### **Karachi Institute of Economics and Technology College of Computing & Information Sciences**

### Accept and Reject Arithmetic Expression by Using CFG

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For

Automata 108983 – Thursday Spring 2022

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### 17-May-22

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#### **Abstract**

A grammar can be used to describe the possible hierarchical structure of a program.

Context-free grammars are more expressive than finite automata: if a language L is accepted by finite automata then L can be generated by a context-free grammar.

A context free grammar has 4 components: — A set of tokens, known as terminal symbols. — A set of no terminals. — A set of productions where each production consists of a no terminal, called the left side of the production, an arrow, and a sequence of tokens and/or no terminals, called the right side of the production. — A designation of one of the no terminals as the start symbol.

#### 1. Project Description

- A 4-tuple G = < V, S, P > is a context-free grammar (CFG) if V and are finite sets sharing no elements between them, S V is the start symbol, and P is a finite set of productions of the form X ->, where X V, and (V)\*.
- T describes a finite set of terminal symbols. S is the start symbol. In CFG, the start symbol is used to derive the string. You can derive the string by repeatedly replacing a non-terminal by the right-hand side of the production until all non-terminals have been replaced by terminal symbols.
- A context-free grammar is studied in the fields of theoretical computer science, compiler design, and linguistics. CFGs are used to describe programming languages and parser programs in compilers can be generated automatically from context-free grammar.

#### 2. 3rd party Libraries and purpose which you used in your project

- We used spilt function from library **re** also known as regular expression.
- Earlier we also used **rnd** as Random for generation of an arithmetic expression.

#### **3.** Project Meetings with your teacher (goals + achievements)

- First, we examined the problem, and how we can check the expression is valid or invalid by using CFG rules
- We understand this by looking into the libraries e.g., NLTK that can solve the same problem
- We also added **regex** which check the expression validity through regular expression but it under the scope of this project.
- We made use of python data type dictionary to define the rules for an arithmetic expression.

• The rules made by CFG method used to check for the input expression and returns if an expression is valid or invalid

# 4. Stack info = (Hardware & Software Requirement for your project)

- System should have python installed and should be run on Jupyter Notebook
- Modern Operating System:

```
x86 64-bit CPU (Intel / AMD architecture)
```

4 GB RAM.

5 GB free disk space

## 5. Methodology (explanation of complex code logic or complex variable structure)

- Rules is a python dictionary
- Spilt is used by regex library re
- Used indexing for list to split a num into different variables
- Open and closing are validated by data structure stack (push and pop)
- Conversion is also used in this program for associability of python features.

#### 6. Project Code

```
rules = {
  "S":[
      ["Exp"]
      "Exp":[
           ["BO","N","O","N","BC"]
      "BO":[
          ['(']
          ],
      "N":[
         '0','1','2','3','4','5','6','7','8','9',['N','N']
      "O": [
          ['+'],['-'],['*'],['/']
          ],
      "BC":[
          [')']
          ],
#(4+5) (12+10) (1+22) (55+698)
#final = input("Please enter an expression")
final = "(4+95)"
result = re.split(r" + |-| *| / | / | / |), final)
#print(result)
num1 = result[1]
num2 = result[2]
# print(num1)
# print(num2)
```

```
def is_valid(myStr):
  opening = ['(']
  closing = [')']
  stack = []
  for i in myStr:
     if i in opening:
       stack.append(i)
     elif i in closing:
       pos = closing.index(i)
       if ((len(stack) > 0)) and
             (opening[pos] == stack[-1])):
          stack.pop()
        else:
          stack.append(i)
     else:
       pass
  if len(stack) == 0:
     return True
  else:
     return False
bracketsAccpet = is_valid(final)
# print(bracketsAccpet, "B")
num1Accept= False
if((len(num1)>=2)):
  x = [int(a) \text{ for a in } str(num1)]
  print(x)
  count_num1=0
  for i in range (10):
     for j in range(len(x)):
       #print(rules['N'][i])
       #print('1'==rules['N'][i])
       if((str(x[i])==rules['N'][i])):
          count_num1=count_num1+1
          #print(count_num1)
          if(count_num1>=2):
             num1Accept = True
```

```
else:
  for i in range(10):
     if(str(num1)==rules['N'][i]):
       num1Accept = True
#print(num1Accept)
num2Accept= False
if((len(num2)>=2)):
  y = [int(b) \text{ for } b \text{ in } str(num2)]
  #print(y)
  count_num2=0
  for k in range(10):
     for 1 in range(len(y)):
       #print(rules['N'][i])
       #print('1'==rules['N'][i])
       if((str(y[1])==rules['N'][k])):
          count_num2=count_num2+1
          #print(count_num2)
          if(count_num2>=2):
            num2Accept = True
else:
  for i in range(10):
     if(str(num2)==rules['N'][i]):
       num2Accept = True
#print(num2Accept)
if(num1Accept and num2Accept and bracketsAccpet):
  print("Expression is Accepted")
else:
  print("Expression is Rejected")
```

#### 7. Three Test cases input + output

else:

```
for i in range(10):
        if(str(num2)==rules['N'][i]):
            num2Accept = True
#print(num2Accept)
if(num1Accept and num2Accept and bracketsAccpet):
    print("Expression is Accepted")
else:
    print("Expression is Rejected")
Please enter an expression(14+8)
[1, 4]
Expression is Accpeted
                  count num2=count num2+1
                #print(count_num2)
                if(count_num2>=2):
                   num2Accept = True
else:
    for i in range(10):
        if(str(num2)==rules['N'][i]):
            num2Accept = True
#print(num2Accept)
if(num1Accept and num2Accept and bracketsAccpet):
    print("Expression is Accepted")
else:
    print("Expression is Rejected")
Please enter an expression(4-+6)
Expression is Rejected
```

```
else:
    for i in range(10):
        if(str(num2)==rules['N'][i]):
            num2Accept = True

if(num1Accept and num2Accept and bracketsAccpet):
    print("Expression is Accpeted")

else:
    print("Expression is Rejected")

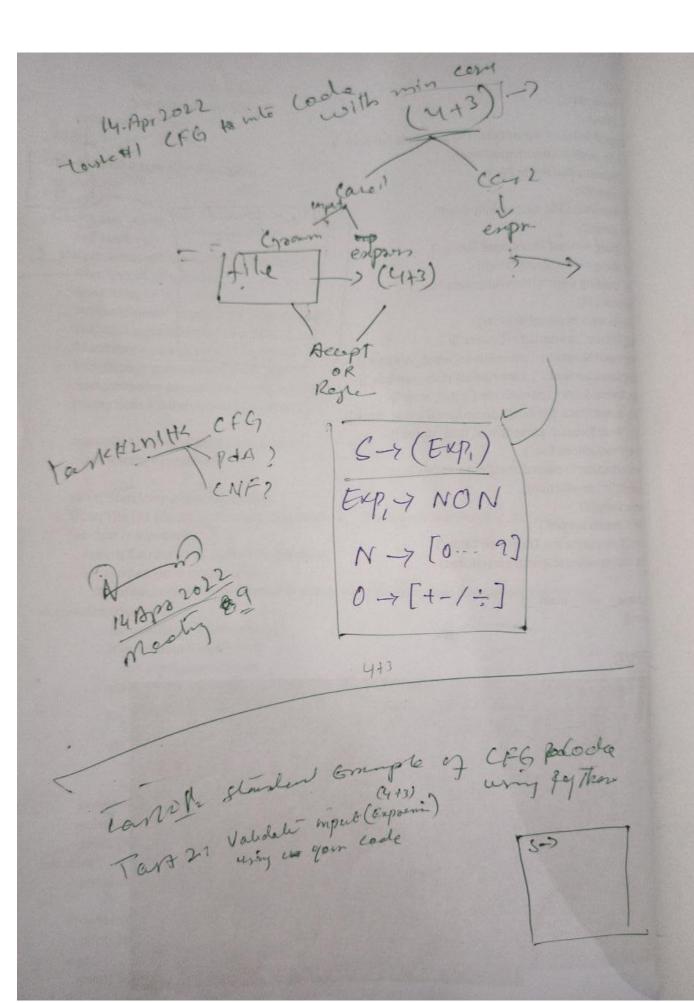
Please enter an expression(10/2)
[1, 0]

Expression is Accpeted
```

#### 8. Conclusion

In future, we want to improve it by adding filing and making it more dynamic by adding UI to it for easy use.

(4+ (3-2)) . Exp -> NON Expy -> Expy S-7 (EXPI) EXP, -> NO EXP2 EXP, -> NON. N->[0--- 9] Exp. -> Exp3 0 -> (+-/ × ] EXP3 -> NON N ->[0---- 9] \* ((2\*6-1)\*(2-1)) EXP,  $0 \rightarrow [+-/+]$ S-7EKP, EXP\_ > EXP\_2 O EXP3 ((5+4)-(3-1)) EXP - EXPY S-> EXP, EXP, > EXP, O EXP3. EXPY -> NONON Exp3 > Exp5 Exp -> Expy Exps -> NON EXPY -> NON N-> [0 --- 9] EXP3 -> EXP5 0-7(+-/\*). Exp. -> NON  $N \rightarrow [0 - - 9]$ 0->[+-/\*] ((3+1)/(1+4)-3(5\*3))S -> EXP, EXP, > EXP O EXP3 ON EXPY XP, > EXPS - MON 43 > EXPG



(4+2-1) (4+2-1) (4+21) (4+21) if(05 1 - Operator one only 2- Number can be (4+3) else SYEXP Exp -> (BO, N, O, N, BC) BO -> C 4 -> (0- 9), [NN) 0-) (+,-1/, 4) BL =>) mbo ( L

\* Task 3 (4+2) STOPENPER CFG

EXP. -> (EXP.)

CXP. -> (EXP.)

CNF

EXP. -> (EXP.)

O -> [+ -1.\*] CFG -> code Tape # 30\$3 xpy -> 1 KP3 -> t → [O--> [+ (3+1) using Lib EXP nitk. CFG