Simple Configuration

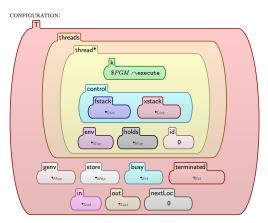
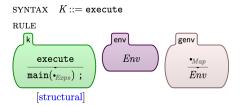
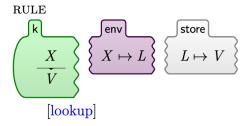


Fig. 1. The K configuration of SIMPLE

Starting execution

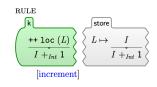


Lookup



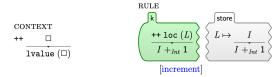
Increment





Ivalue described in the next slide

Increment



- Ivalue described in the next slide
- Wrap increment Exp into an auxilliary Ivalue construct.
- Once Ivalue evaluates to Iocation loc(L), perform increment.

Increment (contd.)

$$\begin{array}{ll} \text{SYNTAX} & \textit{Exp} ::= \text{lvalue} \ (\textit{K}) \\ \\ \text{SYNTAX} & \textit{Val} ::= \text{loc} \ (\textit{Int}) \\ \\ \hline \\ \text{RULE} \\ \hline \\ \text{lvalue} \ (\underbrace{X}_{\text{loc} \ (\textit{L})}) \\ \hline \\ \text{[structural]} \end{array}$$

Arithemic Operators

Boolean Operators

```
RULE \frac{V[N1, N2, Vs]}{V[N1][N2, Vs]} [structural, anywhere]
```

$$\frac{\texttt{array}\;(L,-)[N]}{\texttt{lookup}\;(\check{L}+_{Int}N)}\\ [\texttt{structural}, \texttt{anywhere}]$$

```
 \begin{array}{ccc} \text{RULE} & & \text{RULE} \\ V[NI\,,\,N2\,,\,Vs] & & & \text{array}\,\,(L,-)[N] \\ \hline V[NI][N2\,,\,Vs] & & & \text{lookup}\,\,(L+_{Int}\,\,N) \\ & & & \text{[structural, anywhere]} \end{array}
```

• Lookup in an array - how are expressions handled in array offset?

```
RULE V[N1, N2, Vs]
V[N1][N2, Vs]
[structural, anywhere]
```

```
RULE \frac{\text{array }(L,-)[N]}{\text{lookup }(L+_{Int}N)} [structural, anywhere]
```

- Lookup in an array how are expressions handled in array offset?
- what is anywhere?

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RULE V[N1, N2, Vs]
V[N1][N2, Vs]
[structural, anywhere]
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RULE \frac{\text{array }(L,-)[N]}{\text{lookup }(L+_{Int}N)} [structural, anywhere]
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- Lookup in an array how are expressions handled in array offset?
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RULE \frac{V[N1, N2, Vs]}{V[N1][N2, Vs]} [structural, anywhere]
```

$$\frac{\texttt{array}\;(L,-)[N]}{\texttt{lookup}\;(L+_{Int}\;N)} \\ [\texttt{structural}, \texttt{anywhere}]$$

```
 \begin{array}{lll} \text{RULE} & & \text{RULE} \\ V[N1\,,\,N2\,,\,Vs] & & & \text{array}\;(L,-)[N] \\ \hline V[N1][N2\,,\,Vs] & & & \text{lookup}\;(L+_{Int}\;N) \\ & & & \text{[structural, anywhere]} \end{array}
```

• Lookup in an array - how are expressions handled in array offset?

```
RULE

\frac{V[N1, N2, Vs]}{V[N1][N2, Vs]}

[structural, anywhere]
```

```
\begin{array}{c} \text{RULE} \\ & \underbrace{\text{array} \; (L, -)[N]}_{\text{lookup} \; (L +_{Int} \; N)} \\ & [\text{structural, anywhere}] \end{array}
```

- Lookup in an array how are expressions handled in array offset?
- what is anywhere?

Array Lookup (contd.)

```
\begin{split} \text{SYNTAX} \quad & Exp ::= Int \mid Bool \mid String \mid Id \\ & \mid (Exp) \text{ [bracket]} \\ & \mid ++ Exp \\ & \mid Exp \text{ [Exrs] [strict]} \end{split}
```

Array Lookup (contd.)

```
\begin{split} \text{SYNTAX} \quad & Exp ::= Int \ | \ Bool \ | \ String \ | \ Id \\ & | \ (Exp) \ [\text{bracket}] \\ & | \ ++ Exp \\ & | \ Exp[Exps] \ [\text{strict}] \end{split}
```

SYNTAX
$$Exps ::= List\{Exp, ", "\}$$
 [strict]

Array Lookup (contd.)

```
\begin{split} \text{SYNTAX} \quad & Exp ::= Int \mid Bool \mid String \mid Id \\ & \mid (Exp) \text{ [bracket]} \\ & \mid ++ Exp \\ & \mid Exp \text{ [Exrs] [strict]} \end{split}
```

SYNTAX
$$Exps ::= List\{Exp, ", "\}$$
 [strict]

RULE
$$\frac{\left(\operatorname{lookup}\left(L\right)\right)}{\operatorname{loc}\left(L\right)}$$

$$\underbrace{\left[\operatorname{structural}\right]}$$

Array Size

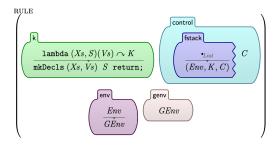
Size of an array

The size of the array is stored in the array reference value, and the **sizeOf** construct was declared strict, so:

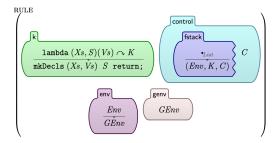
RULE

$$\frac{\texttt{sizeOf}\;(\texttt{array}\;(--,N))}{\check{N}}$$

Function Calls

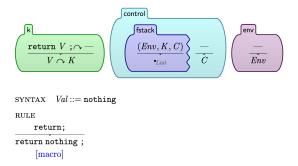


Function Calls

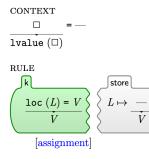


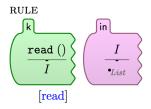
• Switch to *global environment*, where free variables in function are looked up.

Function Calls



Assignments & Reads





Sequential Composition, Expressions & Conditionals

RULE
$$\frac{S1 \quad S2}{S1 \quad \sim S2}$$
[structural]
$$\frac{V ;}{\bullet_{K}}$$

$$\frac{\text{RULE}}{\text{if } (\text{true})S \text{ else} --}{\overset{\bullet}{S}}$$

$$\frac{\text{RULE}}{\text{if (false)} - \text{else } S}$$

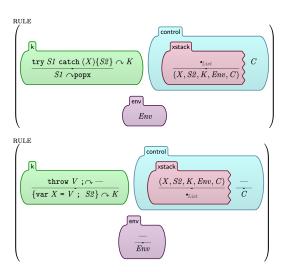
While & Print

```
RULE  \begin{array}{c} \text{print} \ (\underbrace{V,Es}) \ ; \\ \text{[print]} \\ \\ \text{RULE} \\ \\ \text{print} \ (\underbrace{V_{Valls}}) \ ; \\ \\ \vdots \\ \\ \text{Structural]} \\ \\ \text{RULE} \\ \\ \text{while} \ (E)S \\ \\ \text{if} \ (E)\{S \ \text{while} \ (E)S\} \\ \\ \text{[structural]} \\ \\ \text{[structural]} \\ \end{array}
```

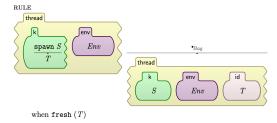
Exception Handling

```
ListItem ::= (Id, Stmt, K, Map, Bag)
SYNTAX
           K ::= popx
RULE
                                               control
                                                   xstack
      \mathtt{try}\;S1\;\mathtt{catch}\;(X)\{S2\}\curvearrowright K
                                                   (X, S2, K, Env, C)
             S1 →popx
                                        Env
RULE
                      RULE
                                           xstack
                           popx
                             •K
                                               •List
```

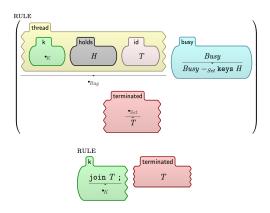
Exception Handling (contd.)



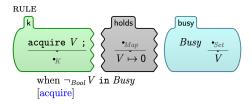
Spawn

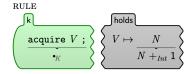


Termination & Joining

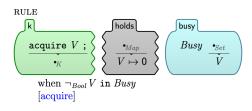


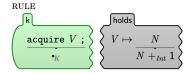
Acquire Locks





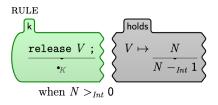
Acquire Locks

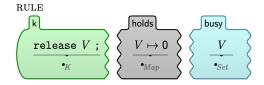




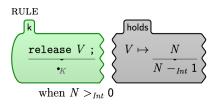
• Assume re-entrace. Same thread can acquire a lock again.

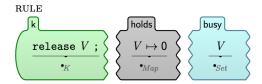
Release Locks





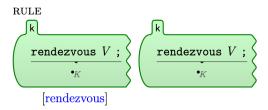
Release Locks





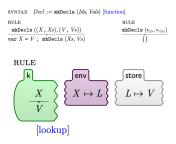
 A lock is considered release only when n release calls match with n-acquires calls.

Rendezvous



Auxilliary Constructs

Declarations, Lookups & Restoring Environments



Auxilliary Constructs

Restoring Environments

