# AzureML로 시작하는 Machine Learning

**Microsoft Student Partner** 

KAIST 정태영

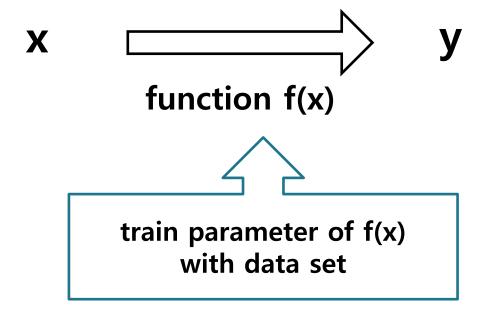
#### Session

Session 1: Introduction to Introduction to ML
 머신러닝의 기초의 기초에 대해 배웁니다.

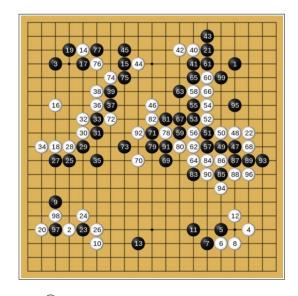
- Session 2 : Azure Machine Learning 맛보기
  - Azure ML studio를 통해 붓꽃을 분류하는 머신 러닝 프로그램을 만들어봅니다.

#### What is ML?

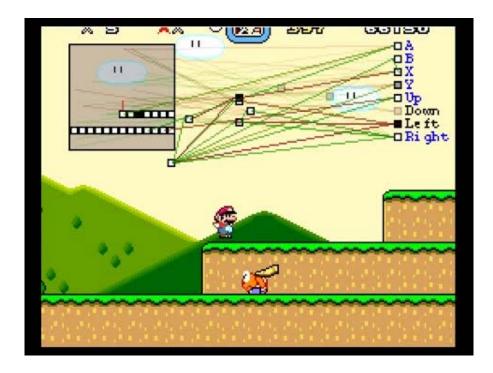
- 명시적으로 프로그래밍 되지 않은 것을 컴퓨터가 학습할 수 있게 하는 학문
- 데이터를 통해 학습하는 알고리즘에 관한 학문



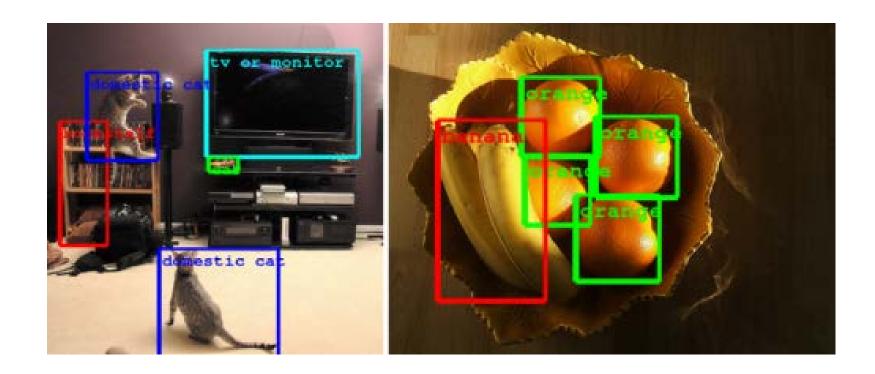
• 인공지능





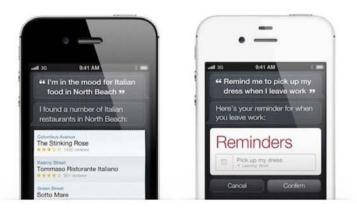


• 시각 정보 처리



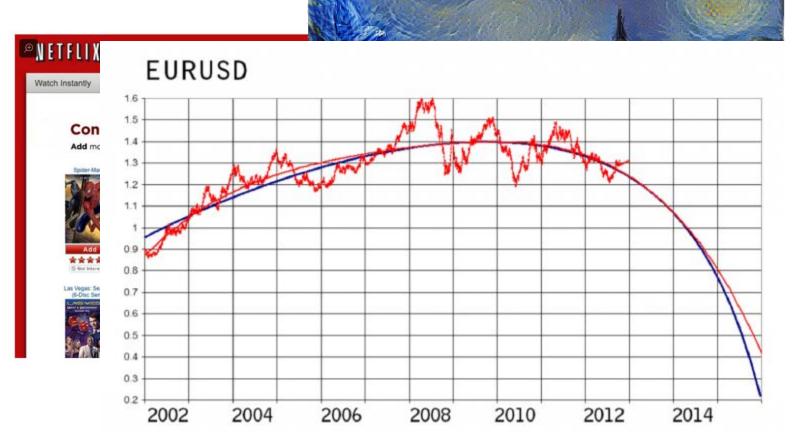
언어







기타

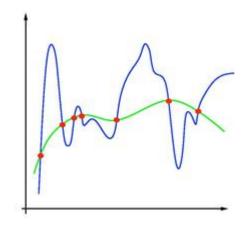


### Why ML?

- 특히, 왜 deep learning이 떠오를까?
- 과거 Deep learning의 문제들



**Computation Cost** 



**Overfitting** 

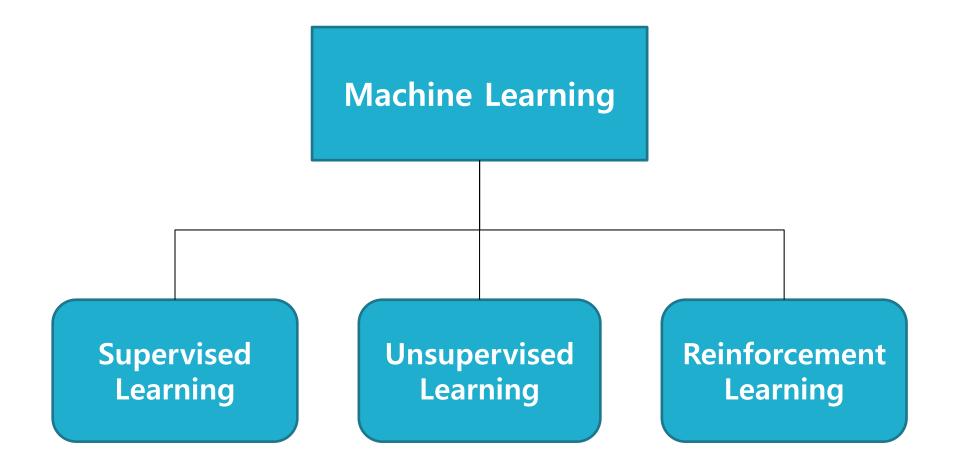


No theorical base

### Why ML?

- 현재
- Overfitting: Regularization(dropout, ReLU..)
- 연산 능력: 하드웨어의 발전, 병렬 처리
- 이론적 배경 : 실제 결과가 좋음

### **Category of ML**



### Category of ML

- 지도학습: Label된 데이터를 통해 훈련
- 회귀분석(regression), 분류(classification)
- 집값 예측, 개/고양이 사진 판단
- 비지도학습: Label이 없는 데이터를 통해 훈련
- Clustering, Dimensionality reduction
- 강화학습: (상태, 행동)에 대한 보상으로 훈련
- Markov Decission Process
- 게임 AI(알파고)

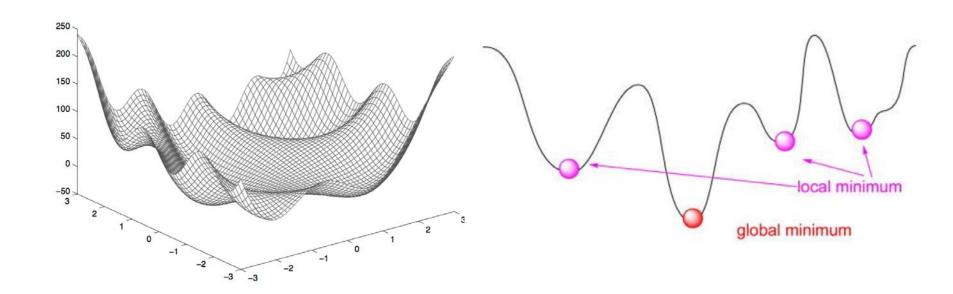


- 확률, 통계, 선형대수, 최적화 이론
- Optimization

minimize 
$$f_0(x)$$
  
s.t  $f_i(x) \le 0, i = 1, 2, ..., k$   
 $h_j(x) = 0, j = 1, 2, ..., l$ 

 복잡한 그래프에서는 최적의 해를 찾기 힘듦

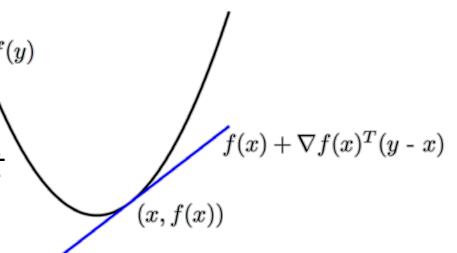






 Convex Optimization (볼록최적화)

볼록한 함수에서 최적해를 찾는 것은 훨씬 쉽다.



minimize 
$$f_0(x)$$
  
s.t  $f_i(x) \le 0, i = 1, 2, ..., k$   
 $h_j(x) = 0, j = 1, 2, ..., l$ 

Lagrangian function with 
$$\lambda_i \geq 0, v_j \in R$$

$$L(x,\lambda,v) = f_0(x) + \sum_{i=1}^k \lambda_i f_i(x) + \sum_{j=1}^l v_j h_j(x)$$

$$maximize_{\lambda \geq 0,v} \left[\inf_{v} L(x,\lambda,v)\right]$$

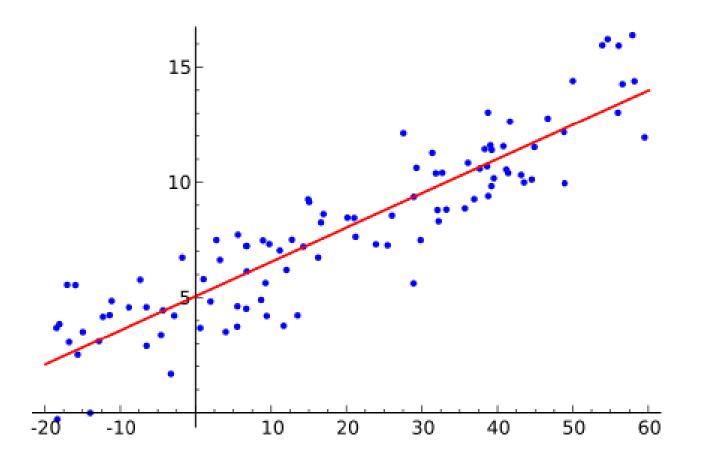


Gradient Descent Method

현재 기울기를 계산해 더 좋은 방향으로 이동  $J(\theta_0,\theta_1)$ -10 -20 -10 t is too small t is too large appropriate 16



• 회귀분석: 주어진 데이터에 가장 적합한 함수 찾기

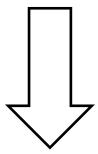


• 선형회귀 : 선형 함수를 찾는 회귀분석  $X \in \mathbb{R}^{n \times p}, y \in \mathbb{R}^n$ 일 때 y = XA에 가장 가까운  $A \in \mathbb{R}^p$ 

$$minimize_A \|y - XA\|_2^2$$

$$X = \begin{bmatrix} x_{11} & \cdots & x_{1p} \\ \vdots & \ddots & \vdots \\ x_{n1} & \cdots & x_{np} \end{bmatrix}, y = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{bmatrix}, A = \begin{bmatrix} a_1 \\ a_2 \\ \vdots \\ a_p \end{bmatrix}$$

$$\nabla \|y - XA\|_2^2 = 0$$

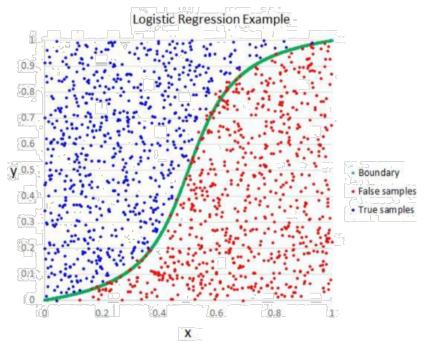


$$A = (X^T X)^{-1} X^T y$$



• 로지스틱회귀 : 분류를 위한 모델

$$P(y = 1|X; A) = \frac{e^{XA}}{1 + e^{XA}}$$
Logistic Regression Example



Multi-class logistic regression(softmax)

$$P(y = k | X; A_1, ..., A_k) = \frac{e^{XA_i}}{\sum_{i=0}^k e^{XA_i}}, A_0 = \vec{0}$$

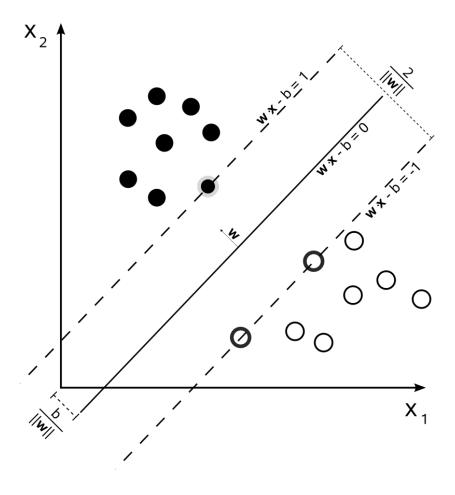
Binary class logistic regression objective

$$\log \left( \prod_{i=1}^{n} P(y_i|x_i;A) \right)$$

$$= \sum_{i=1}^{n} y_i log p_i(A) + (1-y_i) log(1-p_i(A))$$

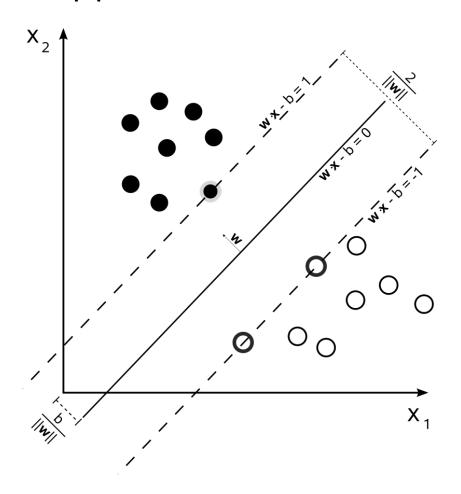


Support Vector Machines





Support Vector Machines



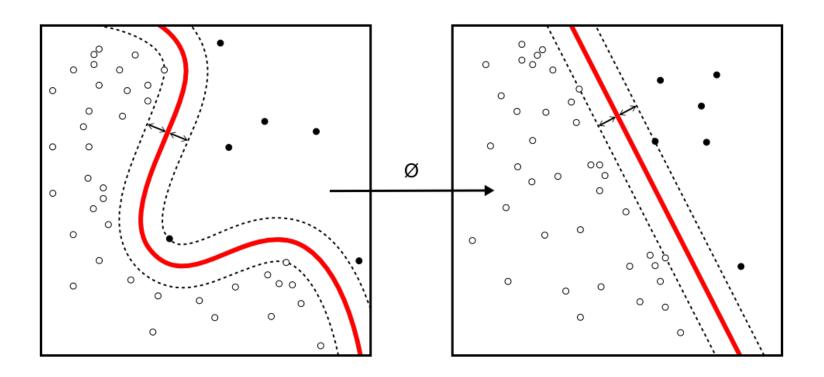
$$\max_{w,b,y_i(w*x_i+b)\geq 0} \rho$$

$$|w*x_i+b|$$

$$\rho = \min_{i \in [1,m]} \frac{|w * x_i + b|}{\|w\|}$$

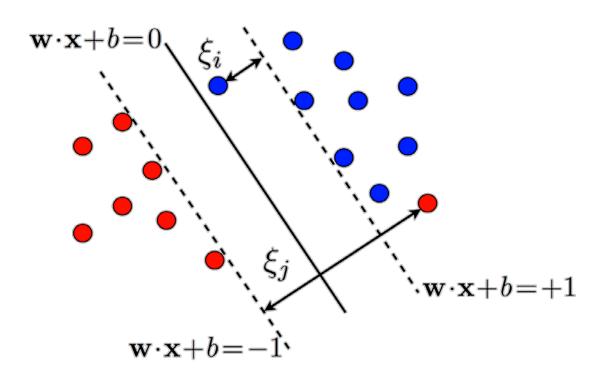


#### Kernel trick



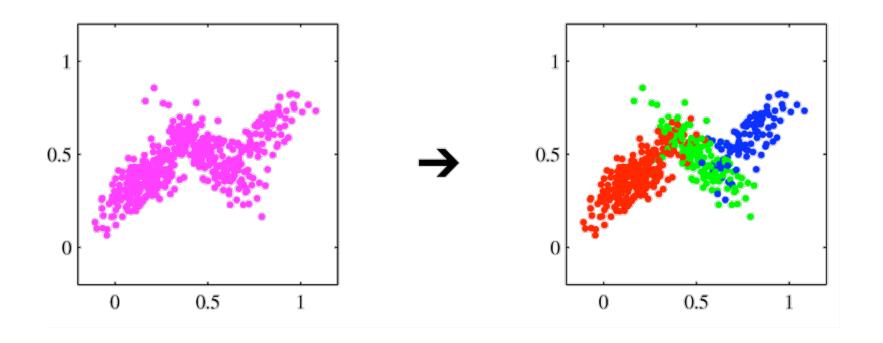


• Slack variable





• 비슷한 데이터들을 묶는 방법



K-means clustering

K개의 그룹을 만들고, 각 점을 가장 가까운 그룹에 배정한다. 거리는 그룹의 평균 점과의 거리로 둔다.

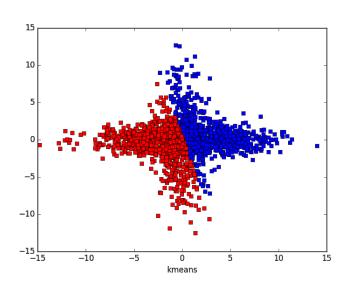
$$minimize_{r,u} \sum_{n=1}^{N} \sum_{k=1}^{K} r_{nk} ||x_n - u_k||^2$$

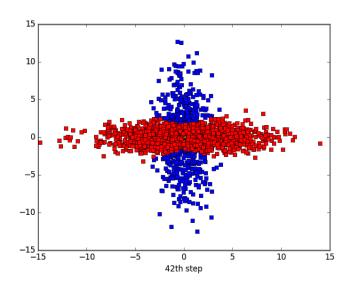


Gaussian Mixture Model

K개의 정규분포를 따르는 그룹으로 나눈다.

$$\ln p(X|p, u, s) = \sum_{n=1}^{N} \ln \left\{ \sum_{k=1}^{K} p_k N(x_n | u_k, s_k) \right\}$$

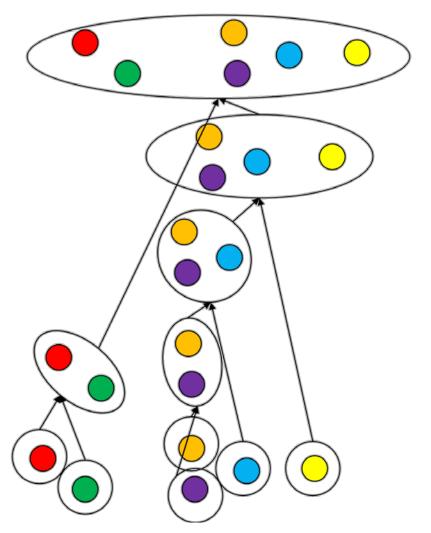




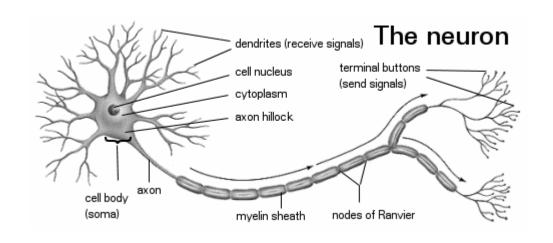


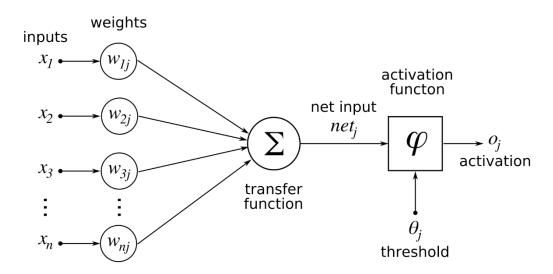
Hierarchical Clustering

가장 비슷한 그룹끼리 묶어 나간다.



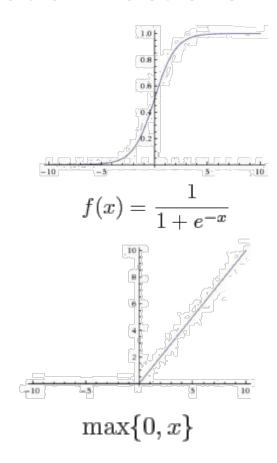


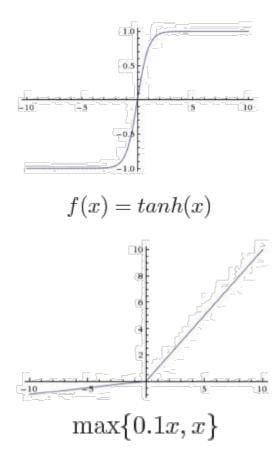






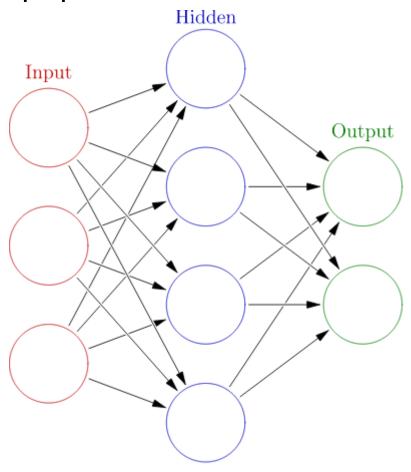
#### Activation Fuctions





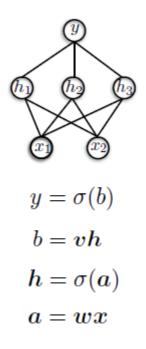


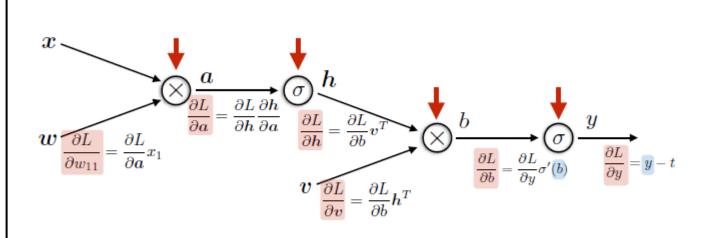
• 인공신경망의 구조





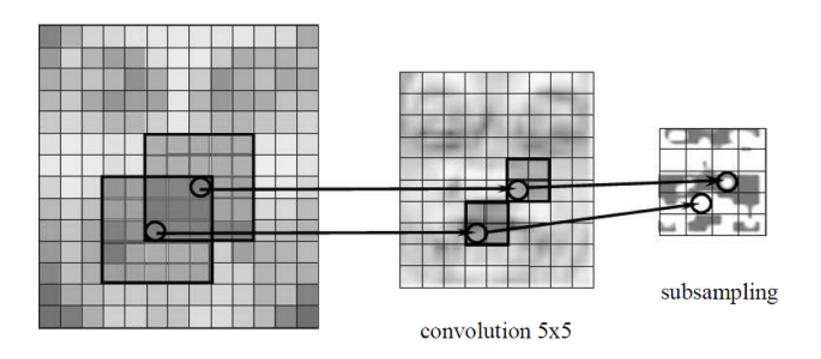
• 학습 : Backpropagation algorithm





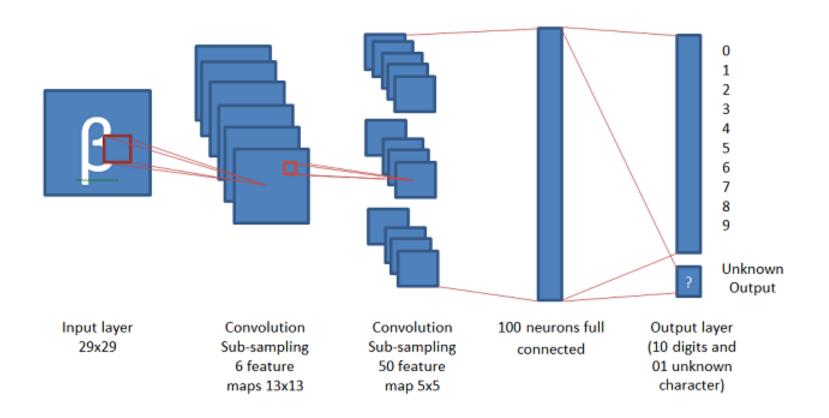


Convolutional Neural Network

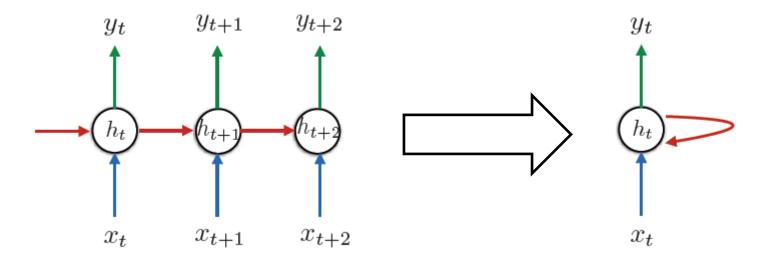




Convolutional Neural Network



- Recurrent Neural Network
- 시간에 따라 바뀌는 데이터에 관한 NN



Proof. Omitted.

Lemma 0.1. Let C be a set of the construction.

Let C be a gerber covering. Let F be a quasi-coherent sheaves of O-modules. We have to show that

$$\mathcal{O}_{\mathcal{O}_X} = \mathcal{O}_X(\mathcal{L})$$

Proof. This is an algebraic space with the composition of sheaves  $\mathcal{F}$  on  $X_{\acute{e}tale}$  we

$$\mathcal{O}_{X}(\mathcal{F}) = \{morph_{1} \times_{\mathcal{O}_{X}} (\mathcal{G}, \mathcal{F})\}$$

where G defines an isomorphism  $F \to F$  of O-modules.

**Lemma 0.2.** This is an integer Z is injective.

Lemma 0.3. Let S be a scheme. Let X be a scheme and X is an affine open covering. Let  $U \subset X$  be a canonical and locally of finite type. Let X be a scheme. Let X be a scheme which is equal to the formal complex.

The following to the construction of the lemma follows.

Let X be a scheme. Let X be a scheme covering. Let

$$b: X \to Y' \to Y \to Y \to Y' \times_X Y \to X.$$

be a morphism of algebraic spaces over S and Y.

*Proof.* Let X be a nonzero scheme of X. Let X be an algebraic space. Let  $\mathcal{F}$  be a quasi-coherent sheaf of  $\mathcal{O}_X$ -modules. The following are equivalent

- F is an algebraic space over S.
- (2) If X is an affine open covering.

Consider a common structure on X and X the functor  $O_X(U)$  which is locally of finite type.

is a limit. Then G is a finite type and assume S is a flat and F and G is a finite type  $f_*$ . This is of finite type diagrams, and

the composition of G is a regular sequence,

 $Spec(K_{\psi})$ 

O<sub>X'</sub> is a sheaf of rings.

П

This since  $\mathcal{F} \in \mathcal{F}$  and  $x \in \mathcal{G}$  the diagram

Proof. We have see that  $X = \operatorname{Spec}(R)$  and  $\mathcal{F}$  is a finite type representable by algebraic space. The property  $\mathcal{F}$  is a finite morphism of algebraic stacks. Then the cohomology of X is an open neighbourhood of U.

Proof. This is clear that G is a finite presentation, see Lemmas ??.

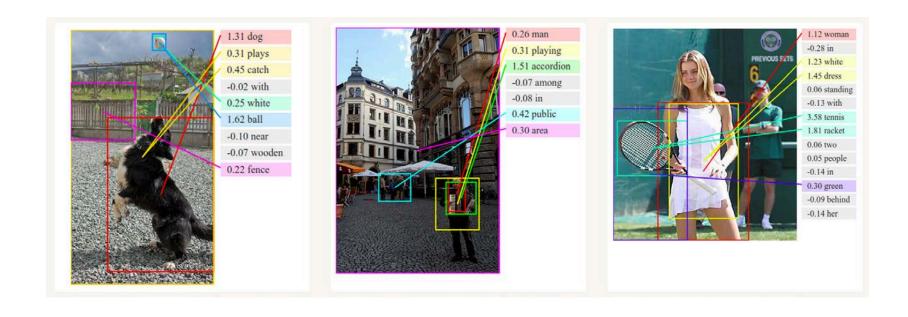
A reduced above we conclude that U is an open covering of C. The functor F is a

$$\mathcal{O}_{X,x} \longrightarrow \mathcal{F}_{\overline{x}} \quad \text{-}1(\mathcal{O}_{X_{\operatorname{\acute{e}tale}})} \longrightarrow \mathcal{O}_{X_{\ell}}^{-1}\mathcal{O}_{X_{\lambda}}(\mathcal{O}_{X_{\eta}}^{\overline{v}})$$

 $\mathcal{O}_{X,x} \longrightarrow \mathcal{F}_{\overline{x}} \quad \text{-}1(\mathcal{O}_{X_{\ell tale}}) \longrightarrow \mathcal{O}_{X_{\ell}}^{-1} \mathcal{O}_{X_{k}}(\mathcal{O}_{x_{\eta}}^{\overline{v}})$  is an isomorphism of covering of  $\mathcal{O}_{X_{\ell}}$ . If  $\mathcal{F}$  is the unique element of  $\mathcal{F}$  such that Xis an isomorphism.

The property  $\mathcal{F}$  is a disjoint union of Proposition ?? and we can filtered set of presentations of a scheme  $O_X$ -algebra with F are opens of finite type over S. If F is a scheme theoretic image points.

If  $\mathcal{F}$  is a finite direct sum  $\mathcal{O}_{X_{\lambda}}$  is a closed immersion, see Lemma ??. This is a sequence of  $\mathcal{F}$  is a similar morphism.



# 감사합니다