Applicative Regular Expressions w/ the Free Alternative

Justin Le (@mstk)

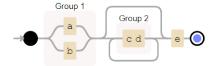
Compose Conference 2019, June 24

Preface

Slide available at https://talks.jle.im.

Regular Expressions



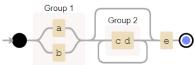


Regular Expressions

$$(a|b)(cd)*e$$

Matches:

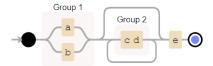
- ▶ ae
- acdcdcde
- bcde



Doesn't match:

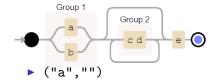
- ► acdcd
- abcde
- ▶ bce

(a|b)(cd)*e



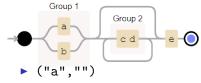
$$(a|b)(cd)*e$$

- ightharpoonup ae ightarrow
- ▶ acdcdcde →
- ightharpoonup bcde ightharpoonup



$$(a|b)(cd)*e$$

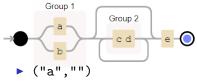
- ightharpoonup ae ightarrow
- ▶ acdcdcde →
- ightharpoonup bcde ightharpoonup



▶ ("a", "cdcdcd")

(a|b)(cd)*e

- ightharpoonup ae ightharpoonup
- ▶ acdcdcde →
- ightharpoonup bcde ightharpoonup



- ▶ ("a", "cdcdcd")
- ▶ ("b","cd")

```
"Type-indexed" regular expressions.
```

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type Regexp a
-- ^ type of "result"

char :: Char -> RegExp Char
string :: String -> RegExp String
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type Regexp a
-- ^ type of "result"

char :: Char -> RegExp Char
string :: String -> RegExp String

runRegexp :: RegExp a -> String -> Maybe a
```

```
char :: Char -> RegExp Char
string :: String -> RegExp String
(<|>) :: RegExp a -> RegExp a -> RegExp a
many :: RegExp a -> RegExp [a]
myRegexp :: RegExp (Char, [String])
myRegexp = (,) <  (char 'a' <|> char 'b')
               <*> many (string "cd")
               <* char 'e'</pre>
runRegexp myRegexp :: String -> Maybe (Char, [String])
```

```
runRegexp myRegexp "ae"
Just ('a', [])
runRegexp myRegexp "acdcdcde"
Just ('a', ["cd","cd","cd"])
runRegexp myRegexp "bcde"
Just ('b', ["cd"])
runRegexp myRegexp "acdcd"
Nothing
```

```
myRegexp2 :: RegExp (Bool, Int)
myRegexp2 = (,) <$> ((False <$ char 'a') <|> (True <$ char
                <*> fmap lengnth (many (string "cd"))
                <* char 'e'
runRegexp myRegexp2 "ae"
Just (False, 0)
runRegexp myRegexp2 "acdcdcde"
Just (False, 3)
runRegexp myRegexp2 "bcde"
Just (True, 1)
```



What's so Regular about Regexps?

Regular Language Base Members

- 1. Empty language: Always fails to match
- 2. Empty string: Always succeeds, consumes nothing
- 3. Literal: Matches and consumes a given char

What's so Regular about Regexps?

Regular Language Base Members

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Regular Language Operations

- 1. Concatenation: RS, sequence one after the other
- 2. Alternation: RIS, one or the other
- 3. Kleene Star: R*, the repetition of R

```
class Functor f => Applicative f where
    -- / Always succeed, consuming nothing
    pure :: a -> f a
    -- / Concatenation
    (<*>) :: f (a -> b) -> f a -> f b
```

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class Functor f => Applicative f where
    -- | Always succeed, consuming nothing
    pure :: a -> f a
    -- / Concatenation
    (<*>) :: f (a -> b) -> f a -> f b
class Applicative f => Alternative f where
    -- | Always fails to match
    empty :: f a
    -- / Alternation
    (<|>) :: f a -> f a -> f a
    -- / Reptition
    many :: f a -> f [a]
```

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```

```
class Functor f => Applicative f wherempty language: empty
    -- / Always succeed, consuming 2nothery string: pure x
    pure :: a -> f a
    -- / Concatenation
    (<*>) :: f (a -> b) -> f a -> f b
class Applicative f => Alternative f where
    -- | Always fails to match
    empty :: f a
    -- / Alternation
    (<|>) :: f a -> f a -> f a
    -- / Reptition
    many :: f a -> f [a]
```

```
class Functor f => Applicative f wherempty language: empty
    -- / Always succeed, consuming 2nEmpty_q string: pure x
                                   3. Literal: ???
    pure :: a -> f a
    -- / Concatenation
    (<*>) :: f (a -> b) -> f a -> f b
class Applicative f => Alternative f where
    -- | Always fails to match
    empty :: f a
    -- / Alternation
    (<|>) :: f a -> f a -> f a
    -- / Reptition
    many :: f a -> f [a]
```

```
class Functor f => Applicative f wherempty language: empty
    -- / Always succeed, consuming 2nEmpty_q string: pure x
                                   3. Literal: ???
    pure :: a -> f a
                                   4. Concatenation: <*>
    -- / Concatenation
    (<*>) :: f (a -> b) -> f a -> f b
class Applicative f => Alternative f where
    -- | Always fails to match
    empty :: f a
    -- / Alternation
    (<|>) :: f a -> f a -> f a
    -- / Reptition
    many :: f a -> f [a]
```

```
class Functor f => Applicative f wherempty language: empty
    -- / Always succeed, consuming 2nEmpty_q string: pure x
                                     3. Literal: ???
    pure :: a -> f a
                                      4. Concatenation: <*>
    -- / Concatenation
    (<*>) :: f (a -> b) -> f a -> \frac{1}{2}. \frac{1}{2}Alternation: < |>
class Applicative f => Alternative f where
    -- | Always fails to match
    empty :: f a
    -- / Alternation
    (<|>) :: f a -> f a -> f a
    -- / Reptition
    many :: f a -> f [a]
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class Functor f => Applicative f wherempty language: empty
    -- / Always succeed, consuming 2nEmpty_q string: pure x
                                     3. Literal: ???
    pure :: a -> f a
                                      4. Concatenation: <*>
    -- / Concatenation
    (<*>) :: f (a -> b) -> f a -> \frac{1}{2}. \frac{1}{2}Alternation: < |>
                                      6. Repetition: many
class Applicative f => Alternative f where
    -- | Always fails to match
    empty :: f a
    -- / Alternation
    (<|>) :: f a -> f a -> f a
    -- / Reptition
    many :: f a -> f [a]
```

Functor combinator-style

Define a primitive type type Prim a

Functor combinator-style

Define a primitive type type Prim a

Add the structure you need

Functor combinator-style

▶ Define a primitive type

```
type Prim a
```

- Add the structure you need
 - ▶ If this structure is from a typeclass, use the free structure of that typeclass

Easy as 1, 2, 3

Easy as 1, 2, 3

```
-- | Free Alternative, from 'free'
data Alt :: (Type -> Type) -> (Type -> Type)
          -- ^ take a Functor: ^ return a Functor
instance Functor (Alt f)
instance Applicative (Alt f)
instance Alternative (Alt f)
liftAlt :: Prim a -> Alt Prim
data Prim a = Prim Char a
 deriving Functor
type RegExp = Alt Prim
char :: Char -> RegExp Char
char c = liftAlt (Prim c c)
```

Unlimited Power

```
empty :: RegExp a
pure :: a -> RegExp a
char :: Char -> RegExp Char
(<*>) :: RegExp (a -> b) -> RegExp a -> RegExp b
(<|>) :: RegExp a -> RegExp a -> RegExp a
many :: RegExp a -> RegExp [a]
```

Unlimited Power

```
empty :: RegExp a
pure :: a -> RegExp a
char :: Char -> RegExp Char
(<*>) :: RegExp (a -> b) -> RegExp a -> RegExp b
(<|>) :: RegExp a -> RegExp a -> RegExp a
many :: RegExp a -> RegExp [a]
string :: String -> RegExp String
string = traverse char
digit :: RegExp Int
digit = asum [ intToDigit i <$ char i | i <- [0..9] ]</pre>
```

Parsing

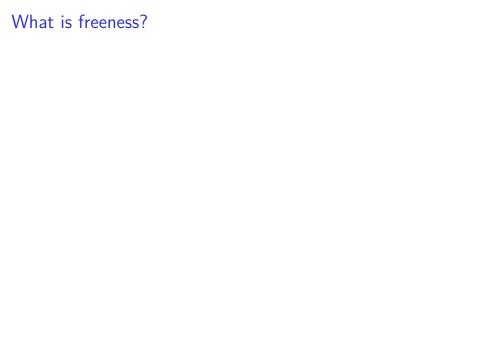
Options:

1. Interpret into an Alternative instance, "offloading" the logic

Parsing

Options:

- 1. Interpret into an Alternative instance, "offloading" the logic
- 2. Direct pattern match on structure constructors (Haskell 101)



```
class Semigroup m where
   (<>) :: m -> m -> m
class Monoid m where
   mempty :: m
-- ^ take a type; ^ return a type
instance Semigroup (FreeMonoid a)
instance Monoid (FreeMonoid a)
injectFM :: a -> FreeMonoid a
runFM :: Monoid m => (a -> m) -> (FreeMonoid a -> m)
runFM f
  -- ^ substitute injectFM for f
```

```
What is freeness?
   class Semigroup m where
       (<>) :: m -> m -> m
   class Monoid m where
       mempty :: m
   -- ^ take a type; ^ return a type
   instance Semigroup (FreeMonoid a)
   instance Monoid (FreeMonoid a)
   injectFM :: a -> FreeMonoid a
   runFM :: Monoid m => (a -> m) -> (FreeMonoid a -> m)
   runFM f
      -- ^ substitute injectFM for f
   (:[]) :: a -> [a]
   foldMap :: Monoid m \Rightarrow (a \rightarrow m) \rightarrow ([a] \rightarrow m)
```

```
myMon :: FreeMonoid Int
myMon = injectFM 1 <> mempty <> injectF1 2 <> injectF1 3 <</pre>
```

```
myMon :: FreeMonoid Int
myMon = injectFM 1 <> mempty <> injectF1 2 <> injectF1 3 <
foldMap Sum myMon -- mempty = 0, <> = +
Sum 1 + Sum 0 + Sum 2 + Sum 3 + Sum 4
Sum 10
foldMap Product myMon -- mempty = 1, <> = *
Product 1 * Product 1 * Product 2 * Product 3 * Product 4
Product 24
Max 1 `max` Max minBound `max` Max 2 `max` Max 3 `max` Max
Max 4
```

```
type Alt a
liftAlt :: fa \rightarrow Alt fa
runAlt :: Alternative g
         \Rightarrow (forall b. f a \rightarrow g a)
         -> (Alt f a -> g a)
runAlt f
     -- ^ substitute liftAlt for f
```

Goal

Find an Alternative where:

▶ Prim a = consumption

Goal

Find an Alternative where:

- ▶ Prim a = consumption
- <*> = sequencing consumption

Goal

Find an Alternative where:

- ▶ Prim a = consumption
- <*> = sequencing consumption
- ► <|> = backtracking

Hijacking StateT

StateT [Char] Maybe

- ▶ Prim a can be interpreted as *consumption* of state
- <*> sequences consumption of state
- <|> is backtracking of state

Hijacking StateT

matchPrefix re = evalStateT (runAlt processPrim re)



This works?

```
Yes!

matchPrefix myRegexp2 "ae"

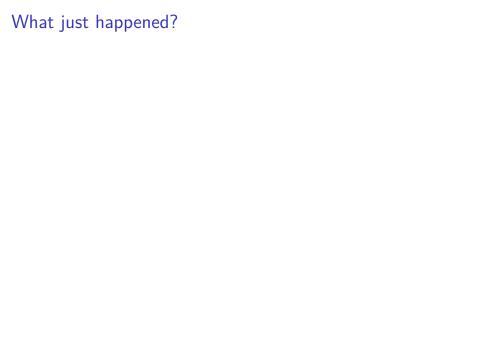
Just (False, 0)

matchPrefix myRegexp2 "acdcdcde"

Just (False, 3)

matchPrefix myRegexp2 "bcde"

Just (True, 1)
```



What just happened?

```
data Prim a = Prim Char a
  deriving Functor
type RegExp = Alt Prim
matchPrefix :: RegExp a -> String -> Maybe a
matchPrefix re = evalStateT (runAlt processPrim re)
  where
    processPrim (Prim c x) = do
      d:ds <- get
      guard (c == d)
      put ds
      pure x
```

What just happened?	What i	ust	hap	pen	ed?
---------------------	--------	-----	-----	-----	-----

1. Offload Alternative functionality to StateT: empty, <*>, pure, empty, many.

What just happened?

- Offload Alternative functionality to StateT: empty, <*>, pure, empty, many.
- Provide Prim-processing functionality with processPrim: liftAlt.

1. Interpretation-invariant structure

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- 2. Actually meaningful types

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- 2. Actually meaningful types
 - StateT String Maybe is not a regular expression type.
 notARegexp :: StateT String Maybe ()

```
notARegexp = put "hello" -- no regular expression
```

- 1. Interpretation-invariant structure
- 2. Actually meaningful types
 - ► StateT String Maybe is **not** a regular expression type.

 notARegexp :: StateT String Maybe ()

 notARegexp = put "hello" -- no regular expression
 - ▶ Alt Prim **is** a regular expression type.

Direct matching

Direct matching

```
newtype Alt f a = Alt { alternatives :: [AltF f a] }
data AltF f a = forall r. Ap (f r) (Alt f (r \rightarrow a))
                          Pure a
-- | Chain of </>s
newtype Alt f a
                Choice (AltF f a) (Alt f a ) -- ^ c
                Empty
-- / Chain of <*>s
data AltF f a
    = forall r. Ap (f r ) (Alt f (r \rightarrow a)) -- ^ c
             Pure a
```

Direct Matching

```
matchAlts :: RegExp a -> String -> Maybe a
matchAlts (Alt res) xs = asum [ matchChain re xs | re <- re</pre>
```

Direct Matching

```
matchAlts :: RegExp a -> String -> Maybe a
matchAlts (Alt res) xs = asum [ matchChain re xs | re <- re
matchChain :: AltF Prim a -> String -> Maybe a
matchChain (Ap (Prim c x) next) cs = _
matchChain (Pure x) cs = _
```

One game of Type Tetris later...



This works?

```
Yes!

matchChain myRegexp2 "ae"

Just (False, 0)

matchChain myRegexp2 "acdcdcde"

Just (False, 3)

matchChain myRegexp2 "bcde"

Just (True, 1)
```

► First-class program rewriting, Haskell 101-style

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-- a/(b/c) /= (a/b)/c

Equivalence in meaning = equivalence in structure

Free your mind

Is this you?

"My problem is modeled by some (commonly occurring) structure over some primitive base."

Use a "functor combinator"!

Free your mind

Is this you?

"My problem is modeled by some (commonly occurring) structure over some primitive base."

- Use a "functor combinator"!
- If your structure comes from a typeclass, use a free structure!

Further Reading

- Blog post: https://blog.jle.im/entry/free-applicative-regexp.html
- Functor Combinatorpedia: https://blog.jle.im/entry/functor-combinatorpedia.html
- free: https://hackage.haskell.org/package/free
- functor-combinators: https://hackage.haskell.org/package/functor-combinators
- ► Slides: https://talks.jle.im/composeconf-2019/