

# Bilinear and quadratic forms

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## 0.1 Bilinear form

$h : V \times V \rightarrow K$  where  $V$  is a vector space and  $h$  is polylinear is called bilinear form.

As usually, bilinear form is characterized by  $a_{i,j} = h(e_i, e_j)$  which is called the matrix of  $h$  in basis  $e$ . Here is how to evaluate  $h$  through  $A$ :

$$h(u, v) = [u]_e^T A [v]_e$$

### Use cases

- o Dot product:  $x, y \mapsto \sum x_i y_i$
- o Convolution?:  $f, g \mapsto \int_0^1 f(x)g(x)w(x) dx$

### Symmetry

$\forall v, u : h(u, v) = h(v, u) \Rightarrow h$  — symmetric. It's equivalent to  $A$  being symmetric.

Examples of asymmetry:

- o Something with  $\mathbb{C}$ :  $x, y \mapsto \bar{x}y$
- o Another Convolution (it's usually antisymmetric):  $f, g \mapsto \int_0^1 f'(x)g(x) dx$

### Positive definite form

$$\forall x \neq 0 : h(x, x) > 0$$

### Use cases

It allows decomposition into a direct sum by taking orthogonal vectors.

**Def.**  $x \perp y \Leftrightarrow h(x, y) = 0$

**Def** (Orthogonal complement).  $U \leq V, U^\perp = \{x \in V \mid \forall y \in U, x \perp y\}$

**Th. 1.**  $h > 0$  — symmetric bilinear,  $U \leq V$  then  $V = U \oplus U^\perp$

*Proof.* 1. Suppose  $0 \neq x \in U \cap U^\perp$ . Then  $x \neq 0 \Rightarrow h(x, x) > 0$ , but  $x \in U \cap U^\perp \Rightarrow h(x, x) = 0$ .

2. Now let's show  $\dim U + \dim U^\perp = \dim V$ . Consider what equations define  $U^\perp$ :

$$\forall y \in U : h(x, e_i) = 0$$

Here  $e$  is the basis of  $U$ , and it's enough to check only with basis elements. So we have not more (they might be dependent, though they aren't) than  $\dim U$  equations. So  $\dim U^\perp \geq \dim V - \dim U$ .  $\square$

### Quadratic form

**Def.**  $q : V \rightarrow K$  which is a homogenous quadratic polynomial in every basis.

For  $q(x) = \sum_{i \leq j} b_{i,j} x_i x_j$  let's define  $A$  this way:

$$a_{i,j} = \begin{cases} \frac{b_{i,j}}{2}, & i < j \\ \frac{b_{j,i}}{2}, & i > j \\ b_{i,j}, & i = j \end{cases}$$

Now, both symmetric bilinear forms and quadratic forms both correspond to symmetric matrices, so they correspond to each other, and mapping is very simple:  $h(x, y) = \frac{q(x+y) - q(x) - q(y)}{2}$ .

### Change of basis

Suppose  $h$  is  $x^T A x$  in  $e$  and  $y^T B y$  in  $f$ , then:

$$x^T A x = (C y)^T A C y = y^T (C^T A C) y, C = [id]_e^f$$

Note that if we find an orthogonal basis than  $A$  will be a diagonal matrix.

**Th. 2.**  $q : V \rightarrow K, \exists e$  in which  $A$  (the matrix of  $q$ ) will be diagonal.

*Proof.* Here  $a_{ij} = a_{ji}$ :

$$q(x) = \sum a_{ij}x_i x_j = a_{11}(x_1 + \sum_{j \geq 2} \frac{a_{1j}}{a_{11}})^2 + q(x_{2\dots n}) - p(x_{2\dots n})$$

Where:

$$p(x_{2\dots n}) = \sum_{j \geq 2} \frac{a_{1j}^2}{a_{11}}$$

Then we continue to extract squares from  $q - p$  on  $x_{2\dots n}$  and the first basis element is:

$$e_1 = x_1 + \sum \frac{a_{1j}}{a_{11}}x_j$$

Final transition matrix will be upper-triangular with ones on the main diagonal, so it will be nondegenerate and give a basis. If  $a_{11} = 0$  we can swap this row with another. If all  $a_{ii} = 0$ , take some  $i, j : a_{ij} \neq 0$  and go to  $x'_i, x'_j = x_i + x_j, x_i - x_j$ , there will appear a difference of squares.  $\square$

We want to compute the main diagonal of matrix after this process.

**Th. 3.**  $q(x) = x^T A x, d_i \neq 0$  — *principal minors* (i.e.  $\det A[:i, :i]$ ).

The main diagonal will be  $\frac{d_i}{d_{i-1}}$ , where  $d_0 \stackrel{\text{def}}{=} 0$  (i.e. they won't change).

*Proof.* It's enough to proof that multiplying  $A \in M_n(K)$  by  $C \in UT_n(K)$  with ones on the main diagonal doesn't change it's principal minors.

Actually, we consider only top-left  $k \times k$  blocks and find the determinant of their product. But the determinant of  $UT_n$  with ones on the diagonal is 1, so the product is equal to the original determinant.

Now we need to show, that on each iteration  $a_{ii} \neq 0$  so that only the main branch of the algorithms works, otherwise transition matrices wouldn't have the right form.

Consider the  $k$ -th step of the algorithm:

$$A^{(k)} = \begin{pmatrix} c_1 & & & 0 & 0 \\ & \dots & & 0 & 0 \\ & & c_k & 0 & 0 \\ 0 & 0 & 0 & x & \dots \\ 0 & 0 & 0 & \dots & \dots \end{pmatrix}$$

Here  $x = a_{k+1,k+1} \neq 0$ , otherwise the  $(k+1)$ -th minor would be zero. More precisely,  $x = \frac{d_{k+1}}{d_k}$ .  $\square$

**Th. 4** (LU decomposition from the diagonalized form). *After the process from the previous theorem we get  $A = C^T D C$ , where  $C \in UT_n(K)$  with ones on the diagonal,  $D$  — diagonal.*

## Real case

By scaling the basis elements we can get  $q(x) = \sum_{i=1}^k x_i^2 - \sum_{i=k+1}^l x_i^2$ .

**Def.**  $(k, l)$  or simply  $k - l$  is called signature

**Th. 5.**  $k = \max\{\dim U \mid U \leq V, q|_U > 0\}$

*Proof.* If  $q(x) = \sum_{i=1}^k x_i^2 - \sum_{i=k+1}^l x_i^2$  is true for some  $(k, l)$ , then  $\max\{\dim U \mid U \leq V, q|_U > 0\} \geq k$ , because we can take  $\langle e_1 \dots e_k \rangle$ .

Now suppose we found  $U \leq V : \dim U \geq k+1$  and  $q|_U > 0$ . Also consider  $W = \langle e_{k+1} \dots e_n \rangle$ . We know  $q|_W < 0$  and  $U \cap W \neq \{0\}$  (from Grassman's formula). So we have a nonzero vector lying in the both forms, which is a contradiction.  $\square$