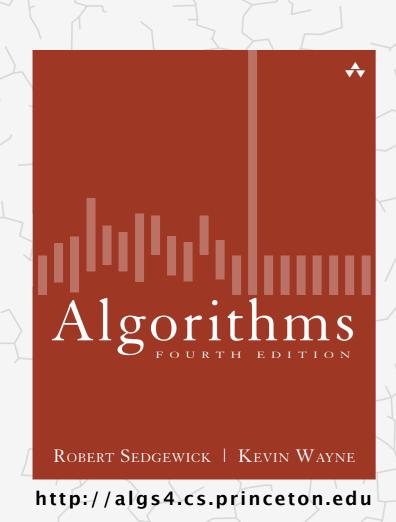
# Algorithms



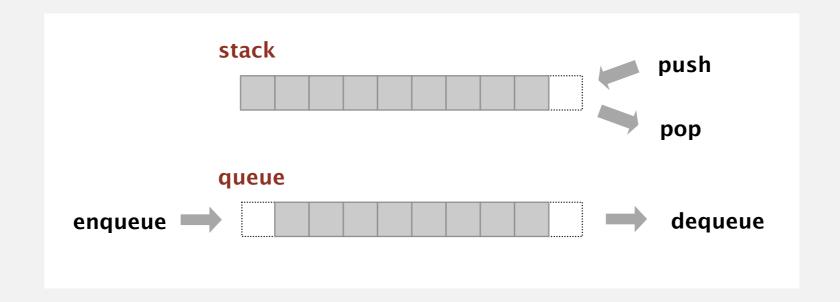
## 1.3 BAGS, QUEUES, AND STACKS

- stacks
- resizing arrays
- queues
- generics
- iterators
- applications

## Stacks and queues

#### Fundamental data types.

- Value: collection of objects.
- Operations: insert, remove, iterate, test if empty.
- Intent is clear when we insert.
- Which item do we remove?



Stack. Examine the item most recently added. ← LIFO = "last in first out"

Queue. Examine the item least recently added. ← FIFO = "first in first out"

### Client, implementation, interface

#### Separate interface and implementation.

Ex: stack, queue, bag, priority queue, symbol table, union-find, ....

#### Benefits.

- Client can't know details of implementation ⇒
   client has many implementation from which to choose.
- Implementation can't know details of client needs ⇒
  many clients can re-use the same implementation.
- Design: creates modular, reusable libraries.
- Performance: use optimized implementation where it matters.

Client: program using operations defined in interface.

Implementation: actual code implementing operations.

Interface: description of data type, basic operations.

# Algorithms

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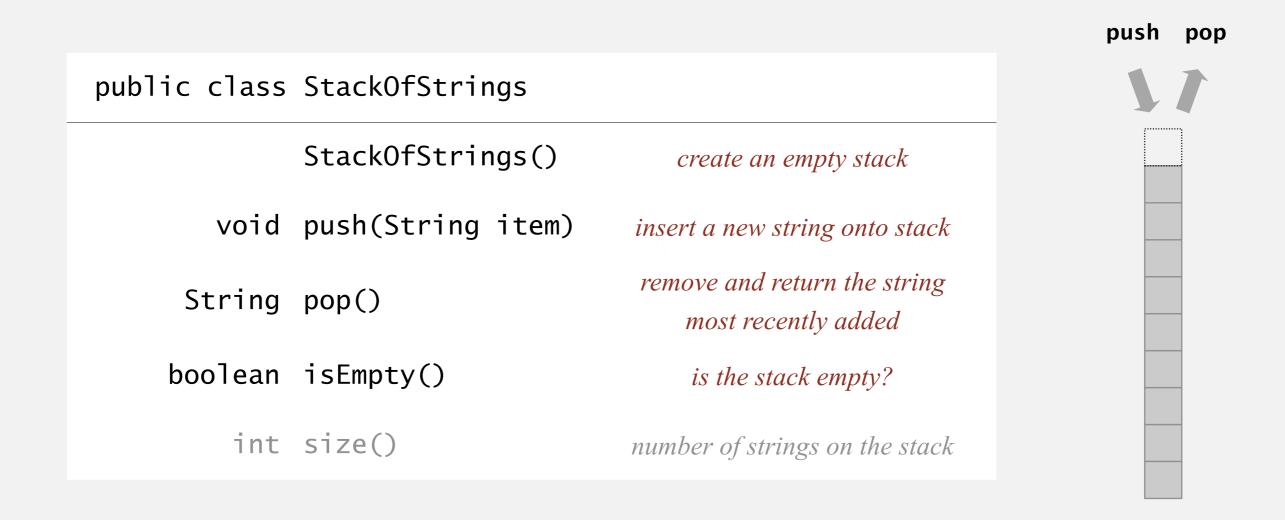
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# 1.3 BAGS, QUEUES, AND STACKS

- stacks
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#### Stack API

Warmup API. Stack of strings data type.

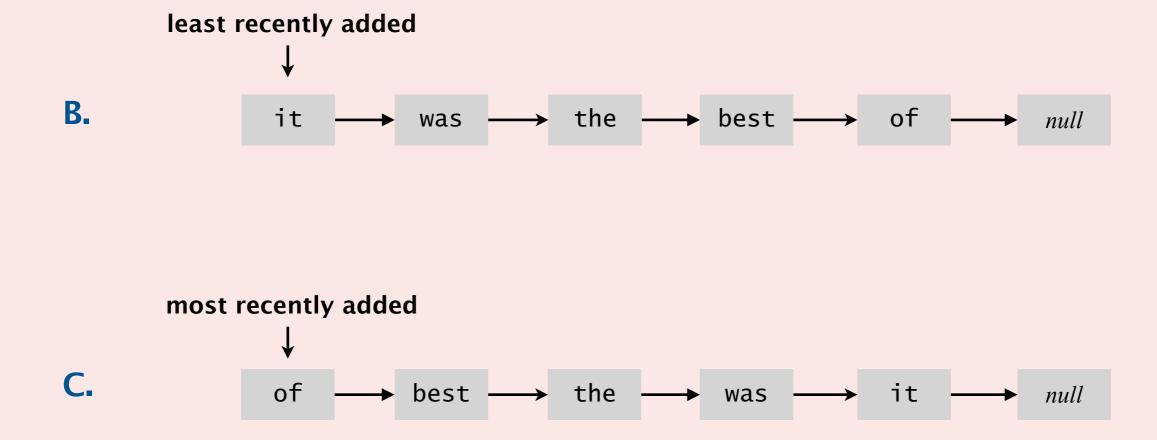


Warmup client. Reverse sequence of strings from standard input.

## Stack quiz

### How to implement a stack with a singly-linked list?

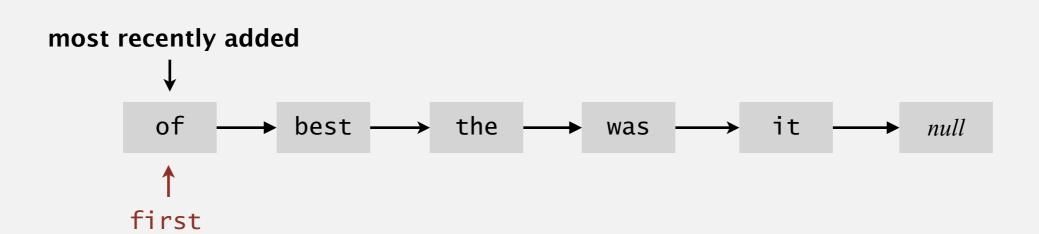
A. Can't be done efficiently with a singly-linked list.



D. I don't know.

## Stack: linked-list implementation

- Maintain pointer first to first node in a singly-linked list.
- Push new item before first.
- Pop item from first.



## Stack pop: linked-list implementation

#### inner class

```
private class Node
{
    String item;
    Node next;
}
```

```
save item to return
   String item = first.item;
delete first node
   first = first.next;
     first -
                                         to
                                         null
     first ~
                                         to
return saved item
   return item;
```

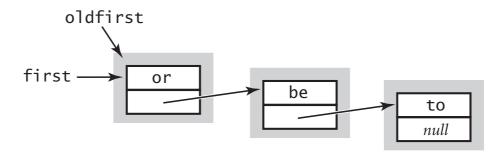
## Stack push: linked-list implementation

#### inner class

```
private class Node
{
    String item;
    Node next;
}
```

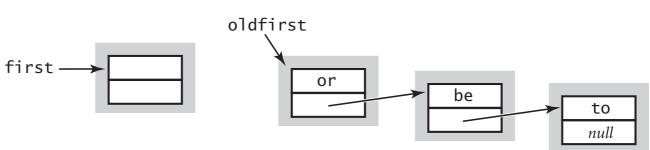
#### save a link to the list

Node oldfirst = first;



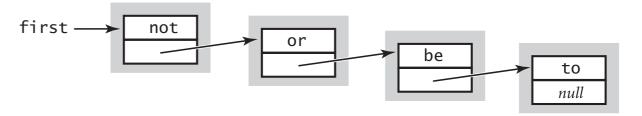
#### create a new node for the beginning

first = new Node();



#### set the instance variables in the new node

first.item = "not";
first.next = oldfirst;



## Stack: linked-list implementation in Java

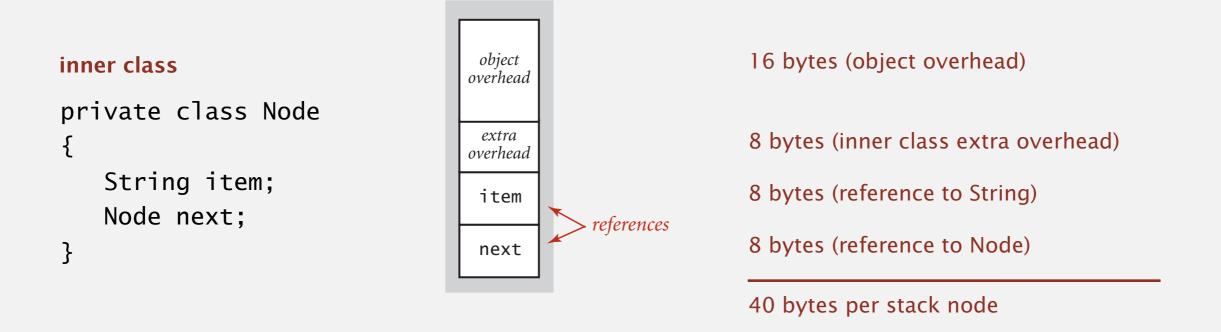
```
public class LinkedStackOfStrings
   private Node first = null;
   private class Node
      String item;
      Node next;
   public boolean isEmpty()
   { return first == null; }
   public void push(String item)
      Node oldfirst = first;
      first = new Node();
      first.item = item;
      first.next = oldfirst;
   }
   public String pop()
      String item = first.item;
      first = first.next;
      return item;
```

private inner class (access modifiers for instance variables don't matter)

## Stack: linked-list implementation performance

Proposition. Every operation takes constant time in the worst case.

Proposition. A stack with N items uses  $\sim 40 N$  bytes.



Remark. This accounts for the memory for the stack (but not the memory for strings themselves, which the client owns).

## Stack quiz

### How to implement a fixed-capacity stack with an array?

A. Can't be done efficiently with an array.

least recently added

B.

it	was	the	best	of	times	null	null	null	null
0	1	2	3	4	5	6	7	8	9

most recently added

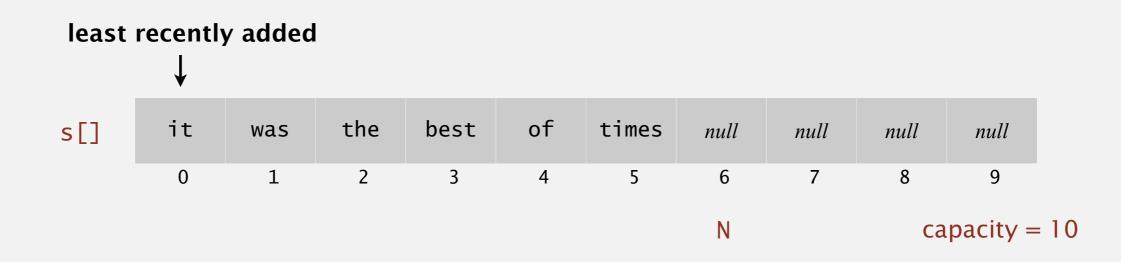
C

times	of	best	the	was	it	null	null	null	null
0	1	2	3	4	5	6	7	8	9

D. I don't know.

## Fixed-capacity stack: array implementation

- Use array s[] to store N items on stack.
- push(): add new item at s[N].
- pop(): remove item from s[N-1].



**Defect.** Stack overflows when N exceeds capacity. [stay tuned]

## Fixed-capacity stack: array implementation

```
public class FixedCapacityStackOfStrings
                                            a cheat
   private String[] s;
                                          (stay tuned)
   private int N = 0;
   public FixedCapacityStackOfStrings(int capacity)
   { s = new String[capacity]; }
   public boolean isEmpty()
   { return N == 0; }
   public void push(String item)
     s[N++] = item; 
   public String pop()
   { return s[--N]; }
```

use to index into array; then increment N

> decrement N; then use to index into array

#### Stack considerations

#### Overflow and underflow.

- Underflow: throw exception if pop from an empty stack.
- Overflow: use resizing array for array implementation. [stay tuned]

Null items. We allow null items to be inserted.

Loitering. Holding a reference to an object when it is no longer needed.

```
public String pop()
{ return s[--N]; }
```

loitering

```
public String pop()
{
    String item = s[--N];
    s[N] = null;
    return item;
}
```

this version avoids "loitering": garbage collector can reclaim memory for an object only if no outstanding references

# Algorithms

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## 1.3 BAGS, QUEUES, AND STACKS

- stacks
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  - applications

Problem. Requiring client to provide capacity does not implement API! Q. How to grow and shrink array?

#### First try.

- push(): increase size of array s[] by 1.
- pop(): decrease size of array s[] by 1.

#### Too expensive.

infeasible for large N

- Need to copy all items to a new array, for each operation.
- Array accesses to insert first N items =  $N + (2 + 4 + ... + 2(N-1)) \sim N^2$ .

per push

1 array access 2(k-1) array accesses to expand to size k (ignoring cost to create new array)

Challenge. Ensure that array resizing happens infrequently.

- Q. How to grow array?
- A. If array is full, create a new array of twice the size, and copy items.

"repeated doubling"

```
public ResizingArrayStackOfStrings()
{ s = new String[1]; }
public void push(String item)
   if (N == s.length) resize(2 * s.length);
   s[N++] = item;
}
private void resize(int capacity)
{
   String[] copy = new String[capacity];
   for (int i = 0; i < N; i++)
      copy[i] = s[i];
   s = copy;
}
```

```
Array accesses to insert first N = 2^i items. N + (2 + 4 + 8 + ... + N) \sim 3N.

1 array access per push k array accesses to double to size k (ignoring cost to create new array)
```

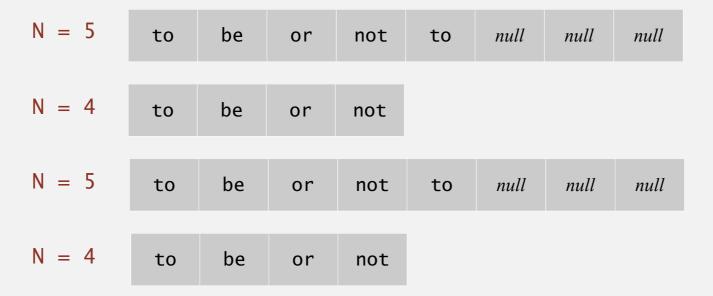
Q. How to shrink array?

#### First try.

- push(): double size of array s[] when array is full.
- pop(): halve size of array s[] when array is one-half full.

#### Too expensive in worst case.

- Consider push-pop-push-pop-... sequence when array is full.
- Each operation takes time proportional to N.



Q. How to shrink array?

#### Efficient solution.

- push(): double size of array s[] when array is full.
- pop(): halve size of array s[] when array is one-quarter full.

```
public String pop()
{
    String item = s[--N];
    s[N] = null;
    if (N > 0 && N == s.length/4) resize(s.length/2);
    return item;
}
```

Invariant. Array is between 25% and 100% full.

## Stack resizing-array implementation: performance

Amortized analysis. Starting from an empty data structure, average running time per operation over a worst-case sequence of operations.

Proposition. Starting from an empty stack, any sequence of M push and pop operations takes time proportional to M.

	best	worst	amortized	
construct	1	1	1	
push	1	N	1	
pop	1	$N \leftarrow$	1	doubling and
size	1	1	1	halving operati

order of growth of running time for resizing array stack with N items

## Stack resizing-array implementation: memory usage

Proposition. Uses between  $\sim 8 N$  and  $\sim 32 N$  bytes to represent a stack with N items.

- $\sim 8 N$  when full.
- $\sim 32 N$  when one-quarter full.

Remark. This accounts for the memory for the stack (but not the memory for strings themselves, which the client owns).

## Stack implementations: resizing array vs. linked list

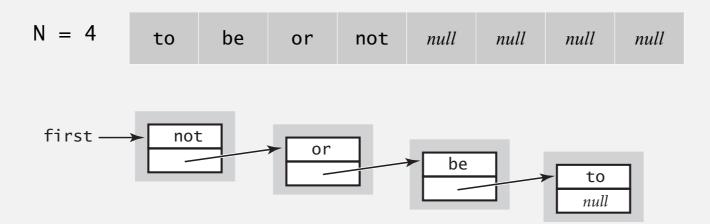
Tradeoffs. Can implement a stack with either resizing array or linked list; client can use interchangeably. Which one is better?

#### Linked-list implementation.

- Every operation takes constant time in the worst case.
- Uses extra time and space to deal with the links.

#### Resizing-array implementation.

- Every operation takes constant amortized time.
- Less wasted space.



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## 1.3 BAGS, QUEUES, AND STACKS

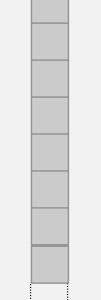
- stacks
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## Queue API

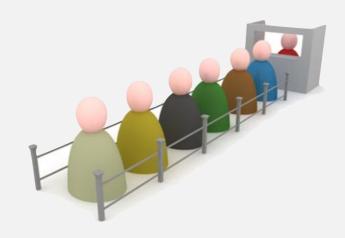
public class	QueueOfStrings	
	QueueOfStrings()	create an empty queue
void	<pre>enqueue(String item)</pre>	insert a new string onto queue
String	dequeue()	remove and return the string least recently added
boolean	isEmpty()	is the queue empty?
int	size()	number of strings on the queue

#### enqueue





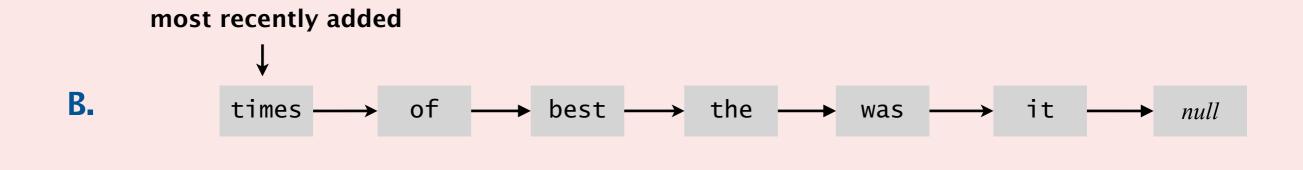


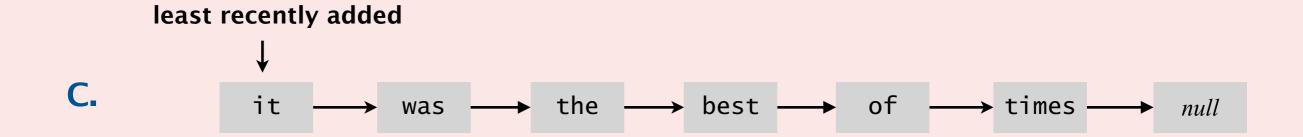


## Queue quiz 1

#### How to implement a queue with a linked list?

A. Can't be done efficiently with a singly-linked list.

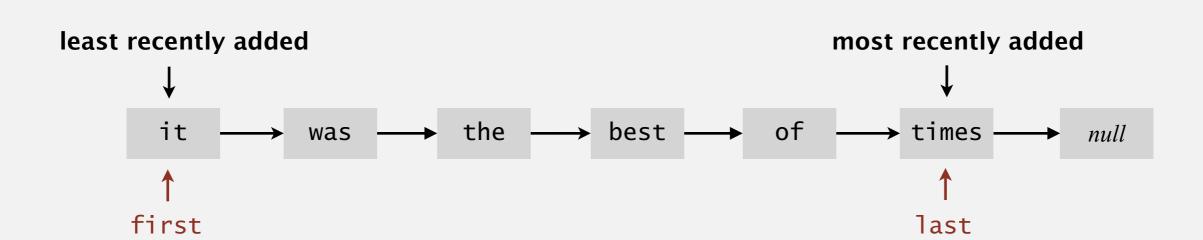




D. I don't know.

## Queue: linked-list implementation

- Maintain one pointer first to first node in a singly-linked list.
- Maintain another pointer last to last node.
- Dequeue from first.
- Enqueue after last.



## Queue dequeue: linked-list implementation

inner class

private class Node

String item;

Node next;

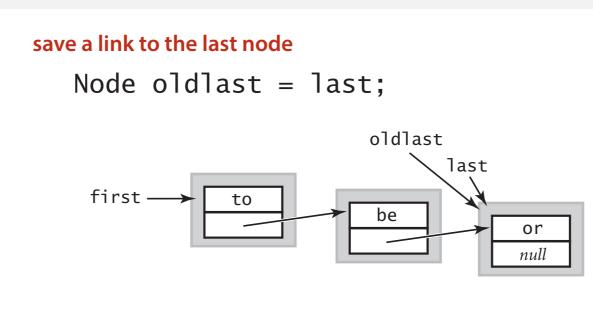
```
save item to return
   String item = first.item;
delete first node
   first = first.next;
                                  last
     first.
                                  last
     first -
return saved item
   return item;
```

Remark. Identical code to linked-list stack pop().

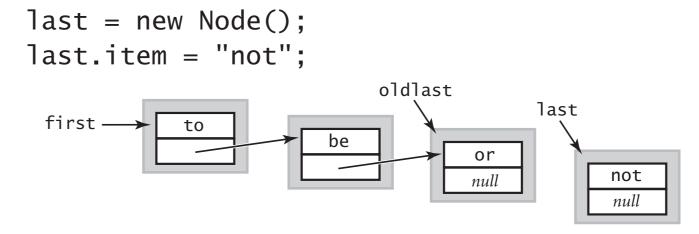
## Queue enqueue: linked-list implementation

#### inner class

```
private class Node
{
    String item;
    Node next;
}
```

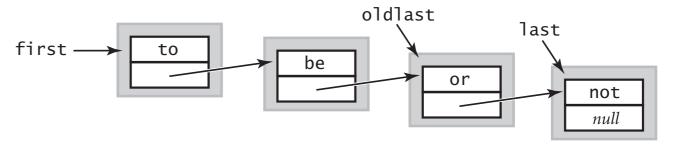


#### create a new node for the end



#### link the new node to the end of the list

oldlast.next = last;



## Queue: linked-list implementation in Java

```
public class LinkedQueueOfStrings
  private Node first, last;
  private class Node
   { /* same as in LinkedStackOfStrings */ }
  public boolean isEmpty()
   { return first == null; }
  public void enqueue(String item)
      Node oldlast = last;
     last = new Node();
     last.item = item;
     last.next = null;
                                                        special cases for
     if (isEmpty()) first = last;
                                                         empty queue
                     oldlast.next = last;
      else
  public String dequeue()
      String item = first.item;
                  = first.next;
     first
      if (isEmpty()) last = null;
      return item;
```

## Queue quiz 2

### How to implement a fixed-capacity queue with an array?

A. Can't be done efficiently with an array.

least recently added



B.

it	was	the	best	of	times	null	null	null	null
0	1	2	3	4	5	6	7	8	9

most recently added



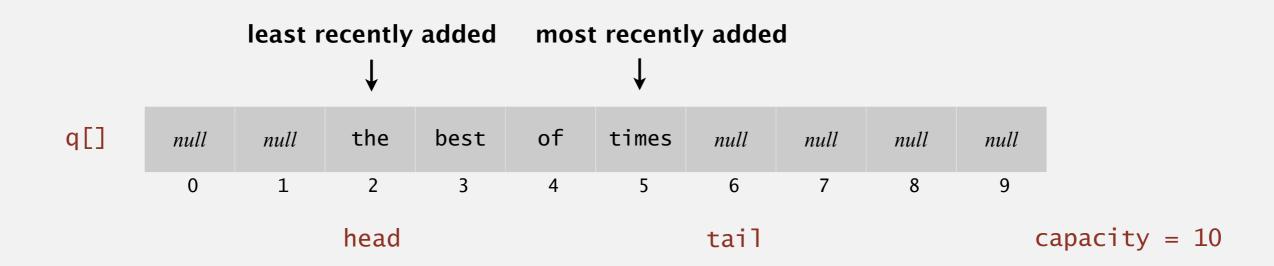
C.

times	of	best	the	was	it	null	null	null	null
0	1	2	3	4	5	6	7	8	9

D. I don't know.

## Queue: resizing-array implementation

- Use array q[] to store items in queue.
- enqueue(): add new item at q[tail].
- dequeue(): remove item from q[head].
- Update head and tail modulo the capacity.
- Add resizing array.



Q. How to resize?

# Algorithms

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## 1.3 BAGS, QUEUES, AND STACKS

- stacks
- resizing arrays
- queues
- generics
  - iterators
- applications

#### Parameterized stack

We implemented: StackOfStrings.

We also want: StackOfURLs, StackOfInts, StackOfVans, ....

Attempt 1. Implement a separate stack class for each type.

- Rewriting code is tedious and error-prone.
- Maintaining cut-and-pasted code is tedious and error-prone.

@#\$\*! most reasonable approach until Java 1.5.



#### Parameterized stack

We implemented: StackOfStrings.

We also want: StackOfURLs, StackOfInts, StackOfVans, ....

Attempt 2. Implement a stack with items of type Object.

- Casting is required in client.
- Casting is error-prone: run-time error if types mismatch.

```
StackOfObjects s = new StackOfObjects();
Apple a = new Apple();
Orange b = new Orange();
s.push(a);
s.push(b);
a = (Apple) (s.pop());
```



#### Parameterized stack

We implemented: StackOfStrings.

We also want: StackOfURLs, StackOfInts, StackOfVans, ....

#### Attempt 3. Java generics.

- Avoid casting in client.
- Discover type mismatch errors at compile-time instead of run-time.

```
Stack<Apple> s = new Stack<Apple>();
Apple a = new Apple();
Orange b = new Orange();
s.push(a);
s.push(b);
a = s.pop();
type parameter

compile-time error
```

Guiding principles. Welcome compile-time errors; avoid run-time errors.

# Generic stack: linked-list implementation

```
public class LinkedStackOfStrings
   private Node first = null;
   private class Node
     String item;
     Node next;
   public boolean isEmpty()
   { return first == null; }
   public void push(String item)
     Node oldfirst = first;
      first = new Node();
      first.item = item;
     first.next = oldfirst;
   public String pop()
      String item = first.item;
      first = first.next;
      return item;
```

```
public class Stack<Item>
   private Node first = nuN;
   private class Node
                                    generic type name
      Item item:
      Node next;
   public boolean is Empty
   { return first == nu////;
   public void push(I/tem item)
      Node oldfirs/t /= first;
      first = new Node();
      first.item /= item;
      first.next/ = oldfirst;
   public/Item pop()
      Item item = first.item;
      first = first.next;
      return item;
```

# Generic stack: array implementation

```
public class FixedCapacityStackOfStrings
   private String[] s;
   private int N = 0;
   public ..StackOfStrings(int capacity)
   { s = new String[capacity]; }
   public boolean isEmpty()
   { return N == 0; }
   public void push(String item)
   \{ s[N++] = item; \}
   public String pop()
   { return s[--N]; }
```

### the way it should be

```
public class FixedCapacityStack<Item>
   private Item[] s;
   private int N = 0;
   public FixedCapacityStack(int capacity)
   { s = new Item[capacity]; }
   public boolean isEmpty()
   { return N == 0; }
   publig void push(Item item)
   \{ s/[N++] = item; \}
   public Item pop()
      return s[--N]; }
```

# Generic stack: array implementation

```
public class FixedCapacityStackOfStrings
   private String[] s;
   private int N = 0;
   public ..StackOfStrings(int capacity)
   { s = new String[capacity]; }
   public boolean isEmpty()
   { return N == 0; }
   public void push(String item)
   \{ s[N++] = item; \}
   public String pop()
   { return s[--N]; }
```

### the way it is

```
public class FixedCapacityStack<Item>
   private Item[] s;
   private int N = 0;
   public FixedCapacityStack(int capacity)
   { s = (Item[]) new Object[capacity]; }
   public boolean isEmpty()
   { return N == 0; }
   public void push(Item item)
   \{ \langle S[N++] = item; \}
   public Item pop()
   { return s[--N]; }
```

### Unchecked cast

Q. Why does Java make me cast (or use reflection)? Short answer. Backward compatibility.



Long answer. Need to learn about type erasure and covariant arrays.

## Generic data types: autoboxing

Q. What to do about primitive types?

### Wrapper type.

- Each primitive type has a wrapper object type.
- Ex: Integer is wrapper type for int.

Autoboxing. Automatic cast between a primitive type and its wrapper.

Bottom line. Client code can use generic stack for any type of data.

# Algorithms

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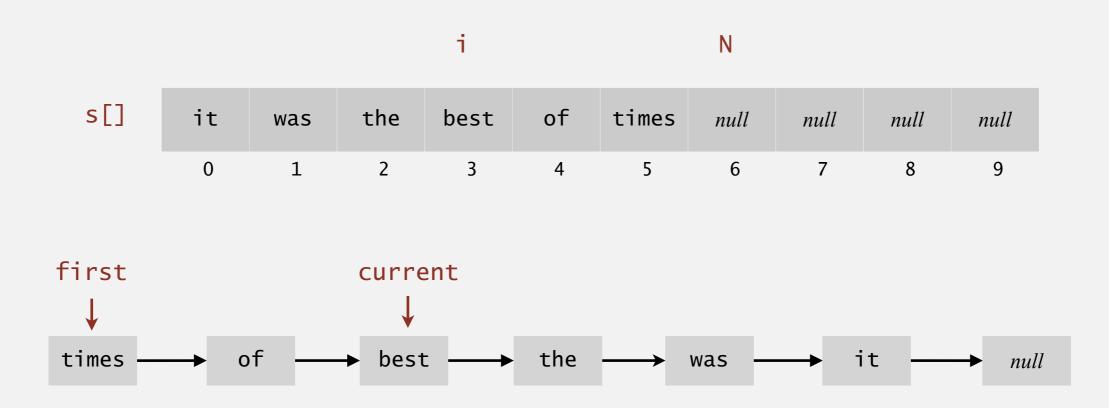
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# 1.3 BAGS, QUEUES, AND STACKS

- stacks
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- iterators
  - applications

### **Iteration**

Design challenge. Support iteration over stack items by client, without revealing the internal representation of the stack.



Java solution. Make stack implement the java.lang.Iterable interface.

### **Iterators**

- Q. What is an Iterable?
- A. Has a method that returns an Iterator.

- Q. What is an Iterator ?
- A. Has methods hasNext() and next().

- Q. Why make data structures Iterable?
- A. Java supports elegant client code.

### java.lang.lterable interface

```
public interface Iterable<Item>
{
    Iterator<Item> iterator();
}
```

### java.util.Iterator interface

### "foreach" statement (shorthand)

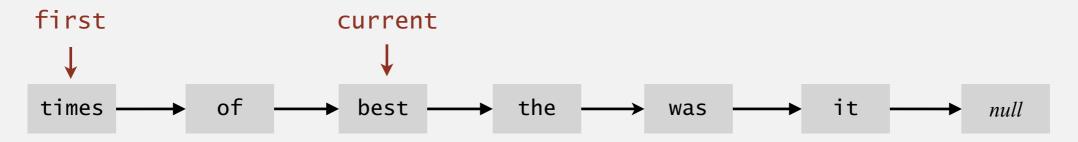
```
for (String s : stack)
   StdOut.println(s);
```

### equivalent code (longhand)

```
Iterator<String> i = stack.iterator();
while (i.hasNext())
{
   String s = i.next();
   StdOut.println(s);
}
```

# Stack iterator: linked-list implementation

```
import java.util.Iterator;
public class Stack<Item> implements Iterable<Item>
    public Iterator<Item> iterator() { return new ListIterator(); }
    private class ListIterator implements Iterator<Item>
        private Node current = first;
        public boolean hasNext() { return current != null; }
        public void remove() { /* not supported */
        public Item next() <</pre>
                                                 throw UnsupportedOperationException
             Item item = current.item:
                                                 throw NoSuchElementException
             current
                       = current.next;
                                                  if no more items in iteration
             return item;
```



## Stack iterator: array implementation

```
import java.util.Iterator;
public class Stack<Item> implements Iterable<Item>
   public Iterator<Item> iterator()
   { return new ReverseArrayIterator(); }
   private class ReverseArrayIterator implements Iterator<Item>
       private int i = N;
       public boolean hasNext() { return i > 0; }
       public void remove() { /* not supported */ }
       public Item next() { return s[--i]; }
```

				1			N			
s[]	it	was	the	best	of	times	null	null	null	null
	0	1	2	3	4	5	6	7	8	9

### Iteration: concurrent modification

- Q. What if client modifies the data structure while iterating?
- A. A fail-fast iterator throws a java.util.ConcurrentModificationException.

#### concurrent modification

```
for (String s : stack)
    stack.push(s);
```

Q. How to detect concurrent modification?

A.

- Count total number of push() and pop() operations in Stack.
- Save counts in \*Iterator subclass upon creation.
- If, when calling either next() or hasNext(), the current counts do not equal the saved counts, throw exception.

# Algorithms

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# 1.3 BAGS, QUEUES, AND STACKS

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## Java collections library

List interface. java.util.List is API for an sequence of items.

```
public interface List<Item> extends Iterable<Item>
                    List()
                                                         create an empty list
          boolean isEmpty()
                                                          is the list empty?
               int size()
                                                          number of items
             void add(Item item)
                                                       append item to the end
             Item get(int index)
                                                      return item at given index
             Item remove(int index)
                                                  return and delete item at given index
          boolean contains(Item item)
                                                  does the list contain the given item?
 Iterator<Item> iterator()
                                                    iterator over all items in the list
```

Implementations. java.util.ArrayList uses resizing array; java.util.LinkedList uses linked list. caveat: only some operations are efficient

## Java collections library

### java.util.Stack.

- Supports push(), pop(), and iteration.
- Extends java.util.Vector, which implements java.util.List interface from previous slide, including get() and remove().
- Bloated and poorly-designed API (why?)

### **Java 1.3 bug report (June 27, 2001)**

The iterator method on java.util.Stack iterates through a Stack from the bottom up. One would think that it should iterate as if it were popping off the top of the Stack.

#### status (closed, will not fix)

It was an incorrect design decision to have Stack extend Vector ("is-a" rather than "has-a"). We sympathize with the submitter but cannot fix this because of compatibility.

## Java collections library

### java.util.Stack.

- Supports push(), pop(), and iteration.
- Extends java.util.Vector, which implements java.util.List interface from previous slide, including get() and remove().
- Bloated and poorly-designed API (why?)



java.util.Queue. An interface, not an implementation of a queue.

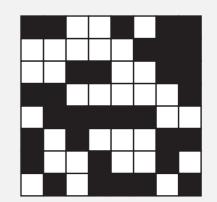
Best practices. Use our implementations of Stack, Queue, and Bag.

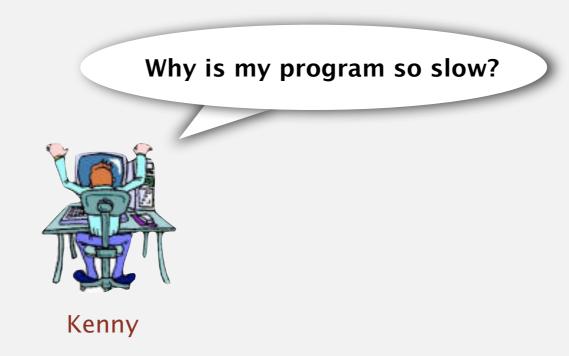
# War story (from Assignment 1)

### Generate random open sites in an N-by-N percolation system.

- Jenny: pick (i, j) at random; if already open, repeat. Takes  $\sim c_1 N^2$  seconds.
- Kenny: create a java.util.ArrayList of  $N^2$  closed sites. Pick an index at random and delete.

Takes ~  $c_2 N^4$  seconds.





Lesson. Don't use a library until you understand its API!

This course. Can't use a library until we've implemented it in class.

# Stack applications

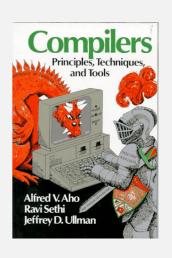
- Parsing in a compiler.
- Java virtual machine.
- Undo in a word processor.
- Back button in a Web browser.
- PostScript language for printers.
- Implementing function calls in a compiler.

• ...









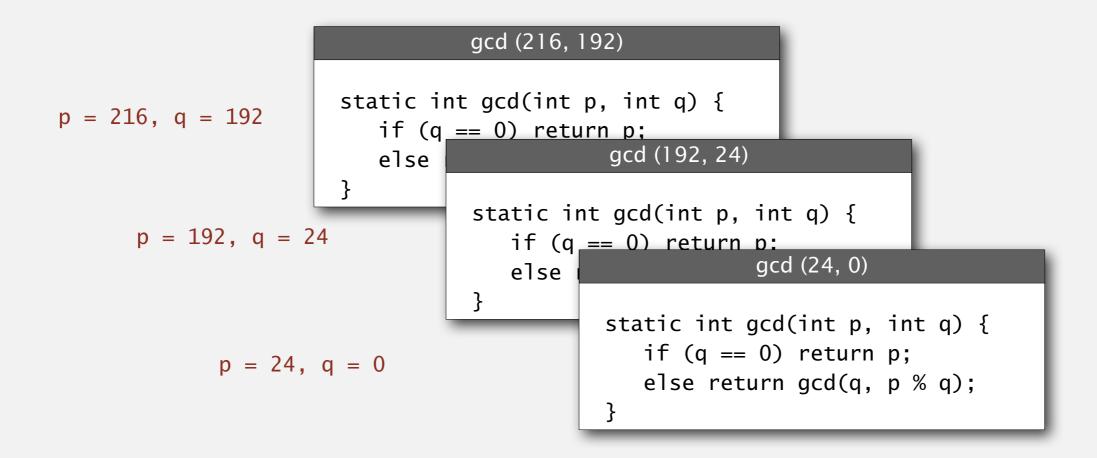
### Function calls

How a compiler implements a function.

- Function call: push local environment and return address.
- Return: pop return address and local environment.

Recursive function. Function that calls itself.

Note. Can always use an explicit stack to remove recursion.



## Arithmetic expression evaluation

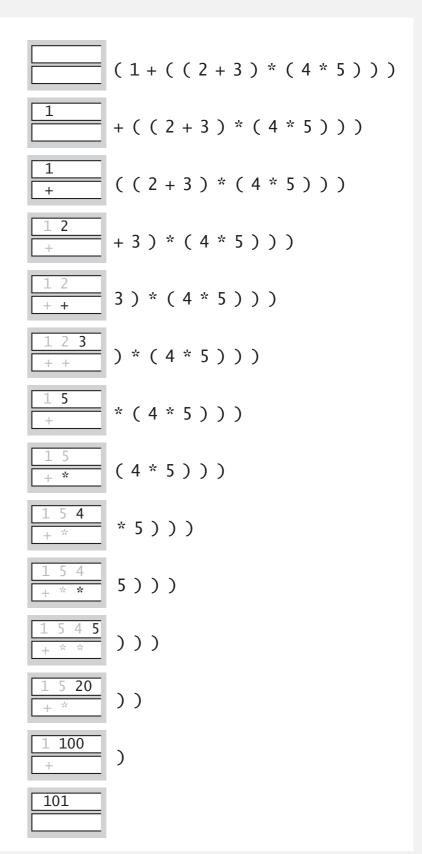
Goal. Evaluate infix expressions.

(1+((2+3)\*(4\*5)))
operand operator

Two-stack algorithm. [E. W. Dijkstra]

- Value: push onto the value stack.
- Operator: push onto the operator stack.
- Left parenthesis: ignore.
- Right parenthesis: pop operator and two values; push the result of applying that operator to those values onto the operand stack.

Context. An interpreter!

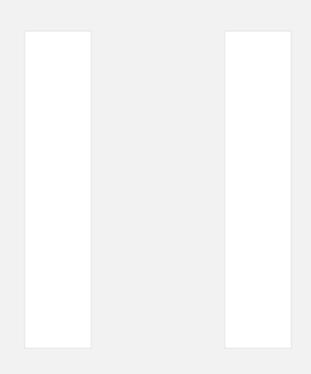


value stack

operator stack

# Dijkstra's two-stack algorithm demo





value stack

infix expression(fully parenthesized)



operator stack

# Arithmetic expression evaluation

```
public class Evaluate
   public static void main(String[] args)
      Stack<String> ops = new Stack<String>();
      Stack<Double> vals = new Stack<Double>();
      while (!StdIn.isEmpty()) {
        String s = StdIn.readString();
        if
            (s.equals("("))
        else if (s.equals("+")) ops.push(s);
        else if (s.equals("*"))      ops.push(s);
        else if (s.equals(")"))
           String op = ops.pop();
           if
              (op.equals("+")) vals.push(vals.pop() + vals.pop());
           else if (op.equals("*")) vals.push(vals.pop() * vals.pop());
        else vals.push(Double.parseDouble(s));
      StdOut.println(vals.pop());
                % java Evaluate
}
                 (1 + ((2 + 3) * (4 * 5)))
                 101.0
```

### Correctness

- Q. Why correct?
- A. When algorithm encounters an operator surrounded by two values within parentheses, it leaves the result on the value stack.

```
(1+((2+3)*(4*5)))
```

as if the original input were:

Repeating the argument:

Extensions. More ops, precedence order, associativity.

# Stack-based programming languages

Observation 1. Dijkstra's two-stack algorithm computes the same value if the operator occurs after the two values.

Observation 2. All of the parentheses are redundant!



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Bottom line. Postfix or "reverse Polish" notation.

Applications. Postscript, Forth, calculators, Java virtual machine, ...