

# Dynamic Reconfiguration of Wireless Sensor Networks to Support Heterogeneous Applications

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# Motivation: Run Two Applications on the Same WSN

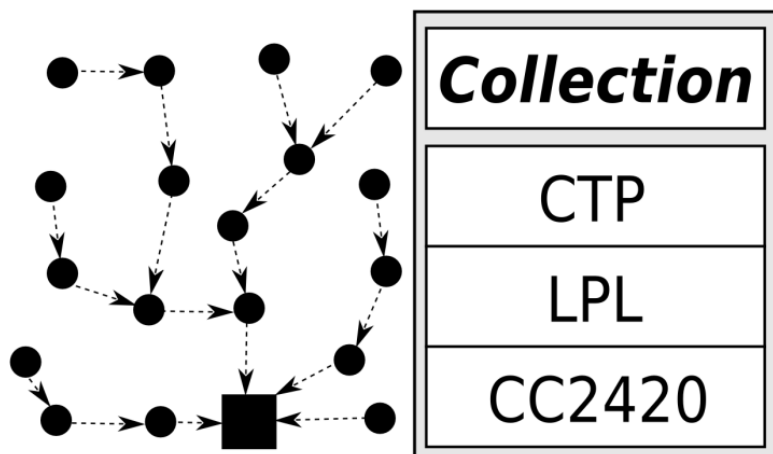
**Collection:** collect building's climate data

- reliable, low-rate
- many-to-one
- energy efficient

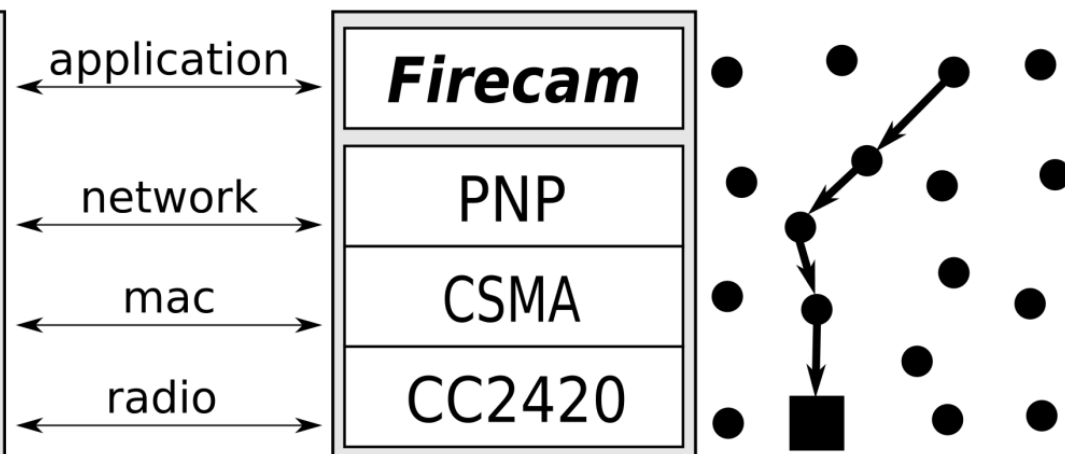
**Firecam:** sends images across a building

- low-latency
- point-to-point
- bulk-data streaming

## COLLECTING CLIMATE DATA



## IMAGE TRANSMISSION

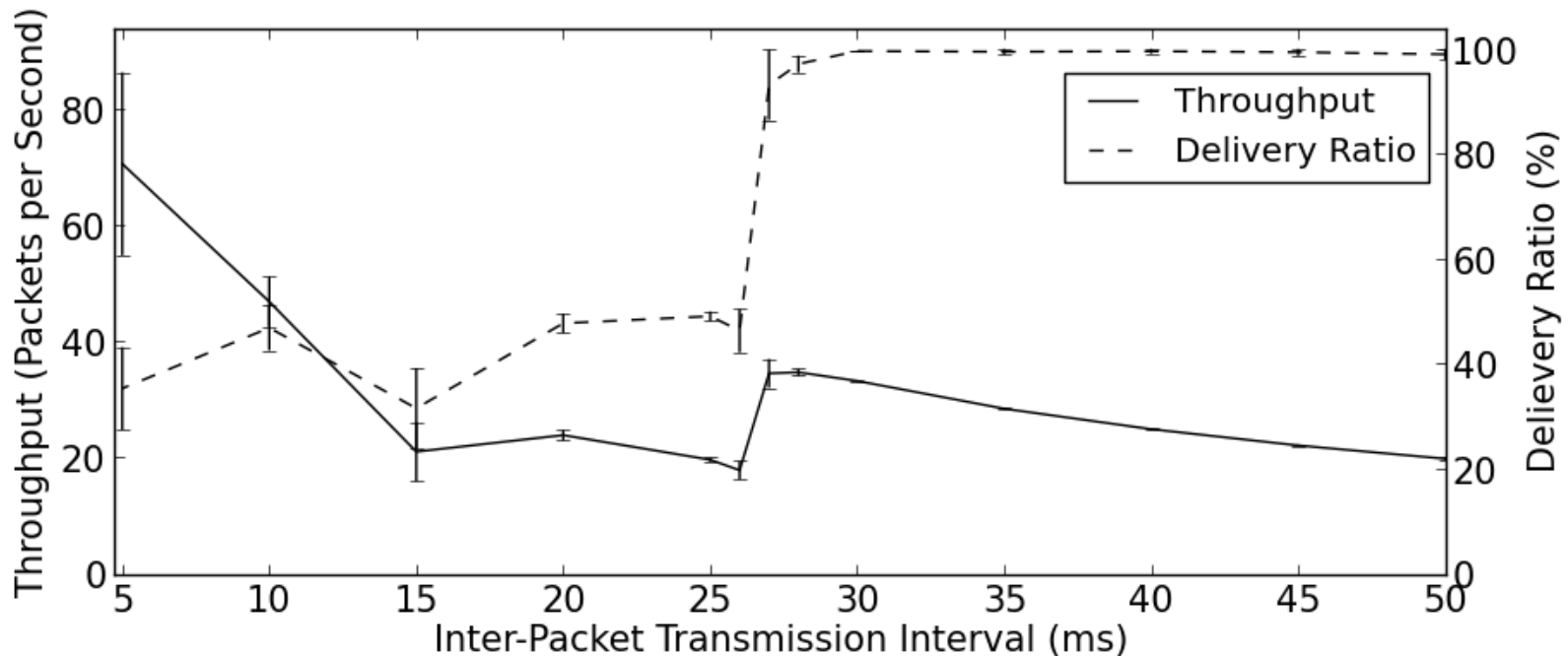


### Emergency Response Requirement:

During normal operation, run **Collection**. When an emergency event occurs (e.g. a smoke sensor goes off), switch to **Firecam**.

# Two Applications, Different Communication Requirements

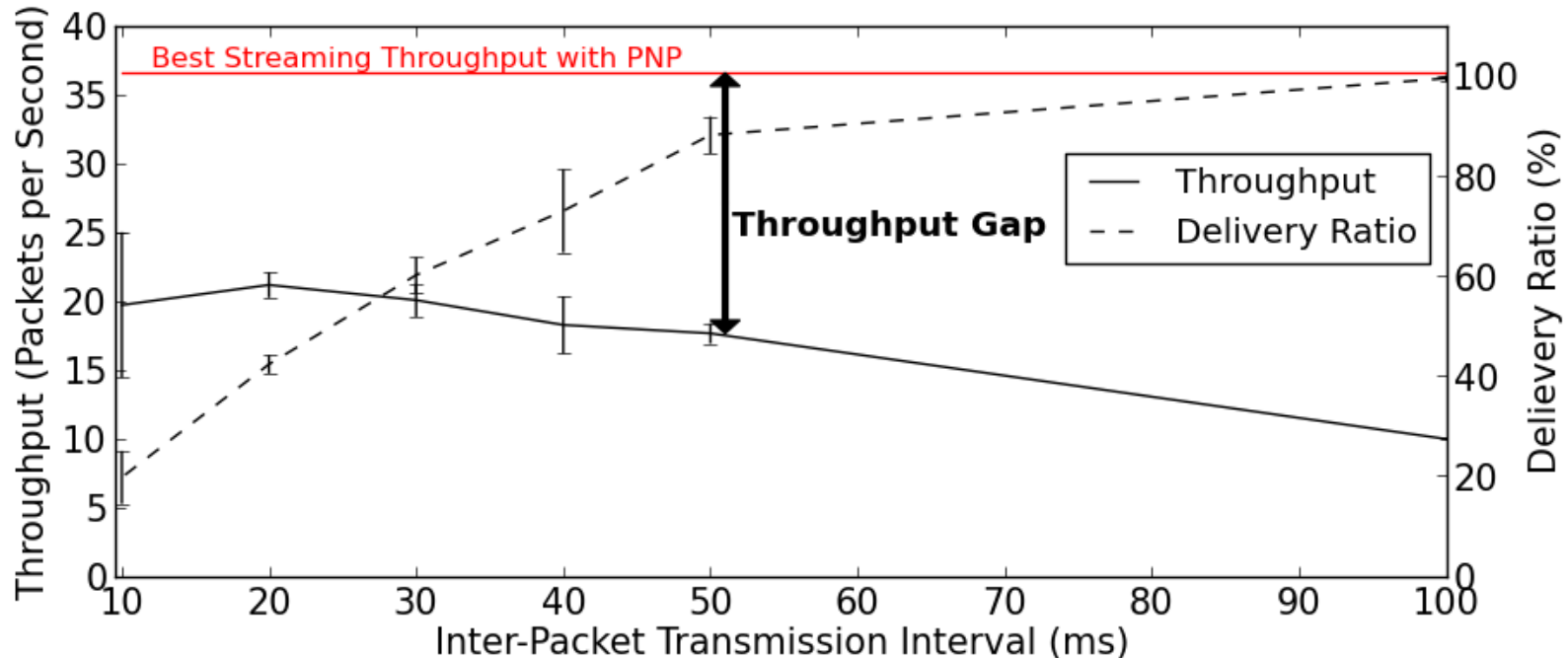
- setup: Indriya testbed, a three-floor building in NUS with 139 TelosB motes
- Transmitting 75 KBytes across the building, over 7-9 hop path
- 76800 bytes takes a 240x320 picture with each pixel encoded as a byte



- Almost 100% of packets delivered when transmission delay was 30ms or more
- For 28ms delay, PNP's delivery ratio was 97.4%
- PNP transmits 36.5 packets per second
- Sending a 75KByte picture takes ~ **21 seconds**

# Two Applications, Different Communication Requirements

- Transmitting 75 KBytes across the building, over 7-9 hop path



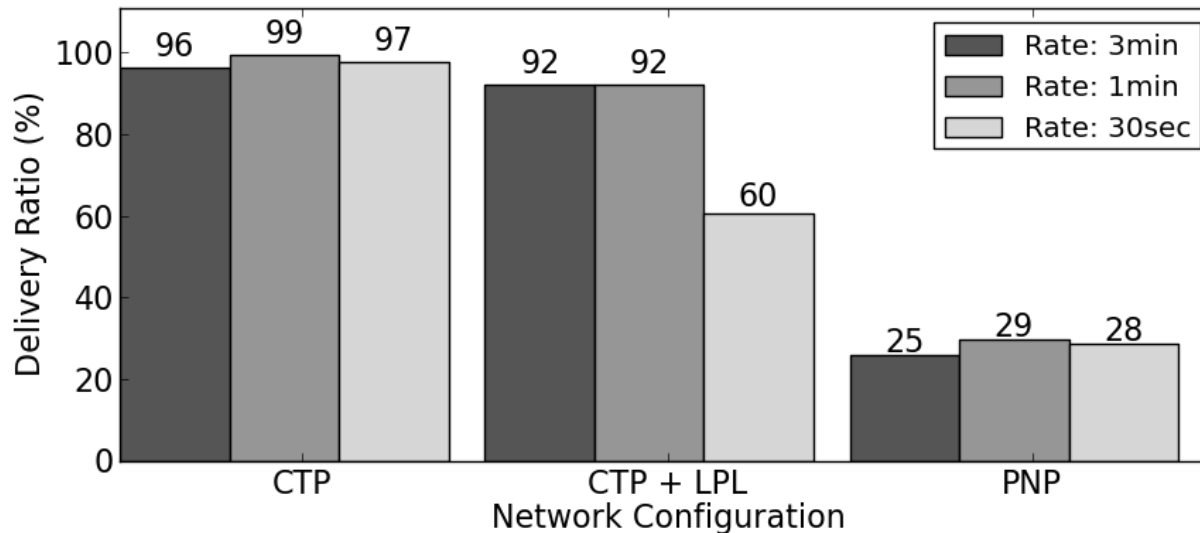
- For 50ms delay CTP delivery ratio is 88%
- This delivers 17.61 packets per second
- Sending 75KBytes picture takes ~ **44 seconds**

## Experimental Observation:

CTP does not support high throughput on a point-to-point path

# Two Applications Need Different Protocol Stacks

- Collecting data from the whole network at one sink node
- Each node periodically sends 4 bytes of data



- Low-Power Listening at 100ms
- CTP performs well with radio always ON or duty-cycled
- PNP drops more than 70% of packets

## Experimental Observation:

PNP cannot support the same many-to-one delivery ratio as CTP

## The Need for Dynamic Protocol Stack Reconfiguration

Each of the protocol stacks supports well the corresponding application, for which it has been optimized, while poorly supporting the other application

# Fennec Fox Framework

## 1. Programming Language – Swift Fox

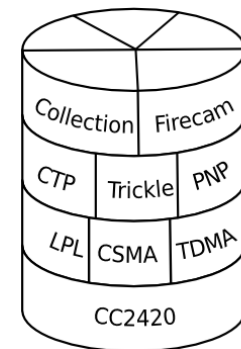
- defines applications and supporting communication infrastructure
- specifies when to context-switch applications

```
configuration Monitoring { collection(1024, NODE, 107)
                           ctp(107) lpl (100, 10), cc2420(1)}
configuration Emergency {firecam(28), pnp(),
                           csma(0,0),cc2420(0)}

event fire {smoke = ON }
event check {timer = 30 sec}
from Monitoring goto Emergency when fire
from Emergency goto Monitoring when check
start Monitoring
```

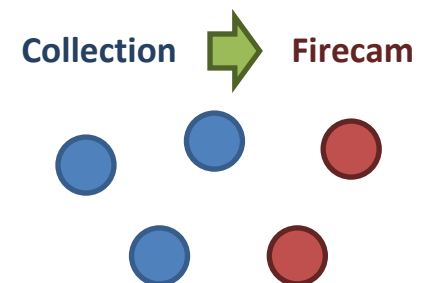
## 2. Protocol Stack

- layered system design of application and communication protocols
- starts and stops modules running across the layers



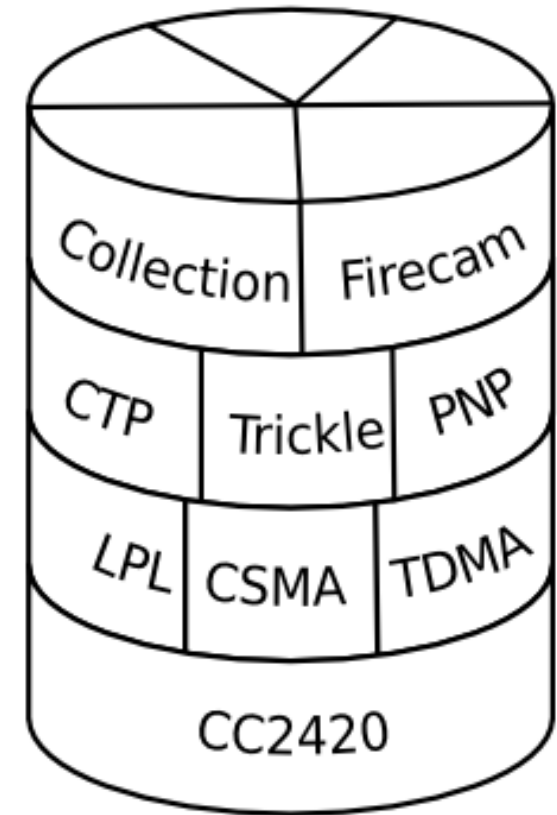
## 3. Network Reconfiguration

- executes application context switch together with communication protocols
- ensures that all nodes run the same modules across the layers of the protocol stack



# Fennec Fox Protocol Stack

- four layer protocol stack
- application, network, MAC, and radio layers
- each layer contains multiple modules implementing its functionality
- Fennec Fox schedules the execution of the modules across the stack
- builds upon existing research work
- runs on top of TinyOS, implemented in nesC
- modules must obey inter-layer API
- reuse and combine already existing protocols
- modular design simplifies programming and validation

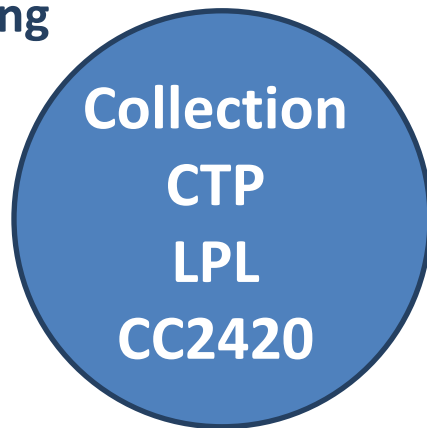


## Protocol Stack Configuration

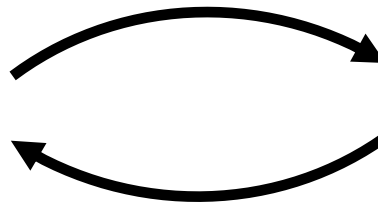
Set of four modules executing on the four-layer stack, one module for each layer. Each configuration has globally unique **id**

# Swift Fox Programming Language

Monitoring

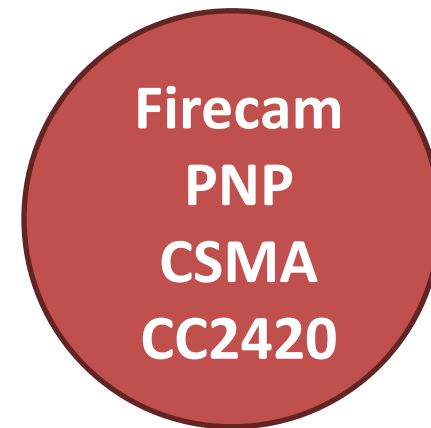


SMOKE = ON



TIMER = 30 SEC

Emergency



- decouple the network reconfiguration from the application implementation
- model network behavior as **Finite State Machine**
- at *design time* a programmer specifies the FSM of the network reconfiguration logic
- programs are compiled to a system that **at run-time automatically reconfigures the network**

```
configuration Monitoring { collection(1024, NODE, 107)
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configuration Emergency {firecam(28) pnp()
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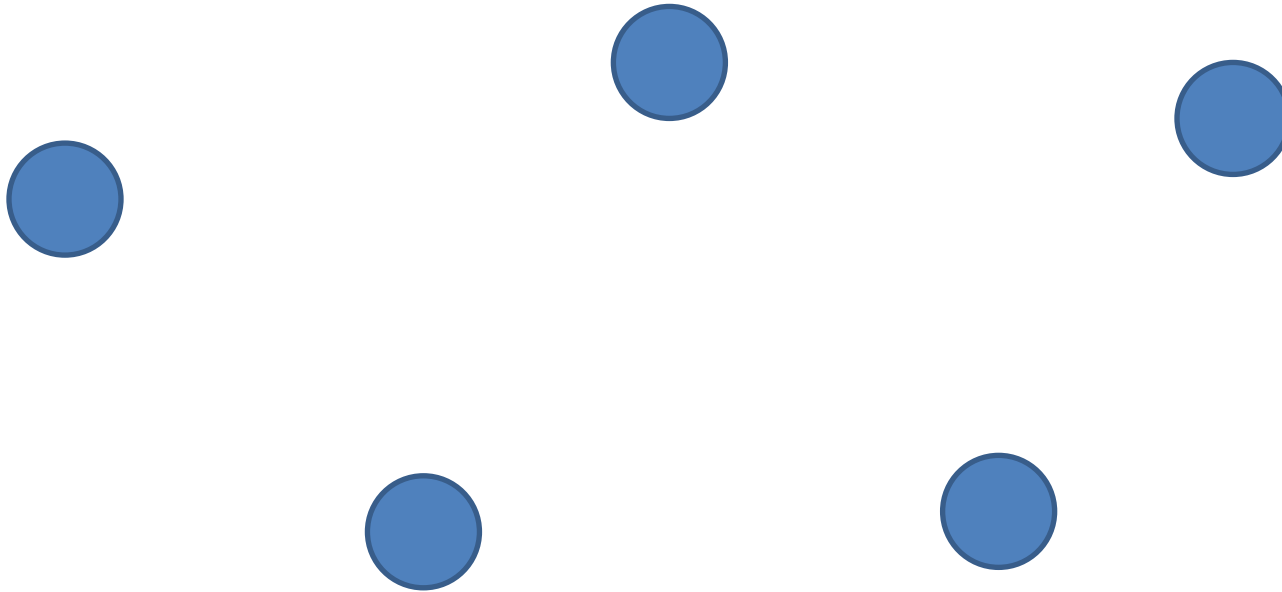
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# Network Reconfiguration

- network is reconfigured by disseminating **control messages**
- each message contains a **sequence number** and the **configuration id**
- control messages are disseminated by following modified Trickle protocol
- new sequence number triggers reconfiguration, but also checks **payload consistency**

## Collection

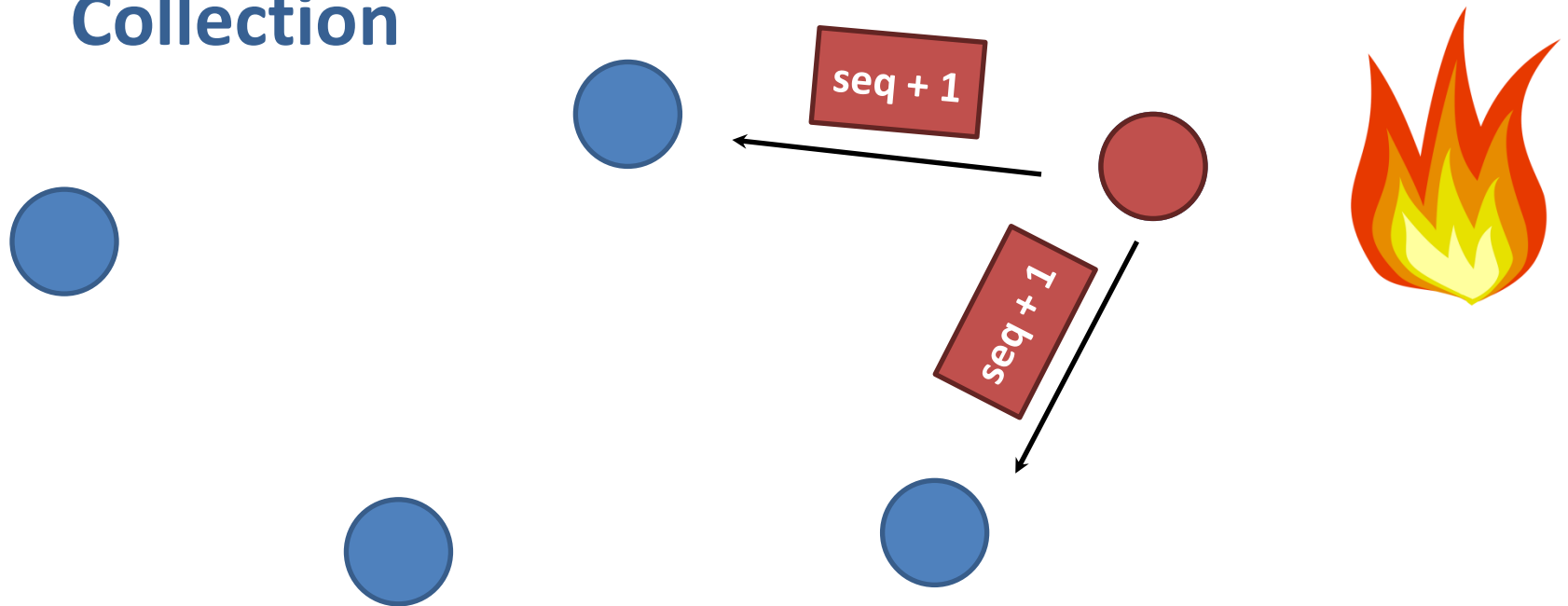


- network is **unsynchronized** when there exist control messages with the same sequence number but different configuration ids
- detects network state inconsistency and synchronizes all nodes

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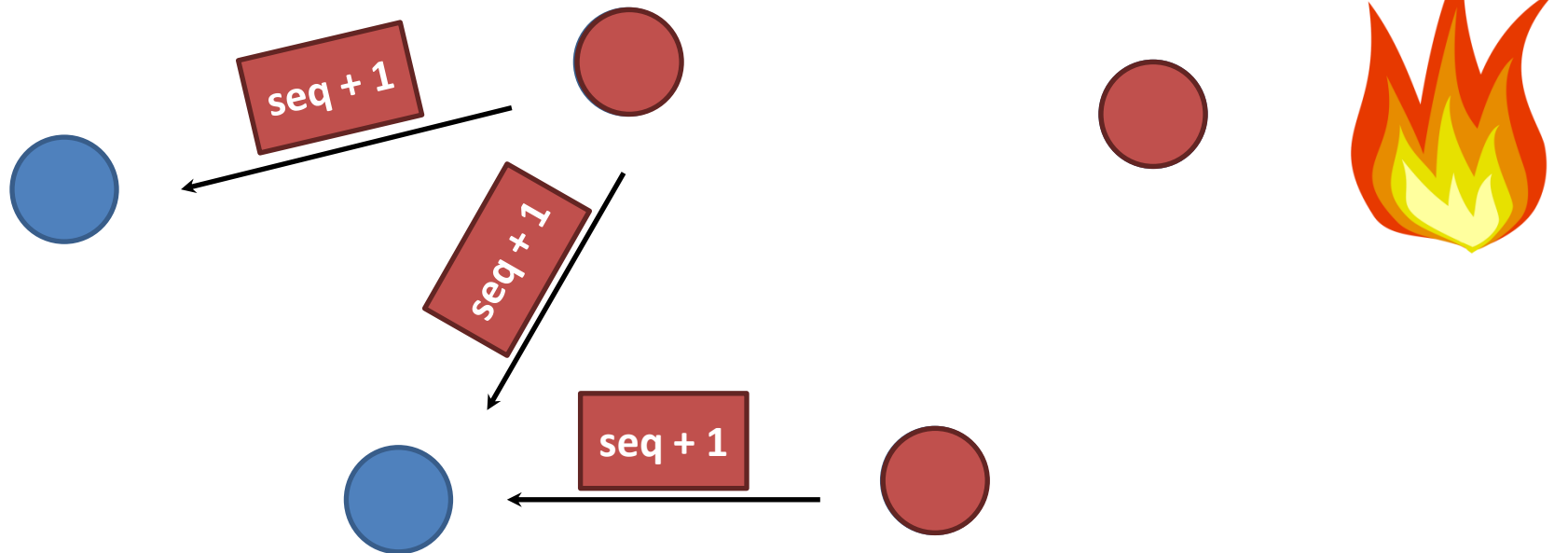


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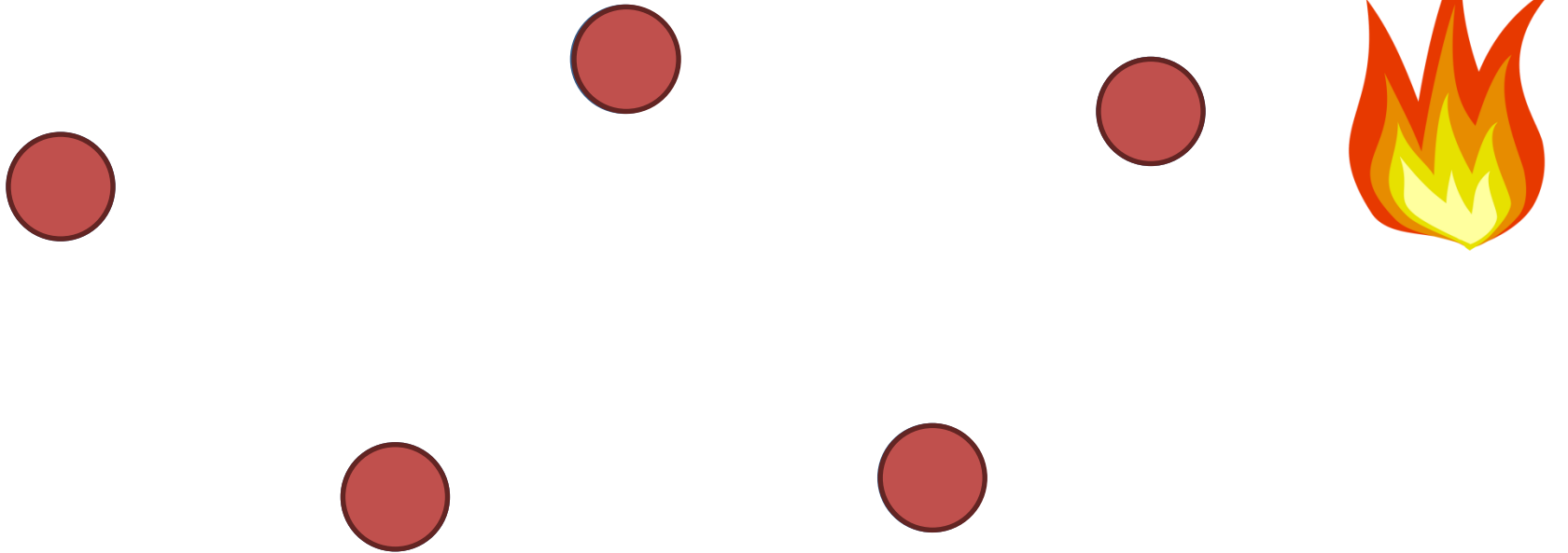


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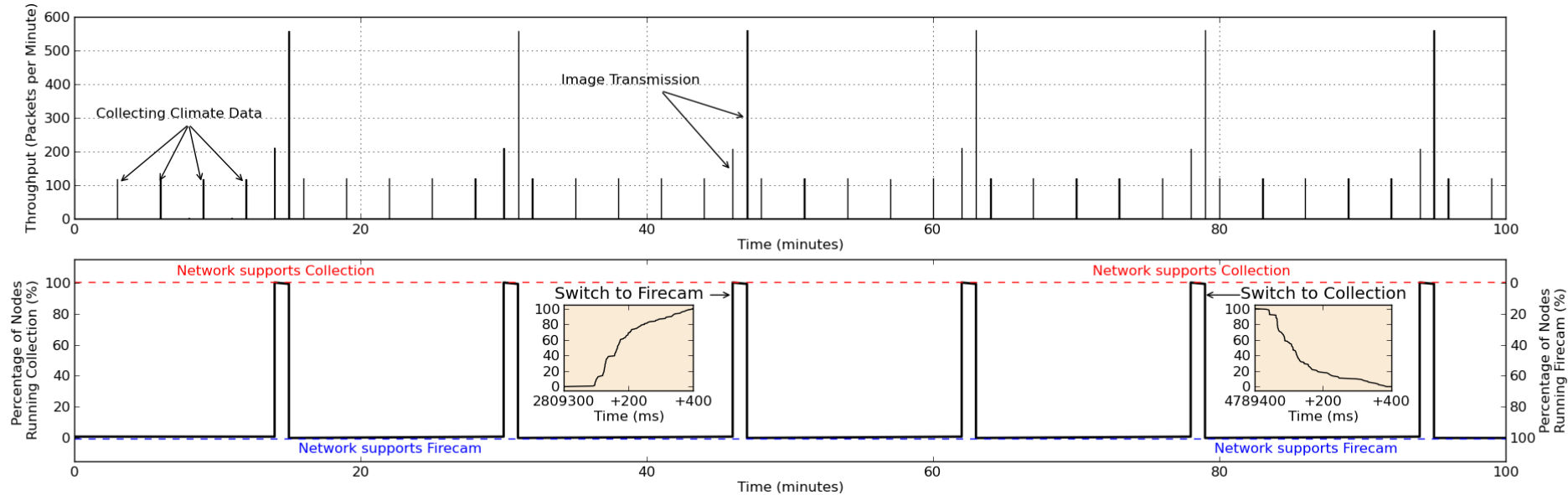
## Firecam



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# Evaluation: Context Switch Robustness

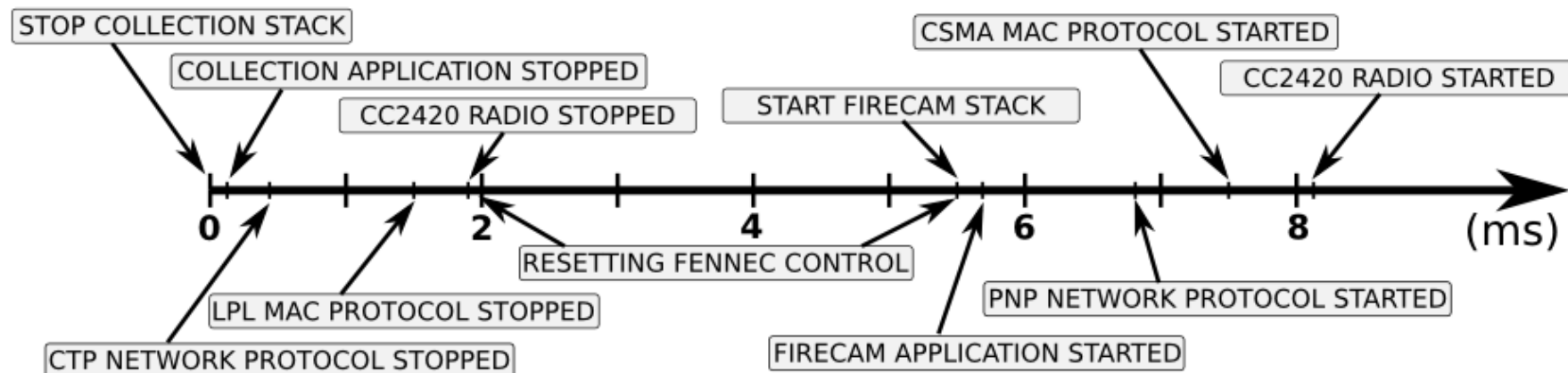
100 minute run of a network reconfiguring between the Collection and Firecam



- on average, 99.91% of network reconfigurations are successful
- 99.5% of nodes reconfigure successfully when radio is duty-cycled
- **network can reconfigure as fast as every 350ms**, with 99.68% nodes successfully following all reconfigurations

setup: Indriya testbed, a three-floor building at the NUS instrumented with 139 TelosB motes

# Evaluation: Reconfiguration Overhead



Protocol stack reconfiguration from Collection to Firecam

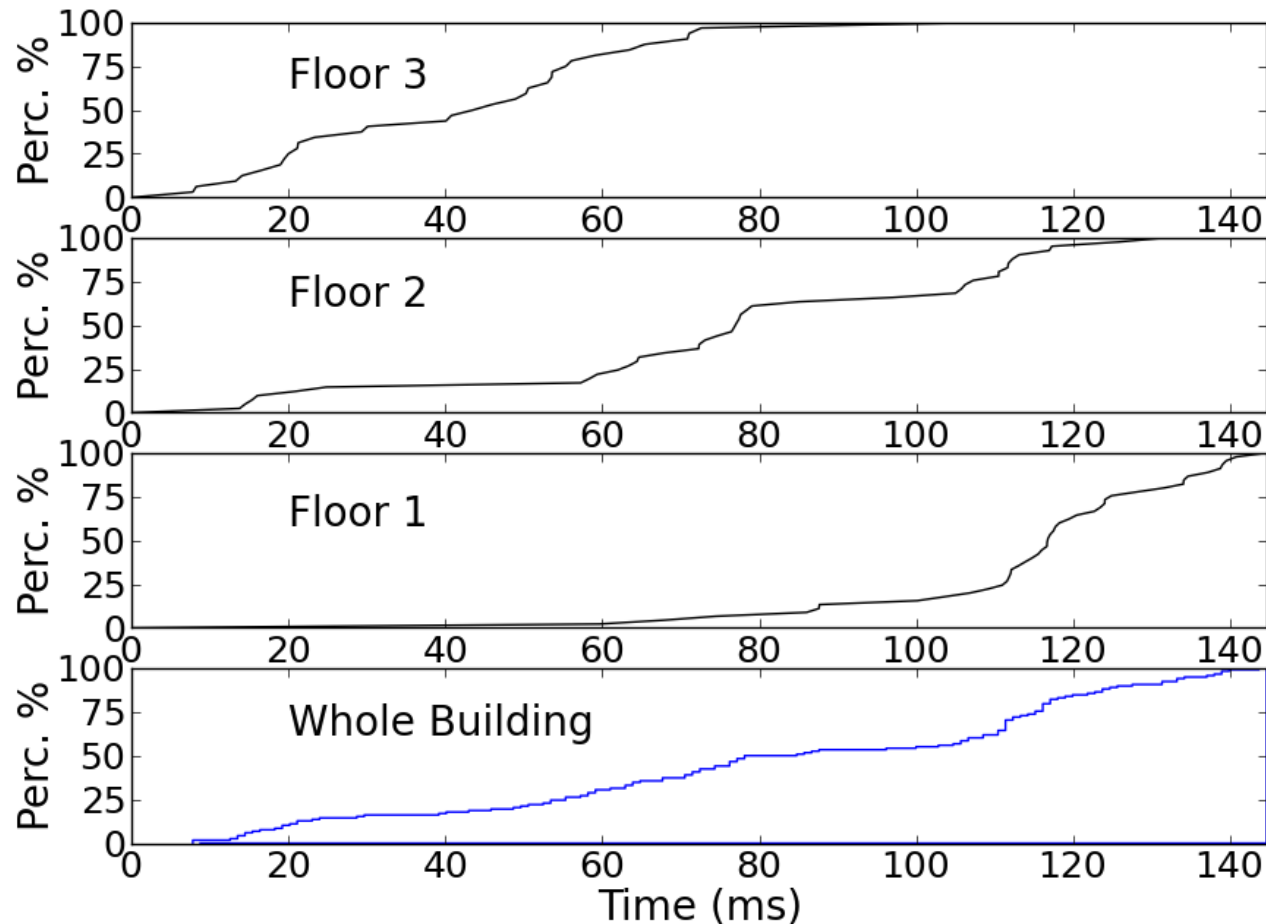
## Delay:

- to switch from Collection to Firecam takes 8.125ms
- to reset the radio itself takes almost 3ms
- the whole stack reconfiguration takes between 8 and 9ms.

## Memory:

- 4.7 KB of ROM and 0.2KB of RAM
- Collection and Firecam require 28.9 KB of ROM and 5.6 KB of RAM

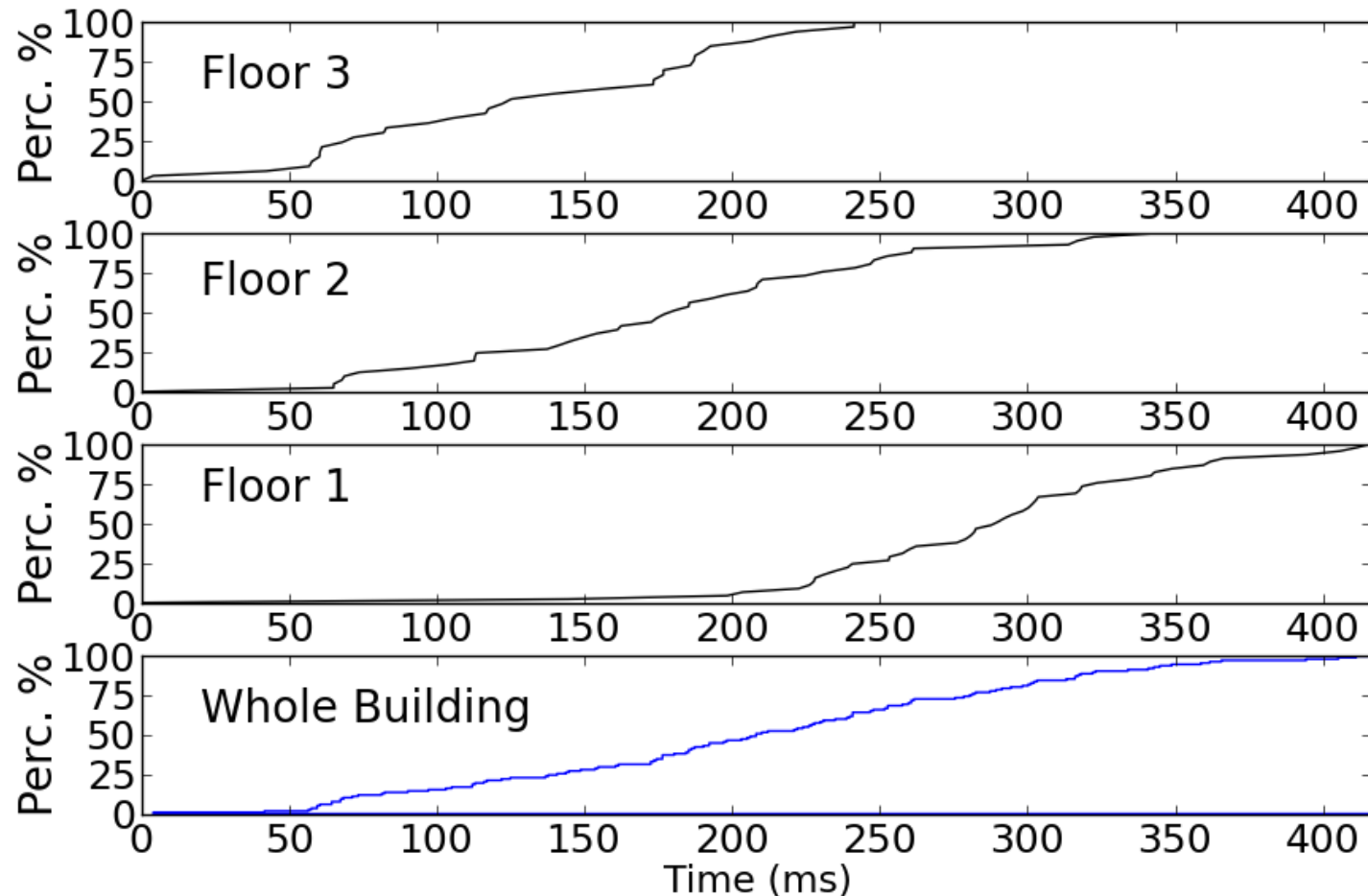
# Evaluation: Reconfiguration Across a Building



## Reconfiguration with radio always turned on

- reconfigurations occur in bursts
- control message broadcast quickly propagates on the same floor
- network reconfiguration is delayed on neighboring floors

# Evaluation: Reconfiguration Across a Building

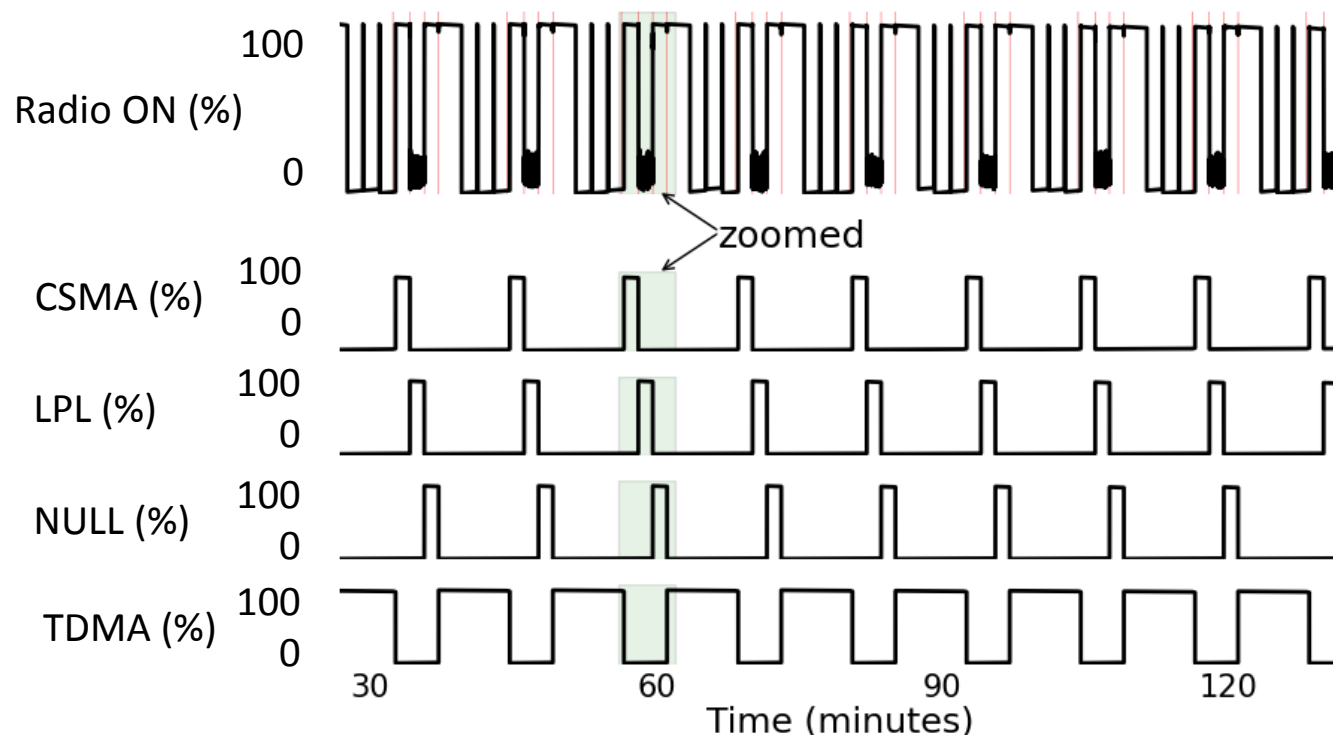


## Reconfiguration with radio duty-cycling

- network reconfigures progressively
- need more control message transmission attempts for smaller duty-cycle times
- 33-40% of control message transmissions attempts are suppressed

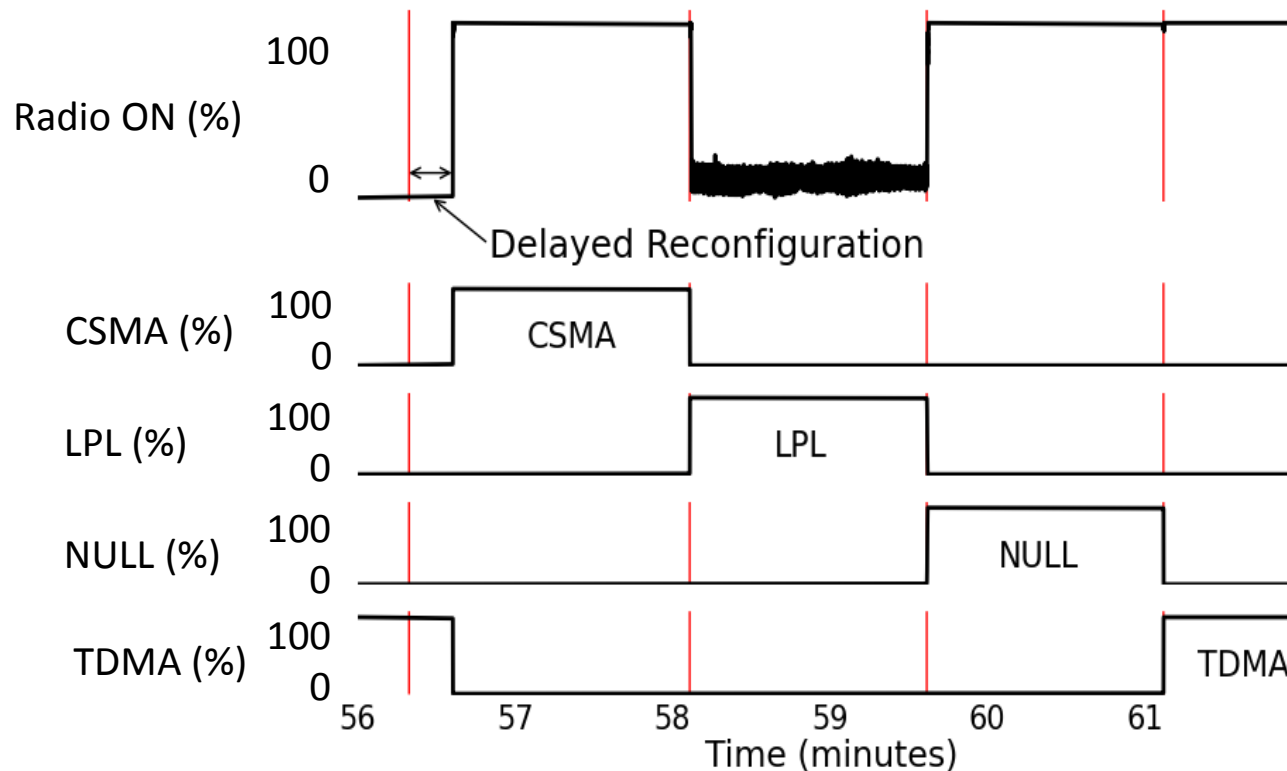


# Evaluation: Reconfiguration Among Various MACs



- reconfigure a WSN among the four MAC protocols:
  - CSMA
  - LPL (unsynchronized duty-cycling)
  - Null (simple MAC that passes packet between the Network and Radio layers)
  - TDMA (synchronized duty-cycling)
- without Fennec Fox these MAC protocols cannot co-exist in the same WSN

# Evaluation: Reconfiguration Among Various MACs



- Fennec Fox stack monitors the status of the radio to guess the state of the network
- when the radio is turned-off, the reconfiguration is delayed until radio is on
- to switch between two configurations with TDMA and with the same duty-cycle period, it requires to insert between these configurations an empty state with CSMA or LPL, to ensure that all motes successfully reconfigure

# Conclusion

## Conclusion:

- studied the problem of scheduling execution of heterogeneous applications with different communication requirements on a single WSN
- developed Fennec Fox, a runtime infrastructure build around a layered protocol stack that executes network-wide application context-switch
- showed the feasibility of Fennec Fox through experimental evaluation on a network deployed across the three-floor building

## Open Source Code:

- Fennec Fox: <https://github.com/mszczodrak/fennecfox>
- Swift Fox: <https://github.com/mszczodrak/swiftfox>

**Thank You,**

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