

CSCE 420 - Fall 2023 Homework 2 (HW2)

due: **Tues, Oct 24, 11:59 pm**

Turn-in answers as a Word document (HW2.docx or .pdf) and commit/push it to your class github repo.

When pushing the final version of your HW1, also run the command:

git tag "HW1" && git push origin "HW1"

This will be used to record time of submission for late penalty when applicable.

All homeworks must be typed, *not* hand-written and scanned/photo-shot.

1a. Prove that "Implication Introduction" (the opposite of Implication Elimination) is a **sound rule of inference** (ROI) using a **truth table**. If you have a Horn clause, with 1 positive literal and $n-1$ negative literals, like $(\neg X \vee \neg Z \vee \neg Y)$, you can transform it into a conjunctive rule by collecting the negative literals as positive antecedents, e.g. $X \wedge Y \rightarrow Z$. It is sufficient to prove this for $n-1=2$ antecedents. (In fact, this is a truth-preserving operation, hence sound.)

X	Y	Z	$\neg X$	$\neg Y$	$\neg X \vee \neg Z \vee \neg Y$	$X \wedge Y$	$X \wedge Y \rightarrow Z$
T	T	T	F	F	T	T	T
T	T	F	F	F	F	T	F
T	F	T	F	T	T	F	T
T	F	F	F	T	T	F	T
F	T	T	T	F	T	F	T
F	T	F	T	F	T	F	T
F	F	T	T	T	T	F	T
F	F	F	T	T	T	F	T

The truth values for these sentences are the same so they are truth preserving meaning it is a sound rule of inference

1b. Prove that $(A \wedge B \rightarrow C \wedge D) \vdash (A \wedge B \rightarrow C)$ ("conjunctive rule splitting") is a **sound rule-of-inference** using a **truth table**.

A	B	C	D	$A \wedge B$	$C \wedge D$	$A \wedge B \rightarrow C \wedge D$	$A \wedge B \rightarrow C$
T	T	T	T	T	T	T	T
T	T	T	F	T	F	F	T
T	T	F	T	T	F	F	F
T	T	F	F	T	F	F	F
T	F	T	T	F	T	T	T
T	F	T	F	F	F	T	T
T	F	F	T	F	F	T	T

T	F	F	F	F	F	T	T
F	T	T	T	F	T	T	T
F	T	T	F	F	F	T	T
F	T	F	T	F	F	T	T
F	T	F	F	F	F	T	T
F	F	T	T	F	T	T	T
F	F	T	F	F	F	T	T
F	F	F	F	F	F	T	T

This rule is sound because all the models that satisfy the premise also satisfy the derived sentence.

1c. Also prove $(A \wedge B \rightarrow C \wedge D) \models (A \wedge B \rightarrow C)$ using **Natural Deduction**. (hint: use 1a above)

- 1 $\neg(A \wedge B) \vee C \wedge D$ implication elimination
- 2 $\neg A \vee \neg B \vee C \wedge D$ demorgans law (1)
- 3 $\neg A \vee \neg B \vee C$ distribution (2)
- 4 $\neg A \vee \neg B \vee D$ distribution (2)
- 5 $A \wedge B \rightarrow C$ implication introduction (4)

1d. Also prove $(A \wedge B \rightarrow C \wedge D) \models (A \wedge B \rightarrow C)$ using **Resolution**.

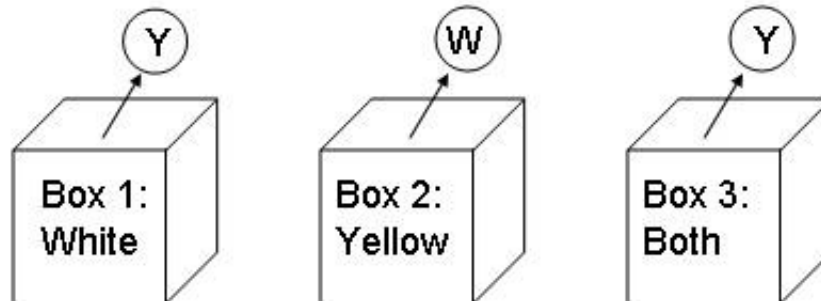
- 1 $\neg(A \wedge B) \vee C$ rewrote derived sentence
- 2 $\neg(\neg(A \wedge B) \vee C)$ negate derived sentence (1)
- 3 $A \wedge B \wedge \neg C$ Demorgans law (2)
- 4 A And elimination (3)
- 5 B And elimination (3)
- 6 $\neg C$ And elimination (3)
- 7 $A \wedge B \rightarrow C \wedge D$ from premise
- 8 $\neg(A \wedge B) \vee C \wedge D$ implication elimination (7)
- 9 $\neg A \vee \neg B \vee C \wedge D$ Demorgans law (8)
- 10 $\neg A \vee \neg B \vee C$ distribution (9)
- 11 $\neg A \vee \neg B \vee D$ distribution (9)
- 12 $\neg B \vee C$ Resolution (10 and 4)
- 13 C Resolution (12 and 6)
- 14 {} Resolution (13 and 6)

By proving the negated derived sentence results in an empty sentence through resolution, the premise is proved to entail the derived sentence.

2. Sammy's Sport Shop

You are the proprietor of *Sammy's Sport Shop*. You have just received a shipment of three boxes filled with tennis balls. One box contains only yellow tennis balls, one box contains only white tennis balls, and one contains both yellow and white tennis balls. You would like to stock

the tennis balls in appropriate places on your shelves. Unfortunately, the boxes have been labeled incorrectly; the manufacturer tells you that you have exactly one box of each, but that **each box is definitely labeled wrong**. You draw one ball from each box and observe its color. Given the initial (incorrect) labeling of the boxes above, and the three observations, use Propositional Logic to infer the correct contents of the middle box.



Use propositional symbols in the following form: O1Y means a yellow ball was drawn (observed) from box 1, L1W means box 1 was initially labeled white, C1W means box 1 contains (only) white balls, and C1B means box 1 actually contains both types of tennis balls. Note, there is no 'O1B', etc, because you can't directly "observe both". When you draw a tennis ball, it will either be white or yellow.

The initial facts describing this particular situation are: {O1Y, L1W, O2W, L2Y, O3Y, L3B}

2a. Using these propositional symbols, write a propositional knowledge base (sammy.kb) that captures the knowledge in this domain (i.e. implications of what different observations or labels mean, as well as constraints inherent in this problem, such as that all boxes have different contents). *Do it in a complete and general way*, writing down *all* the rules and constraints, not just the ones needed to make the specific inference about the middle box. *Do not include derived knowledge* that depends on the particular labeling of this instance shown above; stick to what is stated in the problem description above. Your KB should be general enough to reason about any alternative scenario, not just the one given above (e.g. with different observations and labels and box contents).

// labels are incorrect

- 1 L1Y -> -C1Y
- 2 L1W -> -C1W
- 3 L1B -> -C1B
- 4 L2Y -> -C2Y
- 5 L2W -> -C2W
- 6 L2B -> -C2B
- 7 L3Y -> -C3Y
- 8 L3W -> -C3W
- 9 L3B -> -C3B

// all contents different

- 10 C1W -> -C2W ^ -C3W
- 11 C1Y -> -C2Y ^ -C3Y

```

12    C1B -> -C2B ^ -C3B
13    C2W -> -C1W ^ -C3W
14    C2Y -> -C1Y ^ -C3Y
15    C2B -> -C1B ^ -C3B
16    C3W -> -C1W ^ -C2W
17    C3Y -> -C1Y ^ -C2Y
18    C3B -> -C1B ^ -C2B

```

// observed color implies either box is that color or contains both

```

19    O1Y -> C1Y v C1B
20    O1W -> C1W v C1B
21    O2Y -> C2Y v C2B
22    O2W -> C2W v C2B
23    O3Y -> C3Y v C3B
24    O3W -> C3W v C3B

```

// labels are incorrect (another way to label)

```

25    L1Y -> C2Y v C3Y
26    L1W -> C2W v C3W
27    L1B -> C2B v C3B
28    L2Y -> C1Y v C3Y
29    L2W -> C1W v C3W
30    L2B -> C1B v C3B
31    L3Y -> C1Y v C2Y
32    L3W -> C1W v C2W
33    L3B -> C1B v C2B

```

// cant have multiple labels

```

34    C1Y -> -C1W
35    C1Y -> -C1B
36    C1W -> -C1Y
37    C1W -> -C1B
38    C1B -> -C1Y
39    C1B -> -C1W

```

```

34    C2Y -> -C2W
35    C2Y -> -C2B
36    C2W -> -C2Y
37    C2W -> -C2B
38    C2B -> -C2Y
39    C2B -> -C2W

```

```

34    C3Y -> -C3W
35    C3Y -> -C3B
36    C3W -> -C3Y

```

- 37 C3W \rightarrow \neg C3B
- 38 C3B \rightarrow \neg C3Y
- 39 C3B \rightarrow \neg C3W

2b. Prove that box 2 must contain white balls (**C2W**) using **Natural Deduction**.

- a L2Y observed
- b O2W observed
- c L1W observed
- d O1Y observed
- e L3B observed
- f O3Y observed
- g \neg C2Y (a and 4)
- h \neg C1W (c and 2)
- i \neg C3B (e and 9)
- j C2W \vee C2B (b and 22)
- k C1Y \vee C1B (d and 19)
- l C3Y \vee C3B (f and 23)
- m C3Y resolution (i and l)
- n \neg C1Y \wedge \neg C2Y (m and 17)
- o \neg C1Y and elimination (n)
- p \neg C2Y and elimination (n)
- q C1B resolution (k and o)
- r \neg C2B \wedge \neg C3B (q and 12)
- s \neg C2B and elimination (r)
- t \neg C3B and elimination (r)
- u C2W resolution (j and s)

2c. Convert your KB to CNF.

// labels are incorrect

- 1 \neg L1Y \vee \neg C1Y
- 2 \neg L1W \vee \neg C1W
- 3 \neg L1B \vee \neg C1B
- 4 \neg L2Y \vee \neg C2Y
- 5 \neg L2W \vee \neg C2W
- 6 \neg L2B \vee \neg C2B
- 7 \neg L3Y \vee \neg C3Y
- 8 \neg L3W \vee \neg C3W
- 9 \neg L3B \vee \neg C3B

// all contents different

- 10 \neg C1W \vee \neg C2W
- 11 \neg C1W \vee \neg C3W

12 -C1Y v -C2Y
 13 -C1Y v -C3Y
 14 -C1B v -C2B
 15 -C1B v -C3B
 16 -C2W v -C1W
 17 -C2W v -C3W
 18 -C2Y v -C1Y
 19 -C2Y v -C3Y
 20 -C2B v -C1B
 21 -C2B v -C3B
 22 -C3W v -C1W
 23 -C3W v -C2W
 24 -C3Y v -C1Y
 25 -C3Y v -C2Y
 26 -C3B v -C1B
 27 -C3B v -C2B

// observed color implies either box is that color or contains both

28 -O1Y v C1Y v C1B
 29 -O1W v C1W v C1B
 30 -O2Y v C2Y v C2B
 31 -O2W v C2W v C2B
 32 -O3Y v C3Y v C3B
 33 -O3W v C3W v C3B

// labels are incorrect (another way to label)

34 -L1Y v C2Y v C3Y
 35 -L1W v C2W v C3W
 36 -L1B v C2B v C3B
 37 -L2Y v C1Y v C3Y
 38 -L2W v C1W v C3W
 39 -L2B v C1B v C3B
 40 -L3Y v C1Y v C2Y
 41 -L3W v C1W v C2W
 42 -L3B v C1B v C2B

// cant have multiple labels

43 -C1Y v -C1W
 44 -C1Y v -C1B
 45 -C1W v -C1Y
 46 -C1W v -C1B
 47 -C1B v -C1Y
 48 -C1B v -C1W
 49 -C2Y v -C2W

- 50 $\neg C2Y \vee \neg C2B$
- 51 $\neg C2W \vee \neg C2Y$
- 52 $\neg C2W \vee \neg C2B$
- 53 $\neg C2B \vee \neg C2Y$
- 54 $\neg C2B \vee \neg C2W$

- 55 $\neg C3Y \vee \neg C3W$
- 56 $\neg C3Y \vee \neg C3B$
- 57 $\neg C3W \vee \neg C3Y$
- 58 $\neg C3W \vee \neg C3B$
- 59 $\neg C3B \vee \neg C3Y$
- 60 $\neg C3B \vee \neg C3W$

2d. Prove C2W using **Resolution**.

Facts = { i O1Y,
 ii L1W,
 iii O2W,
 iv L2Y,
 v O3Y,
 vi L3B }

- | | | |
|---|---------------------|-------------------------|
| a | $\neg C2W$ | negate derived sentence |
| b | $\neg L1W \vee C3W$ | resolution (a and 35) |
| c | $\neg L3W \vee C1W$ | resolution (a and 41) |
| d | $\neg O2W \vee C2B$ | resolution (a and 31) |
| e | C3W | resolution (b and iii) |
| f | $C3Y \vee C3B$ | 32 and v |
| g | $\neg C3Y$ | e and 55 |
| h | C3B | f and g |
| i | $\neg C3W$ | h and 60 |
| j | { } | e and i |

Because the negated derived sentence resulted in an empty sentence, the derived sentence is true through resolution refutation.

3. Do **Forward Chaining** for the *CanGetToWork* KB below.

You don't need to follow the formal FC algorithm (with agenda/queue and counts array). Just indicate which rules are triggered (in any order), and keep going until all consequences are generated.

Show the final list of all inferred propositions at the end. *Is CanGetToWork among them?*

KB = { a. CanBikeToWork \rightarrow CanGetToWork
 b. CanDriveToWork \rightarrow CanGetToWork
 c. CanWalkToWork \rightarrow CanGetToWork
 d. HaveBike \wedge WorkCloseToHome \wedge Sunny \rightarrow CanBikeToWork
 e. HaveMountainBike \rightarrow HaveBike
 f. HaveTenSpeed \rightarrow HaveBike
 g. OwnCar \rightarrow CanDriveToWork
 h. OwnCar \rightarrow MustGetAnnualInspection
 i. OwnCar \rightarrow MustHaveValidLicense
 j. CanRentCar \rightarrow CanDriveToWork
 k. HaveMoney \wedge CarRentalOpen \rightarrow CanRentCar
 l. HertzOpen \rightarrow CarRentalOpen
 m. AvisOpen \rightarrow CarRentalOpen
 n. EnterpriseOpen \rightarrow CarRentalOpen
 o. CarRentalOpen \rightarrow IsNotAHoliday
 p. HaveMoney \wedge TaxiAvailable \rightarrow CanDriveToWork
 q. Sunny \wedge WorkCloseToHome \rightarrow CanWalkToWork
 r. HaveUmbrella \wedge WorkCloseToHome \rightarrow CanWalkToWork
 s. Sunny \rightarrow StreetsDry }

Facts: { Rainy, HaveMountainBike, EnjoyPlayingSoccer, WorkForUniversity,
 WorkCloseToHome, HaveMoney, HertzClosed, AvisOpen, McDonaldsOpen }

Facts (agenda):

- 1) Rainy, HaveMountainBike, EnjoyPlayingSoccer, WorkForUniversity,
 WorkCloseToHome, HaveMoney, HertzClosed, AvisOpen, McDonaldsOpen
- 2) Rainy, HaveMountainBike, EnjoyPlayingSoccer, WorkForUniversity,
 WorkCloseToHome, HaveMoney, HertzClosed, AvisOpen, McDonaldsOpen,
 HaveBike
- 3) Rainy, HaveMountainBike, EnjoyPlayingSoccer, WorkForUniversity,
 WorkCloseToHome, HaveMoney, HertzClosed, AvisOpen, McDonaldsOpen,
 HaveBike, CarRentalOpen
- 4) Rainy, HaveMountainBike, EnjoyPlayingSoccer, WorkForUniversity,

WorkCloseToHome, HaveMoney, HertzClosed, AvisOpen, McDonaldsOpen, HaveBike, CarRentalOpen, CanRentCar, IsNotAHoliday

5) Rainy, HaveMountainBike, EnjoyPlayingSoccer, WorkForUniversity, WorkCloseToHome, HaveMoney, HertzClosed, AvisOpen, McDonaldsOpen, HaveBike, CarRentalOpen, CanRentCar, IsNotAHoliday, CanDriveToWork

6) Rainy, HaveMountainBike, EnjoyPlayingSoccer, WorkForUniversity, WorkCloseToHome, HaveMoney, HertzClosed, AvisOpen, McDonaldsOpen, HaveBike, CarRentalOpen, CanRentCar, IsNotAHoliday, CanDriveToWork, CanGetToWork

7) Rainy, HaveMountainBike, EnjoyPlayingSoccer, WorkForUniversity, WorkCloseToHome, HaveMoney, HertzClosed, AvisOpen, McDonaldsOpen, HaveBike, CarRentalOpen, CanRentCar, IsNotAHoliday, CanDriveToWork, CanGetToWork

CanGetToWork is generated from the given facts

4. Do **Backward Chaining** for the *CanGetToWork* KB.

In this case, you should follow the BC algorithm closely (the pseudocode for the propositional version of Back-chaining is given in the lecture slides).

Important: when you pop a subgoal (proposition) from the goal stack, you should systematically go through all rules that can be used to prove it **IN THE ORDER THEY APPEAR IN THE KB**. In some cases, this will lead to *back-tracking*, which you should show.

Also, the sequence of results depends on order in which antecedents are pushed onto the stack. If you have a rule like $A \wedge B \rightarrow C$, and you pop C off the stack, push the antecedents in reverse order, so B goes in first, then A; in the next iteration, A would be the next subgoal popped off the stack.

CanGetToWork	goal
CanBikeToWork	pop CanGetToWork
HaveBike, WorkCloseToHome, Sunny	pop CanBikeToWork
HaveMountainBike, WorkCloseToHome, Sunny	pop HaveBike
HaveTenSpeed, WorkCloseToHome, Sunny	backtrack; not provable

CanDriveToWork	backtrack; not provable
OwnCar	pop CanDriveToWork
CanRentCar	backtrack; not provable
HaveMoney, CarRentalOpen	pop CanRentCar
CarRentalOpen	pop HaveMoney (fact)
HertzOpen	pop CarRentalOpen
AvisOpen	backtrack; not provable
{}	pop AvisOpen

CanGetToWork is true since it resulted in an empty set through back chaining.