

Week - 5 Time-cost trade-off

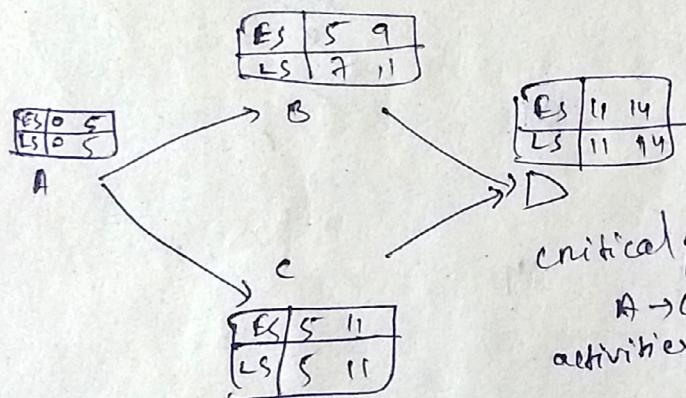
1. Fast tracking vs crashing.

Methods to reduce the time of ^{overall} project.

Fast tracking - Performing tasks in parallel to reduce time

Crashing - Reducing the time of individual activities for overall project time.

Activity	Duration	Pred. duration	(minimum duration for an activity)		ES	EF	CP	C/D	C/day	Cost
			Normal	Critical						
A	5	-	3	250	300	350	2	25		
B	4	A	3	300	375	425	1	75		
C	6	A	3	350	825	875	3	125		
D	3	B, C	2	300	350	350	1	50		



critical path
 $A \rightarrow C \rightarrow D$
 activities $\rightarrow A, C, D$

expenditure of
 crashing a single day

expected dur' = 14
 current comp'd
 dur' = 14 d

expected = 10 d
 How to do?

* Crashing - Reducing activity time by expending additional resources.

Normal dur' - usual/normal activity dur'

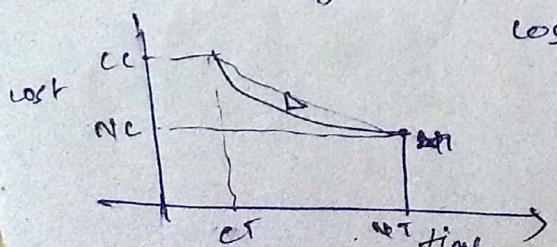
Crash dur' - minimized/reduced time dur'

completing activity in crash dur'

Normal cost - cost incurred in normal dur'

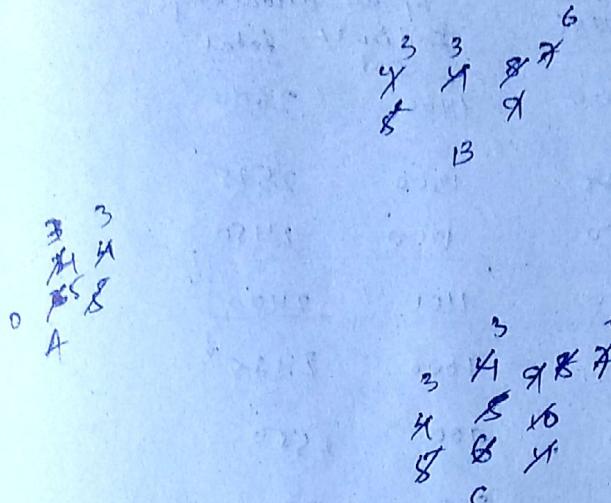
Crash cost - " " to complete activity in crash dur'

Dur' $\propto \frac{1}{\text{cost}}$ less dur' - more cost (additional resources)
 more " - less cost.



$$\text{slope} = \frac{Nc - Cc}{Nt - Ct}$$

$$\text{Slope} = \frac{y}{n} = \frac{\text{crash cost} - \text{normal cost}}{\text{Normal time} - \text{crash time}}$$



	Crash Cost	Normal Cost
14d - basic	14/-	25/-
13d - +25/- ex	14/-	25/-
12 + 25/-	14/-	25/-
11 + 50/-	14/-	25/-
10 + 125/-	14/-	25/-
9 + 125/-	14/-	25/-
8 + 250/-	14/-	25/-

Crashing details

- 1. A by 1 @ 25/-
- 2. A by 1 "
- 3. O by 1 @ 50/-

Task - determine

Determine minimized cost

when deen" is reduced.

Critical activities - A, C, D

13 days expected.

Task - 1. Choose the critical activity that has min. cost/day.

Begin with that.

Hence, A = 25/day - 125/day, D = 50/day.

Step 1. Undercrash A.

that A is chosen. A is crashed till crash deen" is reached,

Cost is calculated for each day. As $CD_A = 3$, we can't

be crashed any more.

task - 2 choose the next minimum costing crashing activity.

Hence, A: $CD = 2$.

task - 3 By 1

(Crashing details).

Now, B is also critical. It is fully crashed out.

- 1. A by 1 @ 25/-
- 2. "
- 3. G by 1 @ 50/-
- 4. C by 1 @ 125/-
- 5. C by 1 @ 125/-
- 6. B by 1 @ 250/-

<u>Days</u>	<u>Add'l cost/day</u>	<u>Direct cost</u>	<u>Indirect cost</u>	<u>Total</u>
14	0	1200	1400	2600
13	25	1225	1300	2525
12	25	1250	1200	2450
11	50	1300	1100	2400
10	125	1425	1000	2425
9	125	1450	900	2550
8	250	1650	800	2500
7		1900		

↓
constant.
@ 100/day.

• this gives the min. possible cost.

Indirect costs

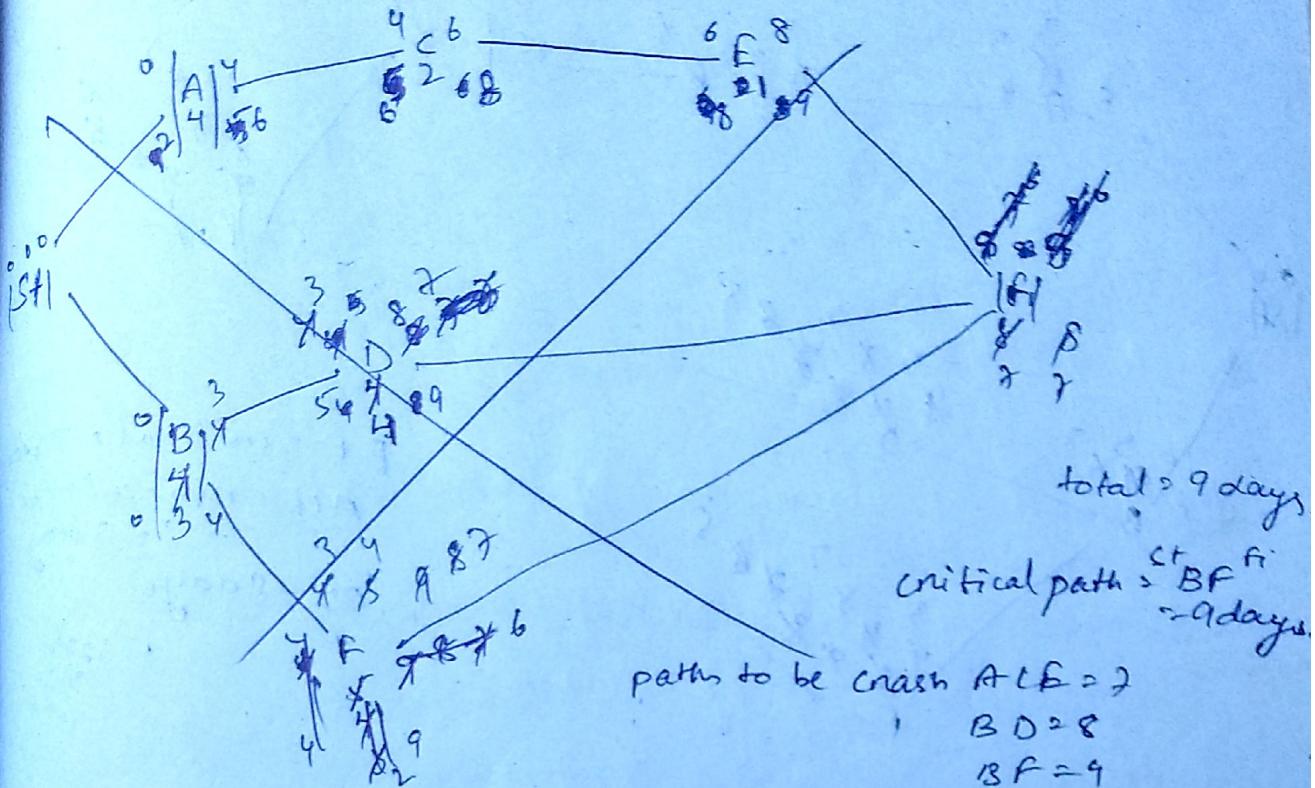
- Salaries, temporary utilities like electricity, phone.
- Site equipment, vehicle, office essentials. Rent, maintenance, etc.
- # When "project day" is reduced, indirect cost will also decrease. They are not specific to any activity.
- Constant for each day, unlike direct costs that vary with activities.

decr' d I., decr' d indirect cost.
direct cost

Total cost

Analysing the total cost shows that, initially with decr' in "decr' cost" decreases. Then it starts inc. with further decrease.

Here, 11 days is the min. total cost point. By Reducing below it will incur larger costs. They reducing decr' from 14 to 11 will give optimum results.



<u>Activity</u>	<u>ND</u>	<u>CD</u>	<u>NC</u>	<u>CC</u>	<u>CF-NC</u> <u>ND-NC</u>	
A	4	3	100	125	25	9
B	4	3	250	400	150	8 + 50/day
C	2		2			.7 + 50
D	4	1	150	300	150	5 + 150
E		1	450	900	150	
F	5	0.5	200	400	100	
		2	200	350	50	

Cost table indirect - 125/day

1. F by 1 @ 50/day → After this F & D have same values, 4 & 8, hence both are critical.

2. F by 1 @ 50/day

3. F by 1 @ 50/day

4. B by 1 @ 150/day

5. B by 1 @ 150

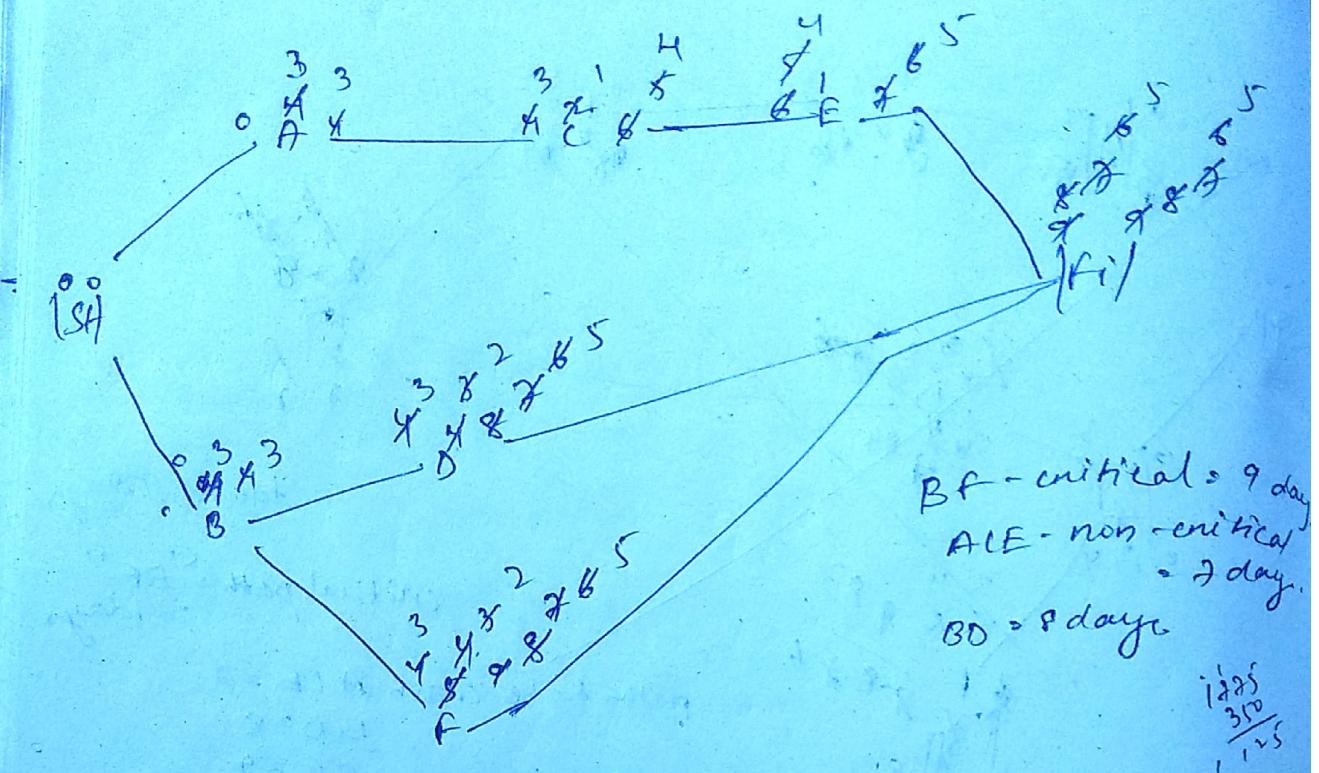
3.

Reducing B - 150

Reducing D & P - 200

so crash B.

9
8 + 50
7 + 150
6 + 225
5 → 350



1. Crash F by 1 @ 50/day

Now both D & F are critical - Crashing B or crashing D & F
 crashing B is cheaper than D & F combined.

2. B by 1 @ 150/day.

E, D, F end at the same time. So A, C, E is also critical now.

3. Possible crashes - DF along with ACE.

DF will be definitely crashed as both are critical.
 Form combinations of DF with A, C, E.

Crash DF & A by 1 - 225/-

DF & C - 350

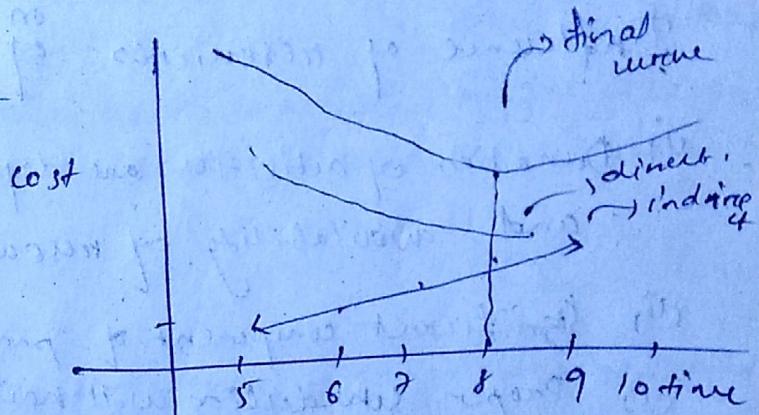
DF & E - 300

• Crash DFA by 1 @ 225/-

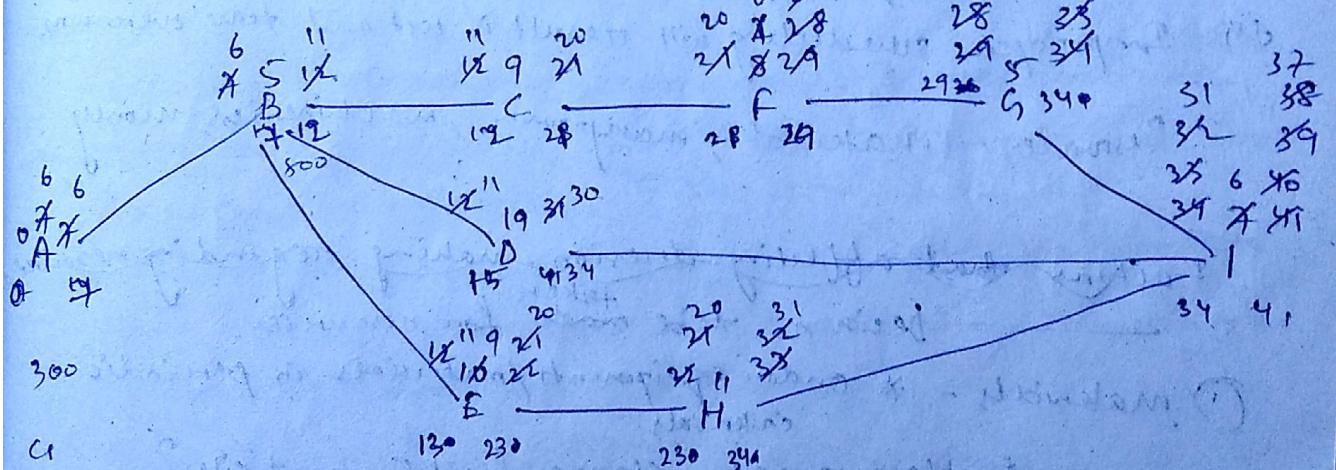
4. DCE & DF by 1 @ 350/-

Duration	Direct cost	Indirect cost (Column \times cost/day)	Total
9	1350	7125	2490
8	1400	1000	2400 ✓
7	1550	875	2425
6	1775	750	2525
5	2125	625	2750

8 days - optimum



2.



Critical path - none. ABCFHI

FBCGI - 41 days.

ABDI - 38 days

ABEHI - 40 days

Critical path = ABCFHI

= 41 days

Other path = ABDI = 38 days

ABEHD = 40 days

Critical path will change while crashing only when activities on the parallel lines are crashed.

Crash activities on ABCFHI.

1. I by 1 @ 200

2. F by 1 @ 250

BST E, H also critical.

Combinations possible - A, B, CE, CH, FB, FH
300 300 525 1025 400 850

3. A by 1

Procurement management

Failure mode analysis.

various possibilities of failure - probabilities of failure,
how severe is the failure.

Week-6 Resource Scheduling

Influence of resources ^{on} schedule.

(i) Duration of activities are dependant on the usage
and availability of resources.

(ii) Significant component of project cost.

(iii) Proper scheduling will have the impact on time and
cost of projects.

(iv) Improper scheduling will result in cost and time overruns.
Resources - material, manpower, machineries, money.

Factors affecting decision making regarding resource
decisions to be taken for resources.

① materials - * Order equipments/materials in periodic intervals.

* Have proper storage facilities and sites?

* Negotiate with multiple vendors. Large quantities
may give a discount?

* Custom delivery of equipment/delivery?

② manpower - * Skill requirements. Find highly-skilled
manpower for current work.

* Welders, masons, carpenters.

If skilled people are unavailable, no substitutes have to
be found or work has to be relocated.

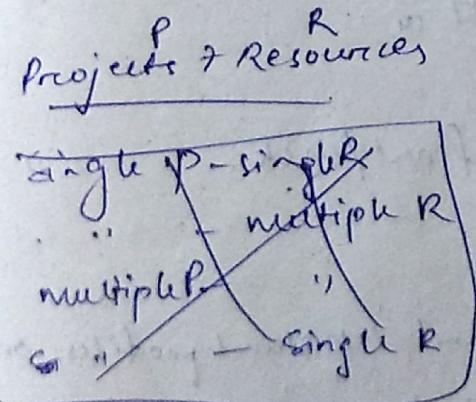
2. Machinery - When special equipment is needed on site,
- certain equipments like cranes, customized
crane, etc aren't easily or widely
available.

- * Preparatory work for equipment requirements.
 - For a ~~to~~ move a heavy-lift crane into a site, foundation
money has to be prepared.
- * Sharing equipment with other sites.

4. Money - * Cash flow prediction.
- Whatever amount of money is invested, it has
to be received back. There has to be returns for the invest-
ment. This is cash flow. ~~But~~ Lack of cash flow will screw
up the project.

- * Credit planning.
 - Borrowing money for project progress. Proper planning should
be done in order to avoid negative cash flow.

* To visually understand requirement of resource per
each activity per day, gantt chart, Histogram and other
Gantt charts are required. Time and cumulative
resource per day will vary with early start and late
start times. All the bars and charts show cumulative
resources over ~~over~~ time.



single P - multiple R
multiple P - multiple R
multiple P - single R
single P - multiple R
large companies

In the gantt chart, for every activity 2 resources are
required. Resources are added up separately.

In histogram, for every activity/day two bars will be plotted
to show two diff resources.

Exercise.

Equipment resource

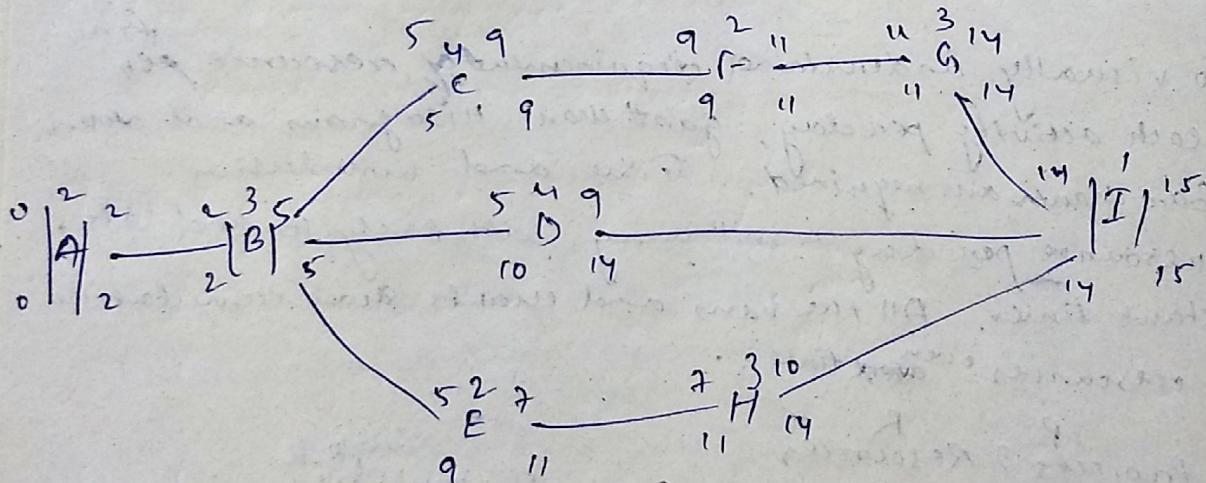
Activity	Preceded by	Duri ⁿ trucks	
		2	2
A			
B	A	3	1
C	B	4	6
D	B	4	4
E	B	2	4
F	C	2	2
G	F	3	2
H	E	3	1
I	D, G, H	1	1

resource.

1. Plot histogram.

critical - ABCFGI.

2. What if only 10 trucks are available?



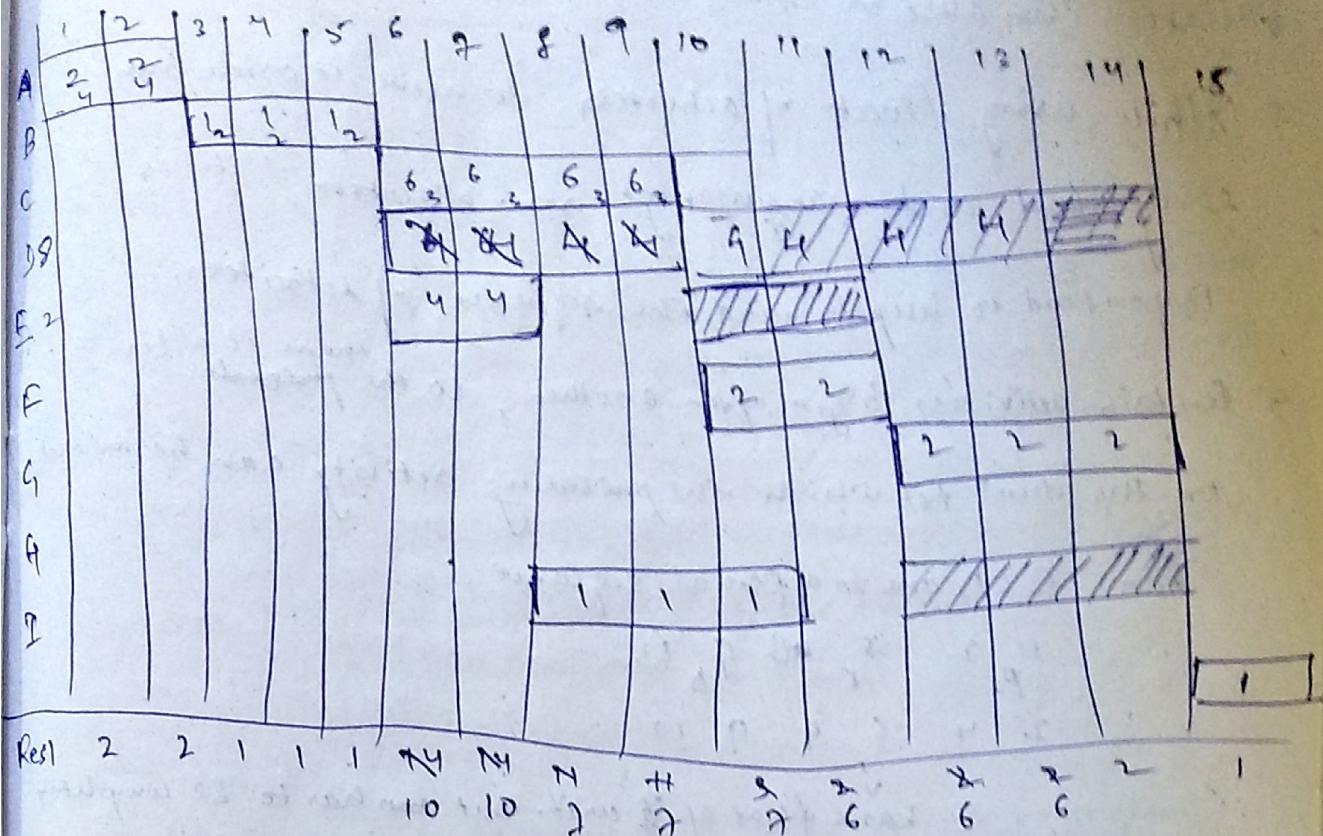
* Critical activities do not have float/slack.

Here - A, B, C, F, G, I.

To have separate early start and late start profiles, only D, F and H can be varied in days.

For the graph - 111118 → Late start profile of activities after considering their floats.

* Over allocation of resources can be resolved by curing float of an activity.



Moved D out of C's space.

This is required because the question states a constraint of 10 trucks available. As the sum is 10, no day can have resources exceeding that value.

For days 6, 7, 8, 9, 10 resources summed up to 14 and 11.

As D has float time, it was moved out to bring resources within constraint. Two possible solutions with D more stand from day 10 or 11.

Cash as a Resource

Total value/cash required for an activity is given for the entire duration - maybe days/months/years. Find a value for a single day/month/year. Divide it equally.

How to manage cash overflow?

By there is always a constraint upon usage of resource (total) for each day. If other resources allocated to activities exceed the given limit, it's called over-allocation of resources. As the resource here is only cash, even-

allocation can also be called cash overflow.

* While using float of activities to move it outside from early start and manage over-allocation, it is important to keep in mind the sequence of activities.

* Certain activities begin after another, so there is a limit on the day till which the preceding activity can be moved.

For BCD are in a serial sequence.

B 2 5 8 9, 13
2 4 6 9 13

has a float of 3rd unit. But C has to be completed before 9th day i.e. before early start of D.

* Here, there is only one decision variable — how much to move each activity?

Two Resources in a single Project

Actv	Preed	Dur ⁿ	Truck R-1	Workers R-2
A	-	2	2	4
B	A	3	1	2
C	B	4	6	3
D	B	4	4	8
E	B	2	4	2
F	C	2	2	3
G	DF	3	2	8
H	E	3	1	1
I	DGH	1	1	3

S - 10 truck, 15 workers constraint

* For multiple resources, same method of chart is followed. Draw alloc' charts separately for each of the resources. Then draw a single histogram for both. * Then look for over-allocation.

* Use float of activities to find a balanced use and allocation of both resources so that overall allocation of both resources neither does not exceed limit.

Resolving over-allocation of one resource might ~~result~~ lead to over-allocation of the other.

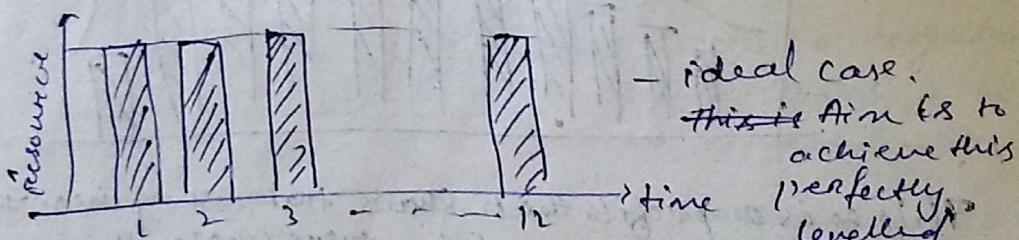
Levelling Resources

* Requirements Resource level requirements on projects.

* Simplifying resource profile.

Levelling approaches - minimum movement algo.

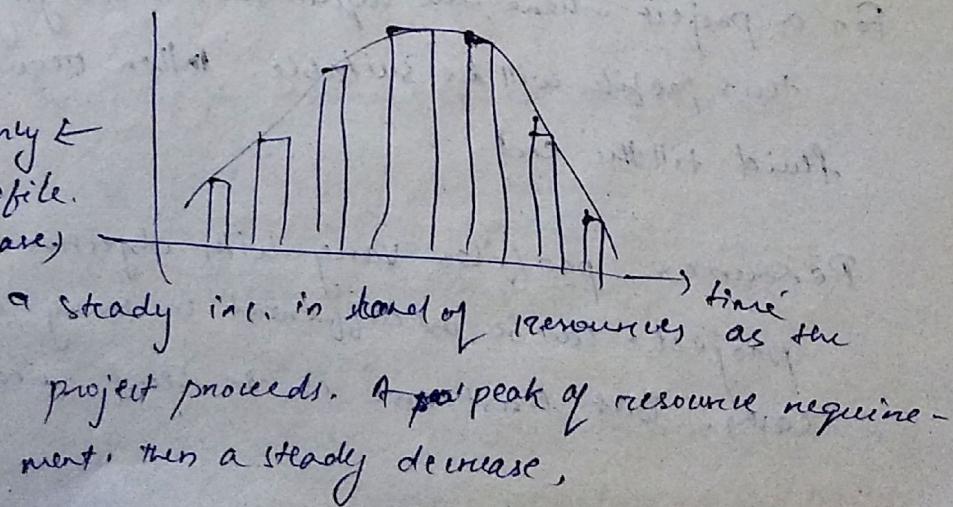
✓ Perfectly leveled resource / ~~Ideal case~~ case - Fixed amount of resources for the entire duration.



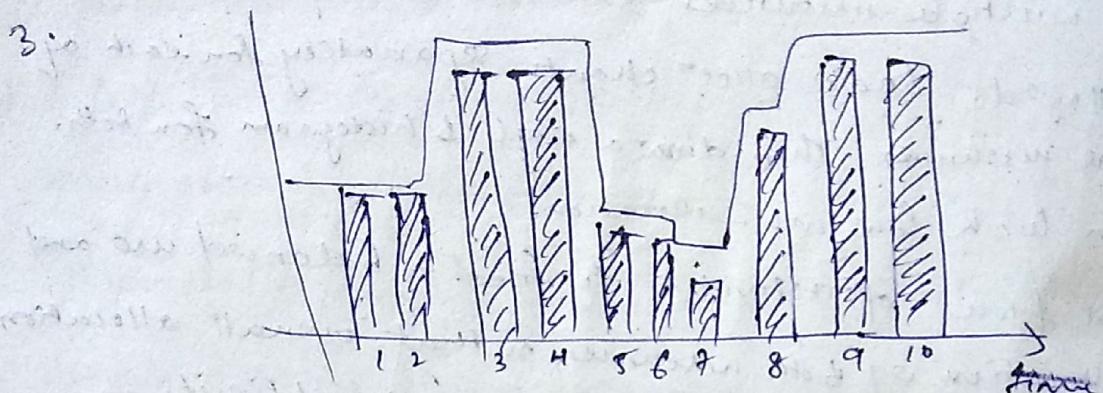
- ideal case.
This is aim is to achieve this perfectly leveled resource histogram.

✓ another possible histogram -

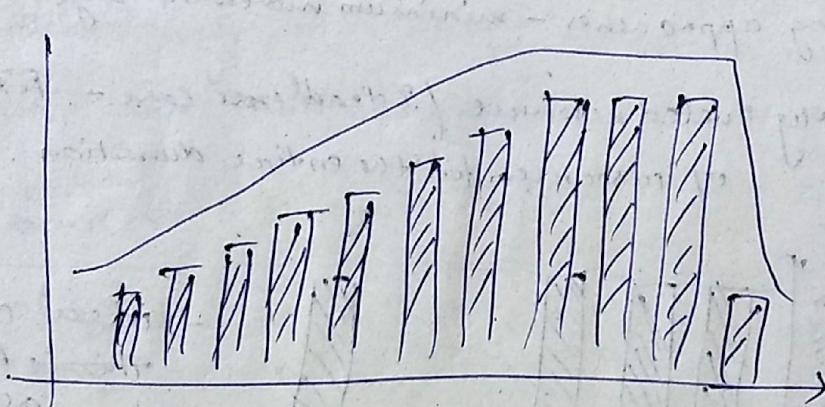
A very commonly observed profile.
(A practical case)



* Steady inc. in demand of resources as the project proceeds. A peak of resource requirement, then a steady decrease,



This can be resource profile or progress graph of a project where weather/climatic conditions affect the working. Let's say it's for each month. In the initial months of a year - Jan/Feb, ~~there are~~ weather condn isn't very favourable. Then around March and April, weather is favourable and project progress is high. Proceeding to summer season, project progress reduces further.



This is a profile that shows that use of resources increases and stays there constantly reaching a peak point for several till the project completes. For a project where the requirement is changing continuously, this profile will be suitable. When requirements are fluid till the end.

Resource profiles vary with objective & demand of project. Can be as different in many different cases. But first one is the ideal case.

Reasons behind resource leveling - manage over-allocation.

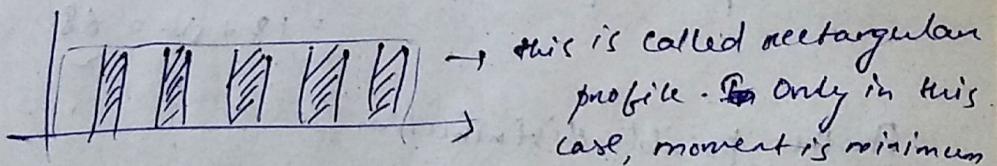
→ Apart from resolving over-allocation, resource leveling is required to remove abrupt random variations in need of resources. The variation over time should be smooth and not abrupt. Hence, to achieve that also resource leveling is required.

→ Resource leveling can be done by exploiting the float of non-critical activities. Forecast the possible shifts in the " " , find the best combi combo, that results in total usage of resources, which looks leveled and stays within constraints.

Minimum Moment concept

→ A mathematical method to find out the leveled profile.

→ The moment of the shape about horizontal axis is minimum only when the profile of the resources is rectangular.
The ideal case of leveled resources is -



→ This is called rectangular profile. Only in this case, moment is minimum.

Objective - Given any profile of resources, ~~converting~~ bring it as close as possible to this rectangular shape, so that there is less variation in use of resources.

Moment of a resource histogram is calculated as -

$$M = \sum (Y * Y/2) \text{ on } x\text{-axis.}$$

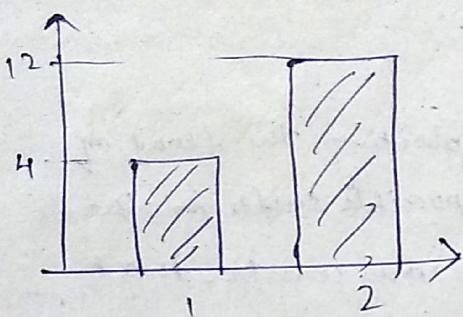
Sum of $(Y * Y/2)$ of all bars given in a histogram.

Y is the height of the bar.

In a resource histogram, x -axis denotes duration of project in days/months, y -axis denotes units/no. of

resources utilized i.e. weight of a bar in the histogram denotes that.

lets say, a resource profile of 2 days that utilizes 16 resource units is as below.



1 \rightarrow 4 units

2 \rightarrow 12 units

Moment of the profile

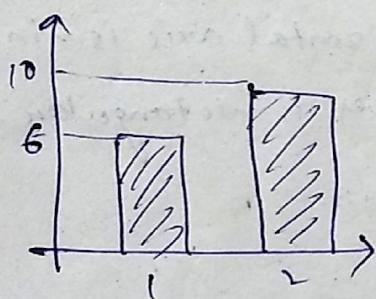
$$= (4 \times 4/2) + (12 \times 12/2)$$

$$= (4 \times 2) + (12 \times 6)$$

$$= 8 + 72 = 80$$

The same profile is modified as below

the same no. of resource units with a different distribution



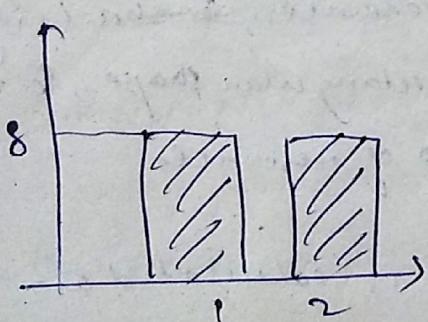
1 \rightarrow 6 2 \rightarrow 10 units

$$M = (6 \times 6/2) + (10 \times 10/2)$$

$$= (6 \times 3) + (10 \times 5)$$

$$= 18 + 50 = 68$$

Another possible distribution -



1 - 8 2 - 8

$$M = (8 \times 8/2) + (8 \times 8/2)$$

$$= 2 \times 8 \times 4$$

$$= 64$$

Hence the minimum moment

Minimum moment is obtained in the third case, where resources are equally utilised on both days. With the same no. of resource units and days, moment can vary widely with distribution and utilising.

From the above it is also observed that rectangular profile truly gives the least moment.

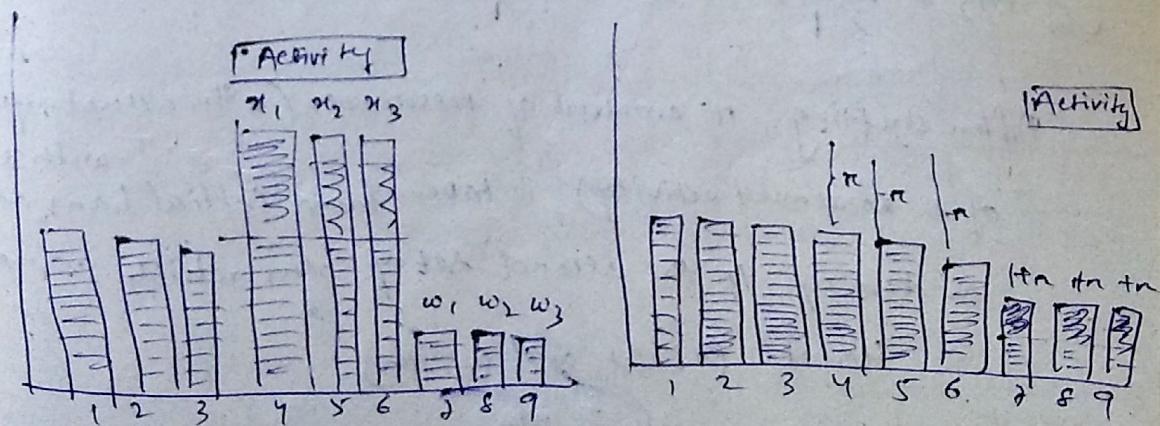
- * The concept is to reduce the moment of histogram in a systematic way and that profile obtained is the most leveled and as closest to possible to a rectangular profile.
- * Possible procedure to calculate the optimum level -

i) Find all possible combinations of floats/shifts of non-critical activities. Find moment for each of one of them. And figure out the min. moment. ~~But this~~ But this method gets tedious when there are higher no. of resources and complex networks of activities.

ii) An improvement function.

Each time an activity shifts towards late start profile by certain days/time units, whether there is an inc. in moment of histograms or a decrease in moment (i.e. improvement) in histogram.

→ denotes shifts/extrn.



When activity x_n happened on days 4, 5, 6 there was a certain extra load of resources on that day. Moving it to days 7, 8, 9 curing its float has created a certain kind of balance. Extra of day 4, 5, 6 was removed and added up to day - 7, 8, 9.

x_i - resource level from which resources are moved
for m days.

w_i - resource level to which the m days of resources
are moved.

m - no. of days when resources are removed/added.
~~(Resource of activity affecting those resources)~~
~~not necessary to shift the entire activity.~~ x_i - initial level w_i - final level of a histogram
 r - resource load of the activity.

For Histogram 1,

$$\checkmark M_1 = \frac{1}{2} \sum_{i=1}^m (x_i - r)^2 + \frac{1}{2} \sum_{i=1}^m (w_i - r)^2 \quad \left[\begin{array}{l} \text{moment} = \frac{\sum x_i w_i}{2} \\ = \frac{r^2}{2} \end{array} \right]$$

* Each bar
Before shifting, the initial bars which had extra resources
are denoted by x_1, x_2, \dots, x_m . And the days which
lacked resources are w_1, w_2, \dots, w_m .

For histogram 2,

$$\checkmark M_2 = \frac{1}{2} \sum_{i=1}^m (x_i - r)^2 + \frac{1}{2} \sum_{i=1}^m (w_i + r)^2$$

After shifting, r amount of resources (the actual requirement
of the concerned activity) is taken from initial bars and
added up to the second set of bars which have resources
Hence $(x_i - r)$ and $(w_i + r)$.

* Now compare M_1 , $\neq M_2$.

$M_1 = M_2 \rightarrow$ shifting did not bring up any benefit.

$M_1 > M_2 \rightarrow$ shifting increased the moment. Should not
be shifted then.

$M_1 < M_2 \rightarrow$ shifting reduced the moment. Hence,
desirable.

$$IF_{A,d} = n \left(\sum_{i=1}^m w_i - \sum_{i=1}^m w_i - mn \right)$$

$IF_{A,d}$ = improvement factor for shifting activity A, d days out of time.

(If > 0, implies reduction in moment due to shift.)

m = min. # days that the activity is shifted, & n = "deem" of activity

w_i = daily resource sum of activity for the current time frame, over which resources will be removed.

w_i = daily resource sum of activity for the time frame over which resources will be added.

n = total daily resource load for the activity.

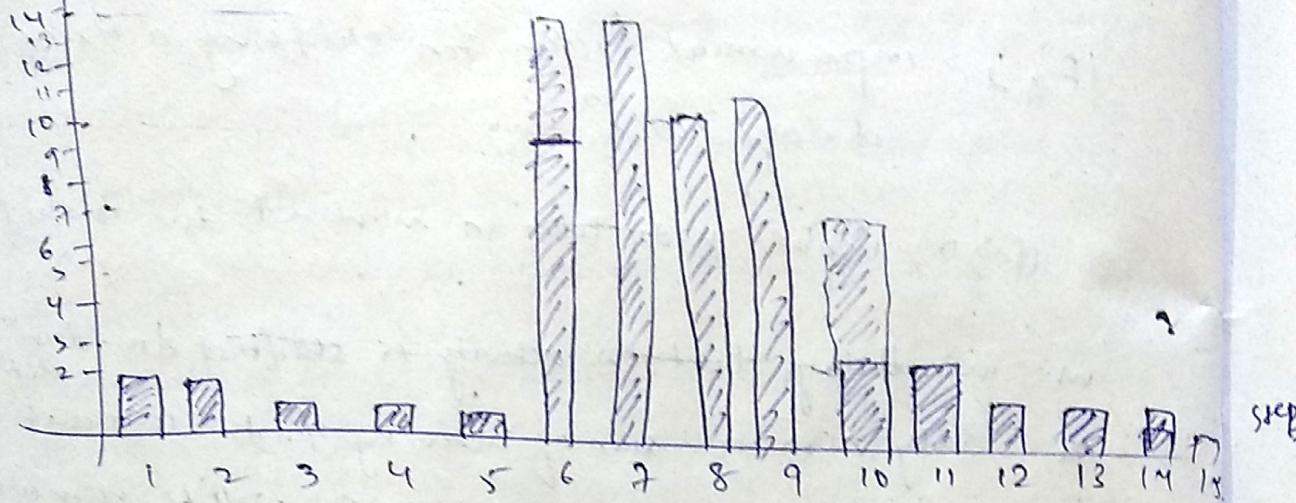
Ex- Problem on minimum moment concept.

A	2	2												
B			3	1	1	1								
C							6	6	6	6				
-D								8	4	4	4	4		
-E									4	4				
F														
G														
-H														
I														
	Days	1	2	3	4	5	6	7	8	9	10	11	12	13
	res	2	2	1	1	1	10	14	10	11	7	3	2	2

If Activities D, E, H can be shifted. Calculate moment at each step.

Start - (i) shift D by 1.

$IF(D, 1) =$



Step-18 Shift D by f.. It will move from Day 6 to Day 10.

γ_j = resource at Day 6 = 14

$$w_1 = w - u - v \rightarrow 10 > 3$$

$$n = 4$$

m = 1

$$\textcircled{1} f(3,1) = 4(\cancel{-6} - 10 - 4 \times 1) = 4(-14 - 4) \\ \rightarrow 4 \times 2 = 28 > 0$$

moment reduced. This means move B by f. Now it

Ranges from Day ¹⁰⁰ 7 to Day 10. Da

change in histogram - Day 6 = 70 (earlier 14)
Day 10 = 7 (earlier 3)

Step - 2

$$\text{① } \cancel{f(0,2)} = f(3,2)$$

$$r=9, m=2, x_1=14, x_2=14, w_1=3 w_2=3$$

$(+9)$	$(+4)$	$(0-10)$	$(0+14)$
$0-6$	$0-7$		

$$f(5,2) = 4((4+4) - (3+3) - 2 \times 4)$$

$$= 4(28 - 6 - 8) = 4 \times 14 = 56 > 0$$

\therefore moment reduced again. As value is higher than
b. step 1, this is a better option.

shifting by 3 days - $n=4, m=3, w_1=4, w_2=3, w_3=2$
~~Step 3~~ $w_1=3, w_2=3, w_3=2$
 010 011 012

$$1F \ (5 \times 3) = 4((14+10+10)-(3+3+2)-4 \times 3)$$

$$= 4 \{ 38 - 8 - 12 \} = 4(38 - 20) = 4 \times 18 = \underline{\underline{72 > 0}}$$

∴ moment reduced even better.

Step 4 shifting D by 4 days - $n=4, m=4, \alpha_1 = 14, \alpha_2 = 14, \alpha_3 = 10$
 $m_4 = 11$

$$\overline{PF}(D,4) = \overline{f}^4(-38)$$

$$\rightarrow \cancel{4} \left[(14+14+10+11) - (3+3+2+2) - (4 \times 14) \right]$$

$$= 4 \left[\cancel{59} - 10 - \cancel{56} \right] = 4 \times (59 - 66) \\ = 4 \times -7 = \underline{-28} \quad 0$$

\therefore Not a ~~g~~ desirable stuff.

$$DP(5,4) = 4 \left[(14+14+10+11) - (3+3+2+2) - (4 \times 4) \right]$$

$$= 4[49 - 10 - 16] = 4[49 - 26] = 4 \times 23 = \underline{\underline{92}}$$

Even better.

Step-5 shifting Δ by 6 days

$IF(5,5) = RE=4, \boxed{m=5}, n_i = 14, 14, 10, 11, w_i \geq 3 \text{ 222}$
 $m=4, \text{ will remain } 4.$

Because 4 days of resources are being shifted. Dur" of O is 4 days, it assigned resources for 4 days. Hence m can never exceed total dur" of an activity.

$$IF = 4 \left(319 - \frac{9}{100} - 16 \right) = 4(49 - 25) = 4 \times 24 = \underline{\underline{96}} > 0$$

\therefore Best option.

Hence, ~~we~~ ~~try~~ to level the resource histogram closer to a rectangular profile. It should be moved by 5 days.

Shifting Activity H

1. $IF(H, 1) = n=1, m=1, x_1=4, w_1=2$

$$M_f = 1(4-2-1) = 11-3 = \underline{8} > 0$$

2. $IF(H, 2) = n=1, m=2, x_1=11, 11, w=3, 2$

$$1. [(11+11)-(3+2)-2 \times 1] = 1[22-15-2] \\ \rightarrow 22-17 = \underline{\cancel{5}} > 0$$

3. $IF(H, 3) = n=1, m=3, x_1=11, 11, 3, w=13, 13$

$$[(11+11+3)-(2+13+13)-3] \\ = 25-28-3 = 25-31 = \underline{-6} < 0 \rightarrow \text{moment increased.}$$

∴ option 1 ~~was~~ is the best. Shifting H by 1 will give best result.

Shifting Activity E

1. $IF(E, 1) = n=4, m=1, x_1=14, w_1=10$

$$4(14-10-4) = 4(0) = 0$$

2. $IF(E, 2) = n=4, m=2, x_1=14, 14, w=10, 11$

$$\therefore 4(28-21-8) = 4(28-29) = 4 \times -1 = \underline{-4} < 0 \\ \text{moment increased.}$$

3. $IF(E, 3) = n=4, m=2, x_1=14, 14, w=11, 3$

$$4(28-14-8) = \cancel{28}(4(28-22)) = 4 \times 6 = \underline{24} > 0$$

4. $IF(E, 4) = n=4, m=2, x_1=14, 14, w=3, 3$

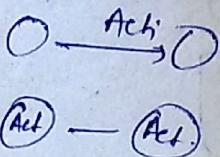
$$4(28-6-8) = 4(28-14) = 4 \times 14 = \underline{56} > 0$$

∴ moving E by 4 days will give least moment.

Week 7

Precedence Diagramming method (PDM) - Module-2

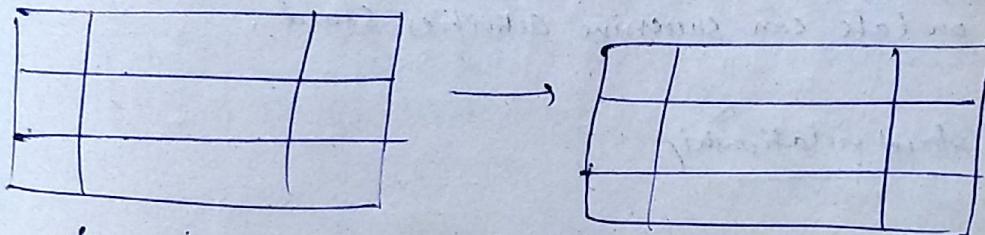
AOA \Rightarrow Activity on Arrow
AON \rightarrow " " Node



\Rightarrow PDM utilises event-based approach of AOA but flow of AON.
 \Rightarrow It is a popular method of representation in almost every scheduling method as given by project management software.

PDM Notation

1. AON Type Notation.



These boxes will contain information about activities. Activity name, duration, ~~float~~ and ~~may be resources info like~~ early & late start time, early & late finish time, float.

2. 4 types of relationships:

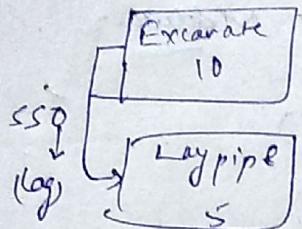
- Start to start "
- Finish to start " ✓ - so far only this representation is being used. Notation shows sequence as end of activity
- Start to finish "
- Finish to finish "

3. Leads and lags.

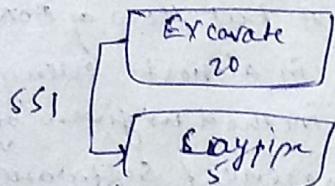
This shows the gap b/w start or finish of 2 following activities in a sequence.

Finish to start with 2 day lag.
 that means, an activity will start 2 days after the preceding activity has finished its execution.

Start to start



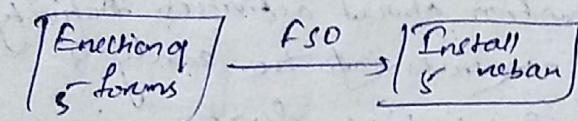
Zero lag: Lay pipe cannot start before Excavate starts.



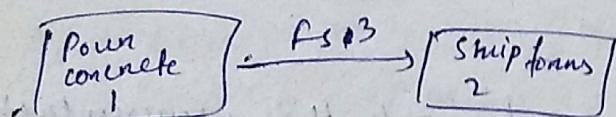
Positive lag: Lay pipe can start no earlier than one day after excavation starts.

Here the diagram shows start of activities only. And how early or late can successive activities start.

Finish to start relationship

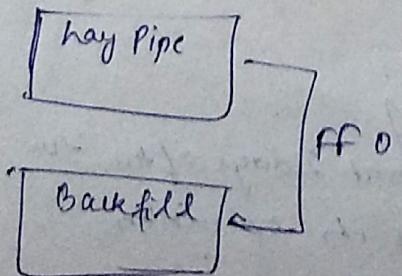


Zero lag: Installation of forms can start only after the erection of forms (1st activity) is completed.



Positive lag: Strip forms activity will start 3 days after completion of pour concrete activity. Cannot start work for 3 days after completion of Activity 1, then start Act 2.

Finish-finish relationship



Zero lag: Backfill (Act 2) can be completed only after Lay pipe (Act 1) is completed.

Act 2 cannot be completed before Act 1 is completed.

Lay pipe

Backfill

FF1

Positive lag: Backfilling can be completed no earlier than 1 day ~~before~~ after Lay pipe activity is completed.

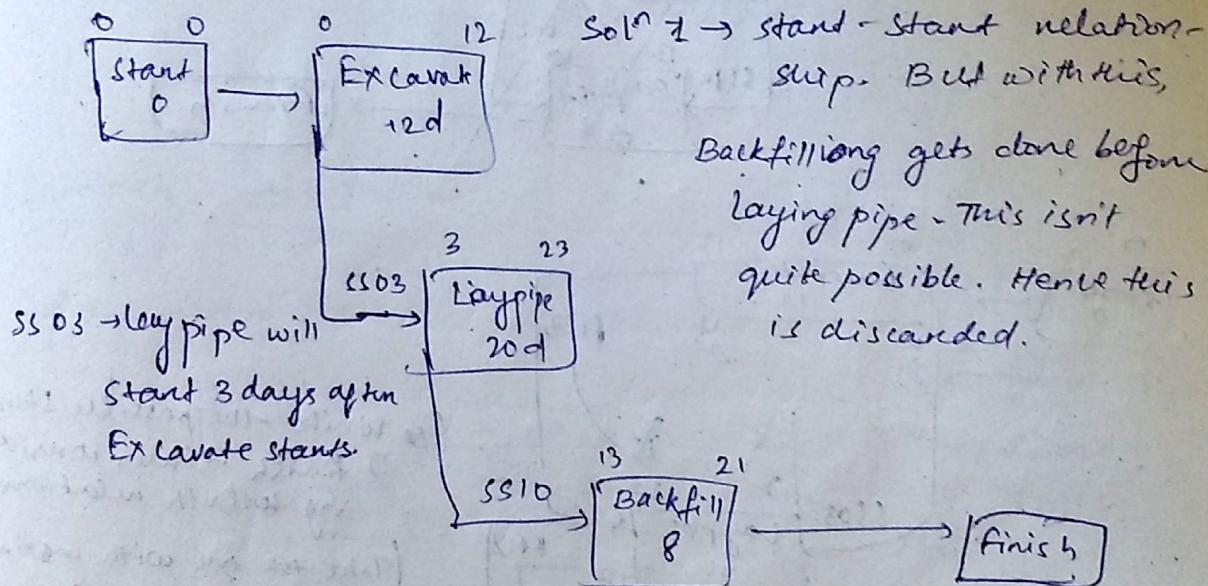
~~It can be finished only at 1 day after laying~~

Alt 2 can be finished 1 day after Alt 1 is finished. Not before 1 day.

+ → delay (lag)

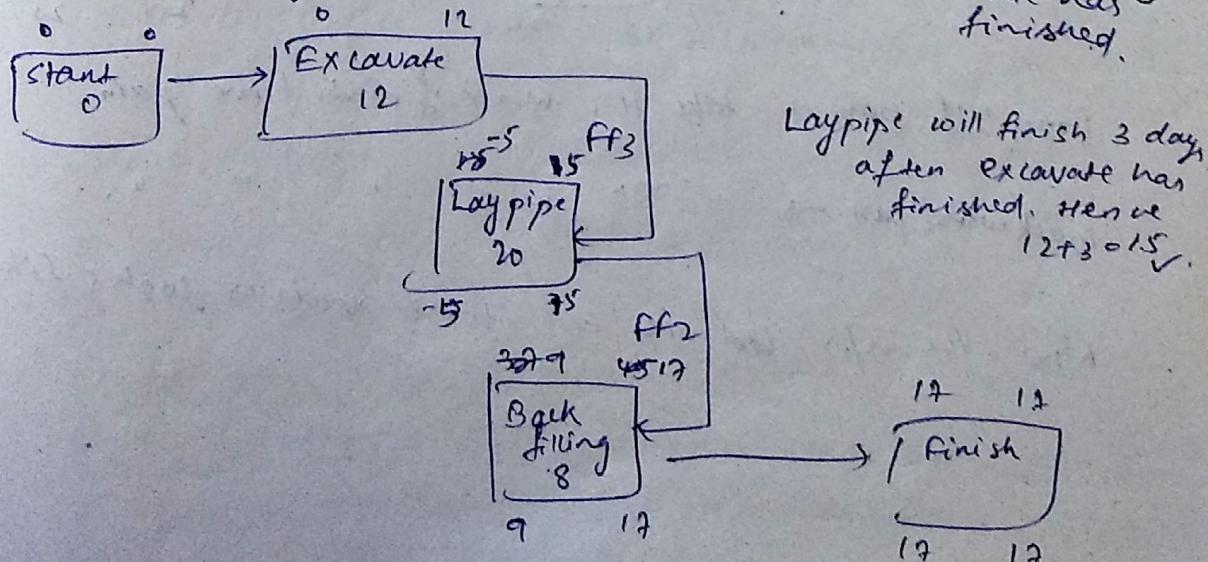
- → lead

PDM analysis



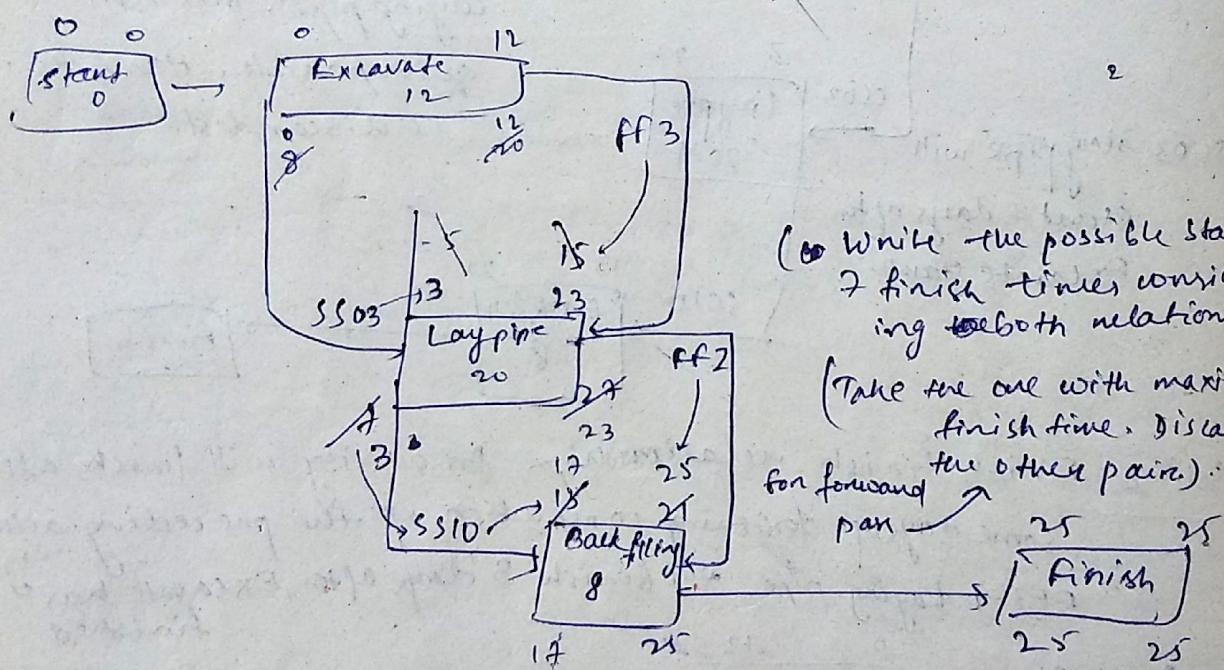
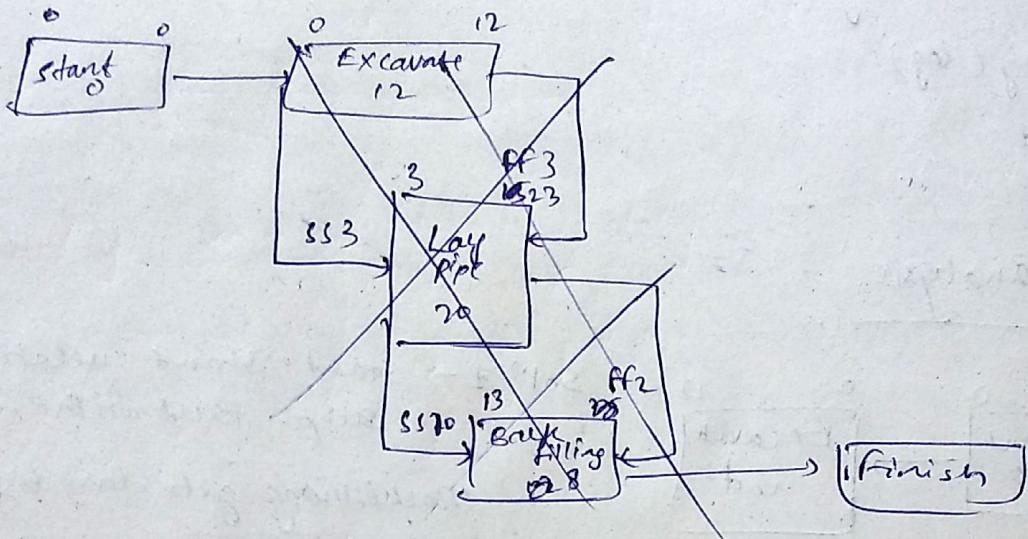
Solⁿ 2 Finish-finish relationship. An activity will finish after some days of finishing completion of the preceding activity.

FF3 → Laying pipe will finish 3 days after Excavate has finished.



As the early finish of Lay pipe is negative which is practically not possible, FF relationships failed.

Solⁿ 3 → Using both SS & FF relationships. - Multiple relationships on activities.



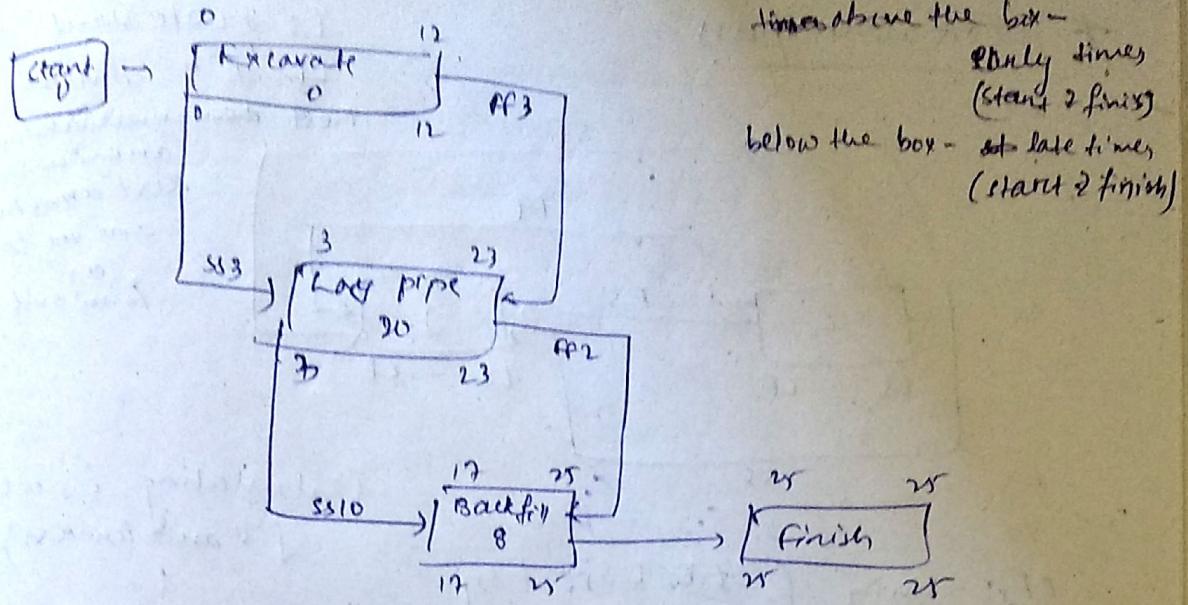
(*) Write the possible start & finish times considering both relationships.

(Take the one with maximum finish time. Discard for forward pass →)

Forward pass → take the max finish time pair.

Backward pass → " " min " " "

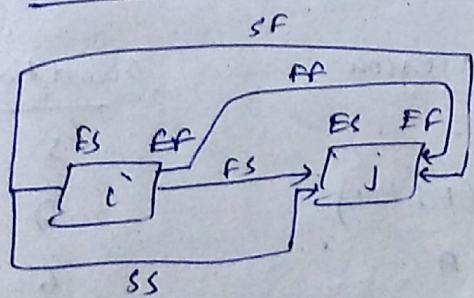
Now the new Gantt chart with all time values looks like -



* PDAM may require multiple relationships between activities. But they should form a loop.

Network Calculations

* Forward pass



$$ES_j = \max \left\{ \begin{array}{l} E_{Sj} + FS_{ij} \quad (\text{FS relationship}) \\ E_{Si} + SS_{ij} \quad (\text{SS relationship}) \\ E_{Fi} + FF_{ij} - D_j \quad (\text{FF relationship}) \\ E_{Si} + SF_{ij} - D_j \quad (\text{SF " }) \end{array} \right.$$

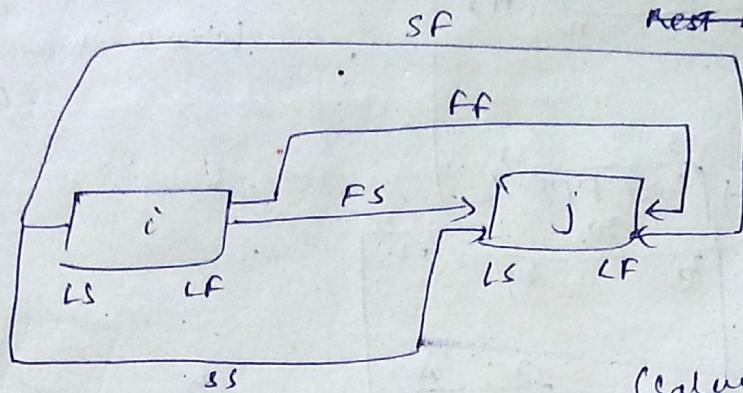
$$EF_j = ES_j + D_j$$

FS_{ij} , FF_{ij} , SF_{ij} , SS_{ij} are the lead/lag values for the different possible relationships.

times above the box -
early times
(start & finish)
below the box - late times
(start & finish)

D_j → duration of activity j
 ES → early start j
 EF → " finish
 SS → start start
 SF → " finish
 FF → finish "
 FS_{ij} → " - start .
lead/lag of
activity i, j

* Backward pass



LS → Late start
LF → " finish.

Rest annotations
are the same.
Rest annotation
are the same
as forward pass.

(calculating 'i's when
j's are known).

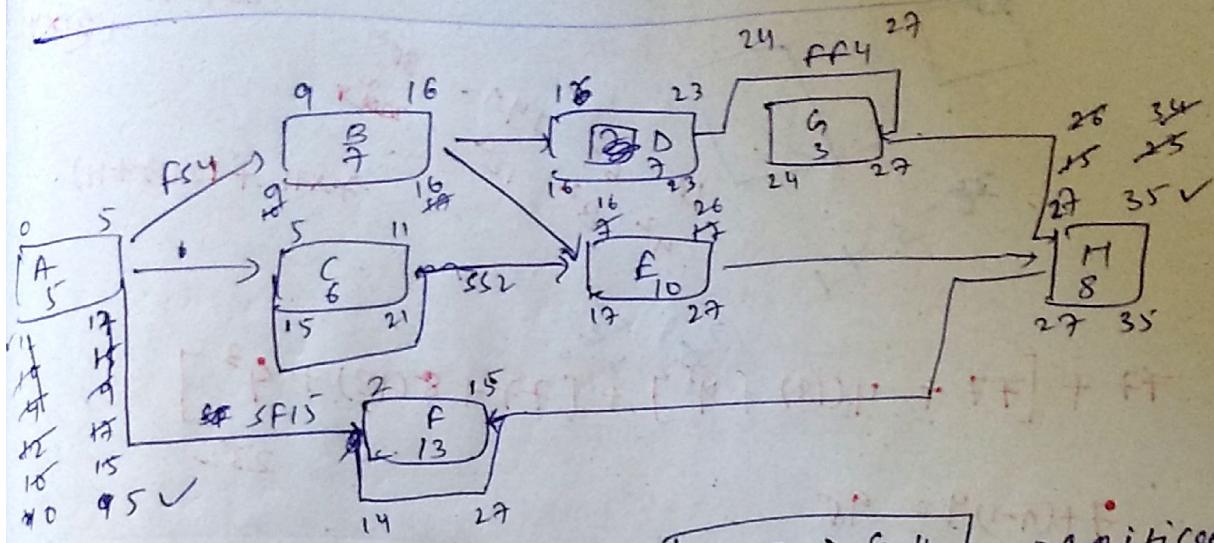
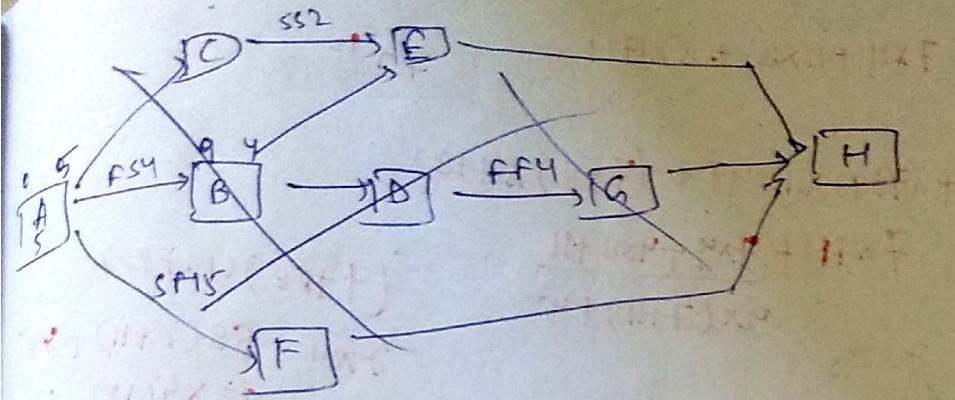
$$LF_i = \min \left\{ \begin{array}{l} LS_j - FS_{ij} \\ LF_j - FF_{ij} \\ LS_j - SS_{ij} + D_i \\ LF_j - SF_{ij} + D_i \end{array} \right.$$

$$LS_i = \max LF_i - D_i$$

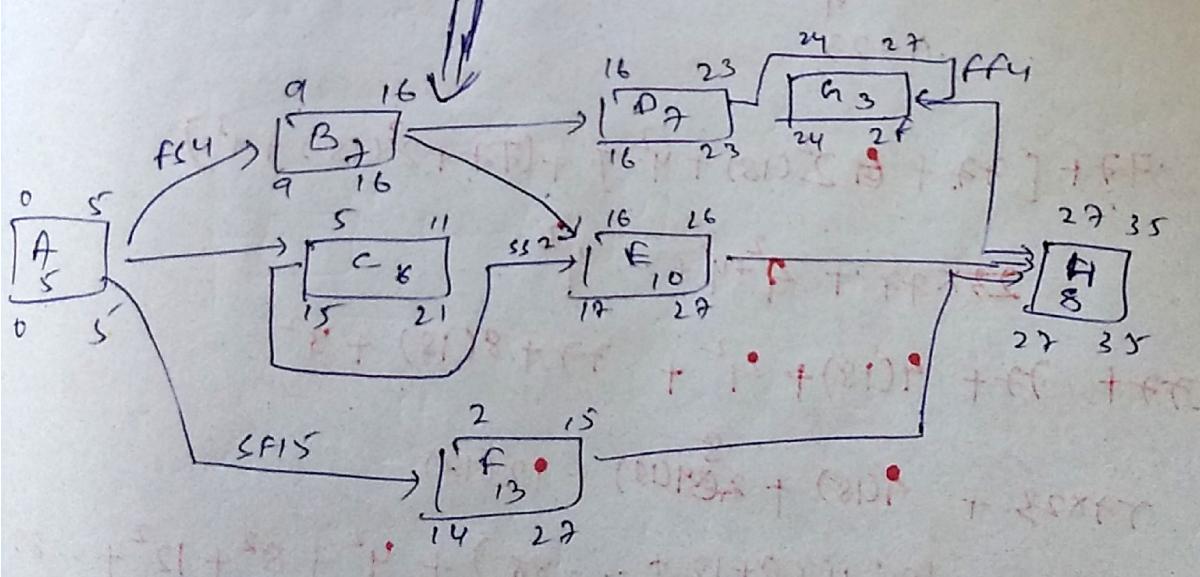
Problem

<u>Activity</u>	<u>Predecessor</u>	<u>Duration</u>
A	-	5
B	A (FS+4)	7
C	A	6
D	B	7
E	C(SS+2), B	10
F	A(CF+15)	13
G	D(FF+4)	3
H	E, F, G	8

Q - Draw PDM network. Find ES, EF, LS, LF & critical path.

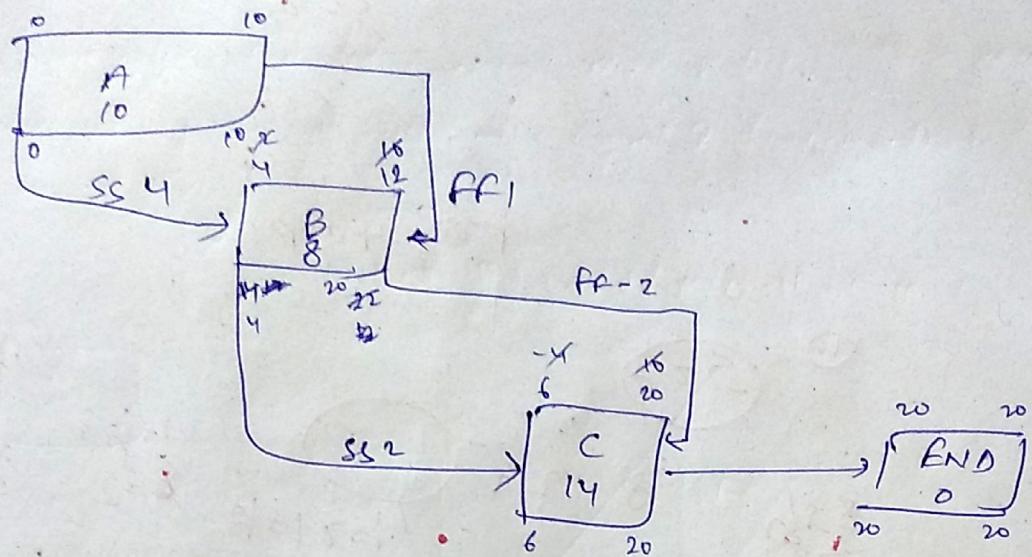


[A B D G H] → critical path



final network

ADM network

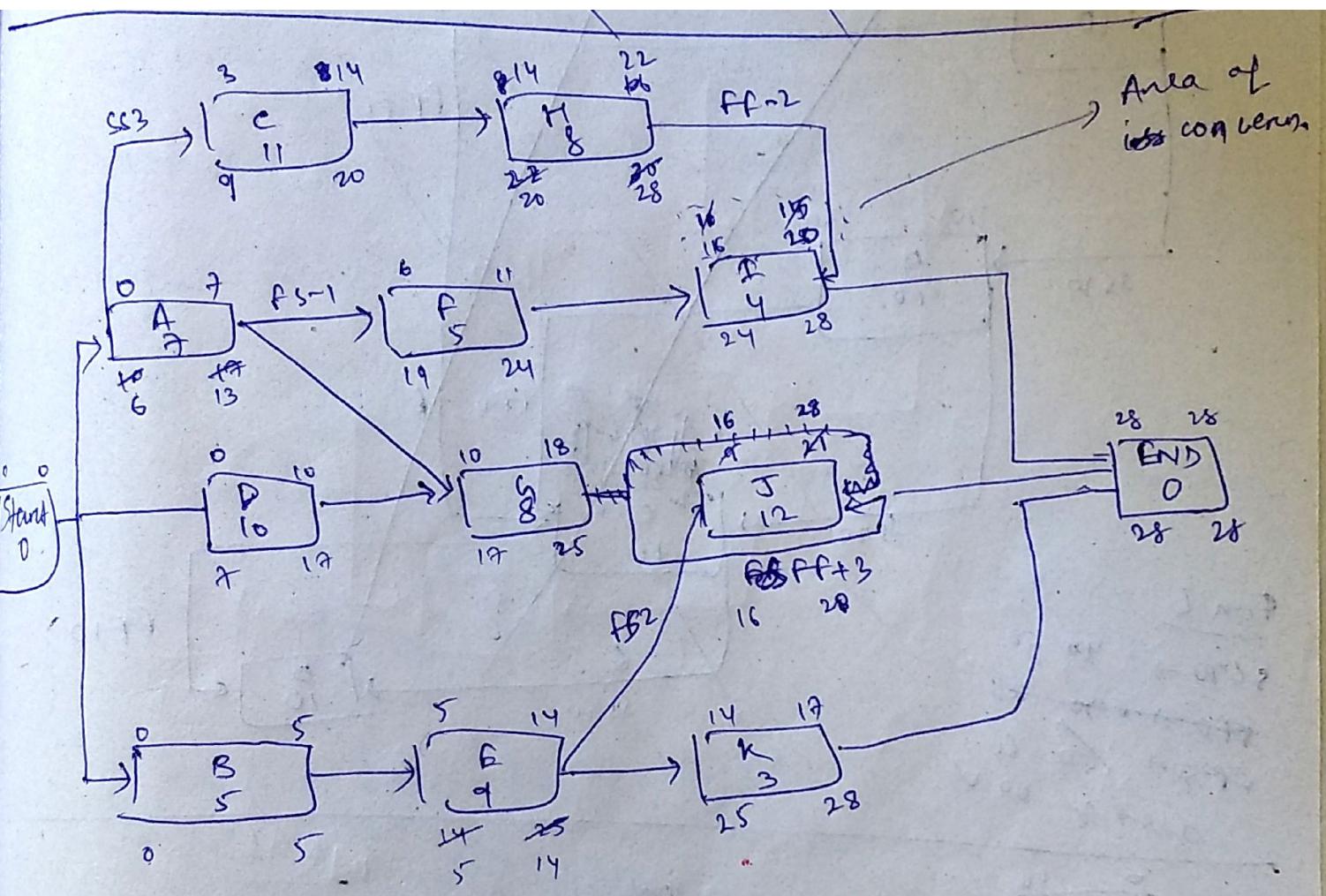


Due to negative lag, there is an inconsistency in the late start & finish times of B. Taking FF-2 from C, late start of B comes as 22.

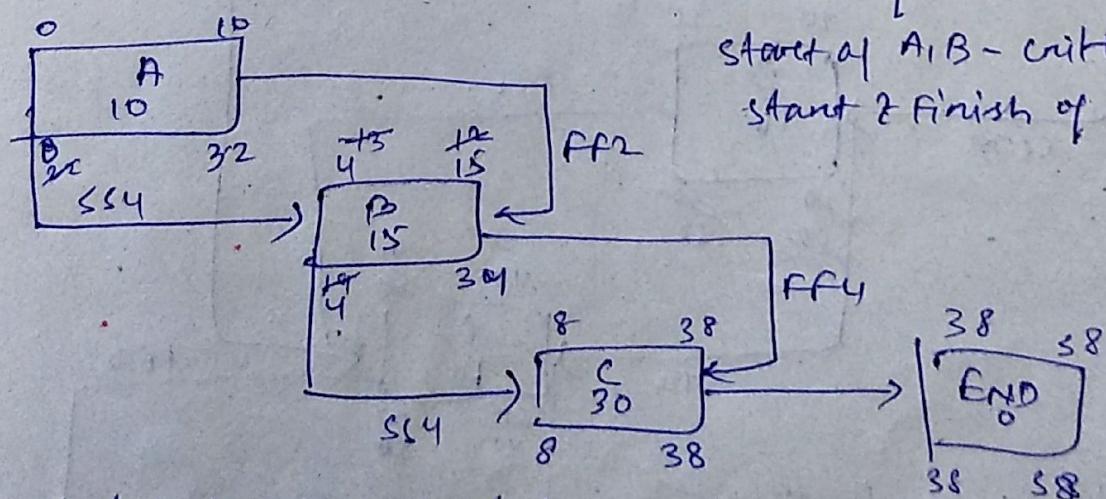
But as the entire dueⁿ of project is 20, 22 isn't possible. So truncated to 20. With SS2 netⁿ, early start of B = 4. This creates an inconsistent situation.

Situation. B, ~~late start = 4~~ duration = 8, ?? Not quite possible.
" finish = 20

<u>Activity</u>	<u>Predecessor</u>	<u>Duration</u>
Start	-	0
A	Start	7
B	"	5
C	A(SS+3)	11
D	Start	10
E	B	9
F	A(FS-1)	5
G	A, D	8
H	C	8
I	F, H(FF-2)	4
J	F(FS+2)	12
K	G(FF+3)	3
End	I, J, K	0



Another Example.



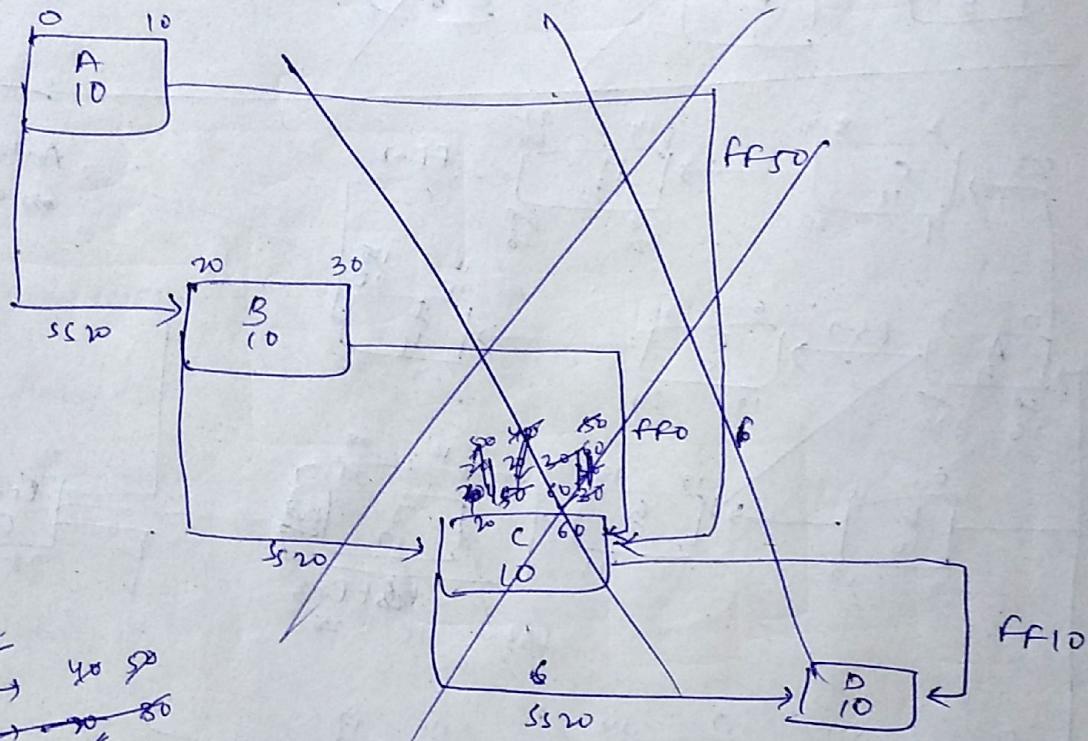
finish of A,B - flexible
start of A,B - critical
start & finish of C → within
crit.

Hence, we calculate the late start & finish times of activities taking both relationships. Then we are considering min start & max finish time (late). This results in critical start times but flexible finish times.

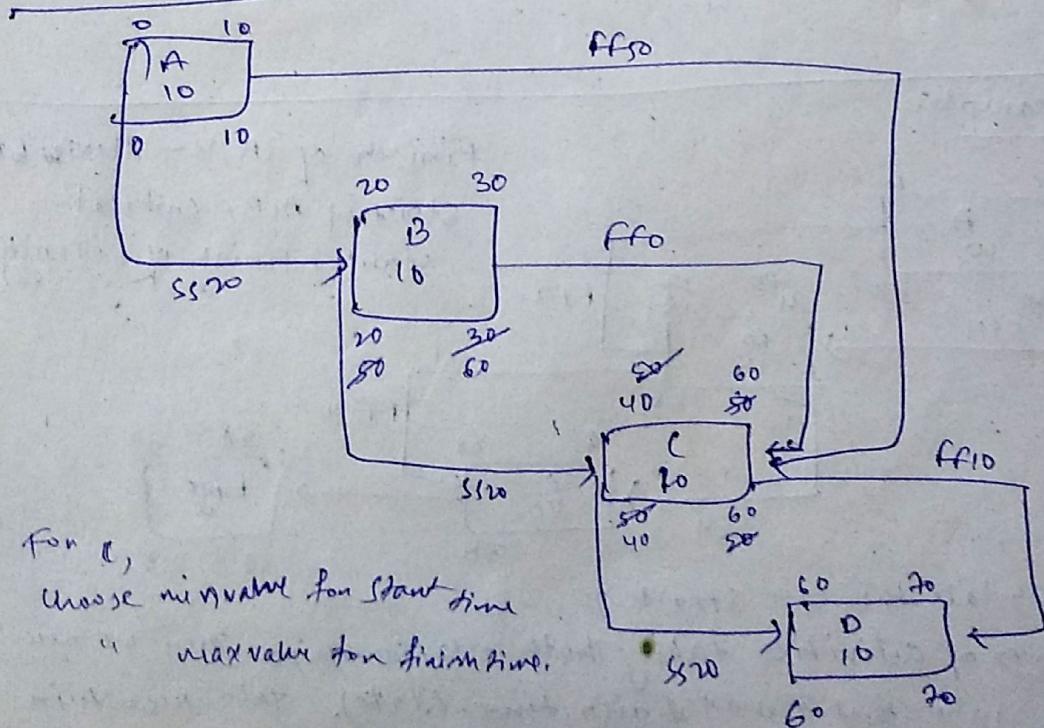
~~min of both start~~ When you have a couple of start & finish time pair options, take; for the starts, choose min value, for finish choose max value. This ~~will~~ min-max pair will form the start & finish times for the activity.

26 30
50 60

where Extract(month from "Date") = 6



For C
 $SS20 \rightarrow 40\ 50$
 $FF0 \rightarrow 20\ 80$
 $FF50 \rightarrow 60\ 60$
 ans $\rightarrow 20\ 60$



For C,
 choose min value for start time
 max value for finish time.

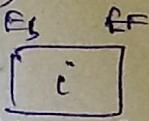
With this, Activity A is critical, Start of B is critical, C & D are critical.

Alternate floats are defined for PDM - start float, finish float, Relation float, Activity float.

~~Floats in PDM (Non-continuous deer)~~

1. Task Float

$$\text{i) total float } (i) = \text{LF}(i) - \text{ES}(i) - \text{Dur}(i)$$

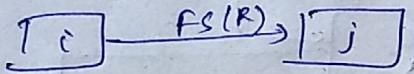


$$\text{ii) start float } (i) = LS(i) - ES(i)$$

$$\text{iii) finish float } (i) = LF(i) - EF(i)$$

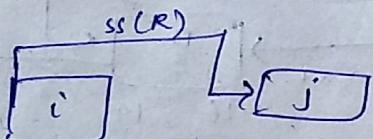
2. Relationship Float

i) FS



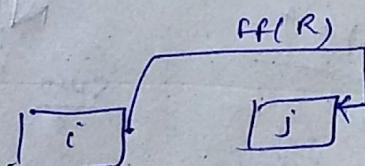
$$RF(FS) = LS(j) - EF(i) - \text{lag}(R) \checkmark$$

ii) SS



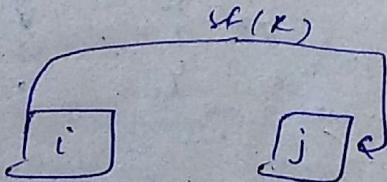
$$RF(SS) = LS(j) - ES(i) - \text{lag}(R) \checkmark$$

iii) FF

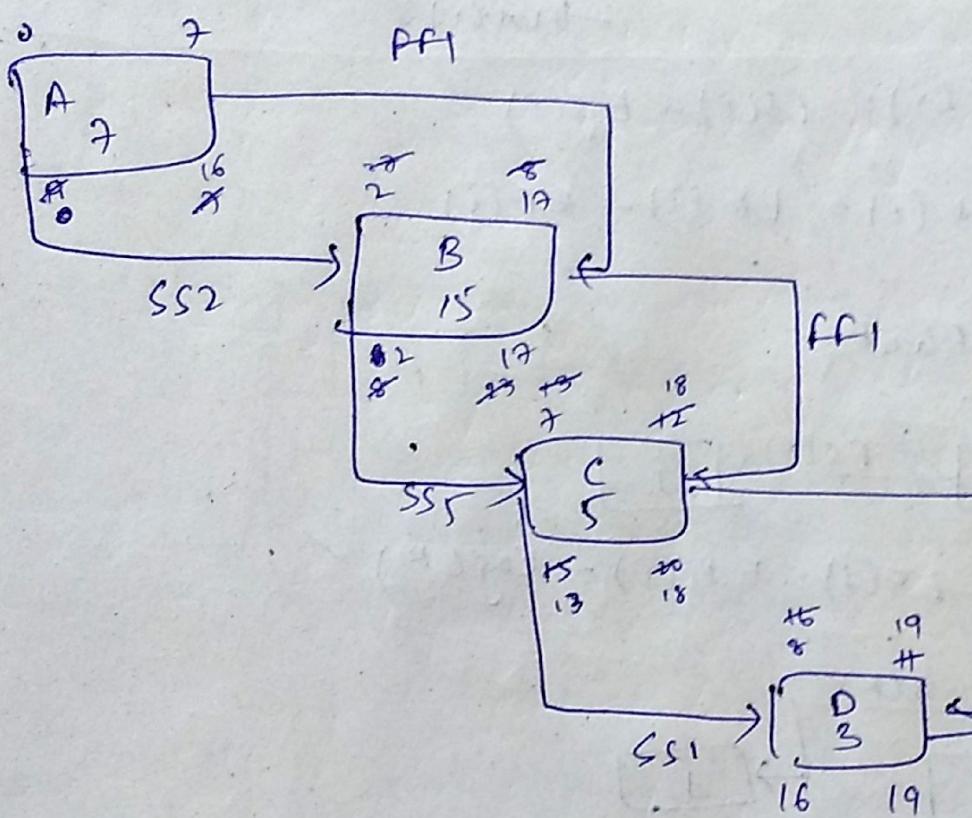
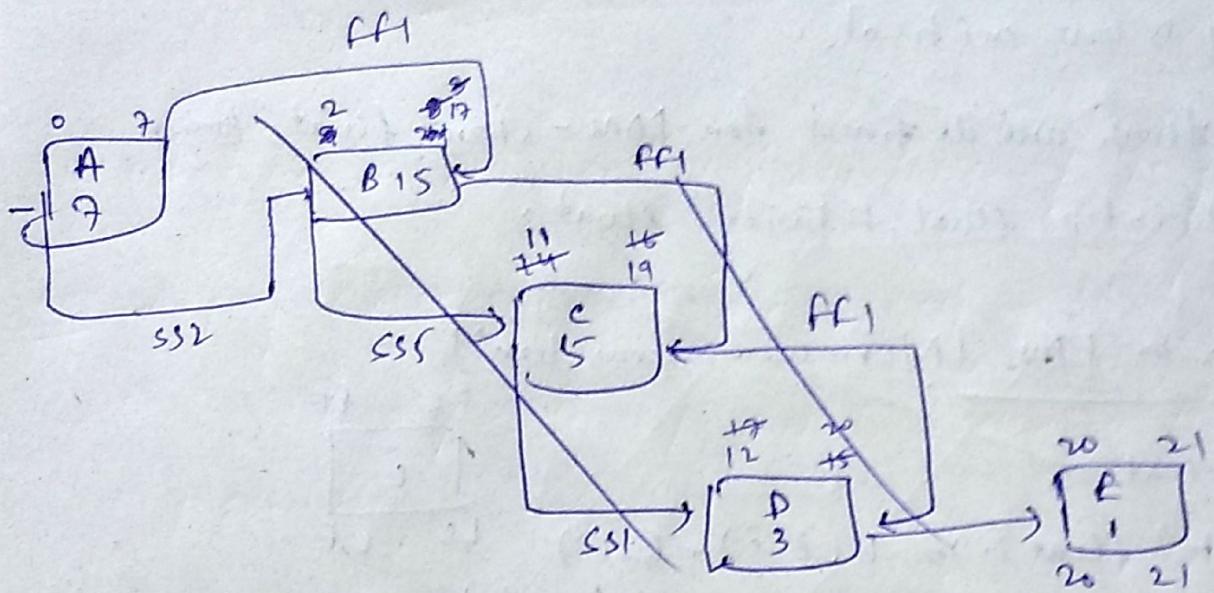


$$RF(FF) = LF(j) - EF(i) - \text{lag}(R) \checkmark$$

iv) SF



$$RF(SF) = LF(j) - ES(i) - \text{lag}(R) \checkmark$$



while ($m \sim$)

$i = s$

if ($i < n$)
 append(i)

else if ($i = n$)
 append(i)

$i = i + 1$