



# **SPINLOCK** [Report]

**Group Members:**

Sohaib Ashraf (20k-0488)

Khawaja Abdullah (20k-0385)

M.Shahzaib (20k-1067)

Teacher: **Sir Ghufraan**

Lab: **Sir Amin Sadiq**

**Date: 24th May 2022**

## Introduction:

Spinlock is a lock that causes a thread trying to acquire it to simply wait in a loop ("spin") while repeatedly checking whether the lock is available. Since the thread remains active but is not performing a useful task, the use of such a lock is a kind of busy waiting.

In our project, we are using a type of spinlock called **Read-Write** spinlock. For e.g we have 4 reader threads (thread-2 to thread-5). Now all four threads want to read the data from the variable at the same time. What will happen in this situation if we use a simple spinlock? Let's assume thread-2 got the lock while other reading threads are trying hard for the lock. Those other reading threads are simply wasting time to take lock since the variable won't change. That means performance will be reduced.

Hence with **read-write** spinlock, let's assume if thread-2 will take the read lock and read the data. And other reading threads also will take the read lock without spinning (blocking) and read the data. Because no one is writing. So what about the performance now? There is no waiting between reading operations. Thus, in this case, a read-write spinlock is useful. So, If multiple threads require read access to the same data, there is no reason why they should not be able to execute simultaneously. Spinlocks don't differentiate between read and read/write access.

**In Read Write spinlock multiple readers are permitted at the same time but only one writer. (i.e) If a writer has the lock, no reader is allowed to enter the critical section. If only a reader has the lock, then multiple readers are permitted in the critical section.**

## Where to use Read Write Spinlock?

- If you are only reading the data then you take read lock
- If you are writing then go for a write lock

## How we implemented:

We made a system call to show implementation of spin lock. We used two kernel threads. In the first thread, Writer is locked and the global variable increments to 1. No reader is allowed to enter. In thread2, readers keep the lock and the writer can not enter.

## STEPS:

1- Create a directory named final/ and made a file final.c

2- We wrote our **read-write spinlock code** in the final.c file.

3- Create a “Makefile” in the final directory and add the following line to it:  
**obj-y := final.o**

This is to ensure that the hello.c file is compiled and included in the kernel source code.

4- Go back to the parent dir (linux-4.17.4) and open “Makefile”

Add final/’ to the end of this line:

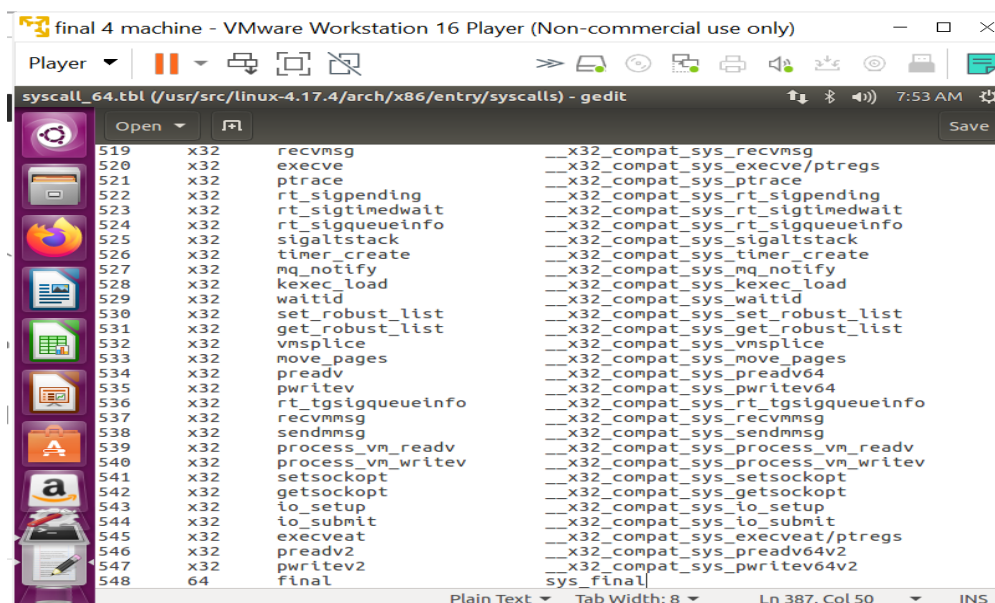
**core-y += kernel/ mm/ fs/ ipc/ security/ crypto/ block/ final/**

Edit the extraversion to our roll name

5- Now go to arch/x86/entry/syscalls/ and open **syscall\_64.tbl**

Go to the last of the document and add a new line like so:

548    64    final    sys\_final



```
final 4 machine - VMware Workstation 16 Player (Non-commercial use only)
Player
syscall_64.tbl (/usr/src/linux-4.17.4/arch/x86/entry/syscalls) - gedit
519 x32 recvmmsg _x32_compat_sys_recvmmsg
520 x32 execve _x32_compat_sys_execve/ptregs
521 x32 ptrace _x32_compat_sys_ptrace
522 x32 rt_sigpending _x32_compat_sys_rt_sigpending
523 x32 rt_sigtimedwait _x32_compat_sys_rt_sigtimedwait
524 x32 rt_sigqueueinfo _x32_compat_sys_rt_sigqueueinfo
525 x32 sigaltstack _x32_compat_sys_sigaltstack
526 x32 timer_create _x32_compat_sys_timer_create
527 x32 mq_notify _x32_compat_sys_mq_notify
528 x32 kexec_load _x32_compat_sys_kexec_load
529 x32 waitid _x32_compat_sys_waitid
530 x32 set_robust_list _x32_compat_sys_set_robust_list
531 x32 get_robust_list _x32_compat_sys_get_robust_list
532 x32 vmsplince _x32_compat_sys_vmsplince
533 x32 move_pages _x32_compat_sys_move_pages
534 x32 preadv _x32_compat_sys_preadv64
535 x32 pwritev _x32_compat_sys_pwritev64
536 x32 rt_tgsigqueueinfo _x32_compat_sys_rt_tgsigqueueinfo
537 x32 recvmmsg _x32_compat_sys_recvmmsg
538 x32 sendmmsg _x32_compat_sys_sendmmsg
539 x32 process_vm_readv _x32_compat_sys_process_vm_readv
540 x32 process_vm_writev _x32_compat_sys_process_vm_writev
541 x32 setsockopt _x32_compat_sys_setsockopt
542 x32 getsockopt _x32_compat_sys_getsockopt
543 x32 io_setup _x32_compat_sys_io_setup
544 x32 io_submit _x32_compat_sys_io_submit
545 x32 execveat _x32_compat_sys_execveat/ptregs
546 x32 preadv2 _x32_compat_sys_preadv64v2
547 x32 pwritev2 _x32_compat_sys_pwritev64v2
548 64 final sys_final
```

6- Go to the linux-4.17.4/ directory and type the following commands:

**cd include/linux/**

**gedit syscalls.h**

Add the following line to the end of the document before the #endif statement:

**asmlinkage long sys\_final(void);**

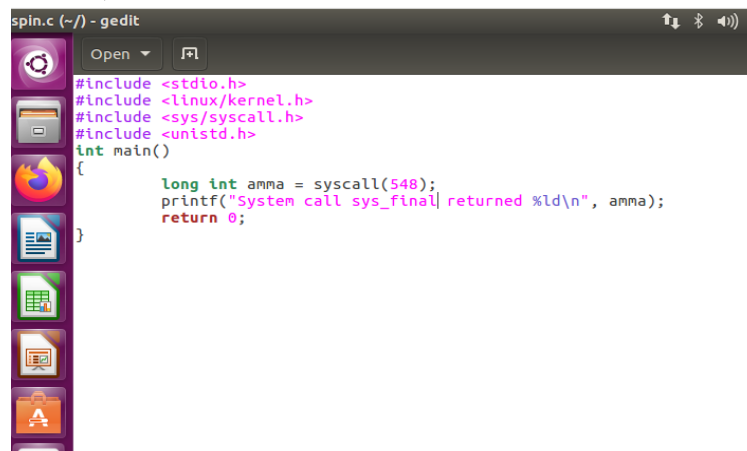
7- Now we compile the kernel using **sudo make**.

8- Run the following command in your terminal:

**sudo make modules\_install install**

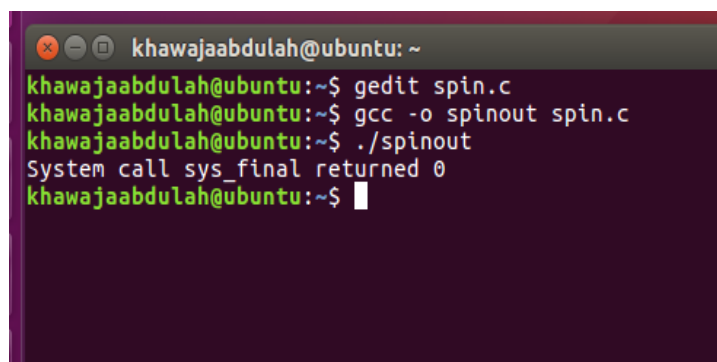
9- Now we restart our system. Go to your home(~) directory using the following commands and create a spin.c file.

We will call our system call from there.

A screenshot of a gedit text editor window titled 'spin.c (~/) - gedit'. The editor shows the following C code:

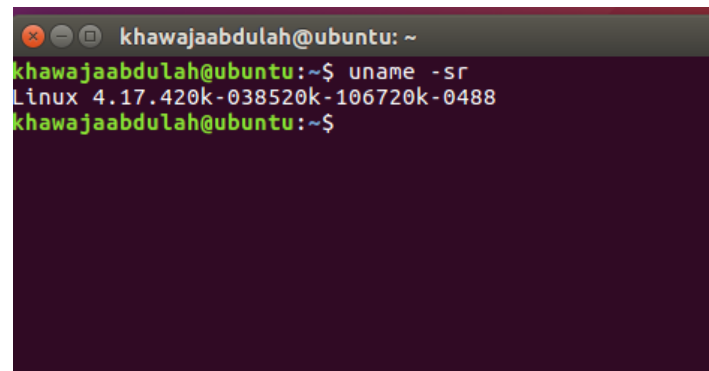
```
#include <stdio.h>
#include <linux/kernel.h>
#include <sys/syscall.h>
#include <unistd.h>
int main()
{
    long int amma = syscall(548);
    printf("System call sys_final returned %ld\n", amma);
    return 0;
}
```

10- After compile our run, spin.c programs returns 0 which means our spinlock system call is working

A screenshot of a terminal window with the title 'khawajaabdulah@ubuntu: ~'. It shows the following commands and output:

```
khawajaabdulah@ubuntu:~$ gedit spin.c
khawajaabdulah@ubuntu:~$ gcc -o spinout spin.c
khawajaabdulah@ubuntu:~$ ./spinout
System call sys_final returned 0
khawajaabdulah@ubuntu:~$
```

## 11- Uname -sr command shows our roll numbers



```
khawajaabduh@ubuntu: ~  
khawajaabduh@ubuntu:~$ uname -sr  
Linux 4.17.420k-038520k-106720k-0488  
khawajaabduh@ubuntu:~$
```

12- Write dmesg command. It will return our results. When read is locked, the value of the shared variable is read. When write is locked, value of shared variable is increased by 1, preventing any race conditions.



```
khawajaabduh@ubuntu: ~  
1929.793021] In Thread Function1 while write is locked: Globalvar is: 3017  
1930.046757] In Thread Function2 while read locked: Globalvar is: 3017  
1930.046764] In Thread Function1 while write is locked: Globalvar is: 3018  
1930.207277] In Thread Function2 while read locked: Globalvar is: 3018  
1930.207287] In Thread Function1 while write is locked: Globalvar is: 3019  
1930.815045] In Thread Function1 while write is locked: Globalvar is: 3020  
1930.815060] In Thread Function2 while read locked: Globalvar is: 3020  
1931.071138] In Thread Function2 while read locked: Globalvar is: 3020  
1931.071147] In Thread Function1 while write is locked: Globalvar is: 3021  
1931.231035] In Thread Function1 while write is locked: Globalvar is: 3022  
1931.231050] In Thread Function2 while read locked: Globalvar is: 3022  
1931.839253] In Thread Function2 while read locked: Globalvar is: 3022  
1931.839262] In Thread Function1 while write is locked: Globalvar is: 3023  
1932.094620] In Thread Function1 while write is locked: Globalvar is: 3024  
1932.094627] In Thread Function2 while read locked: Globalvar is: 3024  
1932.254930] In Thread Function2 while read locked: Globalvar is: 3024  
1932.254948] In Thread Function1 while write is locked: Globalvar is: 3025  
1932.863020] In Thread Function1 while write is locked: Globalvar is: 3026  
1932.863037] In Thread Function2 while read locked: Globalvar is: 3026  
1933.118202] In Thread Function1 while write is locked: Globalvar is: 3027  
1933.118209] In Thread Function2 while read locked: Globalvar is: 3027  
1933.279172] In Thread Function1 while write is locked: Globalvar is: 3028  
1933.279179] In Thread Function2 while read locked: Globalvar is: 3028  
khawajaabduh@ubuntu:~$
```