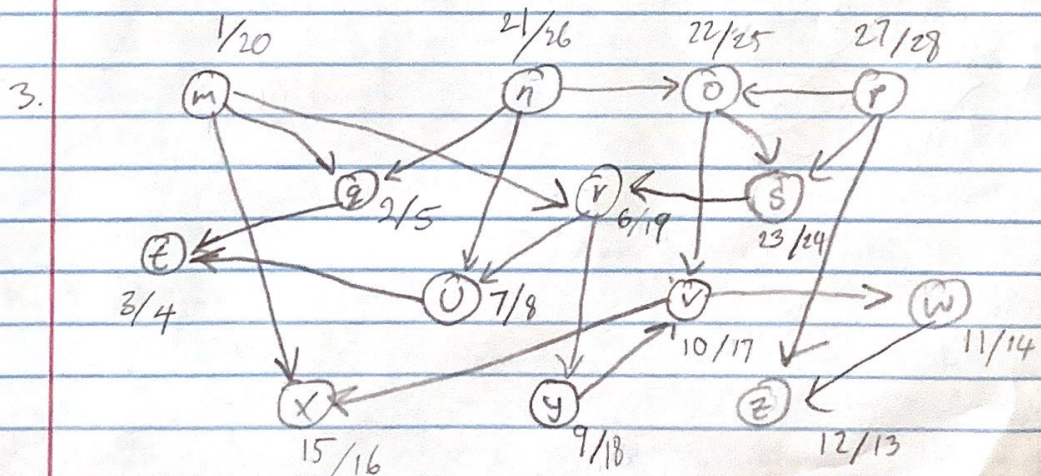
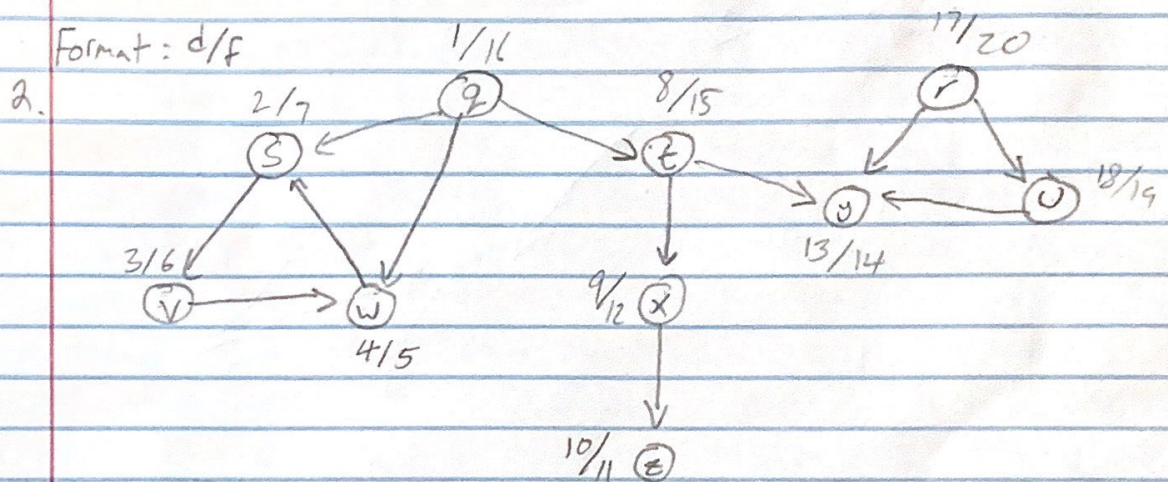
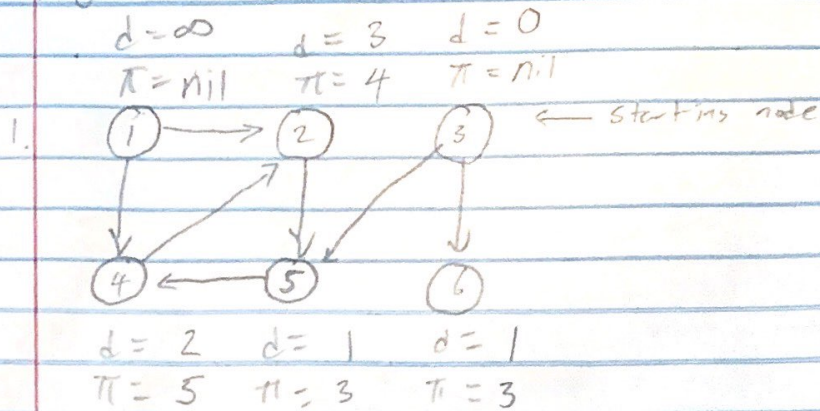


Algorithms HW 4

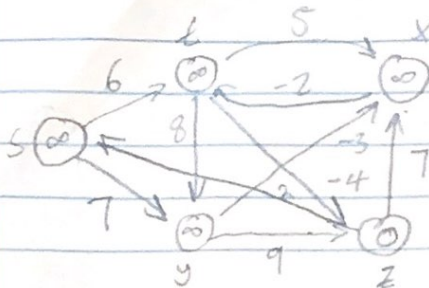
Tae M



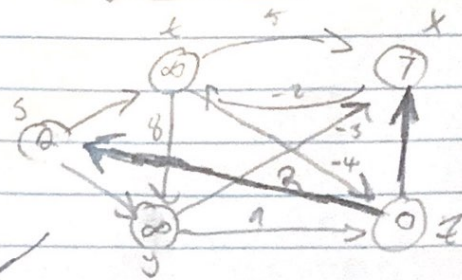
Topological Order: P, n, o, s, m, r, y, v, x, w, z, u, q, t

$|V| = 5 - 1 = 4$ iterations.

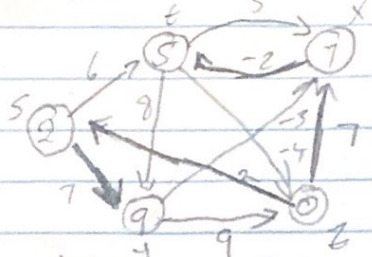
4.



Iteration 1



Iteration 2



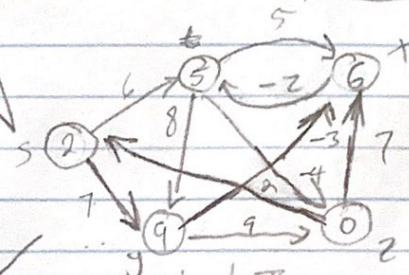
	d	π
s	2	z
t	∞	nil
x	7	z
y	∞	nil
z	0	nil



	d	π
s	2	z
t	5	x
x	7	z
y	9	s
z	0	nil



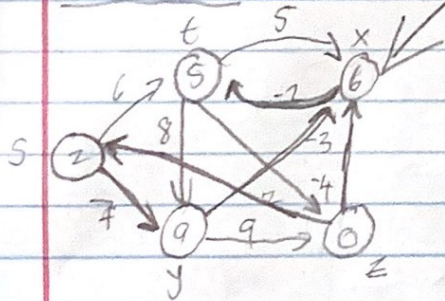
Iteration 3



	d	π
s	2	z
t	5	x
x	6	z
y	9	s
z	0	nil

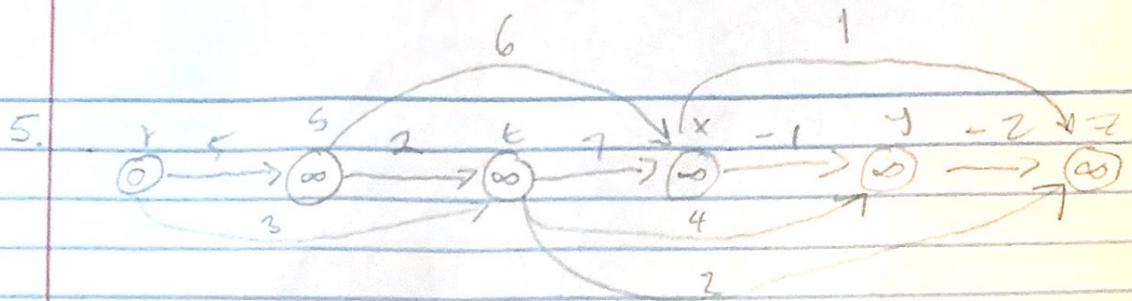


Iteration 4

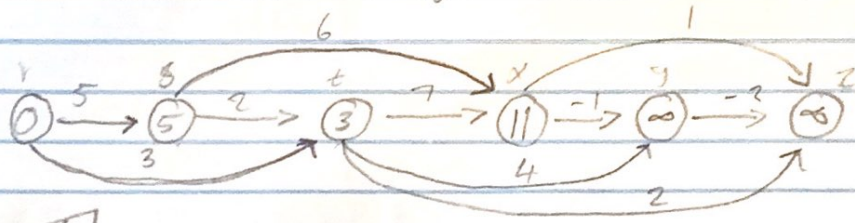
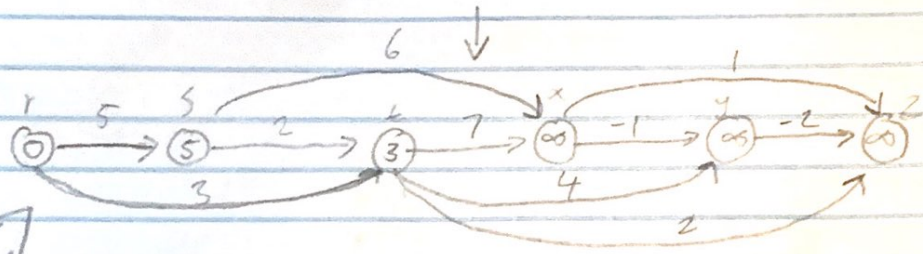


	d	π
s	2	z
t	4	x
x	6	y
y	9	s
z	0	nil

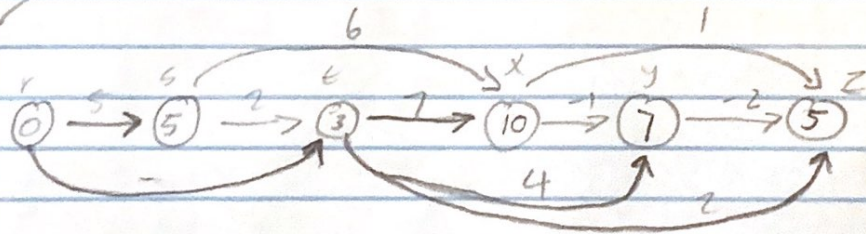




	d	π
r	0	nil
s	5	r
t	3	r
x	∞	nil
y	∞	nil
z	∞	nil

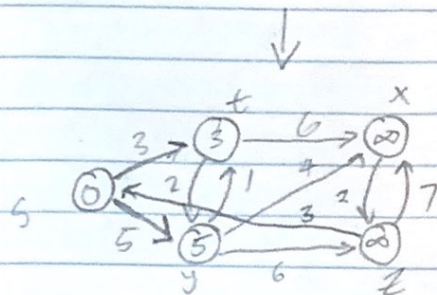
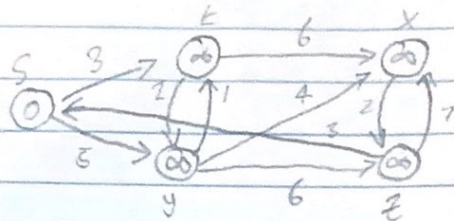


	d	π
r	0	nil
s	5	r
t	3	r
x	10	s
y	∞	nil
z	∞	nil

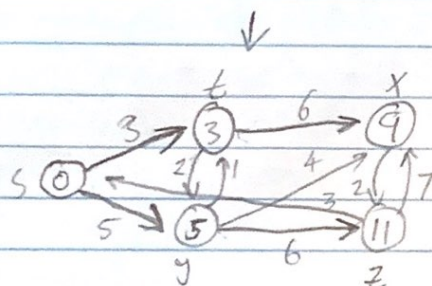


	d	π
r	0	nil
s	5	r
t	3	r
x	10	t
y	7	t
z	5	t

6 Vertex S as Source:

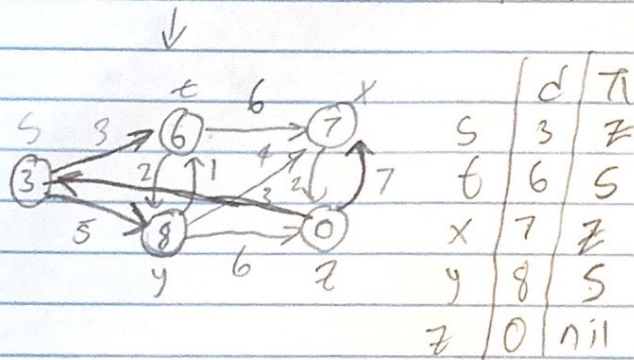
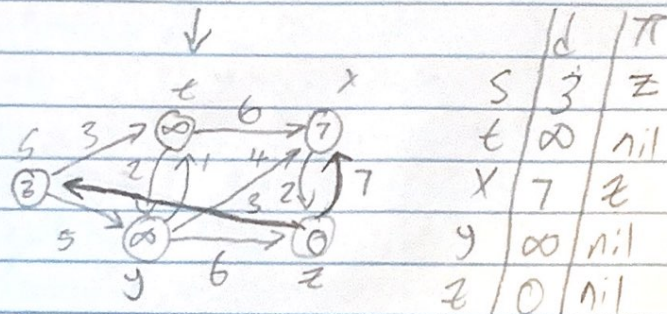
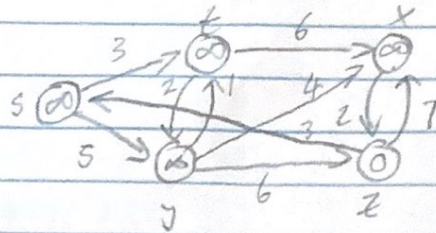


	d	π
s	0	nil
t	3	s
x	5	nil
y	∞	s
z	∞	nil



	d	π
s	0	nil
t	3	s
x	9	t
y	5	t
z	11	y

Vertex z as the source :



7. By adding edges one at a time, cycle is eventually likely to be detected. To deal with this we remove the edge that contains the most-cost edge. Keep iterating this process until all edges have been covered. To detect the most-cost edge we can run through root of the MST to each vertex. Comparing different paths, we are able to find the most-cost edge. Then update MST corresponding to it. The running time is $O(EV)$, since new vertex and incident edge is accounted for each iteration.