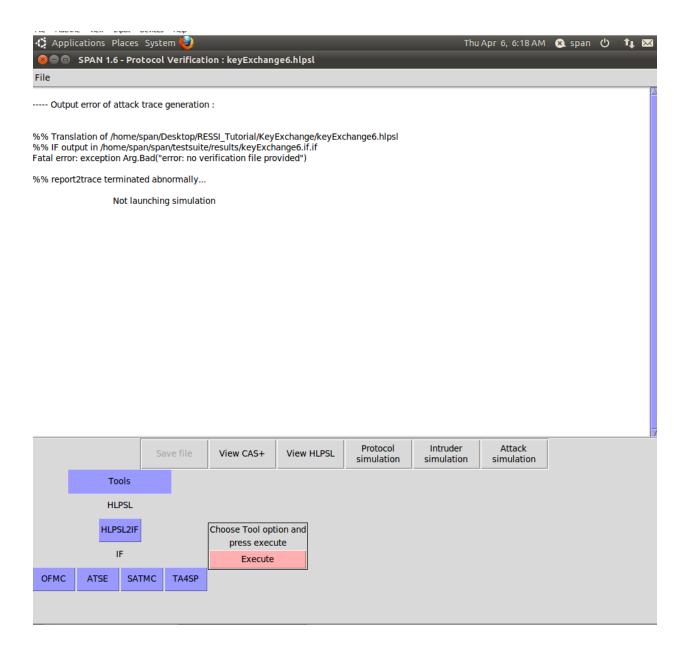
```
goal
secrecy_of sec_1
            authentication_on auth_1
            authentication_on auth_2
end goal
environment()
```

• Attack discovery, strengthening the protocol

this protocol is safe that is optimized version of protol_5 and optimized.



there is no attack simulation and hence safe.



• Tuning and optimizing the protocol

It is the most optimized protocol simulation, which satisfies all steps that are need for the Protocol development using SPAN+AVISPA