

Project # 1

This project is due Thursday, January 29.

1. Warm up (not graded): *Computing RSA by hand*

Let $p = 13, q = 23, e = 17$ be your initial parameters.

- (a) *Key generation:* Compute N and $\phi(N)$. Compute the private key $kp = d = e^{-1} \bmod \phi(N)$ (e.g. by using the extended Euclidean algorithm). Show all intermediate results.
- (b) *Encryption:* Encrypt the message $m = 31$ by applying the square and multiply algorithm (first, transform the exponent to binary representation).
- (c) *Decryption:* Decrypt the ciphertext c computed above by applying the square and multiply algorithm.

2. Warm up (not graded): *RSA Cryptanalysis*

Let $N = pq$ be the product of two distinct primes. Show that if $\phi(N)$ and N are known, then it is possible to compute p and q .

(*Hint:* Derive a quadratic equation (over the integers) in the unknown p)

3. *Algorithms for RSA*

- (a) Implement the extended Euclidean algorithm as a function `xGCD(a,b)` in Python. Use the template given in myWPI as a reference implementation and paste your own code into the given function definition. (for more information on `xGCD` see algorithm 2.107 in HAC)
- (b) Implement the square and multiply algorithm in Python using the provided template.

4. *Textbook RSA*

- (a) Following the example for the RSA encryption function, write a function `RSAdec(c,d,N)` to decrypt ciphertexts. Instead of using the built-in `pow(a,b,N)` function, use your own square and multiply routine.
- (b) Implement a simplified RSA key generation algorithm `RSAGen(p,q)` that outputs public and private key parameters for the given primes p and q . The public exponent e is always chosen as $e = 2^{16} + 1$.

- (c) Compute the plaintext $m = c^d \bmod N$ for the following parameters below.
- (d) Given an RSA signature oracle returning signatures $\sigma' = \text{Sign}_{sk}(m')$ for all $m' \neq m$. Show that the signature of m can be found efficiently, i.e. that *universal* forgeries are possible for RSA signatures. *Hint*: It might be helpful to consider that RSA is *multiplicative*: $m_1^e \cdot m_2^e = (m_1 m_2)^e \bmod N$

Example for the RSA encryption in Python (yes, indentation does matter in Python!):

```
def RSAenc(m,e,N): # return m^e mod N
    """Returns the RSA ciphertext of m encrypted with e and modulus N"""
    return pow(m,e,N)
```

```
p =
117745671967951862648908489379882696196830917954980484704887934600337961
370859639649579886611494760702973713721868261964258172538906822624536069
14957198881,
q =
843508267592143758951725376108546151543072575330048003695892183714052361
704761808357623747291290446435169373648255024227642262144084184562067588
2916844183,
and c = 31
```

5. Straightforward AES

- (a) Implement the AES in Python. A template for this exercise is provided. Your code should support encryption and decryption for keys of 128, 192 and 256 bits. The code should take bytearrays as input (plaintext/ciphertext and key), and also return a bytearray as output.
- (b) Using your implementation of AES, decrypt the following ciphertext (given in hexadecimal representation):

E5 5C D4 A8 EE E5 7D 26 1C 16 CA FE C9 40 A9 44

and the following 128-bit key:

00 01 02 03 04 05 06 07 08 09 10 11 12 13 14 15