CP Algorithms Handbook

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Contents

1 Sorting and Searching

1.1 4SUM

2 Dynamic Programming

2.1 State representaion

2.1.1 All subsets

This way you can generate all subsets in order ready for DP.

```
int sub = whole_set;
while(sub > 0){

sub = (sub-1)&whole_set;
}
```

2.1.2 Permutations to Subsets

3 Graph Algorithms

3.1 Union Find

```
int s_size[MX];
1 int link[MX];
4 int max_size = 0;
7 int find(int x){
8 while(x != link[x])
        x = link[x];
      return x;
11 }
12
bool same(int a, int b){
     return find(a) == find(b);
14
15 }
16
void unite(int a, int b){
    a = find(a);
      b = find(b);
19
20
     if(s_size[a] < s_size[b])</pre>
          swap(a, b);
     s_size[a] += s_size[b];
22
      max_size = max(max_size, s_size[a]);
23
      link[b] = a;
24
      num --;
25
26 }
27
28 int main(){
   for(int i = 0; i < MX; i++){</pre>
          s_size[i] = 1;
link[i] = i;
30
31
32
33
     int n, m;
34
     cin >> n >> m;
35
     num = n;
36
      for(int i = 0; i < m; i++){</pre>
         int a, b;
38
39
          cin >> a >> b;
          if(!same(a, b))
40
          unite(a, b);
cout << num << " " << max_size << endl;
41
43
      return 0;
44
```

3.2 Path through all edges

```
set <int> adj[MX];
3 bool vis[MX];
  void zero(){
      for(int i = 0; i < MX; i++)</pre>
6
           vis[i] = false;
8 }
9
int dfs1(int a){
       int ret = 0;
11
       vis[a] = true;
12
      for(int b : adj[a]){
13
14
           ret += 1;
           if(!vis[b])
15
               ret += dfs1(b);
17
       return ret;
18
19 }
20
21 int main(){
      int n, m;
22
23
       cin >> n >> m;
       for(int i = 0; i < m; i++){</pre>
24
           int a, b;
25
           cin >> a >> b;
26
27
           adj[a].insert(b);
           adj[b].insert(a);
28
29
30
31
      zero();
      if (dfs1(1)/2 != m){
32
           cout << "IMPOSSIBLE";</pre>
33
           return 0;
34
35
36
       for(int i = 1; i <= n; i++){</pre>
37
           if(adj[i].size() & 1){
38
                cout << "IMPOSSIBLE";</pre>
39
40
                return 0;
41
42
43
       vector < int > ans;
44
45
46
       stack<int> s;
       s.push(1);
47
48
       while(!s.empty()){
           int a = s.top();
49
           if(adj[a].size() == 0){
50
51
                ans.push_back(a);
                s.pop();
52
           }
53
54
                int b = *adj[a].begin();
55
                adj[a].erase(adj[a].begin());
56
                adj[b].erase(a);
57
                s.push(b);
58
           }
59
60
61
       for(int a : ans)
62
           cout << a << " ";
63
64
65
       return 0;
66
```

3.3 Flows and Cuts

All problems in this section can be solved using the same basic algorithm defined in 3.3.1.

3.3.1 Maximum Flow

This implementation of Ford-Fulkerson algorithm (known as Edmonds-Karp algorithm) uses BFS to check if it's possible to expand the flow through the graph and assign depths from the source. Later it uses DFS to expand te flow.

```
1 11 flow[MX][MX];
1 used[MX][MX];
vector<int> adj[MX];
5 vector<int> radj[MX];
  bool vis[MX];
9
  bool bfs(int s, int lvl[], const int n){
       queue < int > q;
10
11
       q.push(s);
       lvl[s] = 1;
12
       while(!q.empty()){
13
           int p = q.front();
14
           q.pop();
           for(int v : adj[p]){
                if(lvl[v] == 0 and flow[p][v] - used[p][v] > 0){
    lvl[v] = lvl[p]+1;
17
18
                     q.push(v);
19
20
           }
21
22
            for(int v : radj[p]){
                if(lvl[v] == 0 and used[v][p] > 0){
23
                     lvl[v] = lvl[p]+1;
24
25
                     q.push(v);
                }
26
           }
27
28
       }
       return lvl[n] > 0;
29
30 }
31
  int dfs(int x, ll val, const int lvl[], const int n){
32
       vis[x] = true;
33
       if(x == n){
34
35
           return val;
36
37
       for(int v : adi[x]){
            if(lvl[v] == lvl[x]+1 \text{ and } flow[x][v] - used[x][v] > 0 \text{ and } !vis[v])
38
                int r = dfs(v, min(val, flow[x][v]-used[x][v]), lvl, n);
39
                if(r > 0){
40
41
                     used[x][v] += r;
                     return r;
42
                }
43
           }
44
45
46
       for(int v : radj[x]){
47
           if(lvl[v] == lvl[x]+1 \text{ and } used[v][x] > 0 \text{ and } !vis[v]){
                int r = dfs(v, min(val, used[v][x]), lvl, n);
48
                if(r > 0){
49
                     used[v][x] -= r;
50
51
                     return r;
52
           }
53
       }
54
55
       return 0;
56 }
```

To find the flow:

```
long long int res = 0;

while(bfs(1, lvl, n)){
    res += dfs(1, __INT_MAX__, lvl, n);
    for(int i = 0; i < MX; i++){
        vis[i] = false;
        lvl[i] = 0;
    }
}</pre>
```

3.3.2 Minimum Cut

To find Minimum Cut in a graph we need to find the Maximum Flow and check which edges connect two created disjont sets of nodes. This modification to finds them:

```
long long int res = 0;
         while(bfs(1, lvl, n)){
 3
              res += dfs(1, __INT_MAX__, lvl, n);
for(int i = 0; i < MX; i++) {
    vis[i] = false;
    lvl[i] = 0;
 5
 6
 8
        }
9
         cout << res << endl;</pre>
10
        for(int i = 0 ; i < MX; i++) {
    vis[i] = false;</pre>
11
12
              lv1[i] = 0;
13
14
15
        bfs(1, lv1, n);
16
17
         for(int a = 1; a <= n; a++){</pre>
              if(lvl[a] > 0){
18
                   for(int b : adj[a]){
19
                         if(lvl[b] == 0){
20
                               cout << a << " " << b << endl;
21
22
                    }
23
              }
24
```

3.3.3 Maximum Matching

3.4 2SAT

Given logical formula in the conjunctive normal form:

$$(a_1 \lor b_1) \land (a_2 \lor b_2) \land \dots \land (a_n \lor b_n)$$

we can eliminate disjunction by repacing each $(a_i \vee b_i)$ element with pair:

$$\neg a_i \to b_i \land \neg b_i \to a_i$$

```
#include <iostream>
#include <vector>
3 #include <stack>
5 using namespace std;
7 \text{ const int } MX = 2e5+7;
9 vector < int > adj[MX];
vector<int> Tadj[MX];
stack<int> scc_stack;
12
int scc[MX];
14
bool vis[MX];
int scc_val[MX];
17
18 int not_node(int a){
      if(!(a&1))
19
20
          return a+1;
21
          return a-1;
22
23 }
24
void dfs1(int a){
     vis[a] = true;
      for(int b : adj[a]){
27
         if(!vis[b])
```

```
dfs1(b);
29
30
        scc_stack.push(a);
31
32 }
33
void dfs2(int a, int scc_n){
        vis[a] = true;
for(int b : Tadj[a]){
35
36
             if(!vis[b])
37
38
                  dfs2(b, scc_n);
39
        scc[a] = scc_n;
40
41 }
42
43 void zero(){
       for(int i = 0 ; i < MX; i++)</pre>
44
             vis[i] = false;
45
46 }
47
48 int main(){
        int n, m;
        cin >> n >> m;
50
        for(int i = 0 ; i < n; i++){</pre>
51
52
             char op;
             int a, b;
53
             cin >> op >> a;
if(op == '+'){
   a *= 2;
54
55
56
57
             }
             else{
58
                  a *= 2;
59
             }
61
             cin >> op >> b;
62
             if (op == '+') {
   b *= 2;
63
64
65
66
             else{
                  b *= 2;
67
68
                  b++;
69
70
             Tadj[a].push_back(not_node(b));
71
              Tadj[b].push_back(not_node(a));
             adj[not_node(b)].push_back(a);
72
73
              adj[not_node(a)].push_back(b);
74
75
76
        for(int i = 1 ; i <= 2*m+1; i++){</pre>
             if(!vis[i])
77
                  dfs1(i);
78
        }
79
80
81
        zero();
        int scc_n = 1;
82
83
        while(!scc_stack.empty()){
84
             int a = scc_stack.top();
             scc_stack.pop();
85
             if(scc[a] == 0){
86
                   dfs2(a, scc_n);
87
                   scc_n++;
88
89
             }
90
91
        for(int i = 2; i <= m*2; i+=2){</pre>
             if(scc[i] == scc[i+1]){
   cout << "IMPOSSIBLE" << endl;</pre>
93
94
95
                   return 0;
             }
96
97
        // 1 - false; 2 - true;
for(int i = 2; i <= 2*m; i+=2){
   int a = scc[i];
98
99
100
             int b = scc[i+1];
101
```

```
if(scc_val[a] != 0){
102
                scc_val[b] = (scc_val[a] == 1 ? 2 : 1);
103
            if(scc_val[b] != 0){
                scc_val[a] = (scc_val[b] == 1 ? 2 : 1);
106
            if(a < b){
108
109
                scc_val[a] = 1;
                scc_val[b] = 2;
            }
111
            else{
112
                scc_val[a] = 2;
113
                scc_val[b] = 1;
114
       }
116
117
       for(int i = 2; i <= m*2; i+=2)</pre>
118
            cout << (scc_val[scc[i]] == 2 ? "+" : "-") << " ";
119
120
       return 0;
121
122 }
```

4 Range Queries

4.1 Segment trees

BASE size table:

4.1.1 Point-Range Trees

```
void insert(int tree[], int p, int val){
      p += BASE;
2
      tree[p] = val;
      while (p > 0) {
          p >>= 1;
5
           tree[p] = max(tree[2*p], tree[2*p + 1]);
8 }
int querry(int tree[], int val){
2
      int p = 1;
      while(p < BASE){</pre>
          if(tree[2*p] >= val){
               p = 2*p;
          }
6
          else{
              p = 2*p+1;
9
11
      if(tree[p] >= val)
          return p -= BASE;
12
13
      else
14
          return -1;
15 }
```

4.1.2 Range-Range Trees

1. range ADD insert, range MAX value query

```
void max_add(int a, int b, int val, int k=1, int x=0, int y=BASE-1){

if(a <= x and y <= b){

    max_tree[k][1] += val;

    return;

}

max_tree[2*k][1] += max_tree[k][1];

max_tree[2*k+1][1] += max_tree[k][1];

max_tree[k][0] += max_tree[k][1];</pre>
```

```
9 max_tree[k][1] = 0;
10
       int d = (x+y)/2;
11
       if(a <= d){</pre>
12
           max_add(a, b, val, 2*k, x, d);
13
           }
14
       if(b > d){
15
16
          max_add(a, b, val, 2*k+1, d+1, y);
17
      \max_{\text{tree}}[k][0] = \max(\max_{\text{tree}}[2*k+1][0] + \max_{\text{tree}}[2*k+1][1], \max_{\text{tree}}[2*k][0]
       + max_tree[2*k][1]);
19
20 }
1 ll max_query(int a, int b, int k=1, int x=0, int y=BASE-1){
      if(a <= x and y <= b){</pre>
          return max_tree[k][0] + max_tree[k][1];
3
4
       max_tree[2*k][1] += max_tree[k][1];
5
      max_tree[2*k+1][1] += max_tree[k][1];
6
      max_tree[k][0] += max_tree[k][1];
max_tree[k][1] = 0;
      int d = (x+y)/2;
9
10
      11 ret = -INF;
      if(a <= d){</pre>
11
           ret = max(ret, max_query(a, b, 2*k, x, d));
12
13
       if(b > d){
14
          ret = max(ret, max_query(a, b, 2*k+1, d+1, y));
15
16
17
      return ret;
18 }
2. range ADD insert, range SUM querry
void sum_add(int a, int b, int val, int k=1, int x=0, int y=BASE-1){
     if(a <= x and y <= b){</pre>
          sum_tree[k][1] += val;
           return;
4
5
6
       sum_tree[2*k][1] += sum_tree[k][1];
       sum_tree[2*k+1][1] += sum_tree[k][1];
7
8
       sum_tree[k][0] += sum_tree[k][1]*(y-x+1);
       sum_tree[k][1] = 0;
9
      int d = (x+y)/2;
10
       if(a <= d){
11
          int w = (d <= b ? d : b)+1;</pre>
12
           w -= (x \le a ? a : x);
13
14
           sum_tree[k][0] += val*(w);
           sum_add(a, b, val, 2*k, x, d);
15
16
17
      if(b > d){
          int w = (y <= b ? y : b)+1;
18
           w -= (d+1 \le a ? a : d+1);
19
           sum_tree[k][0] += val*(w);
20
           sum_add(a, b, val, 2*k+1, d+1, y);
21
22
23 }
1 ll sum_query(int a, int b, int k=1, int x=0, int y=BASE-1){
      if(a \le x \text{ and } y \le b)
2
           return sum_tree[k][0] + sum_tree[k][1]*(y-x+1);
3
4
       sum_tree[2*k][1] += sum_tree[k][1];
5
       sum_tree[2*k+1][1] += sum_tree[k][1];
6
       sum_tree[k][0] += sum_tree[k][1]*(y-x+1);
       sum_tree[k][1] = 0;
8
9
       int d = (x+y)/2;
       11 ret = 0;
10
      if(a <= d){
11
12
          ret += sum_query(a, b, 2*k, x, d);
13
```

 $if(b > d){$

14

```
ret += sum_query(a, b, 2*k+1, d+1, y);
}
return ret;
}
```

5 Tree Algorithms

5.1 Binary Lifitng

```
int main(){
   int n, q;
   cin >> n >> q;
   boss[1][0] = 0;
   for(int i = 2; i <= n; i++){
      cin >> boss[i][0];
}

for(int i = 1; i < LOG_MX; i++){
      for(int emp = 1; emp <= MX; emp++){
            boss[emp][i] = boss[boss[emp][i-1]][i-1];
      }
}</pre>
```

5.2 Tree Traversal

```
vector <int> adj[MX];
1 int subtree[MX];
3 int flat[MX];
4 int vals[MX];
6 int idx = 1;
7 int dfs(int a, int p){
      flat[a] = idx;
     idx++;
9
    for(int b : adj[a]){
10
     if(b == p)
11
             continue;
12
         subtree[a] += dfs(b, a)+1;
14
     return subtree[a];
15
```

5.3 LCA

6 String Algorithms

6.1 Trie

```
int trie[TMX][30];
bool stop[TMX];
3 int next_node = 1;
4 int dp[MX];
6 void insert(string s){
    int idx = 0;
       for(char c : s){
           if(trie[idx][c-'a'] == 0){
9
               trie[idx][c-'a'] = next_node;
10
               next_node ++;
11
           }
12
           idx = trie[idx][c-'a'];
13
14
15
       stop[idx] = true;
16 }
17
18 int main(){
    string text;
19
      int n;
21     cin >> text >> n;
22     for(int i = 0; i < n; i++){</pre>
```

6.2 Pattern Finding

6.2.1 KMP

```
int pi[MX];
3 int main(){
       string text, pattern;
cin >> text >> pattern;
       string s = pattern + "+" + text;
       for(int i = 1; i < s.size(); i++){</pre>
           int j = pi[i-1];
8
            while (j > 0 \text{ and } s[i] != s[j])
            j = pi[j-1];
if(s[i] == s[j])
10
11
12
                j++;
            pi[i] = j;
13
14
15
16
     int res = 0;
     for(int i = 0; i < s.size(); i++){</pre>
17
           if(pi[i] == pattern.size())
18
19
                res++;
20
21
      cout << res;
    return 0;
```

6.2.2 Hashing

```
const int MOD = 1e9+9;
const int P = 9973;
3 const int MX = 1e6+7;
5 11 ppow[MX];
6 ll pfx_hash[MX];
8 int res = 0;
10 int main(){
11
      string s, pattern;
      cin >> s >> pattern;
12
      ppow[0] = 1;
13
      for(int i = 1; i < MX; i++) {
    ppow[i] = (ppow[i-1]*P)%MOD;</pre>
14
15
16
17
      11 pattern_hash = 0;
18
      for(int i = 0; i < pattern.size(); i++){</pre>
19
           pattern_hash = (pattern_hash + (pattern[i]-'a'+1)*ppow[i])%MOD;
20
21
22
      for(int i = 0; i < s.size(); i++){</pre>
23
           pfx_hash[i+1] = (pfx_hash[i] + (s[i]-'a'+1)*ppow[i])%MOD;
25
26
       for(int i = 0; i+pattern.size()-1 < s.size(); i++){</pre>
27
           int a = i, b = i+pattern.size();
28
           11 sub_hash = pfx_hash[b] - pfx_hash[a]+MOD;
           sub_hash %= MOD;
30
           if(sub_hash == (pattern_hash*ppow[a])%MOD){
31
               res++;
33
      }
34
35
      cout << res << endl;
    return 0;
36
```

6.3 Palindormes

Fiding longest palindromic substring.

```
int p[MX];
3 int main(){
       string txt1, txt = "#";
        cin >> txt1;
       for(char c : txt1){
6
            string s{c};
             txt += s + "#";
8
9
       txt = "^" + txt + "$";
10
       int l = 1, r = 1;
for(int i = 1; i < txt.size(); i++){
   p[i] = max(0, min(r-i, p[l+r-i]));</pre>
11
12
13
             while(txt[i-p[i]] == txt[i+p[i]])
14
15
                  p[i]++;
             if(i+p[i] > r){
16
                 1 = i-p[i];
17
                  r = i+p[i];
18
19
        }
20
21
        int maxi = 0;
       int imaxi;
for(int i = 0; i < txt.size(); i++){</pre>
22
23
             if(p[i] > maxi){
24
                 maxi = p[i];
imaxi = i;
25
26
27
       }
28
        string res = txt.substr(imaxi-maxi+1, maxi*2-1);
       for(auto c : res){
   if(c != '#')
30
31
                  cout << c;
32
```

7 Mathematics

7.1 Fermat's Theorem

```
x^{\varphi(m)} \mod m = x^{k \mod (m-1)} \mod m
```

```
int exp(int a, int b, int MOD){
    if(b == 0)
        return 1;
    if(b & 1){
        return ((11)a * exp(a, b-1, MOD))%MOD;
    }
    ll tmp = exp(a, b/2, MOD);
    return (tmp*tmp)%MOD;
}
```

7.2 Fast Fibonacci

```
void mul(11 a[][2], 11 b[][2]){
    ll res[2][2];
    res[0][0] = a[0][0]*b[0][0] + a[0][1]*b[1][0];
    res[0][1] = a[0][0]*b[0][1] + a[0][1]*b[1][1];
    res[1][0] = a[1][0]*b[0][0] + a[1][1]*b[1][0];
    res[1][1] = a[1][0]*b[0][1] + a[1][1]*b[1][1];

8    a[0][0] = res[0][0] % MOD;
   a[0][1] = res[0][1] % MOD;
   a[1][0] = res[1][0] % MOD;
   a[1][1] = res[1][1] % MOD;
```

```
12 }
13
14 int main(){
      ll n;
15
       cin >> n;
16
      11 m_pow[2][2] =
17
                {{0, 1},
{1, 1}};
18
19
     11 m[2][2] =
{{1, 0},
{0, 1}};
20
21
22
23
      auto pow = bitset<64> (n);
25
      for(int i = 0 ; i < 64; i++){
   if(pow[i]){</pre>
26
27
                mul(m,m_pow);
28
29
           mul(m_pow, m_pow);
30
31
       cout << m[0][1];
33
34
       return 0;
35
36 }
```

8 Geometry

9 Other Algorithms