

Computational Geometry

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What is CG?

- algorithms w/ geometric input/output

→ I/O is naturally discrete

NOT: taking a cont. thing
+ discretizing to estimate
a solution.

IS: problems you can talk about
+ make progress on in A bar
on a napkin.

e.g., how can I arrange n points
on a plane (\mathbb{R}^2) to maximize

the # of pairs unit distance apart?



④ Forbidden config. 5 pairs (6 total pairs)



Cur best upper bound

$$O(n^{4/3})$$

Open: can we
do tighter UB?

- In CG, we evaluate our solutions

→ asymptotic worst-case RT
(sometimes average-case RT)

→ value efficiency

$\Theta(n)$ preferred over $\Theta(n \log n)$,
which is pref. over $\Theta(n^2)$
 $\Theta(n^3) \dots$

→ correctness: the algo. must do the task at hand.
if approximating, we give bounds on how off our sol'n is.

e.g: Quicksort

worst-case: $\Theta(n^2)$

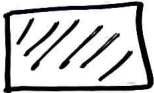
average: $\Theta(n \log n)$

- Dimension: we will focus on 2D / 3D
historically, much of CG is low-d (≤ 10)
curse of dimensionality: complexity often grows
exponentially w/ dimension


More about \mathcal{C}_1

- \mathbb{I}/\mathcal{O} is discrete + geometric, often (not always) flat (= defined by affine sets)

$\rightarrow \mathbb{R}^2$, rectangle in \mathbb{R}^2

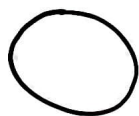


edges = line segments
points, lines, (hyper)planes



our non-flat objects

{ sphere / circle
balls / discs



$$\mathbb{S}^2 := \{x \in \mathbb{R}^3 \mid \|x\|_2 = 1\}$$

$$\mathbb{R}^3 = \mathbb{R} \times \mathbb{R} \times \mathbb{R}$$


$$x = (x_1, x_2, x_3)$$

$$\mathbb{S}^2 := \{x \in \mathbb{R}^3 \mid \sqrt{x_1^2 + x_2^2 + x_3^2} = 1\}$$

"the set of things unit dist from origin"

Not focused on

{ algebraic surfaces



- analysis

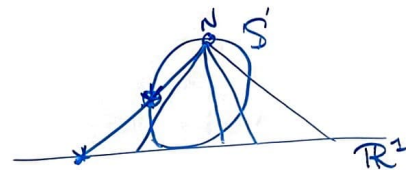
- RAM =
- a given
- memory
- to

- things
- add
- mul
- div
- sub
- √

- * avoid
-

affine sets)

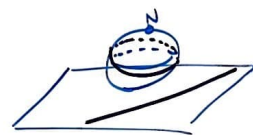
- analysis uses real-RAM MOC
- ↑ ↑
working w/ real #'s random access machine



- RAM $\Rightarrow \Theta(1)$ time to access (R/W)
a given location in memory

- memory stores real #'s

- to store a number in \mathbb{R}^2 : need 2 locations in memory
✓ still $\Theta(1)$



"

in \mathbb{R}^n : need n locations

- things I can do in const. time:

add 2 #'s

multiply

divide

subtract



* avoids dealing w/ floating point arithmetic
 \rightarrow CGAL simulates real-RAM on many ops!

Do 2 line segments intersect?



\rightarrow don't just solve the equation by setting 2 lines -

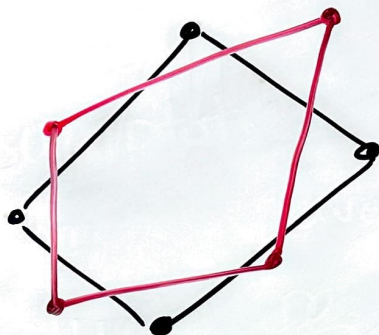
$$\{x \mid \|x\|_2 = 1\}$$

\mathbb{R}

(x_1, x_2, x_3)

$$\{x \mid \sqrt{x_1^2 + x_2^2 + x_3^2} = 1\}$$

things unit dist from origin



Q: are the lengths of these
2 the same or different?

TOP

1. Co



Com
↳

-c

2. In

TOPICS

1. Convex Hulls



CONVEX



NOT

whoops!

↳ def'n: line seg. between any 2 pts
Stays inside the shape.

- CH: given a set of pts P ,
the CH is the smallest convex
shape containing P .

2. Intersections

3. Triangulations + Subdivisions

4. Point Location

5. Linear Programming

6. Duality

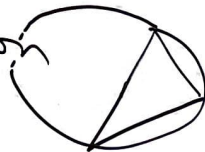
pt - line (in \mathbb{R}^d)
 (a, b) $ax + by = 1$

↳ Sometimes think of
this as $(a, b, 1)$

7. Voronoi Diagrams (post office query data structure)

↑ duals

8. Delaunay Triangulation



9. Line/Plane arrangements

