tramac 10/manual -2+31+0 2-1 Static analysis & C code why study frama C? Contracts: \ nesult nequers, ensures /* (a) requires 2 = int-min; ensures \qualt >=0; ensures $\times \times = = \times$ \quad \qu (alling with we) plugion
frama-c-gui - wp - wp-nte. filenome. ab muspone error - los energos addition vosos calling with (val) plugin frama-c-qui -val filenome.C Some functions to write specification withing specification Valid () (forall lexists) 1*@ *1 intergu Real void swop (int * a, int * b) Spoop Ha requires (valia(a) int Imp; (vola(b)) temp = *a; ensures (*a == \old(*b) *a = *b; 4 b == \ old(40); *b = fem P) */

frame-Coperators ! P P & & 2 p112 P=> 9 -> 17 P then 89 PL=>Q > p if and only if 9 TSwop 2 elements in array void swap (int n, int a[], int ni, int n2) Sperfications 14@ greguires n>=0 880=27127 880=270220; requires |valid (a+(0...n-1)); erques a (a[ni] == \oida [nz] && a[n2] == \odd(n2]); (find() int find (int n, constint all, int v) n>=0 88 \valid(a+(0...n-D)) Specification / + @ requises · mothing; A - motethis. assigns nesult == ensures ==> (Yorall i; O=<i<n ==> a[i]:=v) 0 L= \musult < n ensceres ==> allowellt] == Vs 0=< gresuit < n. ensures 4/

loop invariant, loop assigns, loop variant / * a loopin amant loop invariant 0 1 127) Horall interger ; 0 < j < j ==> a [j] != V; loop assigns i;
loop variant n-i; loop invariant -> Should be true @ start g loop

@ end g loop In (VP) more specifications than code, one it cheeks for all possible values, texhaustiro liso Behaviors: Used to write clean contracts assumes + Les give name for each behavior.
Les conte assumes, requires ensures for each L'all possible ipput la different casses example with find (valid (at(o.m-1)) j Assigns \nothing; behavior found: assumes \inist inleger; 02.1<n & 8 a [i] == V; ensures a[\mull] == V) behavior not/ound: assumes I forall integer 03120 28 a[i]]=v) ensura | qualt ==-1;