

BEAR: sketching BFGS algorithm for ultra-high dimensional feature selection in sublinear memory

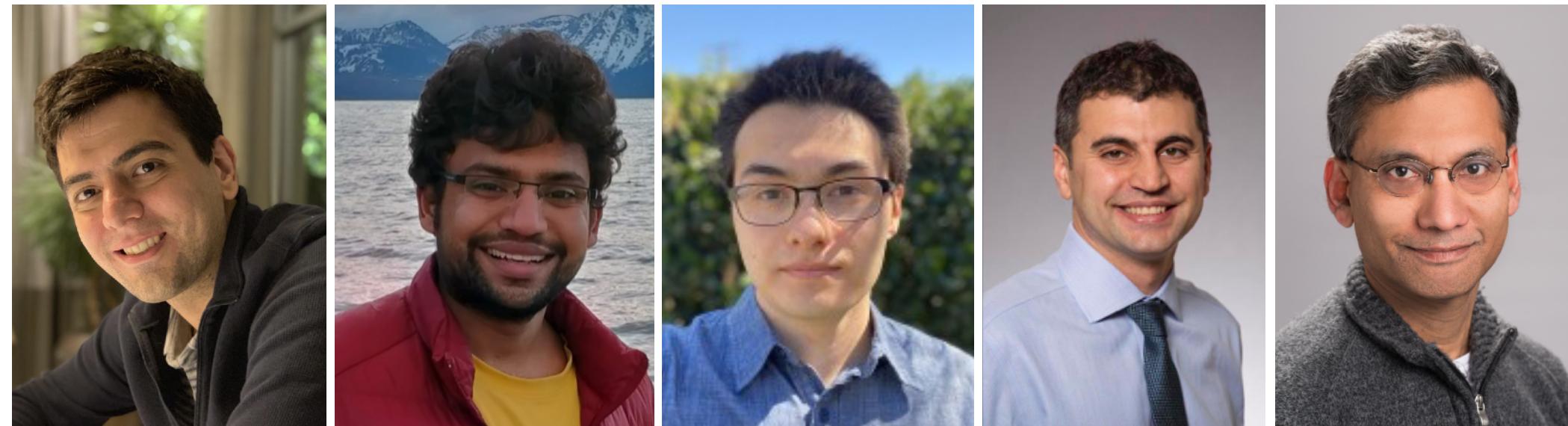
Amirali Aghazadeh

Vipul Gupta

Alex DeWeese

Ozan Koyluoglu

Kannan Ramchandran



Aug 16 - Aug 19th



big and high dimensional data in everyday life

- web services
- language processing
- networking
- genomics/proteomics
- health-care
- critical need for **scalable algorithms** to extract **important features** from the data
- limited computing **resource**



problem setup

- n data points $(\theta_i)_{i=1}^n = (\mathbf{x}_i, y_i)_{i=1}^n$ living in ultra-high dimensional feature space $\mathbf{x}_i \in \mathbb{R}^p$ ($p \gg n$)
goal: find a **small subset of features** best explains the output

>10¹⁵!

- k -**sparse** feature vector $\beta^* \in \mathbb{R}^p$
loss function $f(\beta, \theta) : \mathbb{R}^p \rightarrow \mathbb{R}$
- optimization problem $\min_{\beta} \sum_{i=1}^n f(\beta; \theta_i)$
- stochastic gradient descent (SGD) $\beta_{t+1} = \beta_t - \eta_t \mathbf{g}(\beta_t; \Theta_t)$
minibatch $\Theta_t = \{\theta_{t1}, \theta_{2t}, \dots, \theta_{tb}\}$
with the SGD term defined as $\mathbf{g}(\beta_t; \Theta_t) = \sum_{i=1}^b \nabla_{\beta_t} f(\beta_t; \theta_{ti})$

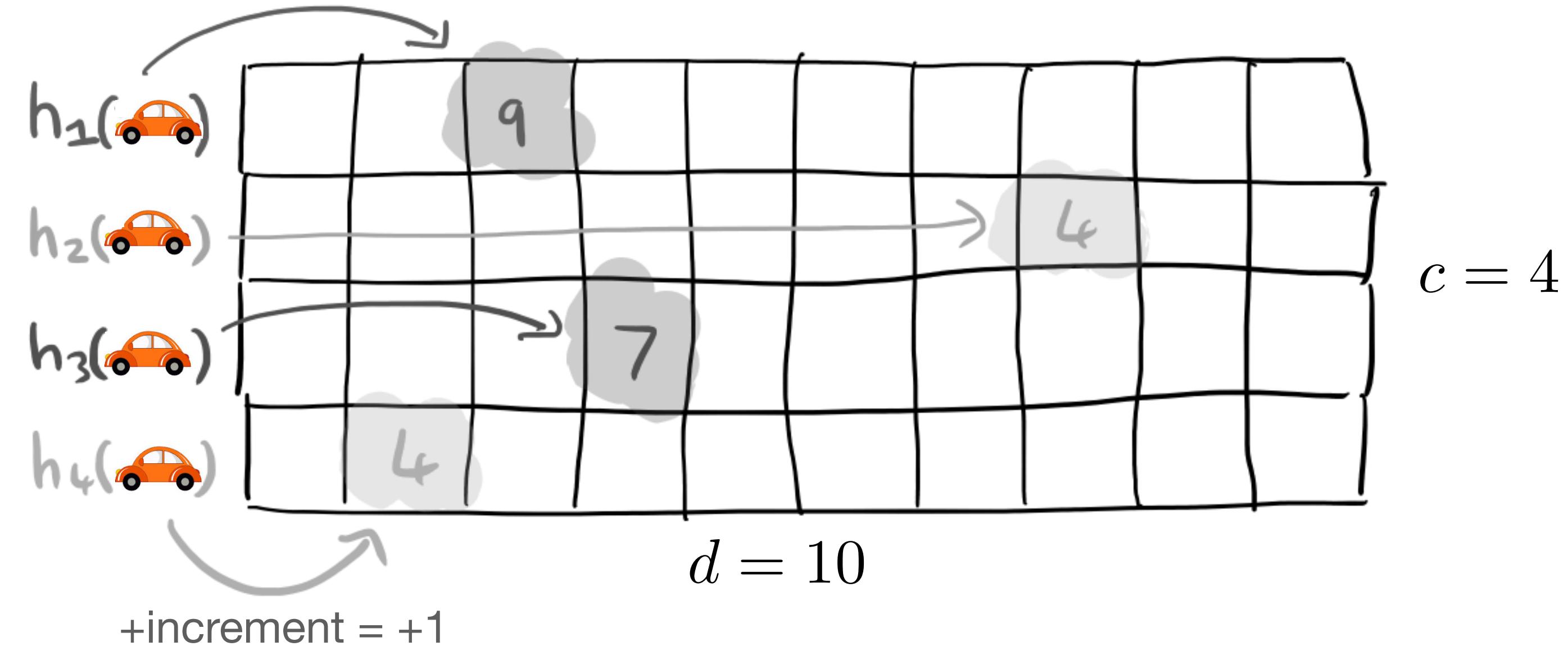
challenge: not enough **memory** to store the **intermediately dense** feature vector β (sublinear alg.)

Count Sketch (CS)

- data structure to **compressively** store the **number of occurrences** of many number of streaming items

random hash
function

$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, d\}$$



- fast operations

- ADD (item, increment)
- QUERY (item)

$$\# \text{ car} \approx \text{median}(\{4, 4, 7, 9\})$$

items (p)

$$m = d \times c$$

frequent items (k)

all colors

memory of CS

top colors



Count Sketch (CS)

- data structure to **compressively** store the **number of occurrences**

random hash
function

$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, d\}$$



Theorem 1 *Charikar et al. (2002) Count Sketch finds top- k items z_i with $\pm\varepsilon\|\mathbf{z}\|_2$ error, with probability at least $1 - \delta$, in space $\mathcal{O}(\log(\frac{p}{\delta})(k + \frac{\|\mathbf{z}^{\text{tail}}\|_2^2}{(\varepsilon\zeta)^2}))$, where $\|\mathbf{z}^{\text{tail}}\|_2^2 = \sum_{i \notin \text{top-}k} z_i^2$ is the energy of the non-top- k items and ζ is the k^{th} largest value in \mathbf{z} .*

- ADD (item, increment)
- QUERY (item)

$$\# \text{ cars} \approx \text{median}(\{4, 4, 7, 9\})$$

items (p)

$$m = d \times c$$

all colors

memory of CS

frequent items (k)

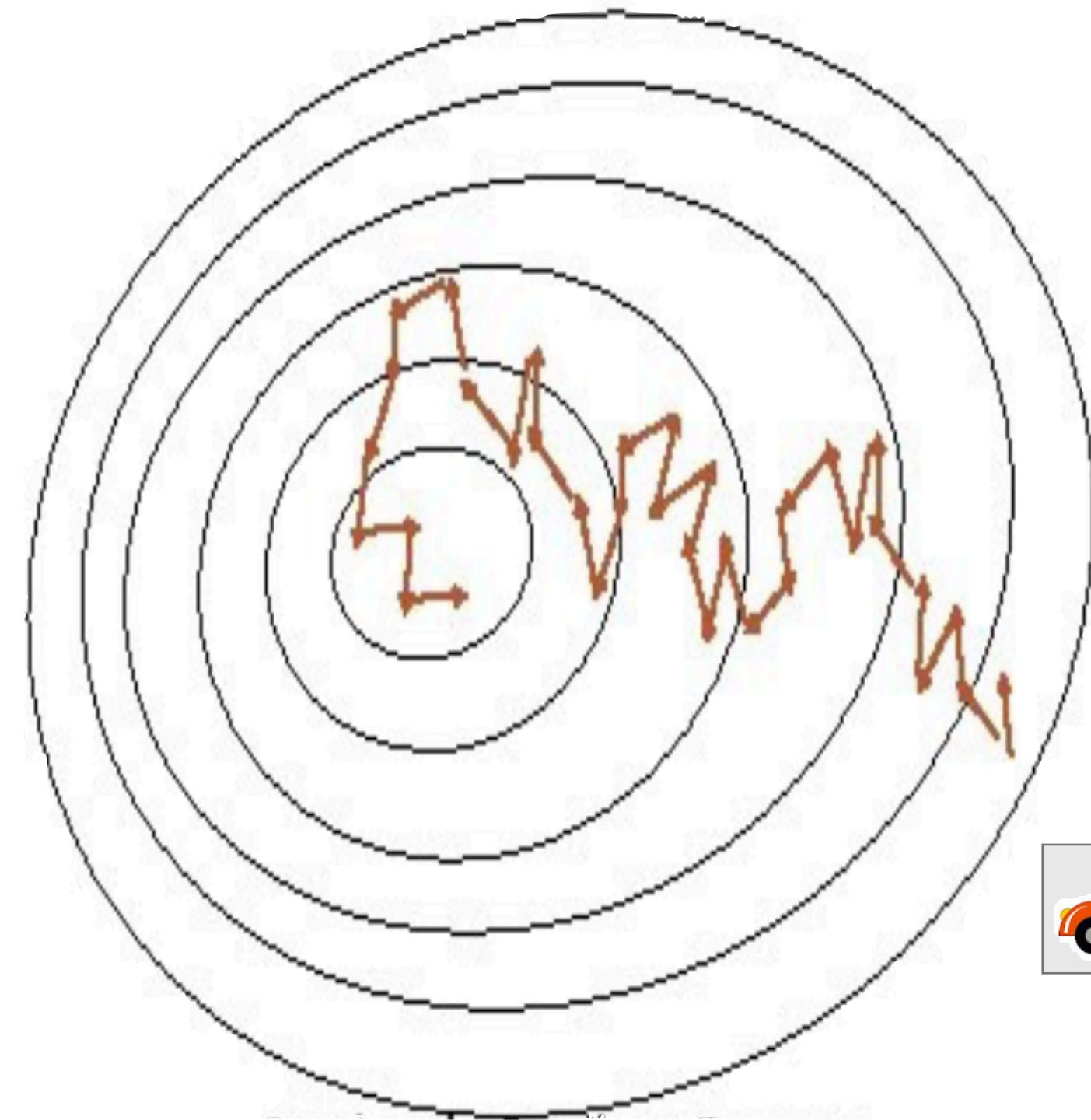
top colors



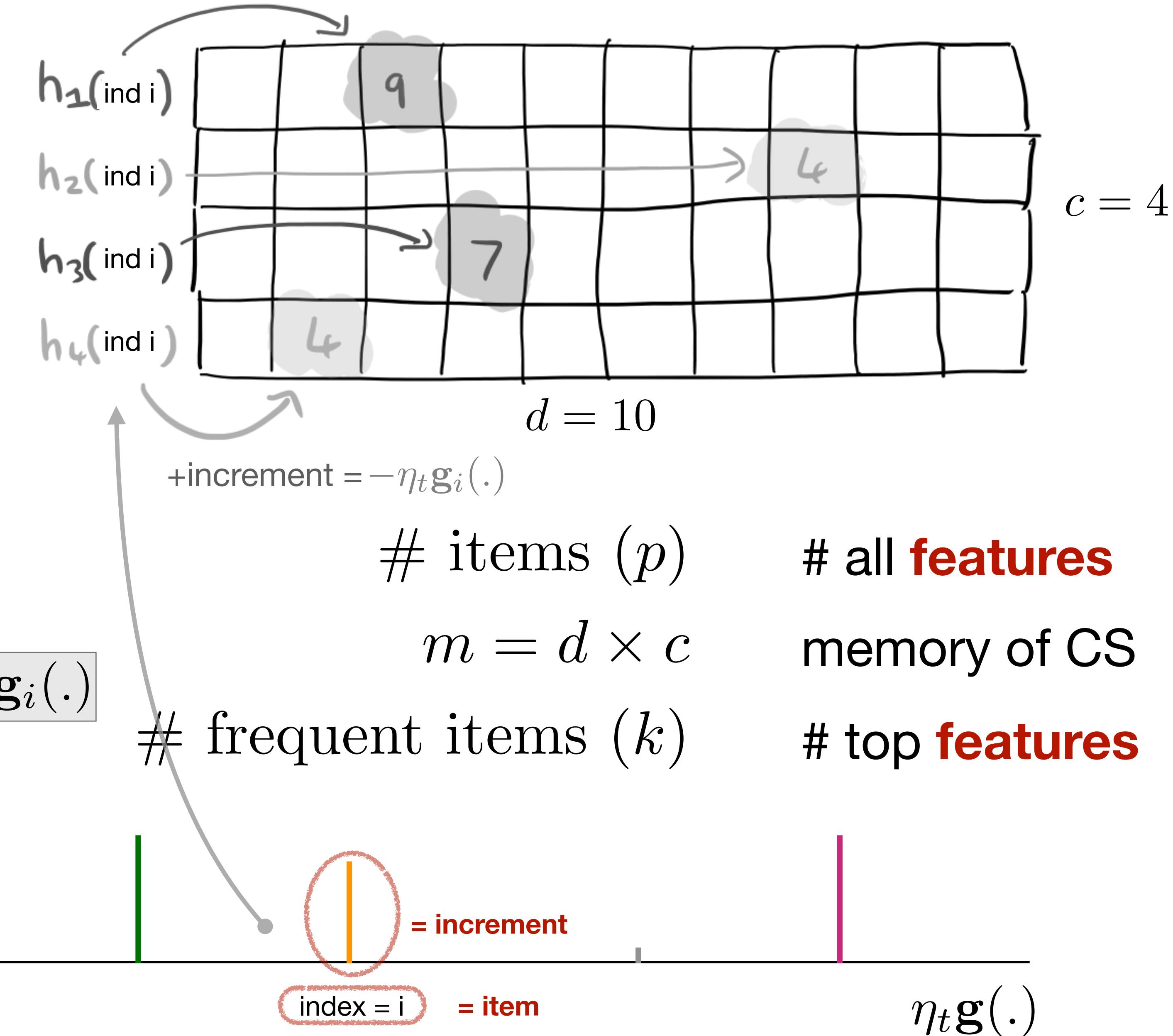
feature selection with CS

$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, d\}$$

$$\beta_{t+1} = \beta_t - \eta_t g(\beta_t; \Theta_t)$$

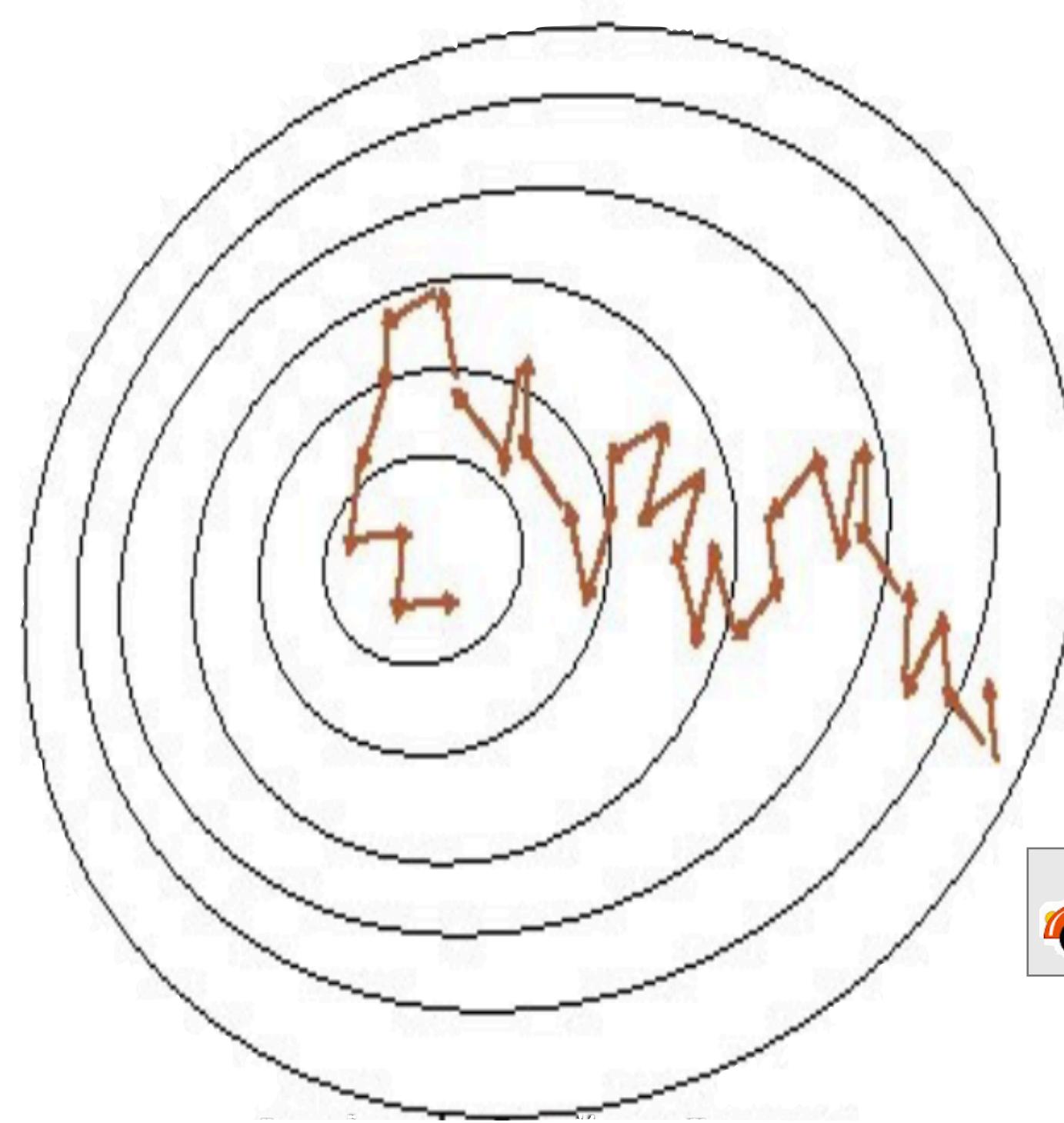


$\equiv \eta_t g_i(\cdot)$



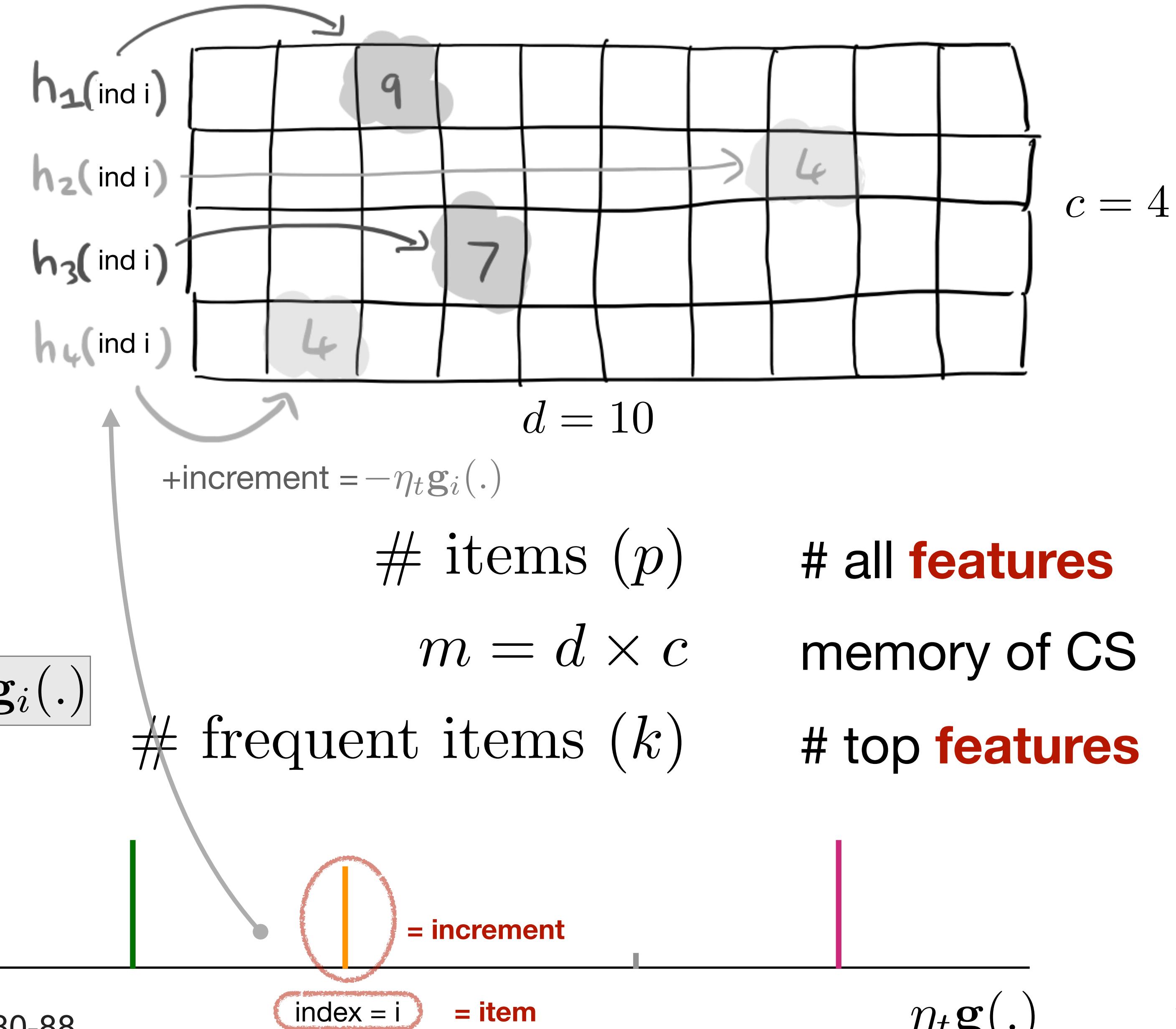
feature selection with CS

MISSION : $\beta_{t+1}^s = \beta_t^s - \eta_t \mathbf{g}^s(\text{Query}_{\text{top-}k}(\beta_t^s); \Theta_t)$



 $\equiv \eta_t \mathbf{g}_i(\cdot)$

$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, d\}$$



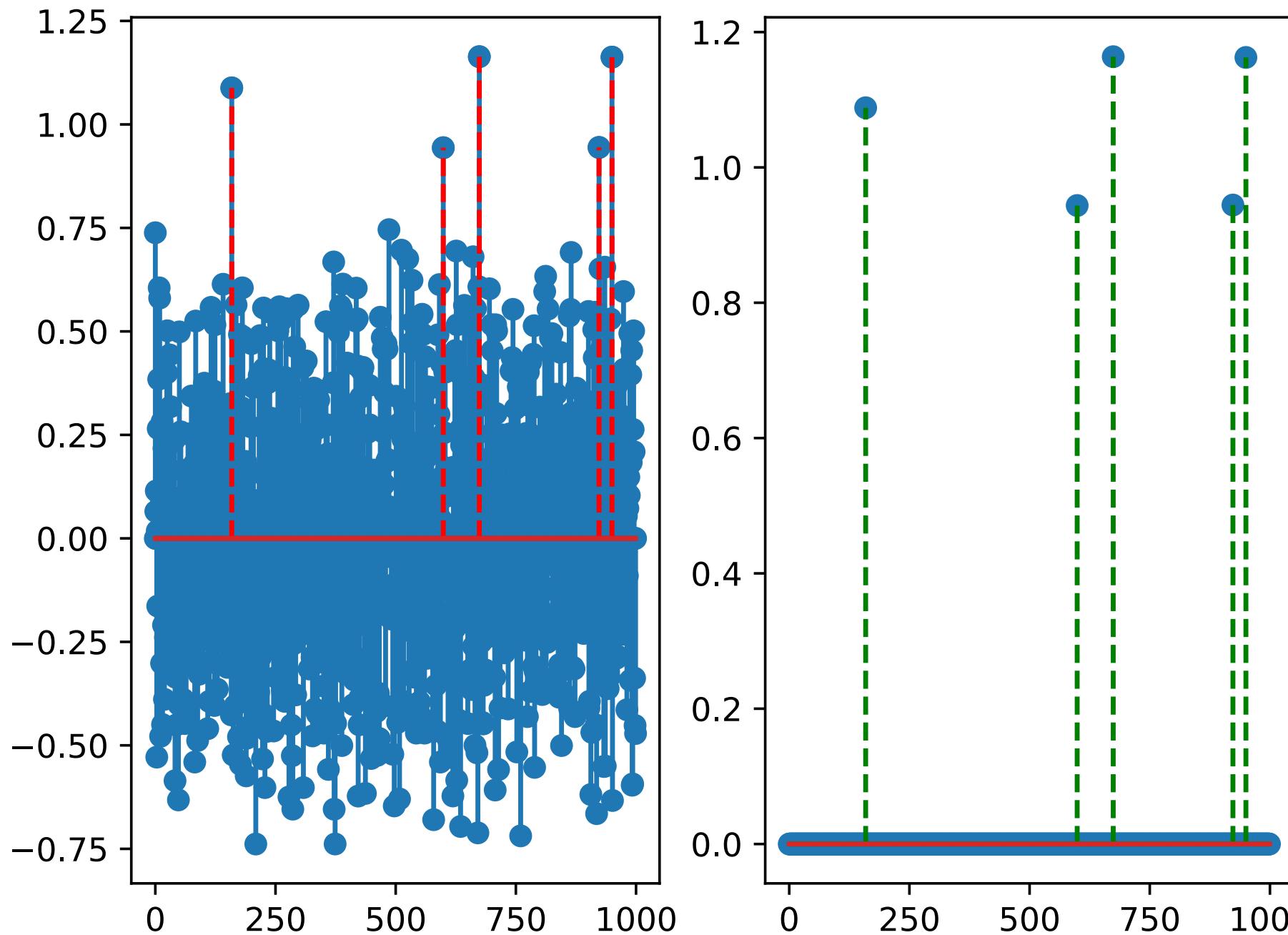
feature selection with CS

MISSION : $\beta_{t+1}^s = \beta_t^s - \eta_t g^s(\text{Query}_{\text{top}-k}(\beta_t^s); \Theta_t)$

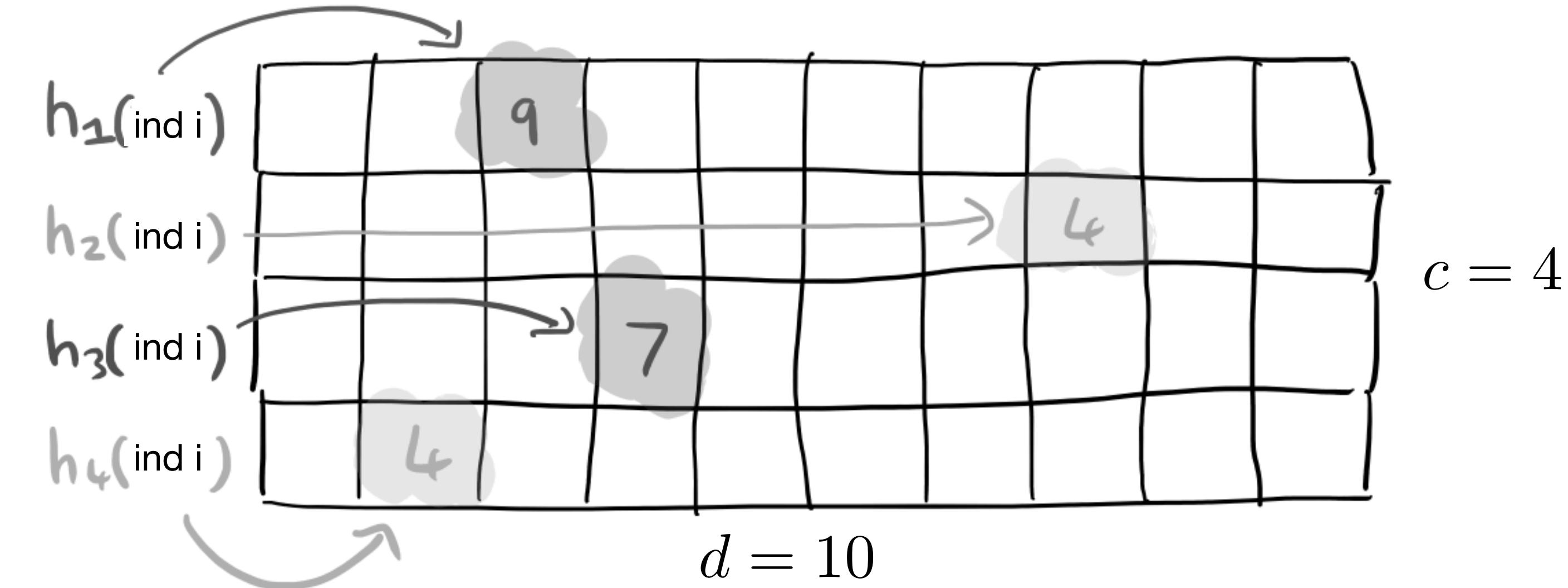
after convergence

content of CS

ground truth



$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, d\}$$



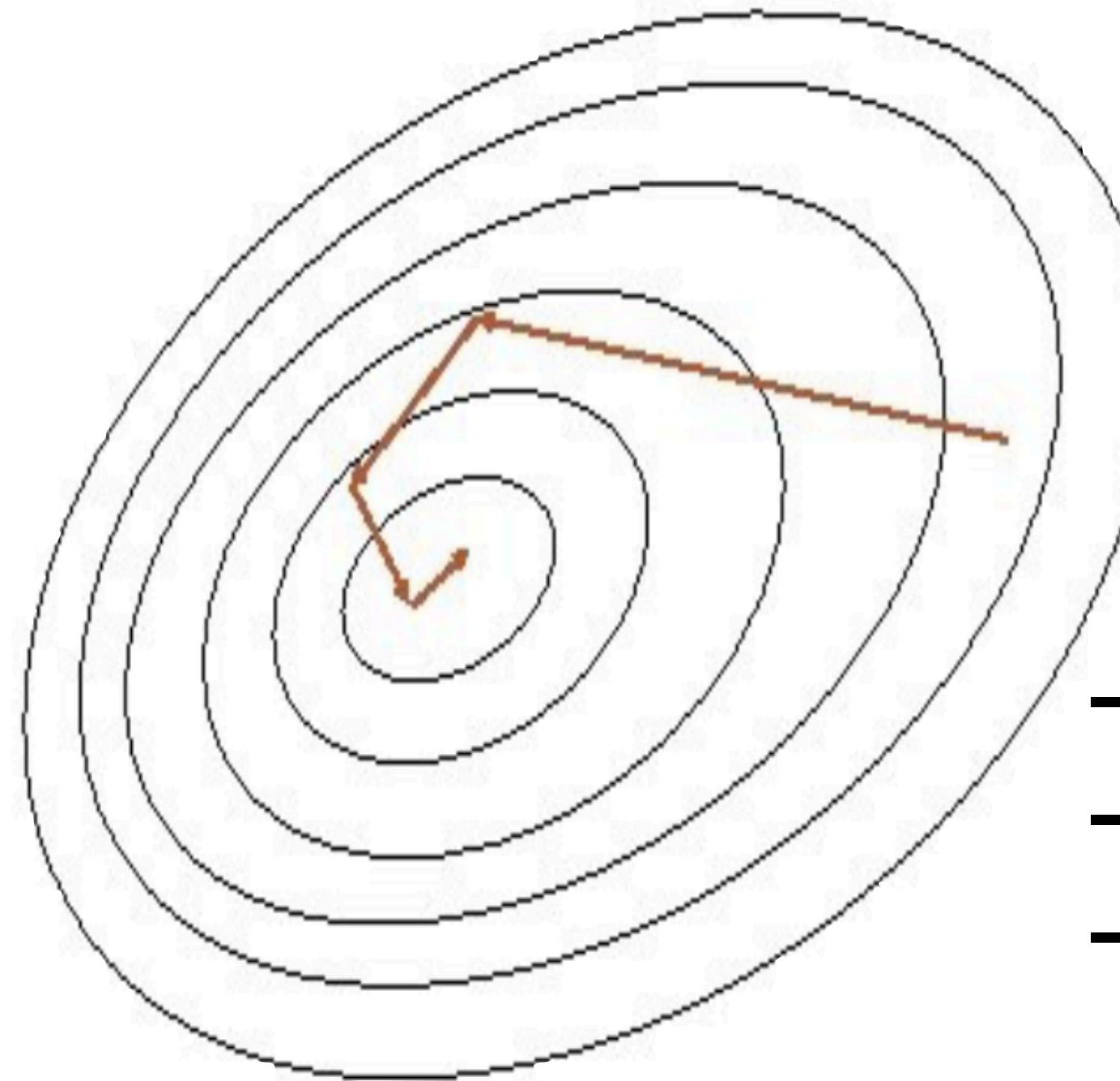
observation: sketch of noisy component of SGD in CS do not cancel out and results in memory wasted to store sketched noise

Theorem 1 Charikar et al. (2002) Count Sketch finds top- k items z_i with $\pm \varepsilon \|z\|_2$ error, with probability at least $1 - \delta$, in space $\mathcal{O}(\log(\frac{p}{\delta})(k + \frac{\|z^{tail}\|_2^2}{(\varepsilon \zeta)^2}))$, where $\|z^{tail}\|_2^2 = \sum_{i \notin \text{top-}k} z_i^2$ is the energy of the non-top- k items and ζ is the k^{th} largest value in z .

idea: second order sketching

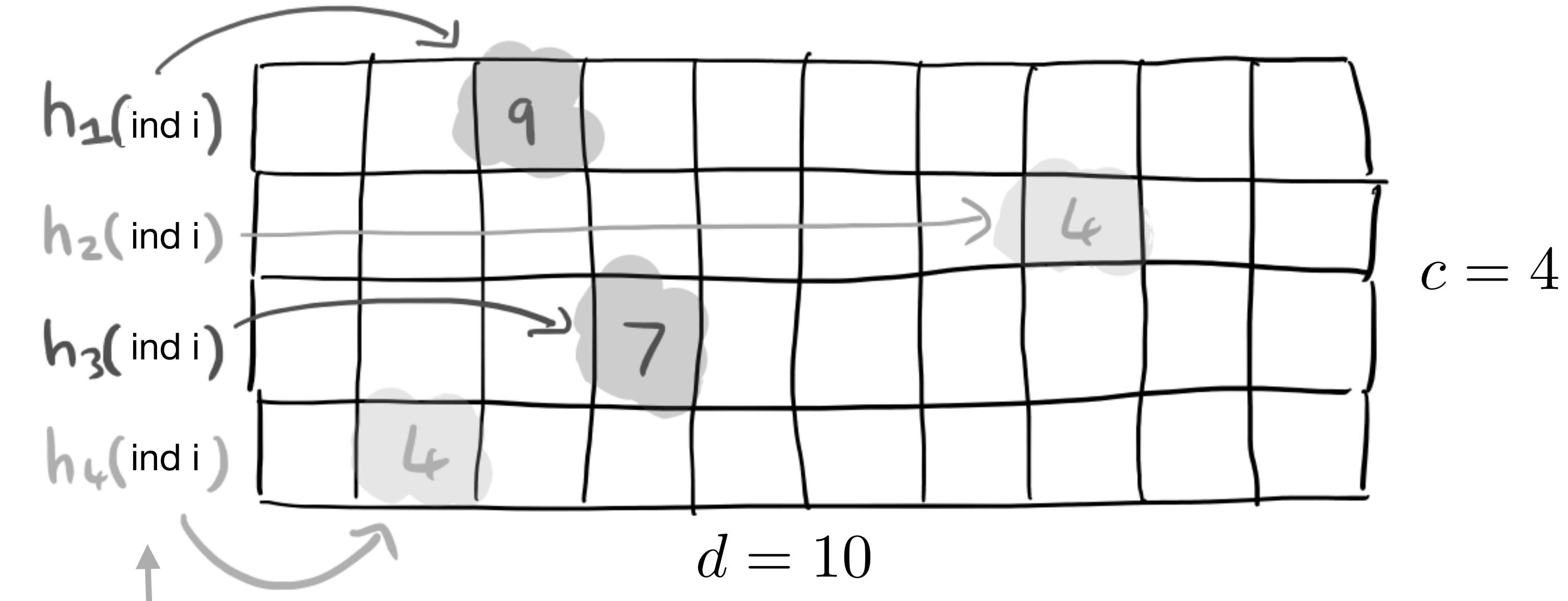
$$\beta_{t+1} = \beta_t - \eta_t \mathbf{B}_t^{-1} \mathbf{g}(\beta_t, \Theta_t)$$

$$\mathbf{B}_t = \nabla_{\beta_t}^2 f(\beta_t, \Theta_t) \in \mathbb{R}^{p \times p}$$

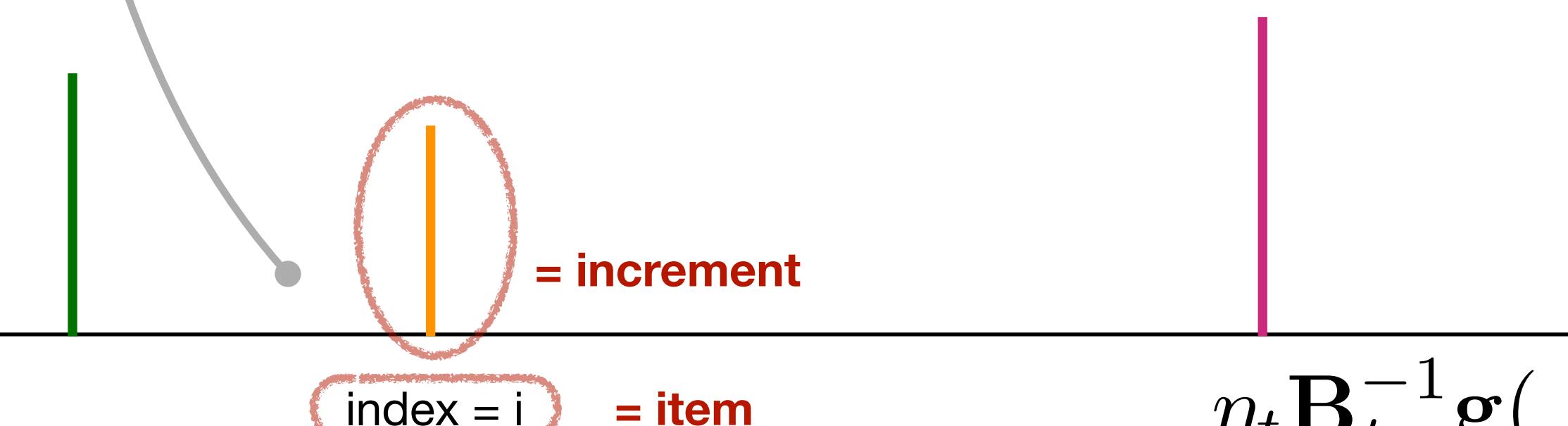


- more comp. cost per iteration
- less noisy gradient
- memory-accuracy tradeoff

$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, d\}$$



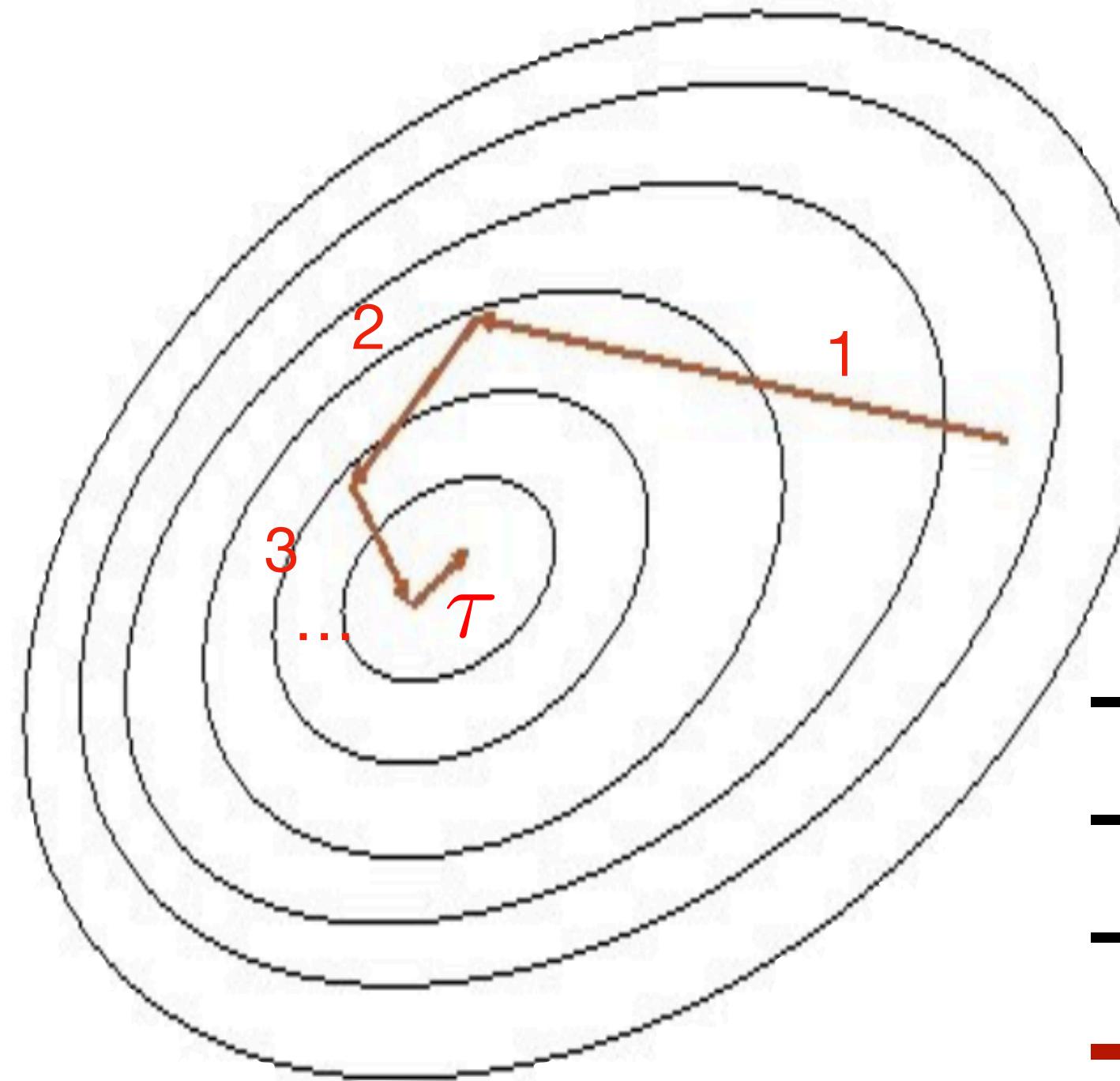
**question: how to compute/
store the Hessian?**



limited-memory BFGS

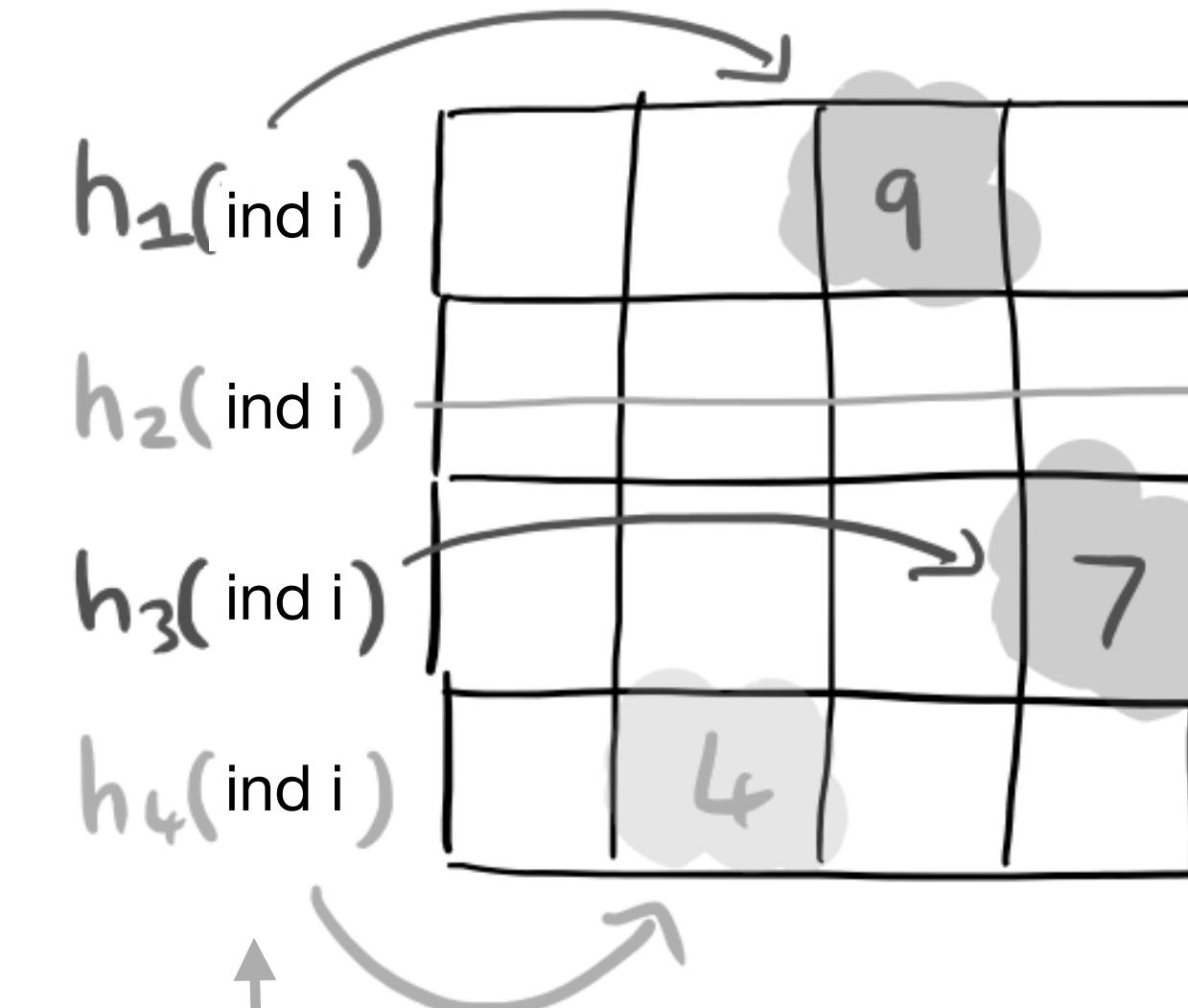
$$\beta_{t+1} = \beta_t - \eta_t \mathbf{B}_t^{-1} \mathbf{g}(\beta_t, \Theta_t)$$

$$\mathbf{B}_t = \nabla_{\beta_t}^2 f(\beta_t, \Theta_t) \in \mathbb{R}^{p \times p}$$



- more comp. cost per iteration
- less noisy gradient
- memory-accuracy tradeoff
- no need to store/compute inverse Hessian**

$$h_j : \{1, 2, \dots, p\} \rightarrow \{1, 2, \dots, c\}$$



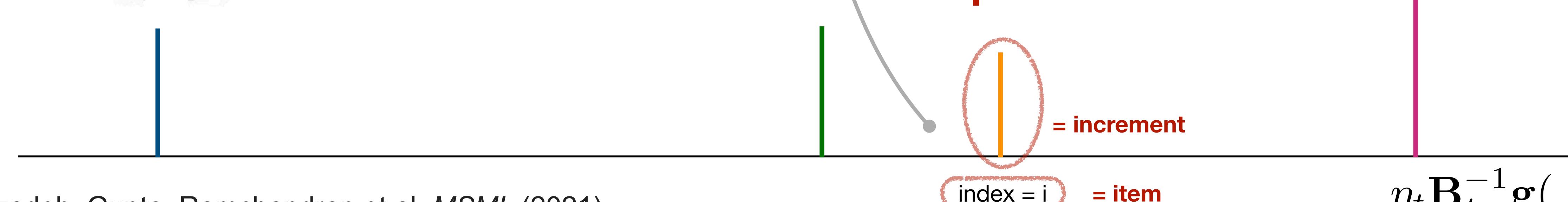
Algorithm 1 Limited-memory BFGS

Input: $\mathbf{g}(\hat{\beta}_t, \Theta_t)$ and $\{\mathbf{s}_i, \mathbf{r}_i\}_{i=t-\tau+1}^t$

1. $\rho_t = \frac{1}{\mathbf{r}_t^T \mathbf{s}_t}$.
2. $\mathbf{q}_t = \mathbf{g}(\hat{\beta}_t, \Theta_t)$,
for $i = t$ to $t - \tau + 1$:
 $\alpha_i = \rho_i \mathbf{s}_i^T \mathbf{q}_i$,
 $\mathbf{q}_{i-1} = \mathbf{q}_i - \alpha_i \mathbf{r}_i$.
3. $\mathbf{z}_{t-\tau} = \frac{\mathbf{r}_t^T \mathbf{s}_t}{\mathbf{r}_t^T \mathbf{r}_t} \mathbf{q}_{t-\tau}$,
for $i = t - \tau + 1$ to t :
 $\gamma_i = \rho_i \mathbf{r}_i^T \mathbf{z}_i$.
 $\mathbf{z}_i = \mathbf{z}_{i-1} + \mathbf{s}_i (\alpha_i - \gamma_i)$.

Return: \mathbf{z}_t

approximate $\mathbf{B}_t^{-1} \mathbf{g}(\cdot)$ using gradients from last few τ iterations



BEAR algorithm: sketch LBFGS gradients using CS

Algorithm 2 BEAR

Initialize: $t = 0$, Count Sketch $\beta_{t=0}^s = 0$, top- k heap.

while stopping criteria not satisfied **do**

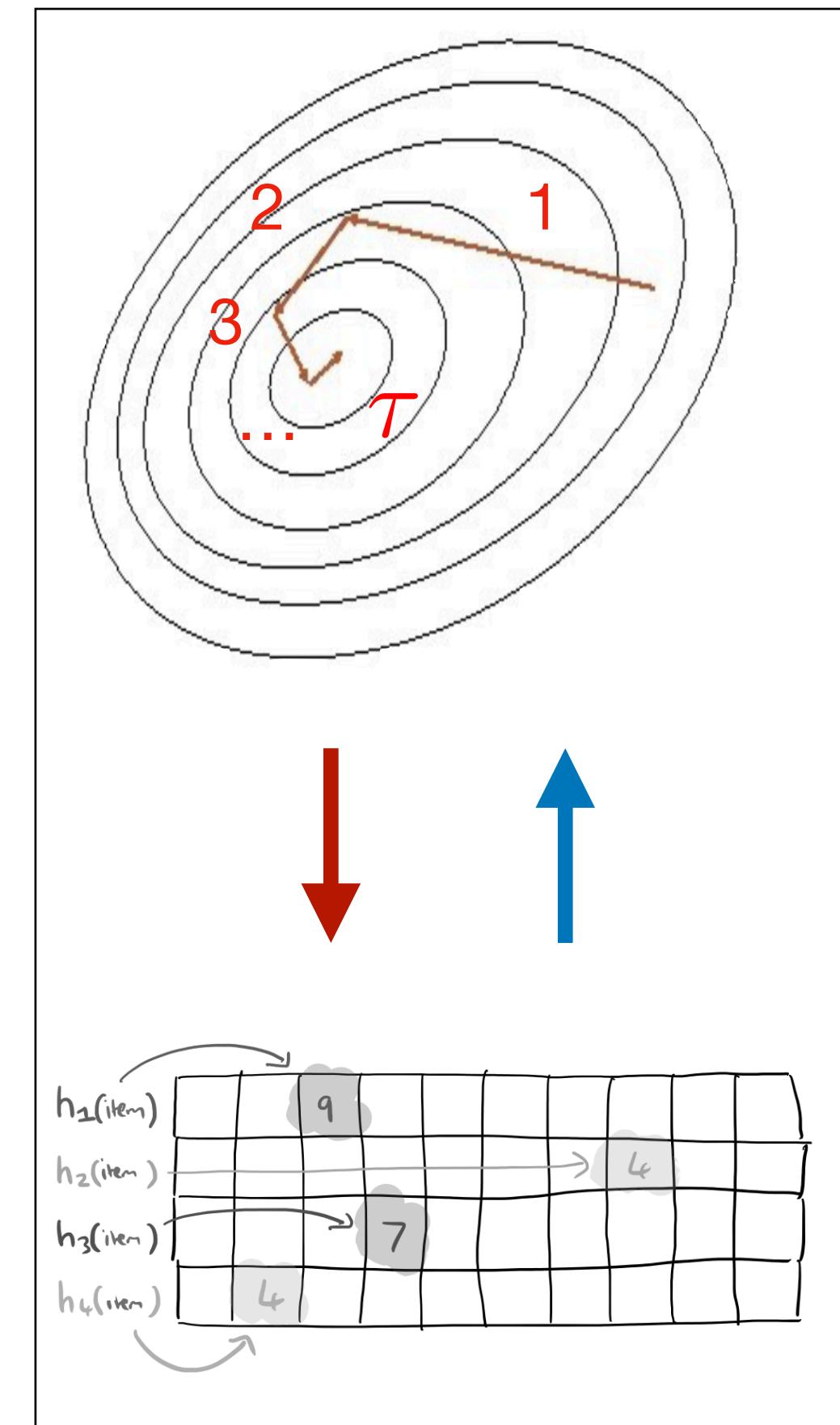
1. Sample b independent data points in a minibatch $\Theta_t = \{\theta_{t1}, \dots, \theta_{tb}\}$.
2. Find the active set \mathcal{A}_t .
3. QUERY the feature weights in $\mathcal{A}_t \cap$ top- k from Count Sketch $\beta_t = \text{query}(\beta_t^s)$.
4. Compute stochastic gradient $\mathbf{g}(\beta_t, \Theta_t)$.
5. Compute the descent direction with Alg. 1 $\mathbf{z}_t = \text{LBFGS}(\mathbf{g}(\beta_t, \Theta_t), \{\mathbf{s}_i, \mathbf{r}_i\}_{i=t-\tau+1}^t)$.
6. ADD the sketch of \mathbf{z}_t at the active set $\hat{\mathbf{z}}_t = \mathbf{z}_t^{\mathcal{A}_t}$ to Count Sketch $\beta_{t+1}^s := \beta_t^s - \eta_t \hat{\mathbf{z}}_t^s$.
7. QUERY the features weights in $\mathcal{A}_t \cap$ top- k from Count Sketch $\beta_{t+1} = \text{query}(\beta_{t+1}^s)$.
8. Compute stochastic gradient $\mathbf{g}(\beta_{t+1}, \Theta_t)$.
9. Set $\mathbf{s}_{t+1} = \beta_{t+1} - \beta_t$, and $\mathbf{r}_{t+1} = \mathbf{g}(\beta_{t+1}, \Theta_t) - \mathbf{g}(\beta_t, \Theta_t)$.
10. Update the top- k heap.
11. $t = t + 1$.

end while

Return: The top- k heavy-hitters in Count Sketch.

find the descent
direction using LBFGS
and update CS

query CS and store the
gradient and feature
difference vectors



convergence

Theorem 2 Let $f(\cdot)$ and the step sizes η_t satisfy the assumptions above. Let the size of Count Sketch be $m = \theta(\varepsilon^{-2} \log 1/\delta)$ with number of hashes $d = \theta(\varepsilon^{-1} \log 1/\delta)$ for $\varepsilon, \delta > 0$. Then, the Euclidean distance between updates β_t^s in the BEAR algorithm and the sketch of the solution of problem (1) converges to zero with probability $1 - \delta$, that is,

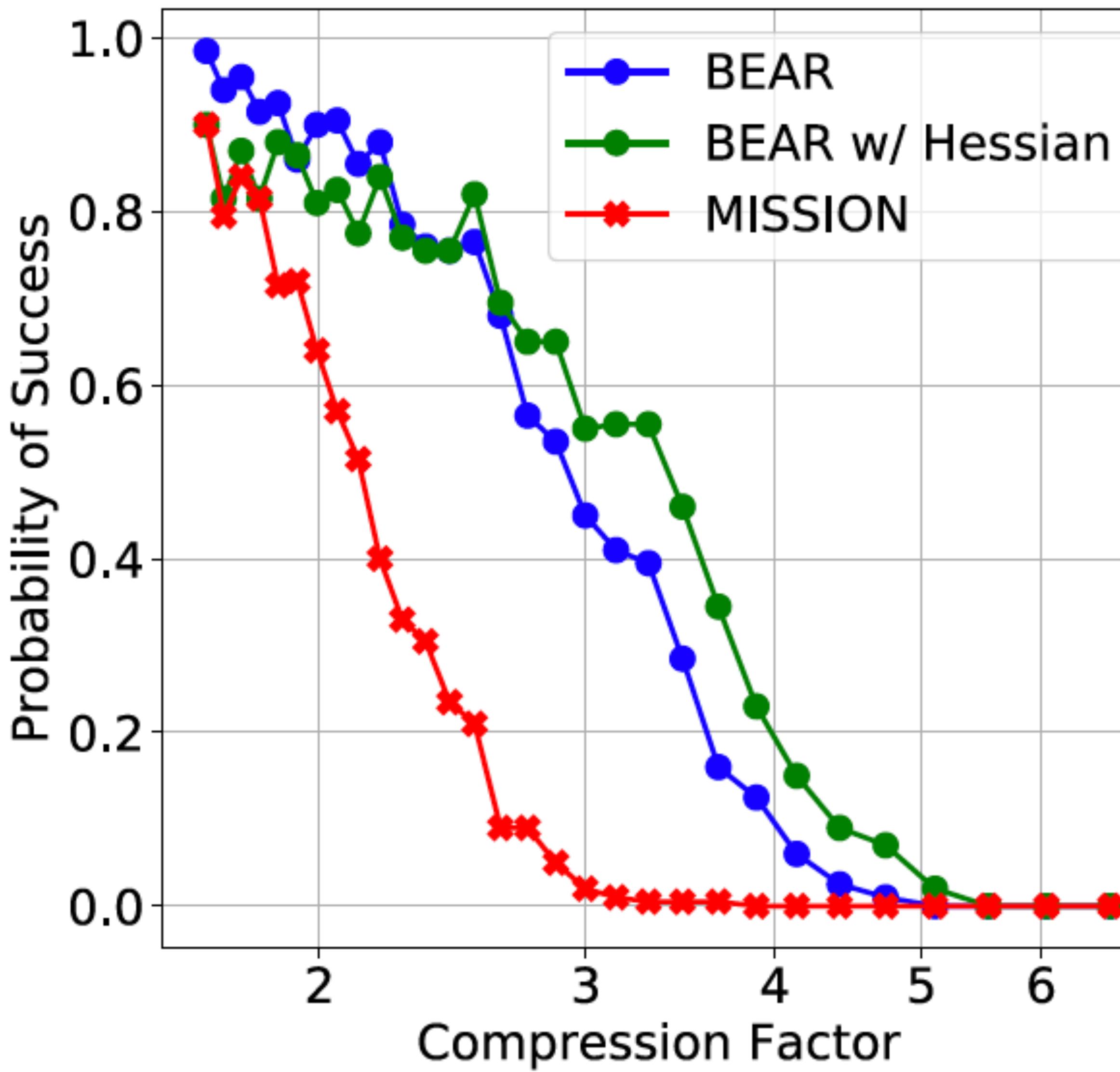
$$\mathbb{P}\left(\lim_{t \rightarrow \infty} \|\beta_t^s - \beta^{s*}\|^2 = 0\right) = 1 - \delta, \quad (2)$$

where the probability is over the random realizations of random samples $\{\Theta_t\}_{t=0}^\infty$. Furthermore, for the specific step size $\eta_t = \eta_0/(t + T_0)$ for some constants η_0 and T_0 , the model parameters at iteration t satisfy

$$\mathbb{E}_\Theta [f(\beta_t^s, \Theta) - \mathbb{E}[f(\beta^{s*}, \Theta)] \leq \frac{C_0}{T_0 + t}, \quad (3)$$

with probability $1 - \delta$. Here, C_0 is a constant depending on the parameters of the sketching scheme, the above assumptions, and the objective function.

simulations

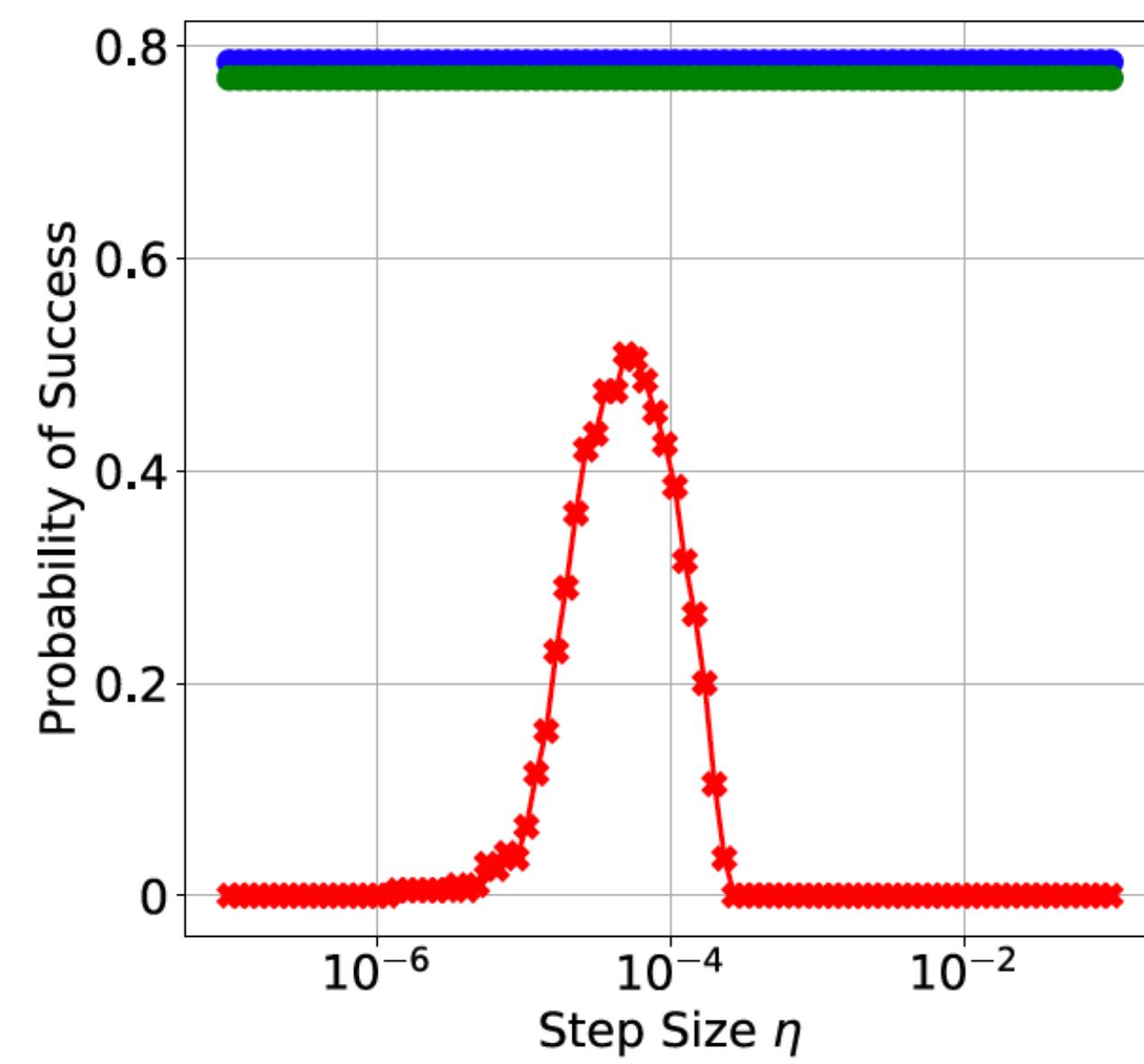
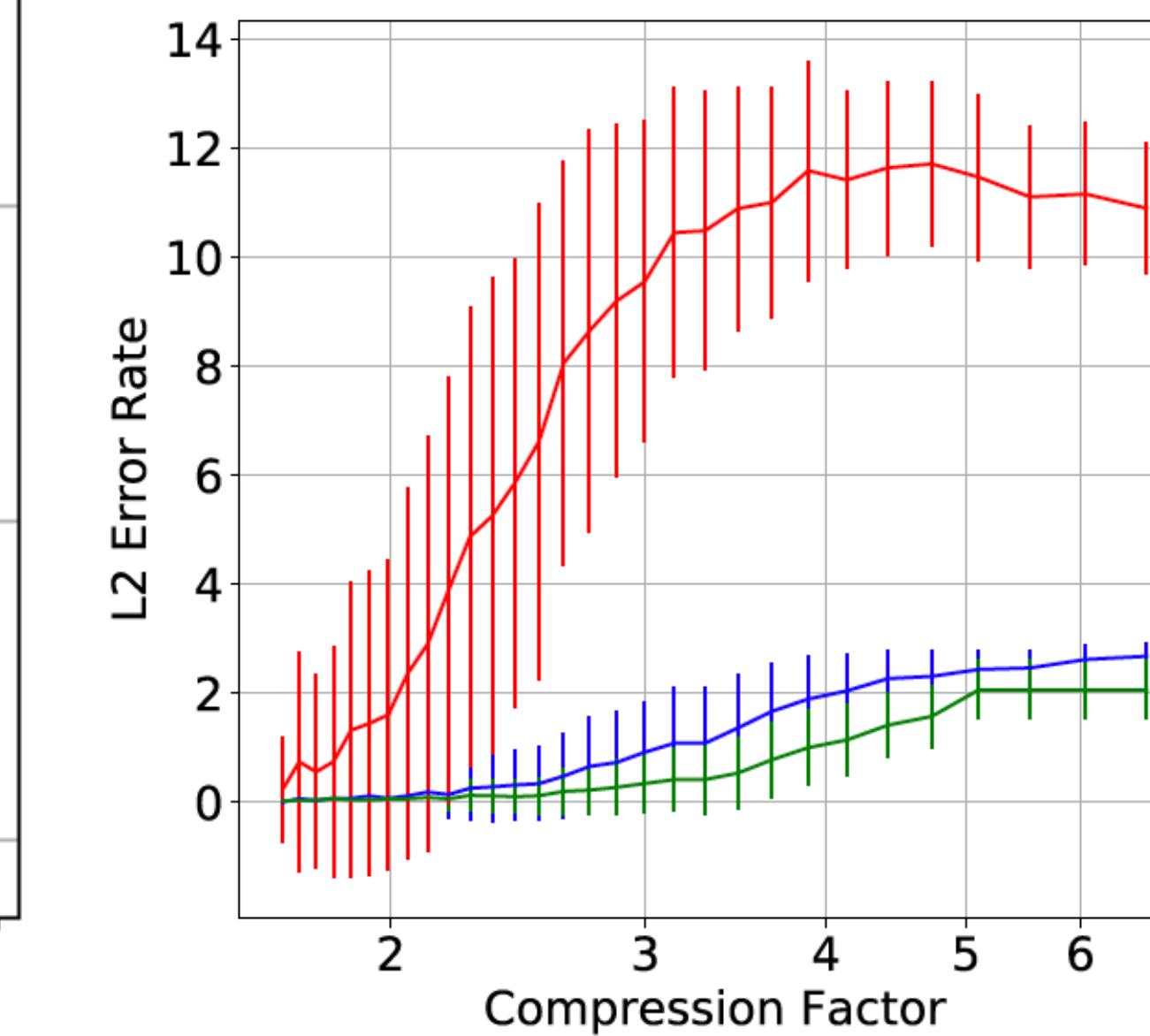


$$\mathbf{x}_i \sim \mathcal{N}(0, 1)$$

$$y_i = \mathbf{x}_i \boldsymbol{\beta}^*$$

$$\boldsymbol{\beta}^* : k - \text{sparse}$$

$$\text{CF} = \frac{p}{\text{size of CS}}$$



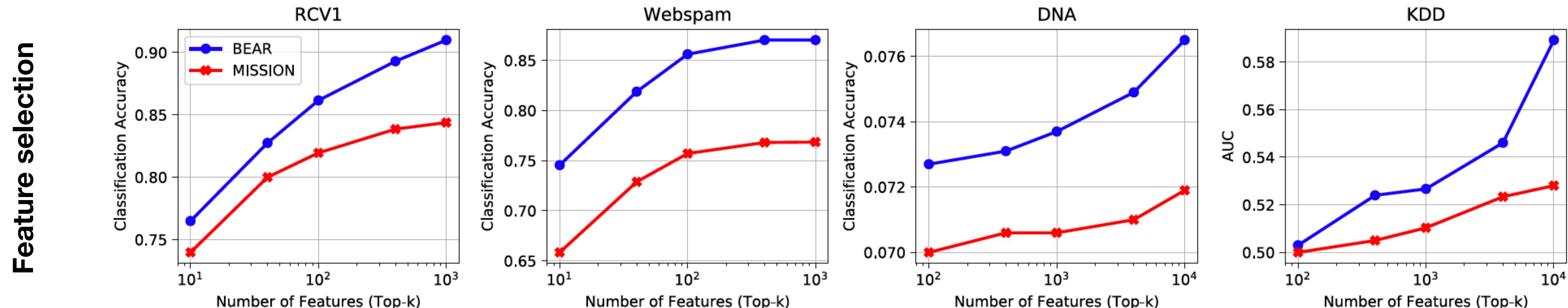
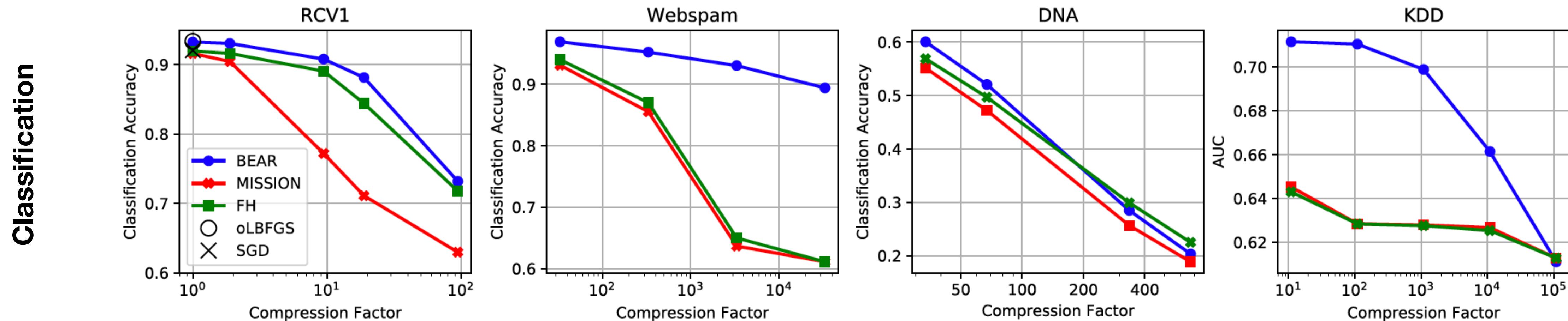
$$p = 1000$$

$$n = 900$$

$$k = 8$$

real-world experiments

Data set	Dim (p)	#Train (n)	#Test	Size	#Act.
RCV1	47,236	20,242	677,399	1.2GB	73
Webspam	16,609,143	280,000	70,000	25GB	3730
DNA	16,777,216	600,000	600,000	1.5GB	89
KDD 2012	54,686,452	119,705,032	29,934,073	22GB	12



summary and future directions

- **adaptively** learn the hashing scheme in the Count Sketch based on the stochastic gradients
- efficient training of massively large **nonlinear models** using LBFGS + sketching (Transforms, etc.)
- **distributed** learning/analysis using LBFGS + sketching explore communication-computation tradeoff

Thanks!

- find the **paper** at <https://arxiv.org/abs/2010.13829>
- find the **code** at <https://github.com/BEAR-algorithm/BEAR>