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Problem 21

Lemma 1.0.1. *In every planar triangulation G on atleast four vertices, there exists a vertex v which does not lie on the outer bound of G*

Proof. For the sake of contradiction let's assume there is no such vertex v. Then all the vertices lie on the outer bound of G forming a cycle with atleast four vertices. This is a contradiction with G being a planar triangulation, because we can add a edge between two not adjacent vertices to G and the result is still planar. Hence, there has to be atleast one vertex v not on the outer bound of G. □

Theorem 1.1. *Every planar triangulation G on atleast four vertices contains a vertex whose neighbourhood induces a cycle.*

Proof. After lemma 1.0.1 there exists a vertex v which does not lie on the outer bound of G. Let $N(v) = p_1, p_2, ..., p_n$ be the neighbourhood of v. $N(v)$ induces a cycle if all p_i are connected to a cycle and if there are no additionally edges between p_i and p_j with $|i - j| > 1$. We first will proof that all p_i form a cycle, hence the induced subgraph $N(v)$ has a cycle as a subgraph. For the sake of contradiction let's assume this is not the case, hence there is p_i and p_{i+1} which are not connected. This is either a contradiction with v not lying on the outer bound of G, or with G being a maximal planar graph. Because if v_i and v_{i+1} are not connected and v is not on the outer bound of G there is a face bounded by v,p_i,p_{i+1} and atleast one additional vertex. This face could be again divided into to smaller faces by adding an edge, hence G is no triangulation. Now we will proof that either $N(v)$ has no additional edges and thus is a cycle or that one of the vertices adjacent to v induces a cycle.

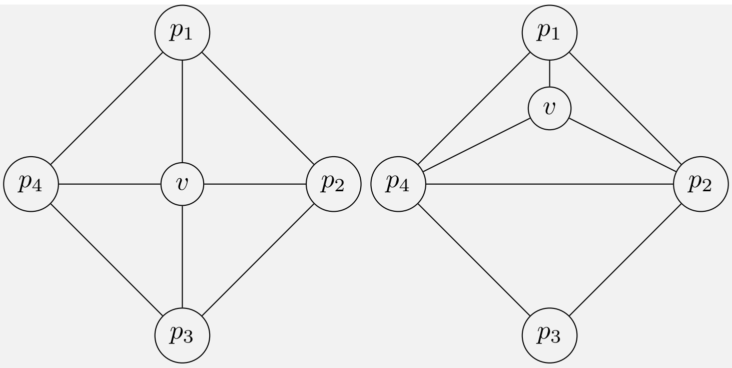
- *There are no additional edges connecting two vertices p_i and p_j with $|i - j| > 1$.*
  In this case $N(v)$ is a cycle and we found a vertex whose neighbourhood induces a cycle.
- *There is an edge p_ip_{i+2}*
  In this case the neighbourhood induces graph of p_{i+1} is a cycle namely $p_ip_{i+1}v$.
- *There is an edge between to vertices p_i and p_j with $|i - j| > 2$ and without loss of generality $i < v$.*
  This case there is a bounded face in G which is not a triangle. This face is the face bounded by at least the edges p_ip_{i+1}, p_jp_{j-1}, p_ip_j and atleast one or more edges forming a path from p_{i+1} to p_{j-1}. Hence this case can never occur in a planar triangulation.

In summary we now that we can always find a vertex v whose neighbourhood induces a cycle or one of the neighbours of v here called p_{i+1} has a neighbourhood inducing a cycle. □

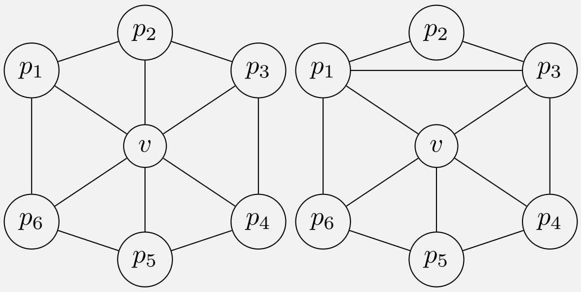
Lemma 1.1.1. *Every planar graph G with a maximal number of triangles and atleast 4 vertices has a vertex of degree 3.*

Proof. Using theorem 1.1 we can find a vertex v whose neighbourhood induces a cycle $p_1p_2...p_n$. In the following we will show by induction over the degree of v that if G has a maximal number of triangles, v cannot have a degree greater than 3.

Basis: v has a degree of 4 If v has a degree 4, then the subgraph with the nodes v, p_1, p_2, p_3, p_4 only has 4 triangles. We can change the edge set of this subgraph to get a subgraph with the same vertices and the same outer structure but with containing 6 triangles. We can achieve this by removing edge vp_3 and adding the edge p_2p_4. Thus we know that if G has a maximal number of triangles v cannot have degree 4.



Induction step Let the degree of v be $n > 3$. We now can again change the edges set without changing the outer structure of the induced subgraph $N(v)$ so that v has degree $n-1$ and the number of triangles in $N(v)$ stays the same. We can achieve this by removing the edge vp_2 and adding the edge p_1p_3. By induction we can repeat this operation until $n = 4$ and have the base case. Thus we can increase the number of triangles in the subgraph without changing the vertex count or the outer structure. Hence, the degree of v cannot be greater than 3 if G has a maximal number of triangles.



□

Theorem 1.2. *Every n-vertex planar graph has at most $3n - 8$ triangles.*

Proof. Let G be a planar graph with a maximal number of triangles. In the following we will show by induction that a planar graph G with atleast 3 vertices with a maximum number of triangles has exactly $3n - 8$ triangles.

Base: $n = 3$ The planar graph with maximum number of triangles is $K3$ which has one triangle. $3n - 8 = 3 * 3 - 8 = 1$ so for the base case the formula is correct.

Induction Step Let G be such a graph with n nodes. By using lemma 1.1.1 we can find a vertex v with degree 3 whose neighbourhood induces a cycle. If we remove this node we decrease the number of nodes by one and the number of triangles by 3. By induction we know that the resulting graph has no more than $3*(n-1)-8$ triangles. Hence, we know that G has no more than $3*(n-1)-8+3 = 3*n-8$ triangles. □

Problem 22