

Free object

In **mathematics**, the idea of a **free object** is one of the basic concepts of **abstract algebra**. It is a part of **universal algebra**, in the sense that it relates to all types of algebraic structure (with **finitary** operations). It also has a formulation in terms of **category theory**, although this is in yet more abstract terms. Examples include **free groups**, **tensor algebras**, or **free lattices**. Informally, a free object over a set A can be thought as being a “generic” algebraic structure over A : the only equations that hold between elements of the free object are those that follow from the defining axioms of the algebraic structure.

1 Definition

Free objects are the direct generalization to **categories** of the notion of **basis** in a vector space. A linear function $u : E_1 \rightarrow E_2$ between vector spaces is entirely determined by its values on a basis of the vector space E_1 . Conversely, a function $u : E_1 \rightarrow E_2$ defined on a basis of E_1 can be uniquely extended to a linear function. The following definition translates this to any category.

Let (C, F) be a **concrete category** (i.e. $F : C \rightarrow \mathbf{Set}$ is a **faithful functor**), let X be a set (called *basis*), $A \in C$ an object, and $i : X \rightarrow F(A)$ a map between sets (called *canonical injection*). We say that A is the **free object on X** (with respect to i) if and only if they satisfy this **universal property**:

for any object B and any map between sets $f : X \rightarrow F(B)$, there exists a unique morphism $\tilde{f} : A \rightarrow B$ such that $f = F(\tilde{f}) \circ i$. That is, the following diagram commutes:

$$\begin{array}{ccc} X & \xrightarrow{i} & F(A) \\ & \searrow f & \downarrow F(\tilde{f}) \\ & & F(B) \end{array}$$

In this way the free functor that builds the free object A from the set X becomes **left adjoint** to the **forgetful functor**.

2 Examples

The creation of free objects proceeds in two steps. For algebras that conform to the **associative law**, the first step is to consider the collection of all possible **words** formed from an **alphabet**. Then one imposes a set of **equivalence relations** upon the words, where the relations are the defining relations of the algebraic object at hand. The free object then consists of the set of **equivalence classes**.

Consider, for example, the construction of the free group in two generators. One starts with an alphabet consisting of the five letters $\{e, a, b, a^{-1}, b^{-1}\}$. In the first step, there is not yet any assigned meaning to the “letters” a^{-1} or b^{-1} ; these will be given later, in the second step. Thus, one could equally well start with the alphabet in five letters that is $S = \{a, b, c, d, e\}$. In this example, the set of all words or strings $W(S)$ will include strings such as *aebecede* and *abdc*, and so on, of arbitrary finite length, with the letters arranged in every possible order.

In the next step, one imposes a set of equivalence relations. The equivalence relations for a **group** are that of multiplication by the identity, $ge = eg = g$, and the multiplication of inverses: $gg^{-1} = g^{-1}g = e$. Applying these relations to the strings above, one obtains

$$aebecede = aba^{-1}b^{-1}$$

where it was understood that c is a stand-in for a^{-1} , and d is a stand-in for b^{-1} , while e is the identity element. Similarly, one has

$$abdc = abb^{-1}a^{-1} = e$$

Denoting the equivalence relation or **congruence** by \sim , the free object is then the collection of **equivalence classes** of words. Thus, in this example, the free group in two generators is the **quotient**

$$F_2 = W(S) / \sim$$

This is often written as

$$F_2 = W(S) / E$$

where

$$W(S) = \{a_1 a_2 \dots a_n \mid a_k \in S; n \text{ finite}\}$$

is the set of all words, and

$$E = \{a_1 a_2 \dots a_n \mid e = a_1 a_2 \dots a_n; a_k \in S; n \text{ finite}\}$$

is the equivalence class of the identity, after the relations defining a group are imposed.

A simpler example are the **free monoids**. The free monoid on a set X , is the monoid of all finite **strings** using X as alphabet, with operation **concatenation** of strings. The identity is the empty string. In essence, the free monoid is simply the set of all words, with no equivalence relations imposed. This example is developed further in the article on the **Kleene star**.

2.1 General case

In the general case, the algebraic relations need not be associative, in which case the starting point is not the set of all words, but rather, strings punctuated with parentheses, which are used to indicate the non-associative groupings of letters. Such a string may equivalently be represented by a **binary tree** or a **free magma**; the leaves of the tree are the letters from the alphabet.

The algebraic relations may then be general **arities** or **finitary relations** on the leaves of the tree. Rather than starting with the collection of all possible parenthesized strings, it can be more convenient to start with the **Herbrand universe**. Properly describing or enumerating the contents of a free object can be easy or difficult, depending on the particular algebraic object in question. For example, the free group in two generators is easily described. By contrast, little or nothing is known about the structure of **free Heyting algebras** in more than one generator.^[1] The problem of determining if two different strings belong to the same equivalence class is known as the **word problem**.

As the examples suggest, free objects look like constructions from **syntax**; one may reverse that to some extent by saying that major uses of syntax can be explained and characterised as free objects, in a way that makes apparently heavy 'punctuation' explicable (and more memorable).

3 Free universal algebras

Main article: **term algebra**

Let S be any set, let \mathbf{A} be an **algebraic structure** of type ρ generated by S . Let the underlying set of this algebraic structure \mathbf{A} , sometimes called universe, be A , and let

$\psi : S \longrightarrow A$ be a function. We say that (A, ψ) (or informally just \mathbf{A}) is a **free algebra** (of type ρ) on the set S of **free generators** if, for every algebra \mathbf{B} of type ρ and function $\tau : S \longrightarrow B$, where B is a universe of \mathbf{B} , there exists a unique homomorphism $\sigma : A \longrightarrow B$ such that $\sigma\psi = \tau$.

4 Free functor

The most general setting for a free object is in **category theory**, where one defines a **functor**, the **free functor**, that is the **left adjoint** to the **forgetful functor**.

Consider the category \mathbf{C} of **algebraic structures**; these can be thought of as sets plus operations, obeying some laws. This category has a functor, $U : \mathbf{C} \rightarrow \mathbf{Set}$, the **forgetful functor**, which maps objects and functions in \mathbf{C} to \mathbf{Set} , the **category of sets**. The forgetful functor is very simple: it just ignores all of the operations.

The free functor F , when it exists, is the left adjoint to U . That is, $F : \mathbf{Set} \rightarrow \mathbf{C}$ takes sets X in \mathbf{Set} to their corresponding free objects $F(X)$ in the category \mathbf{C} . The set X can be thought of as the set of “generators” of the free object $F(X)$.

For the free functor to be a left adjoint, one must also have a **Set**-morphism $\eta : X \rightarrow U(F(X))$. More explicitly, F is, up to isomorphisms in \mathbf{C} , characterized by the following **universal property**:

Whenever A is an algebra in \mathbf{C} , and $g : X \rightarrow U(A)$ is a function (a morphism in the category of sets), then there is a unique \mathbf{C} -morphism $h : F(X) \rightarrow A$ such that $U(h)\eta = g$.

Concretely, this sends a set into the free object on that set; it's the “inclusion of a basis”. Abusing notation, $X \rightarrow F(X)$ (this abuses notation because X is a set, while $F(X)$ is an algebra; correctly, it is $X \rightarrow U(F(X))$).

The **natural transformation** $\eta : \text{id}_{\mathbf{Set}} \rightarrow UF$ is called the **unit**; together with the counit $\varepsilon : FU \rightarrow \text{id}_{\mathbf{C}}$, one may construct a **T-algebra**, and so a **monad**. This leads to the next topic: free functors exist when \mathbf{C} is a monad over \mathbf{Set} .

4.1 Existence

There are general existence theorems that apply; the most basic of them guarantees that

Whenever \mathbf{C} is a **variety**, then for every set X there is a free object $F(X)$ in \mathbf{C} .

Here, a variety is a synonym for a **finitary algebraic category**, thus implying that the set of relations are finitary, and **algebraic** because it is **monadic** over \mathbf{Set} .

4.2 General case

Other types of forgetfulness also give rise to objects quite like free objects, in that they are left adjoint to a forgetful functor, not necessarily to sets.

For example the **tensor algebra** construction on a **vector space** as left adjoint to the functor on **associative algebras** that ignores the algebra structure. It is therefore often also called a **free algebra**.

Likewise the **symmetric algebra** and **exterior algebra** are free symmetric and anti-symmetric algebras on a vector space.

6 See also

- **Generating set**

7 Notes

- [1] Peter T. Johnstone, *Stone Spaces*, (1982) Cambridge University Press, ISBN 0-521-23893-5. (A treatment of the one-generator free Heyting algebra is given in chapter 1, section 4.11)

5 List of free objects

Specific kinds of free objects include:

- **free algebra**
 - **free associative algebra**
 - **free commutative algebra**
- **free category**
 - **free strict monoidal category**
- **free group**
 - **free abelian group**
 - **free partially commutative group**
- **free Kleene algebra**
- **free lattice**
 - **free Boolean algebra**
 - **free distributive lattice**
 - **free Heyting algebra**
- **free Lie algebra**
- **free magma**
- **free module**
- **free monoid**
 - **free commutative monoid**
 - **free partially commutative monoid**
- **free ring**
- **free semigroup**
- **free semiring**
 - **free commutative semiring**
- **free theory**
- **term algebra**
- **discrete space**

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