# Density of States of Hodge Laplacians: Decomposition effects and Sparsification of SC

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#### Abstract:

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#### Introduction

Well, it would be nice to have one. You know, somewhere here, maybe...

## **Graphs and Simplicial Complexes**

Simplicial complex  $\mathcal{K}$  is a higher-order generalization of the classical graph model with pair-wise interactions with a rich set of topological descriptors. Specifically, let  $\mathcal{V}_0(\mathcal{K}) = \{v_1, \ldots v_{m_0}\}$  be a set of nodes; then  $\mathcal{K}$  is a collection of subsets (simplices)  $\sigma$  of nodes from  $\mathcal{V}_0(\mathcal{K})$  such that all the subsets of  $\sigma$  are simplices in  $\mathcal{K}$  too. We refer to a simplex made out of k+1 nodes  $\sigma = [v_{i_1}, \ldots v_{i_{k+1}}]$  as being of order kthe set of all the simplices of order k in the complex  $\mathcal{K}$  is denoted by  $\mathcal{V}_k(\mathcal{K})$ . Thus,  $\mathcal{V}_0(\mathcal{K})$  are the vertices of  $\mathcal{K}$ ,  $\mathcal{V}_1(\mathcal{K})$  are edges between pairs of vertices,  $\mathcal{V}_2(\mathcal{K})$  triangles connecting three vertices, and so on. We let  $m_k = |\mathcal{V}_k(\mathcal{K})|$  denote the cardinality of  $\mathcal{V}_k(\mathcal{K})$ .

Each set of simplices  $\mathcal{V}_k(\mathcal{K}) = \left\{\sigma_1, \ldots \sigma_{m_k}\right\}$  induces a linear space of formal sums over the simplicies  $C_k(\mathcal{K}) = \left\{\sum_{i=1}^{m_k} \alpha_i \sigma_i \mid \alpha_i \in \mathbb{R}\right\}$  referred to as *chain space*; in particular,  $C_0(\mathcal{K})$  is known as the space of vertex states and  $C_1(\mathcal{K})$  as the space of edge flows. Simplices of different orders are related through the boundaries operators  $\partial_k$  mapping the simplex to its boundary; formally,  $\partial_k : C_k(\mathcal{K}) \mapsto C_{k-1}(\mathcal{K})$  is defined through the alternating sum:

$$\partial_k[v_1, v_2, \dots v_k] = \sum_{i=1}^k (-1)^{i-1}[v_1, \dots v_{i-1}, v_{i+1}, \dots v_k]$$

By fixing an ordering for  $\mathcal{V}_k(\mathcal{K})$  we can fix a canonical basis for  $C_k(\mathcal{K})$  and represent each boundary operator as a matrix  $B_k \in \operatorname{Mat}_{m_{k-1} \times m_k}$  with exactly k nonzero entries per each column, being either +1 or -1. For these matrices, the fundamental property of topology holds: the boundary of the boundary is zero, [Lim20, Thm. 5.7]:

$$B_k B_{k+1} = 0 (Eqn. 1)$$

The matrix representation  $B_k$  of the boundary operator  $\partial_k$  requires fixing an ordering of the simplices in  $\mathcal{V}_k(\mathcal{K})$  and  $\mathcal{V}_{k-1}(\mathcal{K})$ . As it will be particularly relevant for the purpose of this work, we emphasize that we order triangles and edges as follows: triangles in  $\mathcal{V}_2(\mathcal{K})$  are oriented in such a way that the first edge (in terms of the ordering of  $\mathcal{V}_1(\mathcal{K})$ ) in each triangle is positively acted upon by  $B_{2i}$  i.e. the first non-zero entry in each column of  $B_2$  is +1, Figure .1.

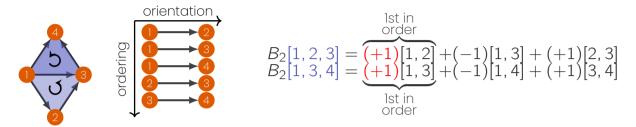


Figure .1: Example of the simplicial complex with ordering and orientation: nodes from  $\mathcal{V}_0(\mathcal{K})$  in orange, triangles from  $\mathcal{V}_2(\mathcal{K})$  in blue. Orientation of edges and triangles is shown by arrows; the action of  $B_2$  operator is given for both triangles.

The following definitions introduce the fundamental concepts of kth homology group and k-th order Laplacian. See |Lim20| e.g. for more details.

(Homology group and higher-order Laplacian) Since im  $B_{k+1} \subset$ Def. 1  $\ker B_k$ , the quotient space  $\mathcal{H}_k = \ker^{B_k}/\ker^{B_{k+1}}$ , known as k-th homology group, is correctly defined and the following isomorphisms hold

$$\mathcal{H}_k \cong \ker B_k \cap \ker B_{k+1}^{\top} \cong \ker \left( B_k^{\top} B_k + B_{k+1} B_{k+1}^{\top} \right).$$

The matrix  $L_k = B_k^ op B_k + B_{k+1} B_{k+1}^ op$  is called the k-th order Laplacian **operator**; the two terms  $L_k^{\downarrow} = B_k^{\top} B_k$  and  $L_k^{\uparrow} = B_{k+1} B_{k+1}^{\top}$  are referred to as the **down-Laplacian** and the **up-Laplacian**, respectively.

The homology group  $\mathcal{H}_k$  describes the k-th topology of the simplicial complex  $\mathcal{K}$ :  $\beta_k = \dim \mathcal{H}_k = \dim \ker L_k$  coincides exactly with the number of k-dimensional holes in the complex, known as the k-th Betti number. In the case k=0, the operator  $L_0=L_0^{\uparrow}$  is exactly the classical graph Laplacian whose kernel corresponds to the connected components of the graph, while  $L_0^{\downarrow}=0$ . For k=1 and k=2, the elements of ker  $L_1$  and ker  $L_2$  describe the simplex 1-dimensional holes and voids respectively, and are frequently used in the analysis of trajectory data, [SBH<sup>+</sup>20, BGL16].

Although more frequently found in their purely combinatorial form, the definitions of simplicial complexes, homology groups, and higher-order Laplacians admit a generalization to the weighted case. For the sake of generality, in the rest of the work, we use the following notion of weighted boundary operators (and thus weighted simplicial complexes), as considered in e.g. GST23.

(Weighted and normalised boundary matrices) For weight functions Def. 2  $w_k: \mathcal{V}_k(\mathcal{K}) \mapsto \mathbb{R}_+ \cup \{0\}$ , define the diagonal weight matrix  $W_k \in$ 

 $\mathrm{Mat}_{m_k \times m_k}$  as  $(W_k)_{ii} = \sqrt{w_k(\sigma_i)}$ . Then the weighting scheme for the boundary operators upholding the Hodge algebras 1 is given by:

$$B_k \mapsto W_{k-1}^{-1} B_k W_k \tag{Eqn. 2}$$

Note that, with the weighting scheme Equation (2), the dimensionality of the homology group is preserved,  $\dim \ker L_k = \dim \ker \widehat{L}_k$ , [GST23] as well as the fundamental property of topology Equation (1).

## Density of States and Hodge Decomposition

Full spectral decomposition of  $L_k^{\uparrow}$  requires  $\mathcal{O}(m_k^3)$  (or, more precisely,  $\mathcal{O}(m_k^{\omega})$  where  $\omega$  is the exponent of matrix multiplication, [BGVKS23]); accounting for the sparsity of  $L_k^{\uparrow}$  the actual computation time is even lower. As a result, even so we are going to use  $\mathcal{O}(m_k^3)$  as a reference point to beat, the actual execution time should be carefully compared.

Instead of the direct computation, one can exploit spectral density functions known as **density of states** which encompass all the spectral information. Specifically,

(Density of States) Fro a given symmetric matrix  $A = Q\Lambda Q^{\top}$  with  $Q^{\top}Q = I$  and diagonal  $\Lambda = \text{diag}(\lambda_1, \dots \lambda_n)$ , the spectral density or density of states (DoS)

$$\mu(\lambda \mid A) = \frac{1}{n} \sum_{i=1}^{n} \delta(\lambda - \lambda_i)$$
 (Eqn. 3)

Additionally, let  $\mathbf{q}_i$  be a corresponding unit eigenvector of A (such that  $A\mathbf{q}_i = \lambda_i \mathbf{q}_i$  and  $Q = (\mathbf{q}_1 \mid \mathbf{q}_2 \mid \cdots \mid \mathbf{q}_n)$ ); then one can define a set of local (entry-wise) densities of states (LDoS):

$$\mu_k(\lambda \mid A) = \sum_{i=1}^n \left| \mathbf{e}_k^{\mathsf{T}} \mathbf{q}_i \right|^2 \delta \left( \lambda - \lambda_i \right)$$
 (Eqn. 4)

with  $\mathbf{e}_k$  being the corresponding versor.

Note that DoS  $\mu(\lambda \mid A)$  only describes the spectrum  $\sigma(A)$  while the family of LDoS  $\{\mu_k(\lambda \mid A)\}$  contain additionally eigenvectors  $\mathbf{q}_i$ . The exact computation of either DoS or LDoS still requires the full spectral decomposition; instead, one can try to approximate densities of states via Kernel Polynomial Method (KPM).

## I. Chebyshev Approximation of DoS/LDoS

add some motivation for this discussion/transition

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Let us assume we transition from symmetric A to a symmetric H such that  $\sigma(H)\subseteq [-1,1]$ , e.g. by  $H=\frac{2}{\lambda_{\max}(A)}A-I$ . This allows us to assume that  $\sup \mu(\lambda\mid H)\subseteq [-1,1]$  and it can be decomposed into an orthonormal basis on [-1,1].

Specifically, let  $T_m(x)$  be a family of Chebyshev polynomials of the first kind iff:

$$T_0(x) = 1$$
,  $T_1(x) = x$ ,  $T_{m+1}(x) = 2xT_m(x) - T_{m-1}(x)$  (Eqn. 5)

Polymonials  $T_m(x)$  are famously orthogonal on [-1,1] in  $L_2$ -scalar product with the weight function  $w(x)=\frac{2}{(1+\delta_{0m})\pi\sqrt{1-x^2}}$ . Then,  $T_m^*(x)=\frac{1}{(1+\delta_{0m})\pi\sqrt{1-x^2}}$ 

 $w_m(x)T_m(x)$  form a dual Chebyshev basis, and one can write an expansion:

figure out how to write here the problem with m=0 for the orthonormality

(Eqn. 6)

$$\mu(\lambda \mid H) = \sum_{m=0}^{\infty} d_m T_m^*(\lambda)$$
$$\mu_k(\lambda \mid H) = \sum_{m=0}^{\infty} d_{mk} T_m^*(\lambda)$$

Note that since  $T_m(x)$  are orthogonal to  $T_m^*(x)$ , so

$$d_{m} = \int T_{m}(\lambda)\mu(\lambda \mid H)d\lambda = \frac{1}{n}\sum_{i=1}^{n}T_{m}(\lambda_{i}) = \frac{1}{n}\operatorname{tr}(T_{m}(H))$$

$$d_{mk} = \int T_{m}(\lambda)\mu_{k}(\lambda \mid H)d\lambda = \sum_{i=1}^{n}\left|\mathbf{e}_{k}^{\top}\mathbf{q}_{i}\right|^{2}T_{m}(\lambda_{i}) = [T_{m}(H)]_{kk}$$
(Eqn. 7)

with

$$\mathbb{E}\left[\mathbf{z}^{\top}H\mathbf{z}\right] = \operatorname{tr}(H) \approx \frac{1}{N_{z}} \sum_{j=1}^{N_{z}} Z_{j}^{\top}HZ_{j}$$

$$\mathbb{E}\left[\mathbf{z} \odot H\mathbf{z}\right] = \operatorname{diag}(H) \approx \frac{1}{N_{z}} \sum_{j=1}^{N_{z}} Z_{j} \odot HZ_{j}$$
(Eqn. 8)

where each entry in  ${\bf z}$  is i.i.d with zero mean and unit variance.

Note that due to the recurrent definition,  $d_m = \frac{1}{n}\mathbb{E}\left[\mathbf{z}^{\top}T_m(H)\mathbf{z}\right] = \frac{1}{n}\mathbb{E}\left[\mathbf{z}^{\top}\left(2HT_{m-1}(H) - T_{m-2}(H)\right)\mathbf{z}\right]$ , the computation of  $d_m$  requires only one new matvec computation for  $HT_{m-1}(H)\mathbf{z}$  since estimations for  $T_{m-1}(H)\mathbf{z}$  and  $T_{m-2}(H)\mathbf{z}$  are inherited from  $d_{m-1}$  and  $d_{m-2}$ . As a result, the computational cost of each new moment  $d_m$  is  $\mathcal{O}(\mathtt{nnz}(H))$ , and, assuming we use M moments in the decomposition, the overall complexity is  $\mathcal{O}\left(N_z M\mathtt{nnz}(H)\right)$ . In the case of  $L_k$ , we get  $\mathcal{O}\left(N_z M(km_k + (k+1)m_{k+1})\right)$ .

#### II. Spike problems and filtration

#### III. Effects of the Hodge Decomposition

## Sparsification of Simplicial Complexes

Simplicial complex  $\mathcal K$  typically has quite an intrinsic structure of the associated Laplacian operators  $L_k$  which becomes more and more complicated as  $\mathcal K$  becomes denser in the sense of complexes of order  $\mathcal K$ . Instead, one can ask a simplifying question: can one find a sparser simplicial complex  $\mathcal L$  with a spectrally close operator  $L_k(\mathcal L)$ ?

The overall idea of the sparsification is built upon the idea that a sparsifier can be subsampled from a given simplificial copmlex with high probability. Specifically:

Def. 4

(Spectral Approximation) The Hermitian matrix A is called spectrally  $\varepsilon$ -close to the Hermitian matrix B,  $A \approx B$ , if  $(1 - \varepsilon)B \preceq A \preceq (1 + \varepsilon)B$ , where  $\succeq$  is the partial ordering induced by the positive definite cone, i.e.  $A \succeq B$  if A - B is positive semi-definite.

Rem III.1

Note that if  $A \approx B$ , then one can directly bound the distance  $\|\mathbf{x}_A - \mathbf{x}_B\|$ , where  $A\mathbf{x}_A = \mathbf{f}$  and  $B\mathbf{x}_B = \mathbf{f}$ . Indeed,  $\mathbf{x}_B = B^{-1}A\mathbf{x}_A$  and  $\|\mathbf{x}_A - \mathbf{x}_B\| = \|(I - B^{-1}A)\mathbf{x}_A\| \le \|I - B^{-1}A\|\|\mathbf{x}_A\| \le \|B^{-1}\| \cdot \|B - A\| \cdot \|\mathbf{x}_A\| \le \varepsilon \|B^{-1}\| \cdot \|B\| \cdot \|\mathbf{x}_A\| = \varepsilon \kappa(B)\|\mathbf{x}_A\|$ , with  $\kappa(B)$  being the condition number of B. Thus, the relative error is controlled by the quality of the approximation  $\varepsilon$  and the condition number  $\kappa(B)$ .

Th IV.III.1

(Simplicial Sparsification, [OPW22]) Let  $\mathcal K$  be a simplicial complex restricted to its p-skeleton,  $\mathcal K=\bigcup_{i=0}^p \mathcal V_i(\mathcal K)$ . Let  $L_k^\uparrow(\mathcal K)$  be its k-th up-Laplacian and let  $m_k=|\mathcal V_k(\mathcal K)|$ . For any  $\varepsilon>0$ , a sparse simplicial complex  $\mathcal L$  can be sampled as follows:

- (1) compute the probability measure  $\mathbf{p}$  on  $\mathcal{V}_{k+1}(\mathcal{K})$  proportional to the generalized resistance vector  $\mathbf{r} = \operatorname{diag}\left(B_{k+1}^{\top}(L_k^{\uparrow})^{\dagger}B_{k+1}\right)$ , where  $\operatorname{diag}(A)$  denotes the vector of the diagonal entries of A;
- (2) sample q simplices  $\tau_i$  from  $\mathcal{V}_{k+1}(\mathcal{K})$  according to the probability measure  $\mathbf{p}$ , where q is chosen so that  $q(m_k) \geq 9C^2m_k\log(m_k/\varepsilon)$ , for some absolute constant C > 0;
- (3) form a sparse simplicial complex  $\mathcal{L}$  with all the sampled simplexes of order k and all its faces with the weight  $\frac{w_{k+1}(\tau_i)}{q(m_k)\mathbf{p}(\tau_i)}$ ; weights of repeated simplices are accumulated.

Then, with probability at least 1/2, the up-Laplacian of the sparsifier  $\mathcal{L}$  is  $\varepsilon$ -close to the original one, i.e. it holds  $L_k^{\uparrow}(\mathcal{L}) \approx L_k^{\uparrow}(\mathcal{K})$ .

The bottleneck of the subsampling above is the construction of the appropriate measure  ${\bf p}$  or, more precisely, generalized effective resistance  ${\bf r}$ . Indeed, one needs a fast pseudo-inverse operator  $\left(L_k^{\uparrow}\right)^{\dagger}$  in order to compute  ${\bf r}$ .

Rem III 2

(Sensitivity of the Sparsification vis-a-vis sampling measure p)

Compare  $\varepsilon$  with  $\frac{1}{m_2}$ : if below, you are fiiiiiine INSERT FIGURE HERE

Th IV.III.2

(GER through DoS) For a given simplicial complex mcK with the k-th order up-Laplacian  $L_k^{\uparrow} = B_{k+1}W_{k+1}^2B_{k+1}^{\top}$ , a generalized effective resistance r can be computed through family of local densities of states  $\{\mu_i(\lambda \mid L_{k+1}^{\downarrow})\}$ :

we need to assume somewhere  $W_k = I$  for simplicity and do not forget about it

$$\mathbf{r}_i = \int_{\mathbb{R}} (1 - \mathbb{1}_0(\lambda)) \mu_i(\lambda \mid L_{k+1}^{\downarrow}) d\lambda$$

(qn. 9

Proof

Let  $B_{k+1}W_{k+1} = USV^{\top}$  where S is diagonal and invertible and both U and V are orthogonal (so it is a truncated SVD decomposition of  $B_{k+1}W_{k+1}$  matrix with eliminated obsolete kernel). Then:

$$(L_k^{\uparrow})^{\dagger} = \left(B_{k+1}W_{k+1}^2B_{k+1}^{\top}\right)^{\dagger} = \left(US^2U^{\top}\right)^{\dagger} = US^{-2}U^{\top}$$
 (Eqn. 10)

$$\mathbf{r} = \operatorname{diag}\left(W_{k+1}B_{k+1}^{\top}(L_{k}^{\uparrow})^{\dagger}B_{k+1}W_{k+1}\right) =$$

$$= \operatorname{diag}\left(VSU^{\top}US^{-2}U^{\top}USV^{\top}\right) = \operatorname{diag}\left(VV^{\top}\right)$$
(Eqn. 11)

As a result,  $\mathbf{r}_i = \|V_i.\|^2 = \sum_j |v_{ij}|^2$ , so the i-th entry of the resistance is defined by the sum of square of i-th components of eigenvectors  $\mathbf{v}_j$  of  $L_{k+1}^{\downarrow} = W_{k+1}B_{k+1}^{\intercal}B_{k+1}W_{k+1}$  operator where  $\mathbf{v}_j \perp \ker L_{k+1}^{\downarrow}$ . Note that

$$\mu_i(\lambda \mid L_{k+1}^{\downarrow}) = \sum_{j=1}^{m_{k+1}} \left| \mathbf{e}_i^{\top} \mathbf{q}_j \right|^2 \delta\left(\lambda - \lambda_j\right) = \sum_{j=1}^{m_{k+1}} \left| q_{ij} \right|^2 \delta\left(\lambda - \lambda_j\right) \quad \text{(Eqn. 12)}$$

SO

$$\begin{aligned} \mathbf{r}_i &= \|V_{i\cdot}\|^2 = \sum_j |v_{ij}|^2 = \int_{\mathbb{R}\backslash\{0\}} \sum_{j=1}^{m_{k+1}} \left|q_{ij}\right|^2 \delta\left(\lambda - \lambda_j\right) d\lambda = \\ &= \int_{\mathbb{R}\backslash\{0\}} \mu_i(\lambda \mid L_{k+1}^{\downarrow}) d\lambda = \int_{\mathbb{R}} (1 - \mathbb{1}_0(\lambda)) \mu_i(\lambda \mid L_{k+1}^{\downarrow}) d\lambda \end{aligned} \tag{Eqn. 13}$$

check for the exact case

error propagation

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