

Optimal Universal Quantum Circuits for Unitary Complex Conjugation

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Abstract—Let U_d be a unitary operator representing an arbitrary d -dimensional unitary quantum operation. This work presents optimal quantum circuits for transforming a number k of calls of U_d into its complex conjugate \overline{U}_d . Our circuits admit a parallel implementation and are proven to be optimal for any k and d with an average fidelity of $\langle F \rangle = \frac{k+1}{d(d-k)}$. Optimality is shown for average fidelity, robustness to noise, and other standard figures of merit. This extends previous works which considered the scenario of a single call ($k = 1$) of the operation U_d , and the special case of $k = d - 1$ calls. We then show that our results encompass optimal transformations from k calls of U_d to $f(U_d)$ for any arbitrary homomorphism f from the group of d -dimensional unitary operators to itself, since complex conjugation is the only non-trivial automorphism on the group of unitary operators. Finally, we apply our optimal complex conjugation implementation to design a probabilistic circuit for reversing arbitrary quantum evolutions.

Index Terms—Information science, quantum information science, quantum channels, quantum circuit.

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I. INTRODUCTION

THE field of quantum information and computation typically describes how physical devices process information encoded into quantum systems. Such devices are often referred to as quantum gates in mathematical frameworks, and can be combined into quantum circuits by defining the order of processing [1]. Conventionally, quantum gates were thought of as operations which take quantum states as inputs, and, after processing, output a quantum state. However, in a paradigm sometimes referred to as “higher-order quantum transformations” quantum gates rather than quantum states may be subject to transformation [2], [3]. Such higher-order quantum transformations can be viewed as a “circuit-board” approach for quantum operations, and has been found to have versatile applications encompassing quantum circuit design [4], [5], [6], control of quantum system dynamics [7], [8], [9], learning/storing of unitary operations [10], [11], [12], identifying cause-effect relations [13], and estimating and discriminating quantum channels [14], [15], [16], among others.

When considering transformations between unitary operations, previous works have considered tasks such as unitary cloning [17], unitary inversion [18], [19] (see also [20]), unitary transposition [19], unitary conjugation [21], unitary iteration [22], unitary controlisation [23], [24], [25], [26], and general homomorphisms and anti-homomorphisms between unitary operators [27], [28]. Unitary complex conjugation was first considered in [18], where non-exact optimal transformations between a single call of an arbitrary d -dimensional unitary operation into its complex conjugation were obtained. Later, [21] presented a parallel quantum circuit which attains exact unitary conjugation by making $k = d - 1$ calls of a d -dimensional unitary. Also, [29] considered the problem of exact unitary conjugation with an allowed rate of failure and showed that the probability of success is exactly zero when $k < d - 1$. Hence, the probability of transforming U_d into its complex conjugation cannot be enhanced by multi-call protocols such as the success-or-draw method [30].

In this work, we present deterministic and universal quantum circuits that transform k calls of an arbitrary d -dimensional unitary operator U_d into its complex conjugate \overline{U}_d . We present an explicit parallel circuit which makes no use of extra auxiliary space or memory for this task. We then prove that our construction attains the maximal performance

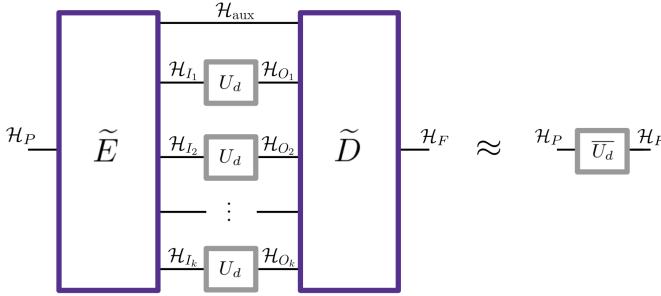


Fig. 1. Parallel superchannel transforming k calls of a unitary operation U_d into its complex conjugation $\overline{U_d}$. Here \tilde{E} stands for an “encoder” quantum channel which is performed before the input operations $U_d^{\otimes k}$ and \tilde{D} stands for a “decoder” quantum channel acting after the input operations. When $k < d - 1$, unitary complex conjugation can only be obtained approximately with an average fidelity of $\langle F \rangle = \frac{k+1}{d(d-k)}$. Here we prove this number and present optimal quantum circuits for this task. The optimal circuit can operate in parallel on input transformations and does not require additional auxiliary spaces (depicted as the extra line labelled by the space \mathcal{H}_{aux} above the gates U_d).

allowed by quantum theory for any number of calls k and any dimension d .

II. TRANSFORMING UNITARY OPERATIONS WITH QUANTUM CIRCUITS

In the following, we formally introduce the problem considered in this manuscript.

We start by pointing out that, with no loss in performance, we can restrict circuits designed for unitary complex conjugation to those that admit a parallel implementation, as depicted in Fig. 1. Indeed, it was shown in [27] that transformations f acting on unitary operators U as $f(U)$ can be implemented optimally by parallel circuits when f is a homomorphism with respect to operator composition, *i.e.*, for f satisfying $f(UV) = f(U)f(V)$ for every unitary U and V ¹. The case of unitary complex conjugation is naturally included, since $\overline{UV} = \overline{U}\overline{V}$.

In quantum theory, states are described by unit-trace, positive semidefinite linear operators² *i.e.*, $\rho \in \mathcal{L}(\mathcal{H})$, $\rho \geq 0$, and $\text{tr}(\rho) = 1$. Deterministic transformations between input-states and output-states are implemented by quantum channels, which are completely positive³ (CP) and trace-preserving⁴ (TP) linear maps⁵ $\tilde{C} : \mathcal{L}(\mathcal{H}_I) \rightarrow \mathcal{L}(\mathcal{H}_O)$, where \mathcal{H}_I and \mathcal{H}_O

¹Interestingly, parallel circuits are optimal even when considering transformations which do not respect a definite causal order, such as the quantum switch [31] or general process matrices [32].

²The symbol \mathcal{H} denotes a finite-dimensional Hilbert space which is isomorphic to \mathbb{C}_d for some dimension $d \in \mathbb{N}$. The symbol $\mathcal{L}(\mathcal{H})$ denotes the set of linear operators acting on \mathcal{H} (endomorphisms).

³A linear map $\tilde{C} : \mathcal{L}(\mathcal{H}_I) \rightarrow \mathcal{L}(\mathcal{H}_O)$ is positive when $\tilde{C}(\rho) \geq 0$ for every positive semidefinite linear operator $\rho \geq 0$, $\rho \in \mathcal{L}(\mathcal{H}_I)$. A map is CP if it is positive for all its trivial extensions, that is, $\tilde{C} \otimes \tilde{\mathbb{1}}(\sigma) \geq 0$ for every $\sigma \geq 0$, $\sigma \in \mathcal{L}(\mathcal{H} \otimes \mathcal{H}_{\text{aux}})$ where $\tilde{\mathbb{1}} : \mathcal{L}(\mathcal{H}_{\text{aux}}) \rightarrow \mathcal{L}(\mathcal{H}_{\text{aux}})$ is the identity map on an arbitrary finite-dimensional space \mathcal{H}_{aux} , *i.e.*, $\tilde{\mathbb{1}}(\rho) = \rho$, $\forall \rho \in \mathcal{L}(\mathcal{H}_{\text{aux}})$.

⁴A linear map $\tilde{C} : \mathcal{L}(\mathcal{H}_I) \rightarrow \mathcal{L}(\mathcal{H}_O)$ is TP when $\text{tr}(\tilde{C}(\rho)) = \text{tr}(\rho)$ for every linear operator $\rho \in \mathcal{L}(\mathcal{H}_I)$.

⁵In this paper, we reserve the word map for linear transformations between linear operators. Also, adopt the convention of writing a tilde on the top of maps to make an easy distinction between operators, *e.g.* $C \in \mathcal{L}(\mathcal{H}_I \otimes \mathcal{H}_O)$ and maps $\tilde{C} : \mathcal{L}(\mathcal{H}_I) \rightarrow \mathcal{L}(\mathcal{H}_O)$. This convention is particularly useful when dealing with Choi representations of linear maps.

are the linear spaces relative to input and output, respectively. A quantum channel $\tilde{U} : \mathcal{L}(\mathcal{H}_I) \rightarrow \mathcal{L}(\mathcal{H}_O)$ is said to be unitary if it can be written as $\tilde{U}(\rho) = U_d \rho U_d^\dagger$ for some d -dimensional unitary operator U_d . Note that unitary operators which only differ by a global phase represent the same unitary channel. Hence, with no loss of generality, we may assume that the unitary operator U_d belongs to the group of special unitary operators $\mathcal{SU}(d)$, that is, unitary operators with determinant one.

Let $\widetilde{C}_{\text{in}} : \mathcal{L}(\mathcal{H}_I) \rightarrow \mathcal{L}(\mathcal{H}_O)$ be an arbitrary input channel to a quantum circuit, which transforms it into an output channel $\widetilde{C}_{\text{out}} : \mathcal{L}(\mathcal{H}_P) \rightarrow \mathcal{L}(\mathcal{H}_F)$. The labels P and F stands for past and future, respectively (see Fig. 1). Quantum circuits designed to obtain the transformation $\widetilde{C}_{\text{in}} \mapsto \widetilde{C}_{\text{out}}$ may be analysed by means of encoder and decoder channels [4], [5], a method which we describe in the following.

- 1) Before performing the input operation $\widetilde{C}_{\text{in}}$, we apply a quantum channel (encoder) $\tilde{E} : \mathcal{L}(\mathcal{H}_P) \rightarrow \mathcal{L}(\mathcal{H}_I \otimes \mathcal{H}_{\text{aux}})$, where \mathcal{H}_{aux} is an arbitrary auxiliary (memory) space which may be used to improve performance.
- 2) Then, the operation $\widetilde{C}_{\text{in}} \otimes \widetilde{\mathbb{1}}_{\text{aux}}$ is applied, where $\widetilde{\mathbb{1}}_{\text{aux}}$ is the identity map, which acts trivially on the auxiliary system, *i.e.*, $\widetilde{\mathbb{1}}_{\text{aux}}(\rho) = \rho$, for every operator ρ .
- 3) Finally, we perform a quantum channel (decoder) $\tilde{D} : \mathcal{L}(\mathcal{H}_O \otimes \mathcal{H}_{\text{aux}}) \rightarrow \mathcal{L}(\mathcal{H}_F)$ to obtain the output

$$\widetilde{C}_{\text{out}} = \tilde{D} \circ (\widetilde{C}_{\text{in}} \otimes \widetilde{\mathbb{1}}_{\text{aux}}) \circ \tilde{E}. \quad (1)$$

The use of k independent calls of a unitary channel $\widetilde{U_d}$ may be mathematically represented by a single channel $\widetilde{U_d}^{\otimes k}$. Then, by identifying $\widetilde{U_d}^{\otimes k}$ as $\widetilde{C}_{\text{in}}$ in the routine described above, a parallel quantum circuit transforms k calls of $\widetilde{U_d}$ as

$$\widetilde{U_d}^{\otimes k} \mapsto \widetilde{D} \circ (\widetilde{U_d}^{\otimes k} \otimes \widetilde{\mathbb{1}}) \circ \widetilde{E}, \quad (2)$$

as illustrated in Fig. 1. When the input channel is composed of k operations, it is convenient to write the input space \mathcal{H}_I explicitly as $\mathcal{H}_I := \mathcal{H}_{I_1} \otimes \dots \otimes \mathcal{H}_{I_k}$, where the bold letter I emphasises that the linear space \mathcal{H}_I has an internal tensor product structure (similarly, we write $\mathcal{H}_O := \mathcal{H}_{O_1} \otimes \dots \otimes \mathcal{H}_{O_k}$). Further, we have $\mathcal{H}_P \cong \mathcal{H}_F \cong \mathcal{H}_{I_i} \cong \mathcal{H}_{O_i} \cong \mathbb{C}_d$ for every $i \in \{1, \dots, k\}$ and $\mathcal{H}_I \cong \mathcal{H}_O \cong \mathbb{C}_d^{\otimes k}$ since the quantum circuit transforms k calls of d -dimensional unitary operations into a d -dimensional unitary operation.

The performance of transformations can be quantified by means of channel fidelity, which is known to respect several properties that are relevant for information processing [33]. The fidelity between quantum channels is usually defined through their Choi operators. The Choi operator of a quantum channel is obtained via the Choi isomorphism is defined as follows. Let $T : \mathcal{H}_I \rightarrow \mathcal{H}_O$ be a linear operator. Its Choi vector $|T\rangle\rangle \in \mathcal{H}_I \otimes \mathcal{H}_O$ is,

$$|T\rangle\rangle := \sum_i |i\rangle \otimes (T|i\rangle), \quad (3)$$

where $\{|i\rangle\}_{i=1}^d$ is the canonical orthonormal basis for \mathbb{C}_d , often referred to as computational basis. Further, let

$\tilde{C} : \mathcal{L}(\mathcal{H}_I) \rightarrow \mathcal{L}(\mathcal{H}_O)$ be a linear map between linear operators. Its Choi operator $C \in \mathcal{L}(\mathcal{H}_I \otimes \mathcal{H}_O)$ is

$$C := \sum_{ij} |i\rangle\langle j| \otimes \tilde{C}(|i\rangle\langle j|). \quad (4)$$

Direct calculation shows that the Choi operator of a unitary channel $\widetilde{U}_d(\rho) = U_d \rho U_d^\dagger$ may be conveniently expressed as $|\widetilde{U}\rangle\langle\widetilde{U}|$.

The fidelity between an arbitrary quantum channel \tilde{C} and a unitary channel acting as $\widetilde{U}_d(\rho) = U_d \rho U_d^\dagger$ on an input ρ is given by [33]

$$F(\tilde{C}, \widetilde{U}_d) := \frac{1}{d^2} \langle\langle U_d | C | U_d \rangle\rangle, \quad (5)$$

where $|U_d\rangle\rangle$ the Choi vector of U_d .

We now turn to the task of transforming k copies of a unitary operator U_d into its complex conjugate \overline{U}_d with respect to the canonical computational basis $\{|i\rangle\}_{i=1}^d$.

Let us decompose $U_d = \sum_{ij} u_{ij} |i\rangle\langle j|$ in the computational basis. The complex conjugate of U_d in said basis reads $\overline{U}_d := \sum_{ij} \overline{u_{ij}} |i\rangle\langle j|$. Note that the definition of \overline{U}_d depends on the choice of basis. To see this, let us express U_d in a different orthonormal basis $\{|i'\rangle\}_{i=1}^d$ as $U_d = \sum_{ij} u'_{ij} |i'\rangle\langle j'|$. Since $u'_{ij} = \langle i'|U_d|j'\rangle$ may be different to $u_{ij} = \langle i|U_d|j\rangle$, we conclude

$$\overline{U}_d := \sum_{ij} \overline{u_{ij}} |i\rangle\langle j| \neq \sum_{ij} \overline{u'_{ij}} |i'\rangle\langle j'| =: \overline{U}'_d. \quad (6)$$

However, \overline{U}_d and \overline{U}'_d are unitarily equivalent: let $W \in \mathcal{L}(\mathbb{C}_d)$ be a unitary operator mapping the $\{|i'\rangle\}_{i=1}^d$ basis into the $\{|i\rangle\}_{i=1}^d$ basis. One simple example is given by $W := \sum_i |i\rangle\langle i'|$. Then, we find $W U_d W^\dagger = \sum_{ij} u'_{ij} |i\rangle\langle j|$, and

$$\begin{aligned} \overline{W U_d W^\dagger} &= \sum_{ij} \overline{u'_{ij}} |i'\rangle\langle j'| \\ &= \overline{U}'_d, \end{aligned} \quad (7)$$

leading to the identity $\overline{U}'_d = \overline{W} \overline{U}_d W^T$, where $\overline{W}^T = W^T$ is the transposition of W on the computational basis⁶.

As discussed earlier, since the complex conjugation function is a homomorphism, *i.e.*, $\overline{UV} = \overline{U} \overline{V}$, we can restrict our analysis to parallel quantum circuits. That is, we focus on circuits which can be as,

$$\widetilde{U}_d^{\otimes k} \mapsto \widetilde{D} \circ (\widetilde{U}_d^{\otimes k} \otimes \widetilde{1}) \circ \widetilde{E} \approx \widetilde{U}_d, \quad (9)$$

where \widetilde{E} and \widetilde{D} are quantum channels, as presented in Eq. (1). We now set the average channel fidelity as our figure of merit. More precisely, we quantify the performance of our transformation via

$$\langle F \rangle := \int_{U_d \in \mathcal{SU}(d)} F\left(\widetilde{D} \circ (\widetilde{U}_d^{\otimes k} \otimes \widetilde{1}) \circ \widetilde{E}, \widetilde{U}_d\right) dU_d, \quad (10)$$

where the fidelity of quantum channels is defined in Eq. (5) and the integral is taken in accordance to the Haar measure. As proven in [28], for the case of homomorphic

⁶Let $A = \sum_{ij} a_{ij} |i\rangle\langle j|$ be an arbitrary operator acting on \mathbb{C}_d , its transposition in the computational basis is defined as $A^T = \sum_{ij} a_{ij} |j\rangle\langle i|$.

transformations (as the complex conjugation), the optimal average fidelity coincides with the optimal worst-case fidelity and there is a one-to-one correspondence between average fidelity and the robustness of the transformation to noise, ensuring that the average fidelity is indeed a relevant figure of merit.

III. OPTIMAL QUANTUM CIRCUIT FOR UNITARY COMPLEX CONJUGATION

In this section, we state our main result and provide an explicit construction for the optimal quantum circuit for complex conjugation of an arbitrary unitary operation. The proof of optimality and performance evaluation of the circuit presented here are given in Sec. IV.

Theorem 1: Let U_d be a unitary operator representing an arbitrary d -dimensional unitary operation $\widetilde{U}_d(\rho) = U_d \rho U_d^\dagger$. When $k \leq d-1$ uses are available, the optimal quantum circuit which transforms k uses of U_d into its complex conjugation \overline{U}_d attains average fidelity $\langle F \rangle = \frac{k+1}{d(d-k)}$. Additionally, the optimal circuit admits a parallel implementation which makes no use of additional auxiliary spaces. Note that when $k = d-1$, we obtain $\langle F \rangle = 1$, hence unitary complex conjugation may be done exactly when the number of calls respects $k \geq d-1$. This optimal and exact transformation for the $k = d-1$ case was first presented at [21] and will be reviewed in this section for convenience.

It is useful to state some facts regarding the antisymmetric subspace, which will play a major role in constructing optimal quantum circuits for unitary complex conjugation. Let S_N be the group of permutations of N elements, and $V(\pi) \in \mathbb{C}_d^{\otimes N}$ be a representation for the permutation $\pi \in S_N$ in $\mathbb{C}_d^{\otimes N}$, explicitly defined via

$$V(\pi)|\psi_1\rangle \otimes |\psi_2\rangle \otimes \dots \otimes |\psi_N\rangle = |\psi_{\pi^{-1}(1)}\rangle \otimes |\psi_{\pi^{-1}(2)}\rangle \otimes \dots \otimes |\psi_{\pi^{-1}(N)}\rangle. \quad (11)$$

The antisymmetric subspace of N d -dimensional parties is a subspace of $\mathbb{C}_d^{\otimes N}$ formed by vectors $|\psi\rangle \in \mathbb{C}_d^{\otimes N}$ satisfying

$$V(\pi)|\psi\rangle = \text{sgn}(\pi)|\psi\rangle, \quad \forall \pi \in S_N \quad (12)$$

where $\text{sgn}(\pi)$ is the sign⁷ of the permutation π . When $d = N$, the antisymmetric subspace of $\mathbb{C}_d^{\otimes N}$ is one-dimensional, and it is spanned by a single vector, known as the totally antisymmetric state,

$$|\psi_{A_d}\rangle := \frac{1}{\sqrt{N!}} \sum_{\pi \in S_N} \text{sgn}(\pi) |\pi(1)\rangle \otimes |\pi(2)\rangle \otimes \dots \otimes |\pi(d)\rangle, \quad (13)$$

where $\{|1\rangle, |2\rangle, \dots, |d\rangle\}$ is the computational basis for \mathbb{C}_d . When $N = d$, the totally antisymmetric state is known to satisfy

$$U_d^{\otimes d} |\psi_{A_d}\rangle = |\psi_{A_d}\rangle, \quad \forall U_d \in \mathcal{SU}(d), \quad (14)$$

which is equivalent to

$$(\mathbb{1}_d \otimes U_d^{\otimes(d-1)}) |\psi_{A_d}\rangle = (U_d^\dagger \otimes \mathbb{1}_d^{\otimes(d-1)}) |\psi_{A_d}\rangle. \quad (15)$$

⁷That is, $\text{sgn}(\pi) = +1$ when π can be written as a composition of even transpositions (pairwise permutation) and $\text{sgn}(\pi) = -1$ when π can be written as a composition of odd transpositions.

The transposition vector $|\psi\rangle \in \mathbb{C}_d$ on the computational basis is a linear map from \mathbb{C}_d to its dual space defined via $|i\rangle^T := \langle i|$. One can then verify that $(U_d|\psi\rangle)^T = \langle\psi| U_d^T$. When $|\psi\rangle_{12} \in \mathbb{C}_d \otimes \mathbb{C}_{d'}$ is a bipartite vector which can be written as $|\psi\rangle_{12} = \sum_{ij} \gamma_{ij} |i\rangle_1 \otimes |j\rangle_2$ in the computational basis, the partial transpose on the first linear space reads

$$|\psi\rangle_{12}^{T_1} := \sum_{ij} \gamma_{ij} |i\rangle_1^T \otimes |j\rangle_2 \quad (16)$$

$$= \sum_{ij} \gamma_{ij} \langle i|_1 \otimes |j\rangle_2. \quad (17)$$

Hence, doing the partial transposition on the first linear space of Eq. (15), we obtain

$$\mathbb{1}_d \otimes U_d^{\otimes(d-1)} |\psi_{A_d}\rangle^{T_1} = \left(\mathbb{1}_d \otimes U_d^{\otimes(d-1)} |\psi_{A_d}\rangle \right)^{T_1} \quad (18)$$

$$= \left(U_d^\dagger \otimes \mathbb{1}_d^{\otimes(d-1)} |\psi_{A_d}\rangle \right)^{T_1} \quad (19)$$

$$= |\psi_{A_d}\rangle^{T_1} \left(\overline{U_d} \otimes \mathbb{1}_d^{\otimes(d-1)} \right). \quad (20)$$

With Eq. (18), we can now recover the result of [21] stating that when $k = d - 1$, it is possible to transform $U_d^{\otimes(d-1)}$ into $\overline{U_d}$ exactly. We now define the linear transformation $V : \mathbb{C}_d \rightarrow \mathbb{C}_d^{\otimes(d-1)}$ as,

$$V := \frac{\sqrt{d!}}{\sqrt{(d-1)!}} |\psi_{A_d}\rangle^{T_1} \quad (21)$$

$$= \frac{1}{\sqrt{(d-1)!}} \sum_{\pi \in S_N} \text{sgn}(\pi) \langle \pi(1)| \otimes |\pi(2)\rangle \otimes \dots \otimes |\pi(d)\rangle. \quad (22)$$

Note that V is an isometry since $V^\dagger V = \mathbb{1}_d$. Now, in order to transform $k = d - 1$ uses of U_d into $\overline{U_d}$, we simply perform V before $U_d^{\otimes(d-1)}$, and the operator V^\dagger after to obtain

$$V^\dagger U_d^{\otimes(d-1)} V = \overline{U_d}. \quad (23)$$

In the following, we move to the case $k < d - 1$. The identification of the optimal circuit for this scenario is the main finding of this work. The case $k < d - 1$ turns out to be more involved, as we need to consider scenarios where the antisymmetric subspace is not one-dimensional. When $k+1 < d$, the antisymmetric subspace of $\mathbb{C}_d^{\otimes(k+1)}$ is spanned by a set of $\binom{d}{k+1} := \frac{d!}{(k+1)!(d-k-1)!}$ orthogonal vectors given by $|\psi_{A_{d,k+1}}^{(i)}\rangle$. These vectors may be explicitly constructed by a simple method: choose $k+1$ vectors from $\{|i\rangle\}_{i=1}^d$ and add the permutation-signed sum over all $k+1$ permutations. For the sake of concreteness we set the vector with index $i = 1$ as

$$|\psi_{A_{d,k+1}}^{(1)}\rangle := \frac{1}{\sqrt{(k+1)!}} \sum_{\pi \in S_{k+1}} \text{sgn}(\pi) |\pi(1)\rangle \otimes |\pi(2)\rangle \otimes \dots \otimes |\pi(k+1)\rangle. \quad (24)$$

Since $k+1 < d$, other orthogonal antisymmetric vectors can be constructed – for instance

$$|\psi_{A_{d,k+1}}^{(2)}\rangle := \frac{1}{\sqrt{(k+1)!}} \sum_{\pi \in S_{k+1}} \text{sgn}(\pi) |\pi(2)\rangle \otimes |\pi(3)\rangle \otimes \dots \otimes |\pi(k+2)\rangle \quad (25)$$

which also respects $V(\pi)|\psi_{A_{d,k+1}}^{(2)}\rangle = \text{sgn}(\pi)|\psi_{A_{d,k+1}}^{(2)}\rangle$ for every permutation $\pi \in S_{k+1}$. The set of $\binom{d}{k+1}$ vectors of this form then span the whole antisymmetric space of $\mathbb{C}_d^{\otimes(k+1)}$. We can now write the projector onto the antisymmetric space as,

$$A(d, k+1) = \sum_{i=1}^{\binom{d}{k+1}} |\psi_{A_{d,k+1}}^{(i)}\rangle \langle \psi_{A_{d,k+1}}^{(i)}|, \quad (26)$$

which may also be equivalently written as [34]

$$A(d, k+1) = \sum_{\pi \in S_{k+1}} \text{sgn}(\pi) V(\pi). \quad (27)$$

We are now in position to present the optimal circuit for unitary complex conjugation. Our circuit does not make use of any auxiliary space. The action of the encoder channel $\widetilde{E} : \mathcal{L}(\mathcal{H}_P) \rightarrow \mathcal{L}(\mathcal{H}_I)$, with $\mathcal{H}_P \cong \mathbb{C}_d$, and $\mathcal{H}_I \cong \mathbb{C}_d^{\otimes k}$ may be equivalently expressed in the Choi or operator-sum representation, which are respectively given by

$$\widetilde{E}(\rho) := \text{tr}_P \left((\rho_P^T \otimes \mathbb{1}_I) E_{PI} \right) \quad (28)$$

$$= \sum_{i=1}^{\binom{d}{k+1}} K_i \rho K_i^\dagger, \quad (29)$$

where tr_P denotes the trace over the past space \mathcal{H}_P , and $E \in \mathcal{L}(\mathcal{H}_P \otimes \mathcal{H}_I)$,

$$E := \frac{d}{\binom{d}{k+1}} A(d, k+1) \quad (30)$$

is the Choi operator of the map \widetilde{E} and

$$K_i := \sqrt{\frac{d}{\binom{d}{k+1}}} |\psi_{A_{d,k+1}}^{(i)}\rangle_{PI}^{T_P} \quad (31)$$

are the Kraus operators of \widetilde{E} . The decoder channel $\widetilde{D} : \mathcal{L}(\mathcal{H}_O) \rightarrow \mathcal{L}(\mathcal{H}_F)$, where $\mathcal{H}_O \cong \mathbb{C}_d^{\otimes k}$ and $\mathcal{H}_F \cong \mathbb{C}_d$, may be equivalently expressed in the Choi or operator-sum representation, given by

$$\widetilde{D}(\rho) := \text{tr}_O \left((\rho_O^T \otimes \mathbb{1}_F) D_{OF} \right) \quad (32)$$

$$= \sum_{i=1}^{\binom{d}{k+1}} K_i^\dagger \rho_O K_i + \text{tr} \left((\mathbb{1}_d^{\otimes k} - A(d, k)) \rho \right) \sigma_F. \quad (33)$$

Here, $D \in \mathcal{L}(\mathcal{H}_O \otimes \mathcal{H}_F)$,

$$D := \frac{\binom{d}{k}}{\binom{d}{k+1}} A(d, k+1)_{OF} + \left(\mathbb{1}_d^{\otimes k} - A(d, k) \right)_O \otimes \sigma_F, \quad (34)$$

is the Choi operator of the map \widetilde{D} and $\sigma \in \mathcal{L}(\mathcal{H}_F) \cong \mathcal{L}(\mathbb{C}_d)$ is an arbitrary quantum state, i.e., $\sigma \geq 0$ and $\text{tr}(\sigma) = 1$.

In Sec. IV, we prove that the circuit with the encoder described in Eq. (28) and the decoder of Eq. (32) achieves optimal unitary complex conjugation. More precisely, we will show that for any unitary operation \widetilde{U}_d the fidelity between $\widetilde{D} \circ \widetilde{U}_d^{\otimes k} \circ \widetilde{E}$ and \overline{U}_d is given by

$$F\left(\widetilde{D} \circ \widetilde{U}_d^{\otimes k} \circ \widetilde{E}, \overline{U}_d\right) = \frac{k+1}{d(d-k)}. \quad (35)$$

Moreover, we show that no other protocol can perform better: for every possible choice of encoder \tilde{E} and decoder \tilde{D} the average channel fidelity satisfies

$$\int_{U_d \in \mathcal{SU}(d)} F\left(\tilde{D} \circ (\tilde{U}_d^{\otimes k} \otimes \tilde{\mathbb{1}}) \circ \tilde{E}, \tilde{U}_d\right) dU_d \leq \frac{k+1}{d(d-k)}. \quad (36)$$

As discussed earlier, the optimal average fidelity coincides with the optimal worst-case fidelity since the complex conjugation is a homomorphism. Also, for homomorphic transformations, protocols which are optimal with respect to average fidelity are also optimal with respect to the robustness to white noise [28].

IV. PROOF OF OPTIMALITY

A. Quantum Superchannels

A useful mathematical framework for analysing transformations between quantum operations is given by the formalism of quantum superchannels, also referred to as 1-slot quantum combs or quantum supermaps [4], [5], [6].

Definition 1: A linear operator $S \in \mathcal{L}(\mathcal{H}_P \otimes \mathcal{H}_I \otimes \mathcal{H}_O \otimes \mathcal{H}_F)$ is a parallel superchannel if it can be written as

$$S = \text{tr}_{\text{aux}} \left((E_{P\mathbf{I}\text{aux}}^{T_{\text{aux}}} \otimes \mathbb{1}_{\mathbf{O}F}) (\mathbb{1}_{PI} \otimes D_{\text{aux}OF}) \right), \quad (37)$$

where \mathcal{H}_{aux} is an arbitrary finite-dimensional linear space, T_{aux} is the partial transpose in the \mathcal{H}_{aux} space, and $E \in \mathcal{L}(\mathcal{H}_P \otimes \mathcal{H}_I \otimes \mathcal{H}_{\text{aux}})$ and $D \in \mathcal{L}(\mathcal{H}_{\text{aux}} \otimes \mathcal{H}_O \otimes \mathcal{H}_F)$ are operators such that,

$$E \geq 0, \quad \text{tr}_{\text{aux}}(E) = \mathbb{1}_P \quad (38)$$

$$D \geq 0, \quad \text{tr}_F(D) = \mathbb{1}_{\text{aux}OF}. \quad (39)$$

That is, E is the Choi operator of a quantum channel from $\mathcal{L}(\mathcal{H}_P)$ to $\mathcal{L}(\mathcal{H}_I \otimes \mathcal{H}_{\text{aux}})$ and D is the Choi operator of a quantum channel from $\mathcal{L}(\mathcal{H}_{\text{aux}} \otimes \mathcal{H}_O)$ to $\mathcal{L}(\mathcal{H}_F)$.

The concept of superchannel presented in Def. 1 is useful because it incorporates every possible quantum circuit based on encoder-decoder schemes [5], [6]: a positive semidefinite operator $S \geq 0$, $S \in \mathcal{L}(\mathcal{H}_P \otimes \mathcal{H}_I \otimes \mathcal{H}_O \otimes \mathcal{H}_F)$ is a parallel superchannel if and only if it there exists a positive semidefinite operator $C \geq 0$ $C \in \mathcal{L}(\mathcal{H}_P \otimes \mathcal{H}_I)$ such that,

$$\text{tr}_F(S) = C_{PI} \otimes \mathbb{1}_O, \quad (40)$$

$$\text{tr}_I(C) = \mathbb{1}_P \quad (41)$$

that is, C is the Choi operator of a quantum channel from $\mathcal{L}(\mathcal{H}_P)$ to $\mathcal{L}(\mathcal{H}_I)$. The decomposition of Eq. (40) can be used to explicitly construct an encoder and decoder channel [5], [6]. For that, we set the auxiliary space as $\mathcal{H}_{\text{aux}} := \mathcal{H}_{P'} \otimes \mathcal{H}_{I'}$ where $\mathcal{H}_{P'}$ and $\mathcal{H}_{I'}$ are isomorphic \mathcal{H}_P and \mathcal{H}_I and define

$$E := \left(\sqrt{C_{PI}} \otimes \mathbb{1}_{P'I'} \right) |\mathbb{1}\rangle\langle\mathbb{1}|_{PP'} \otimes |\mathbb{1}\rangle\langle\mathbb{1}|_{II'} \left(\sqrt{C_{PI}} \otimes \mathbb{1}_{P'I'} \right)^{\dagger} \quad (42)$$

$$D := \left(\sqrt{C_{P'I'}}^{-1} \otimes \mathbb{1}_{OF} \right) S_{P'I'OF} \left(\sqrt{C_{P'I'}}^{-1} \otimes \mathbb{1}_{OF} \right)^{\dagger}, \quad (43)$$

where \sqrt{C} is the unique positive semidefinite square root of C and $\sqrt{C_{P'I'}}^{-1}$ is the Moore-Penrose pseudoinverse defined on the support of $\sqrt{C_{P'I'}}$. By using the relation $C \otimes \mathbb{1}|\mathbb{1}\rangle\langle\mathbb{1}| = \mathbb{1} \otimes C^T|\mathbb{1}\rangle\langle\mathbb{1}|$ and the cyclic property of the trace, we verify that $S = \text{tr}_{\text{aux}} \left((E_{P\mathbf{I}\text{aux}}^{T_{\text{aux}}} \otimes \mathbb{1}_{OF}) (\mathbb{1}_{PI} \otimes D_{\text{aux}OF}) \right)$.

The definition of parallel superchannels prove convenient to investigate the action of parallel superchannels on an arbitrary input channel $\tilde{C} : \mathcal{H}_I \rightarrow \mathcal{H}_O$. In particular, let $C_{IO} \in \mathcal{L}(\mathcal{H}_I \otimes \mathcal{H}_O)$ be the Choi operator of \tilde{C} , [4], [6], [18], and [28] show that

$$\text{tr}_{IO} \left(S \left(\mathbb{1}_P \otimes C_{IO}^T \otimes \mathbb{1}_F \right) \right) \in \mathcal{L}(\mathcal{H}_P \otimes \mathcal{H}_F) \quad (44)$$

is the Choi operator of

$$\tilde{D} \circ \left(\tilde{C} \otimes \widetilde{\mathbb{1}_{\text{aux}}} \right) \circ \tilde{E} : \mathcal{L}(\mathcal{H}_P) \rightarrow \mathcal{L}(\mathcal{H}_F). \quad (45)$$

B. The Performance Operator

A useful tool to evaluate the performance of superchannels for transforming k uses of a unitary operation U_d into $f(U_d)$ is given by the *performance operator*⁸ [18], [28], defined as

$$\Omega := \frac{1}{d^2} \int_{U_d \in \mathcal{SU}(d)} |f(U_d)\rangle\langle f(U_d)|_{PF} \otimes |U_d^{\otimes k}\rangle\langle U_d^{\otimes k}|_{IO}^T dU_d. \quad (46)$$

From Eq. (44) and Eq. (45) we see that using a superchannel S with encoder channel \tilde{E} and decoder channel \tilde{D} , the average fidelity can be written as

$$\text{tr}(S\Omega) = \frac{1}{d^2} \int \text{tr} \left(S |f(U_d)\rangle\langle f(U_d)|_{PF} \otimes |U_d^{\otimes k}\rangle\langle U_d^{\otimes k}|_{IO}^T \right) dU_d \quad (47)$$

$$= \int \frac{1}{d^2} \text{tr}_{PF} \left(|f(U_d)\rangle\langle f(U_d)|_{PF} \right) \quad (48)$$

$$\left[\text{tr}_{IO} \left(S \mathbb{1}_P \otimes |U_d^{\otimes k}\rangle\langle U_d^{\otimes k}|_{IO}^T \otimes \mathbb{1}_F \right) \right] dU_d \quad (49)$$

Hence, when the performance operator Ω is given, by combining Eq. (47) with the superchannel characterisation given by Eq. (40), the problem of finding the optimal average fidelity with a parallel superchannel reads

$$\max \text{tr}(S\Omega) \quad (50)$$

$$\text{s.t.: } S \geq 0 \quad (51)$$

$$\text{tr}_F(S) = C_{PI} \otimes \mathbb{1}_O \quad (52)$$

$$\text{tr}_I(C) = \mathbb{1}_P. \quad (53)$$

References [18] and [28] shows that the dual optimisation problem is given by

$$\min c \quad (54)$$

$$\text{s.t.: } \Omega \leq c\hat{S} \quad (55)$$

$$\hat{S} = W_{P\mathbf{I}\mathbf{O}} \otimes \mathbb{1}_F \quad (56)$$

$$\text{tr}_O(W) = \rho_P \otimes \mathbb{1}_I \quad (57)$$

$$\text{tr}(\rho) = 1. \quad (58)$$

As the notation suggests, $\rho \in \mathcal{L}(\mathcal{H}_P)$ represents a quantum state, $W \in \mathcal{L}(\mathcal{H}_P \otimes \mathcal{H}_I \otimes \mathcal{H}_O)$ is a causally ordered process [6], [32], and \hat{S} lies in the dual affine set of (parallel)

⁸Here we use the definition of [28], which is differs from the definition of [18] by a global transposition. We emphasise that for solve opmisation problems involving superchannels, the definitions from [18] and [28] are fully equivalent since S is a superchannel if and only if S^T is a superchannel.

superchannels [15], [18]. By definition of the dual problem, any feasible point for the dual problem constitutes as an upper bound for the primal. Additionally, for this particular problem we have strong duality, which means that the optimal value for the dual problem coincides with the optimal value for the primal problem.

Strong duality can be shown by making use of the Slater conditions, which state that, if there exists an operator which satisfy the constraints of the primal (or dual) problem strict, the optimisation problem has strong duality. For our particular problem, it is then enough to find an operator $S \in \mathcal{L}(\mathcal{H}_P \otimes \mathcal{H}_I \otimes \mathcal{H}_O \otimes \mathcal{H}_F)$ which is strictly positive $S > 0$, $\text{tr}_F(S) = C_{PI} \otimes \mathbb{1}_O$, and $\text{tr}_I(C) = \mathbb{1}_P$. One such choice which trivially satisfy such strictly properties is $S = \mathbb{1}_P \otimes \frac{\mathbb{1}_O}{d_I} \otimes \mathbb{1}_O \otimes \frac{\mathbb{1}_F}{d_F}$.

Since $|U_d\rangle\langle U_d|^T = |\overline{U}_d\rangle\langle \overline{U}_d|$, the performance operator for transforming k calls of a unitary U_d into its complex conjugate \overline{U}_d reads as

$$\Omega = \frac{1}{d^2} \int_{U_d \in \mathcal{SU}(d)} |\overline{U}_d\rangle\langle \overline{U}_d|_{PF} \otimes |U_d^{\otimes k}\rangle\langle U_d^{\otimes k}|_{IO}^T dU_d, \quad (59)$$

$$= \frac{1}{d^2} \int_{U_d \in \mathcal{SU}(d)} |\overline{U}_d\rangle\langle \overline{U}_d|_{PF} \otimes |\overline{U}_d^{\otimes k}\rangle\langle \overline{U}_d^{\otimes k}|_{IO} dU_d, \quad (60)$$

$$= \frac{1}{d^2} \int_{U_d \in \mathcal{SU}(d)} |\overline{U}_d\rangle\langle \overline{U}_d|_{PF} \otimes |U_d^{\otimes k}\rangle\langle U_d^{\otimes k}|_{IO} dU_d \quad (61)$$

$$= \frac{1}{d^2} \int \left(\mathbb{1}_P \otimes \mathbb{1}_I^{\otimes k} \otimes (U_d^{\otimes k})_O \otimes (U_d)_F \right) \quad (62)$$

$$|\mathbb{1}\rangle\langle \mathbb{1}|_{PF} \otimes |\mathbb{1}\rangle\langle \mathbb{1}|_{IO} \left(\mathbb{1}_P \otimes \mathbb{1}_I^{\otimes k} \otimes (U_d^{\otimes k})_O \otimes (U_d)_F \right)^{\dagger} dU_d. \quad (63)$$

Hence, the performance operator respects the commutation relation

$$[\Omega, \mathbb{1}_P \otimes \mathbb{1}_I^{\otimes k} \otimes (U_d^{\otimes k})_O \otimes (U_d)_F] = 0, \quad \forall U_d \in \mathcal{SU}(d), \quad (63)$$

and since $\mathbb{1} \otimes U_d |\mathbb{1}\rangle\langle \mathbb{1}| = U_d^T \otimes \mathbb{1} |\mathbb{1}\rangle\langle \mathbb{1}|$, we also have

$$[\Omega, (U_d)_P \otimes (U_d^{\otimes k})_I \otimes \mathbb{1}_O^{\otimes k} \otimes \mathbb{1}_F] = 0, \quad \forall U_d \in \mathcal{SU}(d). \quad (64)$$

Now, let $\{B^i\}_i$ be an orthonormal basis for the subspace spanned by operators commuting with $U_d^{\otimes(k+1)}$, i.e., $\text{tr}(B^i B^{j\dagger}) = d_i \delta_{ij}$, and $d_i := \text{tr}(B^i B^{i\dagger})$. Eq. (64) ensures that we can write the performance operator as

$$\Omega = \sum_i B_{PI}^i \otimes \Omega_{OF}^i \quad (65)$$

for some set of operators $\Omega_{OF}^i \in \mathcal{L}(\mathcal{H}_O \otimes \mathcal{H}_F)$. As noticed in [28], we can find such operators by using the orthogonality relation, $\text{tr}(B^i B^{j\dagger}) = d_i \delta_{ij}$. More explicitly, by making use of $\text{tr}_2 |\mathbb{1}\rangle\langle \mathbb{1}| = \mathbb{1}_1$ we can write

$$d_i \Omega_{OF}^i = \text{tr}_{PI} \left[((B_{PI}^i)^\dagger \otimes \mathbb{1}_{OF}) \Omega \right] \quad (66)$$

$$= \frac{1}{d^2} \int \text{tr}_{PI} \left[((B_{PI}^i)^\dagger \otimes \mathbb{1}_{OF}) \right] \quad (67)$$

$$(\mathbb{1}_P \otimes \mathbb{1}_I^{\otimes k} \otimes (U_d^{\otimes k})_O \otimes U_{dF}) \quad (68)$$

$$|\mathbb{1}\rangle\langle \mathbb{1}|_{PF} \otimes |\mathbb{1}\rangle\langle \mathbb{1}|_{IO} \left(\mathbb{1}_P \otimes \mathbb{1}_I^{\otimes k} \otimes (U_d^{\otimes k})_O \otimes U_{dF} \right)^{\dagger} dU_d]$$

$$= \frac{1}{d^2} \int \text{tr}_{PI} \left[(\mathbb{1}_P \otimes \mathbb{1}_I^{\otimes k} \otimes (U_d^{\otimes k})_O \otimes U_{dF}) \right] \quad (68)$$

$$(\mathbb{1}_{PI} \otimes \overline{B}_{FO}^i) |\mathbb{1}\rangle\langle \mathbb{1}|_{PF} \otimes |\mathbb{1}\rangle\langle \mathbb{1}|_{IO} \\ \left(\mathbb{1}_P \otimes \mathbb{1}_I^{\otimes k} \otimes (U_d^{\otimes k})_O \otimes U_{dF} \right)^{\dagger} dU_d \quad (69)$$

$$= \frac{1}{d^2} \int \left((U_d^{\otimes k})_O \otimes U_{dF} \right) \overline{B}_{FO}^i \left((U_d^{\otimes k})_O \otimes U_{dF} \right)^{\dagger} dU_d \quad (70)$$

$$= \frac{1}{d^2} \int \overline{B}_{FO}^i \left((U_d^{\otimes k})_O \otimes U_{dF} \right) \left((U_d^{\otimes k})_O \otimes U_{dF} \right)^{\dagger} dU_d \quad (71)$$

$$= \frac{1}{d^2} \int \overline{B}_{FO}^i dU_d \quad (72)$$

$$= \frac{1}{d^2} \overline{B}_{FO}^i. \quad (73)$$

This yields the following form of the performance operator for unitary complex conjugation

$$\Omega = \frac{1}{d^2} \sum_i \frac{B_{PI}^i \otimes \overline{B}_{FO}^i}{d_i}. \quad (74)$$

C. Evaluating the Performance of the Superchannel Presented in Sec. III

In this section, we show that the superchannel presented in Sec. III attains optimal performance. The proof of its optimality is presented in subsection IV-D.

The superchannel $S \in \mathcal{L}(\mathcal{H}_P \otimes \mathcal{H}_I \otimes \mathcal{H}_O \otimes \mathcal{H}_F)$ described in Eq. (28) and Eq. (32) reads as

$$S := E_{PI} \otimes D_{OF} \quad (75)$$

$$E := \frac{d}{\binom{d}{k+1}} A(d, k+1) \quad (76)$$

$$D := \frac{\binom{d}{k}}{\binom{d}{k+1}} A(d, k+1) + \left(\mathbb{1}_d^{\otimes k} - A(d, k) \right) \otimes \sigma, \quad (77)$$

where $A(d, k)$ is the projector onto the antisymmetric space of $\mathbb{C}_d^{\otimes k}$ and σ is an arbitrary quantum state.

Lemma 1: The quantum circuit described in Sec. III transforms k uses of a d -dimensional unitary operation U_d into its complex conjugation \overline{U}_d with an average fidelity respecting $\langle F \rangle \geq \frac{k+1}{d(d-k)}$.

Proof: We prove the theorem by showing that

$$\text{tr}(\Omega S) = \frac{k+1}{d(d-k)} \quad (78)$$

with Ω being the performance operator for the unitary complex conjugation transformation

$$\Omega = \int_{U_d \in \mathcal{SU}(d)} |U_d\rangle\langle U_d|_{PF} \otimes |U_d^{\otimes k}\rangle\langle U_d^{\otimes k}|_{IO} dU_d, \quad (79)$$

and S the superchannel described in Eq. (75).

As shown in Eq. (74) and [28], the performance operator for unitary complex conjugation may be written as

$$\Omega = \sum_i \frac{1}{d^2} \frac{B_{PI}^i \otimes \overline{B}_{FO}^i}{d_i} \quad (80)$$

where $\{B^i\}_i$ is an orthonormal basis, i.e., $\text{tr}(B^i B^{j\dagger}) = \delta_{ij} d_i$, for the subspace spanned by operators commuting

with all unitary operators $U^{\otimes(k+1)}$. Since the projector onto the antisymmetric space $A(d, k+1)$ respects $A(d, k+1)U^{\otimes(k+1)} = U^{\otimes(k+1)}A(d, k+1)$, we can set $A(d, k+1)$ as B^1 , one of the operators of the orthonormal basis $\{B^i\}_i$. We can then write

$$\Omega = \frac{1}{d^2} \frac{A(d, k+1)_{PI} \otimes A(d, k+1)_{FO}}{\text{tr}(A(d, k+1))} + \frac{1}{d^2} \sum_{i>1} \frac{B^i_{PI} \otimes \overline{B^i_{FO}}}{d_i}. \quad (81)$$

Now, since $E_{PI} = \frac{d}{\binom{d}{k+1}} A(d, k+1)$ and $\text{tr}(A(d, k+1) B^i) = 0$ for $i > 1$, it holds that

$$\text{tr}(\Omega S) = \frac{1}{d^2} \frac{\text{tr}(E A(d, k+1)) \text{tr}(D A(d, k+1))}{\text{tr}(A(d, k+1))}. \quad (82)$$

We now recall from Sec. III that $\text{tr}(A(d, k)) = \binom{d}{k}$, and since $A(d, k)^2 = A(d, k)$, we have that

$$\begin{aligned} \text{tr}(\Omega S) &= \frac{1}{d^2} \frac{d \binom{d}{k} \binom{d}{k+1} \binom{d}{k+1}}{\binom{d}{k+1} \binom{d}{k+1} \binom{d}{k+1}} \\ &+ \frac{1}{d^2} \frac{d \binom{d}{k+1}}{\binom{d}{k+1} \binom{d}{k+1}} \text{tr}(A(d, k+1) (\mathbb{1}_d^{\otimes k} - A(d, k)) \otimes \sigma). \end{aligned} \quad (83)$$

Since $\mathbb{1}_d^{\otimes k} - A(d, k)$ is positive semidefinite, we have that

$$\text{tr}(\Omega S) \geq \frac{1}{d^2} \frac{d \binom{d}{k} \binom{d}{k+1} \binom{d}{k+1}}{\binom{d}{k+1} \binom{d}{k+1} \binom{d}{k+1}} \quad (84)$$

$$= \frac{\binom{d}{k}}{d \binom{d}{k+1}} \quad (85)$$

$$= \frac{d!}{k!(d-k)(d-k-1)!} \frac{(k+1)k!(d-k-1)!}{d d!} \quad (86)$$

$$= \frac{k+1}{d(d-k)}. \quad (87)$$

This finishes the proof. \square

As we are going to show next, this protocol cannot be improved. Its average fidelity is exactly $\langle F \rangle = \frac{k+1}{d(d-k)}$ and we have that $\text{tr}(A(d, k+1) (\mathbb{1}_d^{\otimes k} - A(d, k)) \otimes \sigma) = 0$ for any quantum state σ .

D. A Tight Upper Bound for the Unitary Conjugation Problem

To prove that the superchannel S presented in section III is indeed optimal, we provide here a feasible solution for the dual problem defined by equations (54)-(58), and show that its solution matches with the solution of the respective primal problem from (50)-(53) solved in section IV-C. In particular, we show that for our proposed ansatz, the value of the constant c from (54) matches with the value from (50) – reproducing the value of optimal fidelity from Theorem 1. In order to proceed, we need to introduce basic knowledge from representation theory, including the famous Schur-Weyl duality [34], [35], Young diagrams and their relation to irreducible representation (irrep) of the symmetric group [34], [35]. We give a brief introduction in the following subsection, while the solution of the dual problem for the unitary conjugation problem is derived later.

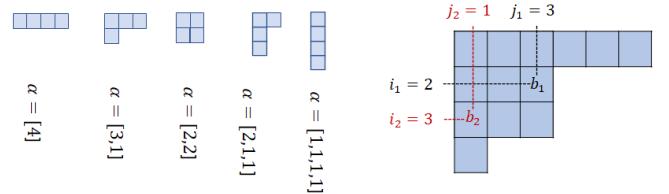


Fig. 2. The left figure presents five possible Young diagrams for $n = 4$. These Young diagrams corresponds to all possible abstract irreducible representations of $S(4)$. Considering representation space $(\mathbb{C}^d)^{\otimes 4}$, there appear only irreps for which height of the corresponding Young diagram is no larger than d . For example, considering qubits ($d = 2$) we have only three Young diagrams: (4), (3, 1), (2, 2). In particular no antisymmetric space exists. The right figure presents graphical example of assigning the coordinates to a box within a Young diagram. Here we have $\mu = [6, 3, 3, 1]$ and the box b_1 has the coordinates $(i_1, j_1) = (2, 3)$, while the box b_2 has the coordinates $(i_2, j_2) = (3, 1)$. The axial distance between these to boxes, according to (91) is equal $d(b_1, b_2) = |i_1 - i_2| + |j_1 - j_2| = 3$.

1) Young Diagrams, Schur-Weyl Duality and Irreducible Representations of the Symmetric Group: A partition λ of a natural number n , which we denote as $\lambda \vdash n$, is a sequence of positive numbers $\lambda = [\lambda_1, \lambda_2, \dots, \lambda_m]$, such that

$$\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_m \geq 0 \quad \sum_{i=1}^m \lambda_i = n. \quad (88)$$

Every partition can be visualised as a *Young diagram* - a collection of boxes arranged in left-justified rows. With every Young diagram we associate its height $H(\lambda)$, which is defined as a number of boxes in the first (the longest one) column. See Figure 2 for illustration of the above definitions.

Now we introduce the notion of a step representation for a given Young diagram [36]. Let k_1, k_2, \dots, k_s be multiplicities of numbers λ_i in the non-increasing sequence (88). Let us also define $p_i := \lambda_{k_i} - \lambda_{k_{i+1}}$ for $i = 1, 2, \dots, s$, where we assume $\lambda_{k_{s+1}} = 0$. Observe that

$$\sum_{1 \leq i \leq j \leq s} k_i p_j = n, \quad (89)$$

and a Young diagram is uniquely determined by the sequence,

$$(k_1, p_1, k_2, p_2, \dots, k_s, p_s). \quad (90)$$

Let us name the sequence a *step representation* of a diagram λ , see Figure 3.

For every two boxes $b_1 = (i_1, j_1)$, $b_2 = (i_2, j_2)$ let us define their *axial distance* (see Figure 2) as:

$$d(b_1, b_2) := |i_1 - i_2| + |j_1 - j_2|. \quad (91)$$

Let c_1, c_2, \dots, c_s be corner boxes

$$c_j = (k_1 + k_2 + \dots + k_j, p_j + p_{j+1} + \dots + p_s), \quad j = 1, 2, \dots, s. \quad (92)$$

For further applications let us observe that for $i < j$,

$$d(c_i, c_j) = p_i + \dots + k_j, \quad (93)$$

where the above notation means the sum of all numbers occurring in the sequence (90) between p_i and k_j including them.

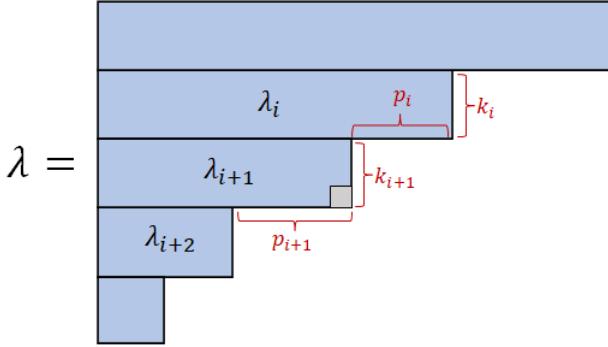


Fig. 3. The graphic presents parameters describing a step representation for an arbitrary Young diagram λ . The numbers k_i, k_{i+1} count the number of appearances of row of lengths λ_i, λ_{i+1} respectively, while the numbers p_i, p_{i+1} are the respective differences $\lambda_{i+1} - \lambda_i, \lambda_{i+1} - \lambda_{i+2}$ between the rows lengths. By a grey square we depicted one corner box c_{i+1} whose coordinates can be calculated through expression (92).

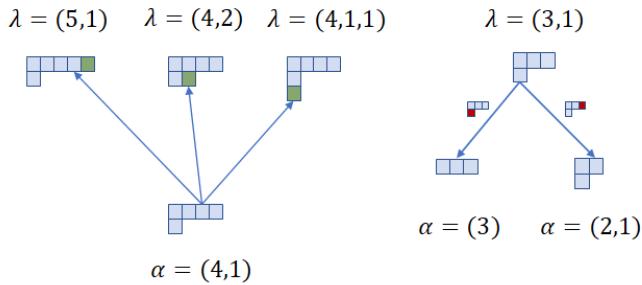


Fig. 4. The left graphic presents possible Young diagrams $\lambda \vdash 6$, which can be obtained from Young diagram $\alpha = [4, 1]$ by adding a single box, depicted here by the green colour. In this particular case, by writing $\lambda \ni \alpha$, we take λ represented only by these three diagrams. In the same manner we define subtracting of a box from a Young diagram. This is depicted on the right graphic where we remove a single box (red colour) from the diagram $\lambda = [3, 1]$ and obtain two three-boxed Young diagrams $\alpha = [3]$ and $\alpha = [2, 1]$. Diagrams α obtained from a diagram λ by removing a single box we denote as $\alpha \in \lambda$.

For a fixed number n , the number of Young diagrams determines the number of nonequivalent irreps of S_n in an abstract decomposition. However, working in the representation space $\mathcal{H} \equiv (\mathbb{C}^d)^{\otimes n}$, for every decomposition of S_n into irreps we take Young diagrams λ with height $H(\lambda)$ of a given λ of at most d . For example, in the case of qubits ($d = 2$), we allow only for diagrams up to two rows.

Suppose we have $\alpha \vdash n - 1$ and $\lambda \vdash n$. Writing $\lambda \ni \alpha$ we consider such Young diagrams λ which can be obtained from α by adding a single box (green colour). Similarly, writing $\alpha \in \lambda$ we consider such Young diagrams α , which can be obtained from λ by removing a single box. The procedure of adding/removing a box from Young diagrams is illustrated in Figure 4.

Having the above definitions and notations we are in position to summarise more results on representation theory used further in this manuscript. We start from the celebrated Schur-Weyl duality relating irreps of the general linear group $GL(d)$ and symmetric group S_n . Namely, it is known that the diagonal action of the general linear group $GL(d)$ of invertible complex matrices and of the symmetric group on $(\mathbb{C}^d)^{\otimes n}$

commute:

$$V(\pi)(X \otimes \cdots \otimes X) = (X \otimes \cdots \otimes X)V(\pi), \quad (94)$$

where $\pi \in S_n$ and $X \in GL(d)$. Due to the above relation, there exists a basis called the Schur basis, producing a decomposition into irreps of $GL(d)$ and S_n simultaneously. Hence, for the whole space $(\mathbb{C}^d)^{\otimes n}$ we have the following

Theorem 2 (Schur-Weyl Duality): The tensor product space $(\mathbb{C}^d)^{\otimes n}$ can be decomposed as

$$(\mathbb{C}^d)^{\otimes n} = \bigoplus_{\substack{\lambda \vdash n \\ H(\lambda) \leq d}} \mathcal{U}_\lambda \otimes \mathcal{S}_\lambda, \quad (95)$$

where the symmetric group S_n acts on the space \mathcal{S}_λ and the general linear group $GL(d)$ acts on the space \mathcal{U}_λ , labelled by the same Young diagram $\lambda \vdash n$.

In fact, for our purposes we need the Schur-Weyl duality not for $GL(d)$ but rather its subgroup $SU(d)$, see [28]. However, it turns out that any irreducible representation of $GL(d)$ on $(\mathbb{C}^d)^{\otimes n}$ is an irreducible representation of $SU(d)$ on $(\mathbb{C}^d)^{\otimes n}$ and vice versa. In particular, it means that that also the linear invariants of $SU(d)$ are those of $GL(d)$, and they belong to the group algebra of S_n denoted by $\mathbb{C}[S_n]$. This however means that we can write the same statement in Theorem 2 for group $SU(d)$ [37]. From the decomposition given in Theorem 2 we deduce that for a given irrep λ of S_n , the space \mathcal{U}_λ is the multiplicity space of dimension m_λ (multiplicity of irrep λ), while the space \mathcal{S}_λ is the representation space of dimension d_λ (dimension of irrep λ). Every operator commuting with the diagonal action of $U^{\otimes n}$, due to the Schur's lemma must be proportional to identity on spaces \mathcal{U}_λ and admits non-trivial parts on spaces \mathcal{S}_λ . Conversely, any operator commuting with action of the permutation group S_n is non-trivially supported on irreducible spaces \mathcal{U}_λ . In the former case, i.e. in very irreducible space \mathcal{S}_λ one can construct an orthonormal irreducible basis $\{|\lambda, i\rangle\}$, where $i = 1, \dots, d_\lambda$, for example by exploiting Young-Yamanouchi construction [38]. With this vectors we associate an irreducible basis operators E_{ij}^λ , for $i, j = 1, \dots, d_\lambda$ admitting the following form:

$$E_{ij}^\lambda := \mathbb{1}^{\mathcal{U}_\lambda} \otimes |\lambda, i\rangle \langle \lambda, j|_{\mathcal{S}_\lambda}. \quad (96)$$

Sine the basis $\{|\lambda, i\rangle\}_{i=1}^{d_\lambda}$ is orthogonal in both indices the above operators fulfil the following properties, which are

$$E_{ij}^\lambda E_{kl}^{\lambda'} = \delta^{\lambda\lambda'} \delta_{jk} E_{il}^\lambda, \quad \text{tr } E_{ij}^\lambda = \delta_{ij} m_\lambda. \quad (97)$$

From the operators E_{ij}^λ , by using the relation $\sum_i |\lambda, i\rangle \langle \lambda, i|_{\mathcal{S}_\lambda} = \mathbb{1}_\lambda^{\mathcal{S}}$, we can construct so-called Young projectors P^λ [38] on the irreducible components λ :

$$P^\lambda := \sum_{i=1}^{d_\lambda} E_{ii}^\lambda = \mathbb{1}_\lambda^{\mathcal{U}} \otimes \mathbb{1}_\lambda^{\mathcal{S}}, \quad (98)$$

$$P^\lambda P^{\lambda'} = \delta^{\lambda\lambda'} P^\lambda, \quad (99)$$

$$\text{tr } P^\lambda = m_\lambda d_\lambda. \quad (100)$$

Let us take $\lambda \vdash n$ and $\alpha, \beta \vdash n - 1$, such that $\alpha, \beta \in \lambda$. Consider an arbitrary unitary irrep $\phi^\lambda(\pi)$ for some $\pi \in S_n$.

Thus it can be always unitarily transformed (reduced) to block form consisting of elements of the subgroup $S_{n-1} \subset S_n$ as

$$\phi^\lambda(\pi) = (\phi_{i_\lambda j_\lambda}^\lambda(\pi)) = (\phi_{i_\alpha j_\beta}^{\alpha\beta}(\pi)), \quad (101)$$

where the indices i_λ, j_λ are the standard matrix indices running from 1 to d_λ . From the above we see the following connection between the indices, namely for fixed i_λ, j_λ one has $i_\lambda = (\alpha \in \lambda, i_\alpha), j_\lambda = (\beta \in \lambda, j_\beta)$, where the indices i_α, j_β run from 1 to d_α, d_β respectively. Please see Figure 5 where the graphical illustration of the above considerations is presented or see [39] for more details. The diagonal blocks of the matrix $\phi^\lambda(\pi)$, i.e. the blocks with $\alpha = \beta$, are of dimension of the corresponding irrep φ^α of S_{n-1} . It means that the both indices i_α, j_α run from 1 to d_α . In particular, when one considers an irrep $\phi^\lambda(\pi')$ for $\pi' \in S_{n-1}$, we have

$$\phi^\lambda(\pi') = (\delta^{\alpha\beta} \varphi_{i_\alpha j_\alpha}^{\alpha\alpha}(\pi')) = \bigoplus_{\alpha \in \lambda} \varphi^\alpha(\pi'). \quad (102)$$

A similar procedure (induction) can be also done in the reverse order, i.e. starting from an irrep φ^α of S_{n-1} we can induce irreps ϕ^λ of a larger group S_n . Furthermore, utilizing the relabelling $i_\lambda = (\alpha \in \lambda, i_\alpha), j_\lambda = (\beta \in \lambda, j_\beta)$ introduced above leads us to the following notation for the operators E_{ij}^λ from (96):

Notation 1: Every operator E_{ij}^λ , where E_{ij}^λ are basis operators of the algebra $\mathbb{C}[S_n]$, can be re-written as

$$E_{i_\lambda j_\lambda}^\lambda := E_{i_\alpha j_{\alpha'}}^{\alpha\alpha'}(\lambda), \quad (103)$$

where $\alpha, \alpha' \in \lambda$.

The above considerations give us a way to represents elements of the algebra $\mathbb{C}[S_{n-1}]$ in terms of elements from the algebra $\mathbb{C}[S_n]$. Concretely, one can formulate the following:

Lemma 2: Operator $E_{ij}^\lambda \otimes \mathbb{1}$, where E_{ij}^λ are basis operators of the algebra $\mathbb{C}[S_{n-1}]$, can be written in terms of basis elements of the algebra $\mathbb{C}[S_n]$ as

$$E_{i_\alpha j_\alpha}^\alpha \otimes \mathbb{1} = \sum_{\lambda \ni \alpha} E_{i_\alpha j_\alpha}^{\alpha\alpha}(\lambda). \quad (104)$$

Notice that the resulting irreducible basis operator $E_{i_\alpha j_\alpha}^{\alpha\alpha}(\lambda)$ from $\mathbb{C}[S_n]$ is block-diagonal in $\alpha \in \lambda$.

Proof: The operator $E_{i_\alpha j_\alpha}^\alpha \otimes \mathbb{1}$ clearly belongs to the algebra $\mathbb{C}[S_n]$, so in general it can be decomposed in terms of operators $E_{i_\lambda j_\lambda}^\lambda := E_{i_\beta j_{\beta'}}^{\beta\beta'}(\lambda)$ as

$$E_{i_\alpha j_\alpha}^\alpha \otimes \mathbb{1} = \sum_{\lambda} \sum_{\beta, \beta' \in \lambda} \sum_{k_\beta, l_{\beta'}} x_{k_\beta l_{\beta'}}^{\beta\beta'}(\lambda) E_{k_\beta l_{\beta'}}^{\beta\beta'}(\lambda), \quad (105)$$

where the numbers $x_{k_\beta l_{\beta'}}^{\beta\beta'}(\lambda)$ are unknown coefficients. To determine the coefficients, let us compute overlap of the left-hand side of (105) with some basis operator $(E_{k_\gamma l_{\gamma'}}^{\gamma\gamma'}(\lambda))^\dagger = E_{l_{\gamma'} k_\gamma}^{\gamma'\gamma}(\lambda)$:

$$\text{tr} [(E_{i_\alpha j_\alpha}^\alpha \otimes \mathbb{1}) E_{l_{\gamma'} k_\gamma}^{\gamma'\gamma}(\lambda)] = \text{tr} [E_{i_\alpha j_\alpha}^\alpha \text{tr}_n (E_{l_{\gamma'} k_\gamma}^{\gamma'\gamma}(\lambda))] \quad (106)$$

$$= \frac{m_\lambda}{m_\gamma} \text{tr} (E_{i_\alpha j_\alpha}^\alpha E_{l_{\gamma'} k_\gamma}^{\gamma'\gamma}) \delta_{\gamma\gamma'} \quad (107)$$

$$= \frac{m_\lambda m_\alpha}{m_\gamma} \delta_{\gamma\gamma'} \delta_{\alpha\gamma} \delta_{j_\alpha l_\gamma} \delta_{i_\alpha k_\gamma}. \quad (108)$$

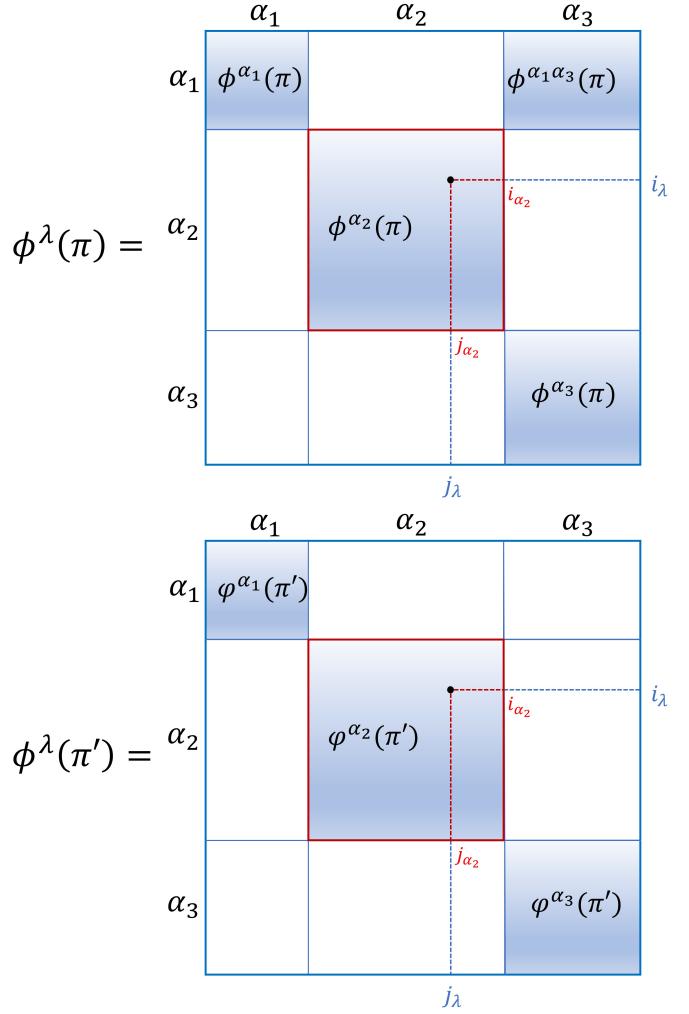


Fig. 5. The left-hand side graphic represents unitary transformed irrep $\phi^\lambda(\pi) = (\phi_{i_\lambda j_\lambda}^\lambda(\pi))$ for some $\pi \in S_n$ and $\lambda \vdash n$ into a block form related to subgroup S_{n-1} . Each block $(\phi_{i_\alpha j_\alpha}^{\alpha\alpha}(\pi))$ is indexed by a pair of indices from $\{\alpha_1, \alpha_2, \alpha_3\}$ which are obtained from λ by removing a single box, so by irreps of S_{n-1} . If we take $\pi' \in S_{n-1}$ then $\phi^\lambda(\pi')$ takes a block diagonal form, i.e. it can be represented as $\varphi^{\alpha_1}(\pi') \oplus \varphi^{\alpha_2}(\pi') \oplus \varphi^{\alpha_3}(\pi')$ - the right-hand side graphic. Moreover, the chain $\dots S_{n-1} \subset S_n \subset \dots$ is multiplicity free, which means that every $\alpha \vdash (n-1)$, for $\alpha \in \lambda$, occurs exactly once in an irrep $\lambda \vdash n$. There is also a one-to-one correspondence between matrix indices of $\phi^\lambda(\pi)$ and every sub-block $(\phi_{i_\alpha j_\alpha}^{\alpha\alpha}(\pi))$. Indeed, every element (black dot) can be indexed by a pair (i_λ, j_λ) , where $1 \leq i_\lambda, j_\lambda \leq d_\lambda$ or by a pair of indices $(i_{\alpha'}, j_{\alpha'})$ from a sub-block $\phi^{\alpha'}(\pi)$, where $1 \leq i_{\alpha'}, j_{\alpha'} \leq d_{\alpha'}$, and we suppressed the double index. The same reasoning works also for off-diagonal blocks. Since the construction is multiplicity free, we have the following relations between dimensions of the irreps: $d_\lambda = d_{\alpha_1} + d_{\alpha_2} + d_{\alpha_3}$.

To get the third equality we use Lemma 7 from [39], to get the fourth equality we directly apply relations from (97). Now, let us apply the same procedure for the right-hand side of (105) but for some $\lambda' \neq \lambda$. In the same lines we get the following:

$$\sum_{\lambda} \sum_{\beta, \beta' \in \lambda} \sum_{k_\beta, l_{\beta'}} x_{k_\beta l_{\beta'}}^{\beta\beta'}(\lambda) \text{tr} (E_{k_\beta l_{\beta'}}^{\beta\beta'}(\lambda) E_{l_{\gamma'} k_\gamma}^{\gamma'\gamma}(\lambda')) \quad (109)$$

$$= \sum_{\lambda} \sum_{\beta, \beta' \in \lambda} \sum_{k_\beta, l_{\beta'}} m_\lambda x_{k_\beta l_{\beta'}}^{\beta\beta'}(\lambda) \delta_{\lambda\lambda'} \delta_{\beta'\beta'} \delta_{l_{\beta'} l_{\gamma'}} \delta_{\beta\gamma} \delta_{k_\beta k_\gamma} \quad (110)$$

$$= m_\lambda x_{k_\gamma l_{\gamma'}}^{\gamma\gamma'}(\lambda). \quad (111)$$

The final expressions in (106) and (109) must be equal giving us the following condition:

$$x_{k_\gamma l_{\gamma'}}^{\gamma\gamma'}(\lambda) = \frac{m_\alpha}{m_\gamma} \delta_{\gamma\gamma'} \delta_{\alpha\gamma} \delta_{j_\alpha l_{\gamma'}} \delta_{i_\alpha k_\gamma}. \quad (112)$$

Inserting the right hand side of (112) into decomposition (105) we have:

$$E_{i_\alpha j_\alpha}^\alpha \otimes \mathbb{1} = \sum_{\lambda} \sum_{\gamma, \gamma' \in \lambda} \sum_{k_\gamma, l_{\gamma'}} \frac{m_\alpha}{m_\gamma} \delta_{\gamma\gamma'} \delta_{\alpha\gamma} \delta_{j_\alpha l_{\gamma'}} \delta_{i_\alpha k_\gamma} E_{k_\gamma l_{\gamma'}}^{\gamma\gamma'}(\lambda) \quad (113)$$

$$= \sum_{\lambda \ni \alpha} E_{i_\alpha j_\alpha}^{\alpha\alpha}(\lambda). \quad (114)$$

The final sum runs only over λ which can be obtained from α by adding a single box. This finishes the proof. \square

2) *Direct Derivations of the Upper Bound for the Unitary Conjugation Problem:* Due to the presentation of the performance operator Ω in (59), we see that it commutes with the diagonal action of the unitary group. This means that its nontrivial part, due to the Schur-Weyl duality in Theorem 2, is supported on the representation space \mathcal{S}_λ connected to the symmetric group. Therefore, the operators from (74) can be chosen to be operators E_{ij}^λ defined in (96). Then the operator Ω from (74) has the following presentation:

$$\Omega = \frac{1}{d^2} \sum_{\substack{\lambda \vdash (k+1) \\ H(\lambda) \leq d}} \sum_{p,q=1}^{d_\lambda} \frac{(E_{pq}^\lambda)_{IP} \otimes (E_{pq}^\lambda)_{OF}}{m_\lambda}. \quad (115)$$

Since the operator Ω is non-trivial only on the space \mathcal{S}_λ , we assume the same property for the operator \hat{S} from (55) given in formulation of the dual problem described in section IV-B. For the time being, except additional symmetries imposed on the operator \hat{S} we see that it can have an arbitrary form. However, the considered operator \hat{S} must also satisfy constraints from (55)-(58), which obviously limits the compatible choices for \hat{S} . Since our goal is to find a tight upper bound which matches with the solution of the primal problem from section IV-C, it is sufficient to present a feasible solution satisfying aforementioned conditions. To this purpose, we choose the operator \hat{S} to be of the following form:

$$\hat{S} = \sum_{\substack{\alpha \vdash k \\ H(\alpha) \leq d}} \sum_{k,l=1}^{d_\alpha} \frac{d}{m_\alpha} (E_{kl}^\alpha)_I \otimes \tau_P \otimes (E_{kl}^\alpha)_O \otimes \tau_F. \quad (116)$$

In the above, the operators $\tau_P = \tau_F = \frac{1}{d}$ are maximally mixed states on the respective systems. It is easy to see that the operator \hat{S} satisfies the constraints (56)-(58) by construction. Indeed, we have the following:

$$W_{PIO} = \sum_{\substack{\alpha \vdash k \\ H(\alpha) \leq d}} \sum_{k,l=1}^{d_\alpha} \frac{1}{m_\alpha} (E_{kl}^\alpha)_I \otimes \tau_P \otimes (E_{kl}^\alpha)_O. \quad (117)$$

Now, exploiting properties of irreducible operator basis $\{E_{ij}^\alpha\}$ from expression (97), we write:

$$\text{tr}_O(W_{PIO}) = \sum_{\substack{\alpha \vdash k \\ H(\alpha) \leq d}} \sum_{k,l=1}^{d_\alpha} \frac{\text{tr}[(E_{kl}^\alpha)_O]}{m_\alpha} (E_{kl}^\alpha)_I \otimes \tau_P \quad (118)$$

$$= \sum_{\substack{\alpha \vdash k \\ H(\alpha) \leq d}} \sum_{k=1}^{d_\alpha} (E_{kk}^\alpha)_I \otimes \tau_P \quad (119)$$

$$= \sum_{\substack{\alpha \vdash k \\ H(\alpha) \leq d}} (P^\alpha)_I \otimes \tau_P \quad (120)$$

$$= \mathbb{1}_I \otimes \tau_P, \quad (121)$$

and finally $\text{tr}(\tau_P) = 1$. Having the above, we are in position to evaluate the minimal value of the constant c satisfying constraint (55).

Lemma 3: The minimal value of the constant c satisfying constraint (55) for the operator \hat{S} from (116) is given by,

$$c = \frac{1}{d} \max_{\lambda \vdash (k+1)} \frac{\sum_{\beta \in \lambda} m_\beta}{m_\lambda}, \quad (122)$$

where m_λ, m_β are multiplicities of irreps $\lambda \vdash k+1, \beta \vdash k$ in the Schur-Weyl duality. The symbol $\beta \in \lambda$ denotes all Young diagrams obtained from a Young diagram λ by removing a single box.

Proof:

Writing the constraint from (55) explicitly by using expressions (115) and (116) we get the following inequality

$$\begin{aligned} & \frac{1}{d^2} \sum_{\substack{\lambda \vdash (k+1) \\ H(\lambda) \leq d}} \sum_{p,q=1}^{d_\lambda} \frac{(E_{pq}^\lambda)_{IP} \otimes (E_{pq}^\lambda)_{OF}}{m_\lambda} \\ & \leq c \cdot \sum_{\substack{\alpha \vdash (k) \\ H(\alpha) \leq d}} \sum_{k,l=1}^{d_\alpha} \frac{d}{m_\alpha} (E_{kl}^\alpha)_I \otimes \tau_P \otimes (E_{kl}^\alpha)_O \otimes \tau_F. \end{aligned} \quad (123)$$

To keep the notation simple, we drop indices the I, O, P, F , since it will be clear from the context which systems are referred to. First, we notice that the right-hand side of (123) can be re-written by applying Lemma 2 with $\tau = \mathbb{1}/d$, leading to

$$\begin{aligned} & \sum_{\alpha} \sum_{k,l=1}^{d_\alpha} \frac{d}{m_\alpha} E_{kl}^\alpha \otimes \tau \otimes E_{kl}^\alpha \otimes \tau \\ & = \sum_{\alpha} \sum_{k,l=1}^{d_\alpha} \frac{1}{dm_\alpha} E_{kl}^\alpha \otimes \mathbb{1} \otimes E_{kl}^\alpha \otimes \mathbb{1} \\ & = \sum_{\alpha} \sum_{k_\alpha, l_\alpha=1}^{d_\alpha} \sum_{\lambda, \lambda' \ni \alpha} \frac{1}{dm_\alpha} E_{k_\alpha l_\alpha}^{\alpha\alpha}(\lambda) \otimes E_{k_\alpha l_\alpha}^{\alpha\alpha}(\lambda') \\ & = \sum_{\lambda, \lambda'} \sum_{\alpha \in \lambda \wedge \lambda'} \sum_{k_\alpha, l_\alpha=1}^{d_\alpha} \frac{1}{dm_\alpha} E_{k_\alpha l_\alpha}^{\alpha\alpha}(\lambda) \otimes E_{k_\alpha l_\alpha}^{\alpha\alpha}(\lambda'). \end{aligned} \quad (124)$$

Here, we introduced symbol $\alpha \in \lambda \wedge \lambda'$, indicating that simultaneously $\alpha \in \lambda$ and $\alpha \in \lambda'$. Next, by Notation 1 every

irreducible operator basis labelled by $\lambda \vdash k+1$ on the left-hand side of equation (123) can be written as

$$E_{pq}^\lambda = E_{p_\beta q_{\beta'}}^{\beta\beta'}(\lambda). \quad (125)$$

The above expression and final result of (124) allows us to re-write inequality (123) in the following form:

$$\begin{aligned} & \frac{1}{d^2} \sum_{\lambda} \sum_{\beta, \beta' \in \lambda} \sum_{p_\beta, q_{\beta'}} \frac{E_{p_\beta q_{\beta'}}^{\beta\beta'}(\lambda) \otimes E_{p_\beta q_{\beta'}}^{\beta\beta'}(\lambda)}{m_\lambda} \\ & \leq c \cdot \sum_{\lambda, \lambda'} \sum_{\alpha \in \lambda \wedge \lambda'} \sum_{k_\alpha, l_\alpha=1}^{d_\alpha} \frac{1}{dm_\alpha} E_{k_\alpha l_\alpha}^{\alpha\alpha}(\lambda) \otimes E_{k_\alpha l_\alpha}^{\alpha\alpha}(\lambda'). \end{aligned} \quad (126)$$

Now, notice that every $\lambda, \lambda' \vdash (k+1)$ label a different nonequivalent irrep. This means that we deal with different orthogonal blocks. From the above inequality it is clear that non-trivial solutions for the constant c can be obtained only for $\lambda = \lambda'$ – for $\lambda \neq \lambda'$ we do not have cross terms on the left-hand side of (126). Consequently, it is sufficient to restrict ourselves to comparing terms for every λ separately:

$$\begin{aligned} & \frac{1}{d^2} \sum_{\beta, \beta' \in \lambda} \sum_{p_\beta, q_{\beta'}} \frac{E_{p_\beta q_{\beta'}}^{\beta\beta'}(\lambda) \otimes E_{p_\beta q_{\beta'}}^{\beta\beta'}(\lambda)}{m_\lambda} \\ & \leq c \cdot \sum_{\alpha \in \lambda} \sum_{i_\alpha, j_\alpha} \frac{1}{dm_\alpha} E_{i_\alpha j_\alpha}^{\alpha\alpha}(\lambda) \otimes E_{i_\alpha j_\alpha}^{\alpha\alpha}(\lambda). \end{aligned} \quad (127)$$

All the operators in the above expression are the irreducible basis operators for the symmetric groups S_{k+1} and S_k . Hence, by Schur-Weyl duality, they act non-trivially only on the symmetric component of the Schur-Weyl decomposition (95). Now, using the explicit forms of the irreducible basis operators in the Schur basis from (96) we rewrite (127) in the following way (ignoring the multiplicity space):

$$\begin{aligned} & \frac{1}{d^2 m_\lambda} \sum_{\beta, \beta' \in \lambda} \sum_{p_\beta, q_{\beta'}} |\beta, p_\beta\rangle \langle \beta', q_{\beta'}| \otimes |\beta, p_\beta\rangle \langle \beta', q_{\beta'}| \\ & \leq \sum_{\alpha \in \lambda} \sum_{i_\alpha, j_\alpha} \frac{c}{dm_\alpha} |\alpha, i_\alpha\rangle \langle \alpha, j_\alpha| \otimes |\alpha, i_\alpha\rangle \langle \alpha, j_\alpha|. \end{aligned} \quad (128)$$

Using the definition $|\Phi_\alpha^+\rangle = \sum_{i_\alpha=1}^{d_\alpha} |i_\alpha\rangle |i_\alpha\rangle$ for an unnormalised maximally entangled state on irrep α and rearranging the irreps, we get

$$\begin{aligned} & \frac{1}{d^2 m_\lambda} \sum_{\beta, \beta' \in \lambda} |\beta\beta'\rangle \langle \beta'\beta'| \otimes |\Phi_\beta^+\rangle \langle \Phi_{\beta'}^+| \\ & \leq \sum_{\alpha \in \lambda} \frac{c}{dm_\alpha} |\alpha\alpha\rangle \langle \alpha\alpha| \otimes |\Phi_\alpha^+\rangle \langle \Phi_\alpha^+|. \end{aligned} \quad (129)$$

Now, we notice that by applying an isometry $|\beta\beta'\rangle \mapsto |\beta\rangle$ we can rewrite the above inequality as:

$$\frac{1}{d^2 m_\lambda} \sum_{\beta, \beta' \in \lambda} |\beta\rangle \langle \beta'| \otimes |\Phi_\beta^+\rangle \langle \Phi_{\beta'}^+| \quad (130)$$

$$\leq \sum_{\beta \in \lambda} \frac{c}{dm_\beta} |\beta\rangle \langle \beta| \otimes |\Phi_\beta^+\rangle \langle \Phi_\beta^+|. \quad (131)$$

Defining a new set of normalised vectors $|\psi_\beta\rangle := \frac{|\tilde{\psi}_\beta\rangle}{\|\tilde{\psi}_\beta\|} = \frac{1}{\sqrt{d_\beta}} |\tilde{\psi}_\beta\rangle$, where $|\tilde{\psi}_\beta\rangle := |\beta\rangle \otimes |\Phi_\beta^+\rangle$, we rewrite the above equation as

$$\frac{1}{d^2 m_\lambda} \sum_{\beta, \beta'} \sqrt{d_\beta d_{\beta'}} |\psi_\beta\rangle \langle \psi_{\beta'}| \leq \frac{c}{d} \sum_{\beta \in \lambda} \frac{d_\beta}{m_\beta} |\psi_\beta\rangle \langle \psi_\beta|. \quad (132)$$

Notice that the matrix $A := \sum_{\beta \in \lambda} \frac{d_\beta}{m_\beta} |\psi_\beta\rangle \langle \psi_\beta|$ is diagonal and of the full rank. These properties allow us to compute easily $A^{-1/2}$, which is $A^{-1/2} = \sum_{\beta \in \lambda} \sqrt{\frac{m_\beta}{d_\beta}} |\psi_\beta\rangle \langle \psi_\beta|$, since blocks labelled by different β are orthogonal. Sandwiching inequality (132) by $A^{-1/2}$ and using fact that $\langle \psi_\beta | \psi_{\beta'} \rangle = \delta_{\beta\beta'}$, gives the following after a few simple steps

$$\frac{1}{dm_\lambda} \sum_{\beta, \beta' \in \lambda} \sqrt{m_\beta m_{\beta'}} |\psi_\beta\rangle \langle \psi_{\beta'}| \leq c \cdot \mathbb{1}, \quad (133)$$

$$\frac{1}{dm_\lambda} |\psi\rangle \langle \psi| \leq c \cdot \mathbb{1}, \quad (134)$$

where we defined $|\psi\rangle := \sum_{\beta \in \lambda} \sqrt{m_\beta} |\psi_\beta\rangle$, with $\|\psi\|^2 = \sum_{\beta \in \lambda} m_\beta$. Expression (134) implies that

$$\forall \lambda \vdash k+1 \quad \frac{1}{d} \frac{\sum_{\beta \in \lambda} m_\beta}{m_\lambda} \leq c. \quad (135)$$

To satisfy the above inequality, one has to choose the constant c at least equal to the maximal possible value of the ratio $(1/d) \sum_{\beta \in \lambda} m_\beta / m_\lambda$ calculated for all possible λ 's. This finally leads to

$$c = \frac{1}{d} \max_{\lambda \vdash (k+1)} \frac{\sum_{\beta \in \lambda} m_\beta}{m_\lambda}. \quad (136)$$

□

Lemma 3 connects the minimal value of the constant c from (55) with the specific choice of the operator \hat{S} in (116) and group theoretical parameters describing irreps of the symmetric groups S_k and S_{k+1} . However, the final expression (122) still involves a maximisation problem. To solve this maximisation explicitly, let us define an auxiliary function of Young diagrams of n boxes:

$$c(\lambda) := \frac{\sum_{\beta \in \lambda} m_\beta}{m_\lambda}. \quad (137)$$

For the above defined function we can formulate the following:

Proposition 1: For a Young diagram $\lambda \vdash n$ with its step representation $(k_1, p_1, k_2, p_2, \dots, k_s, p_s)$, the function from (137) is given by

$$\begin{aligned} c(\lambda) = \sum_{j=1}^s \frac{k_j p_j}{d - \sum_{i=1}^j k_i + \sum_{i=j}^s p_i} \prod_{i=1}^{j-1} \left(1 + \frac{k_i}{d(c_i, c_j)}\right) \\ \prod_{i=j+1}^s \left(1 + \frac{p_i}{d(c_j, c_i)}\right), \end{aligned} \quad (138)$$

where we assume that the product is equal to 1 when the lower range of the products is larger than the upper one.

Proof: Let $(k_1, p_1, \dots, k_s, p_s)$ be a step representation of $\lambda \vdash n$. Note that removing a one box from the Young diagram λ produces a new Young diagram if and only if the box removed is a corner box. Diagrams $\beta \in \lambda$ can therefore

be labelled by $\mu = 1, 2, \dots, s$. We use the convention that β_μ denotes the diagram obtained by removing the box c_μ . According to formula (137)

$$c(\lambda) = \sum_{\mu=1}^s \frac{m_{\beta_\mu}}{m_\lambda} = \sum_{\mu=1}^s \frac{\prod_{b \in \beta_\mu} \frac{d - i_b + j_b}{h_\mu(b)}}{\prod_{b \in \lambda} \frac{d - i_b + j_b}{h(b)}}, \quad (139)$$

where i_b, j_b denote the coordinates of a box b , $h_\mu(b)$ and $h(b)$ is the length of the hook of b in β_μ and λ respectively. Here we used a relation between multiplicity of a given irrep with the number of semisimple Young diagrams [35].

For $\mu = 1, 2, \dots, s$ we define the following subsets R_μ, C_μ of λ

$$R_\mu := \{b \in \lambda : i_b = i_{c_\mu}, j_b < j_{c_\mu}\}, \quad (140)$$

$$C_\mu := \{b \in \lambda : i_b < i_{c_\mu}, j_b = j_{c_\mu}\}. \quad (141)$$

These subsets contain respectively boxes which are in the same row/column as the corner box c_μ not taking into account c_μ itself. Let $G_\mu := R_\mu \sqcup C_\mu \sqcup \{c_\mu\}$. Observe that $h_\mu(b) = h(b)$ if $b \notin G_\mu$. Thus, (139) simplifies to

$$c(\lambda) = \sum_{\mu=1}^s \frac{\prod_{b \in G_\mu \setminus \{c_\mu\}} \frac{d - i_b + j_b}{h_\mu(b)}}{\prod_{b \in G_\mu} \frac{d - i_b + j_b}{h(b)}} \quad (142)$$

$$= \sum_{\mu=1}^s \frac{1}{d - i_{c_\mu} + j_{c_\mu}} \prod_{b \in C_\mu} \frac{h(b)}{h(b) - 1} \prod_{b \in R_\mu} \frac{h(b)}{h(b) - 1}, \quad (143)$$

where we used the fact that $h(c_\mu) = 1$, and $h_\mu(b) = h(b) - 1$ for $b \in G_\mu \setminus \{c_\mu\}$. We define

$$C_{\mu,\nu} := \{b \in C_\mu : i_{c_{\nu-1}} < i_b \leq i_{c_\nu}\} \quad (144)$$

for $\nu = 1, \dots, \mu$, and

$$R_{\mu,\nu} := \{b \in R_\mu : j_{c_{\nu+1}} < j_b \leq j_{c_\nu}\} \quad (145)$$

for $\nu = \mu, \dots, s$, with the convention that $i_{c_0} = 0 = j_{c_{s+1}}$. Observe that $C_\mu = \bigsqcup_{\nu=1}^\mu C_{\mu,\nu}$ and $R_\mu = \bigsqcup_{\nu=\mu}^s R_{\mu,\nu}$, and for $b \in C_\mu \sqcup R_\mu$,

$$h(b) = \begin{cases} i_{c_\mu} - i_b + j_{c_\nu} - j_{c_\mu} + 1, & \text{if } b \in C_{\mu,\nu}, \\ i_{c_\nu} - i_{c_\mu} + j_{c_\mu} - j_b + 1, & \text{if } b \in R_{\mu,\nu}. \end{cases} \quad (146)$$

Hence, exploiting (92), we have

$$\prod_{b \in C_\mu} \frac{h(b)}{h(b) - 1} = \prod_{\nu=1}^\mu \prod_{b \in C_{\mu,\nu}} \frac{h(b)}{h(b) - 1} \quad (147)$$

$$= \prod_{\nu=1}^\mu \prod_{b \in C_{\mu,\nu}} \frac{i_{c_\mu} - i_b + j_{c_\nu} - j_{c_\mu} + 1}{i_{c_\mu} - i_b + j_{c_\nu} - j_{c_\mu}} \quad (148)$$

$$= \prod_{i=i_{c_{\mu-1}}+1}^{i_{c_\mu}-1} \frac{i_{c_\mu} - i + 1}{i_{c_\mu} - i} \prod_{\nu=1}^{\mu-1} \prod_{i=i_{c_{\nu-1}}+1}^{i_{c_\nu}} \frac{i_{c_\mu} - i + j_{c_\nu} - j_{c_\mu} + 1}{i_{c_\mu} - i + j_{c_\nu} - j_{c_\mu}} \quad (149)$$

$$= (i_{c_\mu} - i_{c_{\mu-1}}) \prod_{\nu=1}^{\mu-1} \frac{i_{c_\mu} - i_{c_{\nu-1}} + j_{c_\nu} - j_{c_\mu}}{i_{c_\mu} - i_{c_\nu} + j_{c_\nu} - j_{c_\mu}} \quad (150)$$

$$= k_\mu \prod_{\nu=1}^{\mu-1} \frac{k_\nu + p_\nu + \dots + k_\mu}{p_\nu + \dots + k_\mu} \quad (151)$$

$$= k_\mu \prod_{\nu=1}^{\mu-1} \left(1 + \frac{k_\nu}{d(c_\nu, c_\mu)} \right). \quad (152)$$

Analogously, we have

$$\prod_{b \in R_\mu} \frac{h(b)}{h(b) - 1} = \prod_{\nu=\mu}^s \prod_{b \in R_{\mu,\nu}} \frac{h(b)}{h(b) - 1} \quad (153)$$

$$= \prod_{\nu=\mu}^s \prod_{b \in R_{\mu,\nu}} \frac{i_{c_\nu} - i_{c_\mu} + j_{c_\mu} - j_b + 1}{i_{c_\nu} - i_{c_\mu} + j_{c_\mu} - j_b} \quad (154)$$

$$= \prod_{j=j_{c_{\mu+1}}+1}^{j_{c_\mu}-1} \frac{j_{c_\mu} - j + 1}{j_{c_\mu} - j} \prod_{\nu=\mu+1}^s \prod_{j=j_{c_{\nu+1}}+1}^{j_{c_\nu}} \frac{i_{c_\nu} - i_{c_\mu} + j_{c_\mu} - j + 1}{i_{c_\nu} - i_{c_\mu} + j_{c_\mu} - j} \quad (155)$$

$$= (j_{c_\mu} - j_{c_{\mu+1}}) \prod_{\nu=\mu+1}^s \frac{i_{c_\nu} - i_{c_\mu} + j_{c_\mu} - j_{c_{\nu+1}}}{i_{c_\nu} - i_{c_\mu} + j_{c_\mu} - j_{c_\nu}} \quad (156)$$

$$= p_\mu \prod_{\nu=\mu+1}^s \frac{p_\mu + \dots + k_\nu + p_\nu}{p_\mu + \dots + k_\nu} \quad (157)$$

$$= p_\mu \prod_{\nu=\mu+1}^s \left(1 + \frac{p_\nu}{d(c_\mu, c_\nu)} \right). \quad (158)$$

Coming back to (143), we get

$$c(\lambda) = \sum_{\mu=1}^s \frac{k_\mu p_\mu}{d - i_{c_\mu} + j_{c_\mu}} \prod_{\nu=1}^{\mu-1} \left(1 + \frac{k_\nu}{d(c_\nu, c_\mu)} \right) \prod_{\nu=\mu+1}^s \left(1 + \frac{p_\nu}{d(c_\mu, c_\nu)} \right) \quad (159)$$

what, after taking (92) into account, ends the proof. \square

As an example, let us consider the diagram $\lambda = [1, 1, \dots, 1] = [1^{\times n}]$. Its step representation is $s = 1$, $(k_1, p_1) = (n, 1)$. Direct application of the formula (137) shows that

$$c([1^{\times n}]) = \frac{n}{d - n + 1}. \quad (160)$$

Now, we are ready to prove the main theorem of this subsection

Theorem 3: For every diagram λ with n boxes,

$$c(\lambda) \leq c([1^{\times n}]). \quad (161)$$

Proof: It was shown in [36] that

$$\sum_{j=1}^s k_j p_j \prod_{i=1}^{j-1} \left(1 + \frac{k_i}{d(c_i, c_j)} \right) \prod_{i=j+1}^s \left(1 + \frac{p_i}{d(c_j, c_i)} \right) = n. \quad (162)$$

Observe that $\sum_{i=1}^j k_i \leq n$ and $\sum_{i=j}^s p_i \geq 1$ for every $j = 1, 2, \dots, s$. Hence, $d - \sum_{i=1}^j k_i + \sum_{i=j}^s p_i \geq d - n + 1$. From (138) and (162) we get

$$c(\lambda) \leq \frac{1}{d - n + 1} \sum_{j=1}^s k_j p_j \prod_{i=1}^{j-1} \left(1 + \frac{k_i}{d(c_i, c_j)} \right) \prod_{i=j+1}^s \left(1 + \frac{p_i}{d(c_j, c_i)} \right) \quad (163)$$

$$= \frac{n}{d - n + 1}. \quad (164)$$

According to (160) the latter number is equal to $c([1^{\times n}])$. \square

The above theorem clearly shows that the maximum possible value of the constant from (136) is attained for antisymmetric partitions with $n = k + 1$ boxes:

$$c = \frac{1}{d} \max_{\lambda \vdash (k+1)} c(\lambda) = \frac{k+1}{d(d-k)}. \quad (165)$$

This exactly reproduces the lower bound obtained from the primal problem and proves optimally of chosen ansatz.

V. DISCUSSIONS

We have presented the optimal quantum circuits for transforming k calls of a d -dimensional unitary operation U_d into its complex conjugate \overline{U}_d in a deterministic non-exact manner. When considering probabilistic exact transformations, [29] shows that transforming k uses of U_d into its complex conjugation \overline{U}_d necessarily has probability one or probability zero, and shows that probability one can only be attained when $k \geq d - 1$ – completely solving the probabilistic exact complex conjugation. Hence, the problem of designing circuits for unitary complex conjugation is solved on the probabilistic exact paradigm and following our work now also on the deterministic non-exact one.

A direct application of our circuit for unitary complex conjugation is to design a quantum circuit for reversing quantum operations. In [29], the authors present a simple probabilistic circuit which performs exact unitary transposition $U_d \mapsto U_d^T$ with probability $1/d^2$. Since $U_d^{-1} = U_d^T$, if we perform a unitary transposition protocol after the complex conjugation circuit presented here, we can transform k uses of U_d into its inverse U_d^{-1} with probability $1/d^2$ and an average fidelity of $\langle F \rangle = \frac{k+1}{d(d-k)}$.

As stated earlier, since $(\overline{U_d V_d}) = (\overline{U_d})(\overline{V_d})$, the complex conjugation function $f(U_d) = \overline{U_d}$ is a homomorphism. Also, complex conjugation is the only non-trivial homomorphism f between d -dimensional special unitary operators $f : \mathcal{SU}(d) \rightarrow \mathcal{SU}(d)$. This follows from the fact that there exist only three d -dimensional representations for the group $\mathcal{SU}(d)$ [34]: the trivial representation $f_{\text{trivial}}(U_d) = \mathbb{1}$, the defining representation, $f_{\text{def}}(U_d) = U_d$, and the conjugate representation $f_{\text{conj}}(U_d) = \overline{U_d}$. Hence, we have obtained optimal circuits for transforming k uses of U_d into $f(U_d)$ for all homomorphic functions $f : \mathcal{SU}(d) \rightarrow \mathcal{SU}(d)$.

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