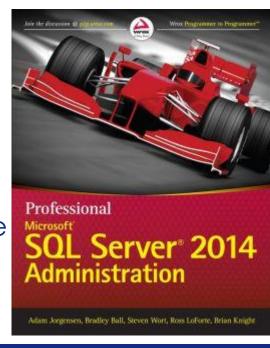


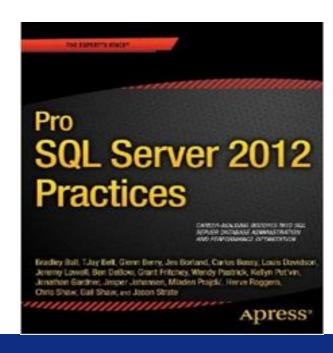


Performance Troubleshooting

Speaker Introduction Bradley Ball

- Almost 15 Years IT Experience
- Previous experience DBA, for the U.S. Army, The Executive Office of the President, Sr. SQL DBA Staff Specialist at Publix
- Currently the Data Platform Management Lead for <u>Pragmatic Works</u>
- Microsoft VTSP for the Greater North East
- MCITP SQL 2005 DBA & SQL 2008 DBA
- Blog: http://www.SQLBalls.com
- Twitter: @SQLBalls @BradleyBall_PW
- Email: <u>bball@PragmaticWorks.com</u>
- Pro SQL Server 2012 Practices Author
- Chapter 14 PAGE & ROW COMPRESSION!
- Managing Author on Pro Admin 2014 Guide





Agenda



- Wait Stats
- Baseline Disks
- Power Configuration
- Instant Database File Initialization
- Locks and Blocking
- Joins
- Execution Plans



SQLOS

Preemptive Scheduling (OS)

Cooperative Scheduling (SQL)



Dynamic Management Views and Functions

- sys.dm_os_wait_stats
 - Server-level
 - Contains aggregated wait statistics by wait type since last restart of SQL
 - DBCC SQLPerf('sys.dm_os_wait_stats',clear)
- sys.dm_os_waiting_tasks
 - Task-level
 - Displays all waiting sessions (waiter list)
- sys.dm_exec_requests
 - Describes each SQL Server Session (spid)

Demo





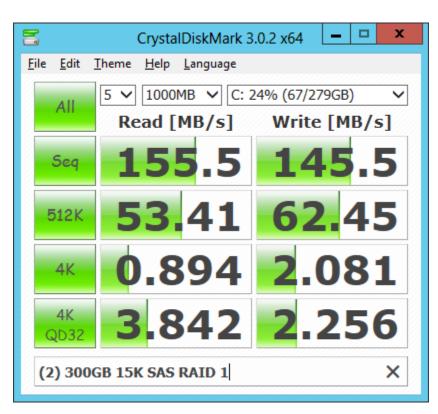
Agenda



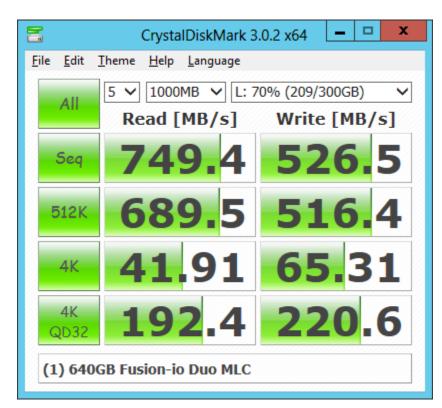
- Wait Stats
- Baseline Disks
- Power Configuration
- Instant Database File Initialization
- Locks and Blocking
- Joins
- Execution Plans



Disk Performance Baseline



Two 300 GB 15K SAS drives in RAID 1



One 640 GB Fusion-IO Duo MLC device



- Balanced *DEFAULT
 - Default Setting. Targets good energy efficiency with minimal performance impact. Matches capacity to demand. Energy-Saving features balance power and performance.

- High Performance
 - Increases performance at the cost of high energy consumption. Power and thermal limitations, operating expenses, and reliability considerations apply

- Power Saver
 - Limits performance to Save Energy and reduce operating costs. Caps processor frequency at a percentage of maximum (if supported), and enables other energy-saving features.



- Bal
 - Default Setting. Targets good energy

 demand. Energy-Saving features balance power and performance impact. Matches capa

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy

 demand. Energy-Saving features balance power and performance impact.

 Targets good energy features balance power and performance impact.

 Targets good energy features balance power and performance impact.

 Targets good energy features balance power and performance impact.

 Targets good energy features balance power and performance impact.

 Targets good energy features balance power and performance impact.

 Targets good energy features balance power and performance impact.

 Targets good energy features balance power and performance impact.

 Targets good energy features balance power and performance impact.

 Targets good energy features good energy featur

- High Performance
 - Increases performance at the cost of high energy consumption. Power and thermal limitations, operating expenses, and reliability considerations apply

Limits performance to Save Energy and reduce.

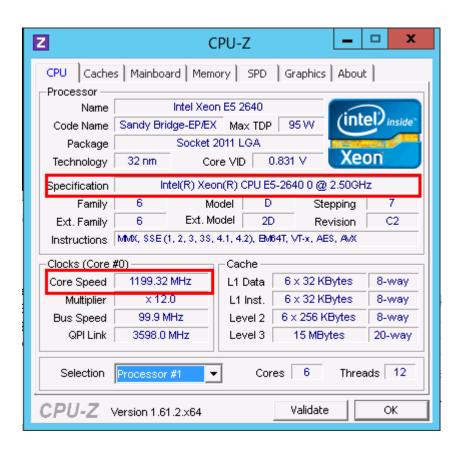
Cans processor frequency and analysis and analysis.

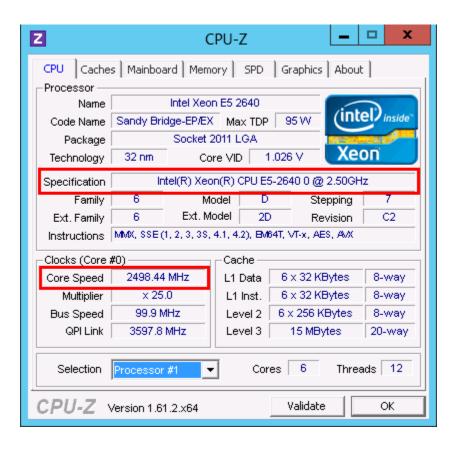






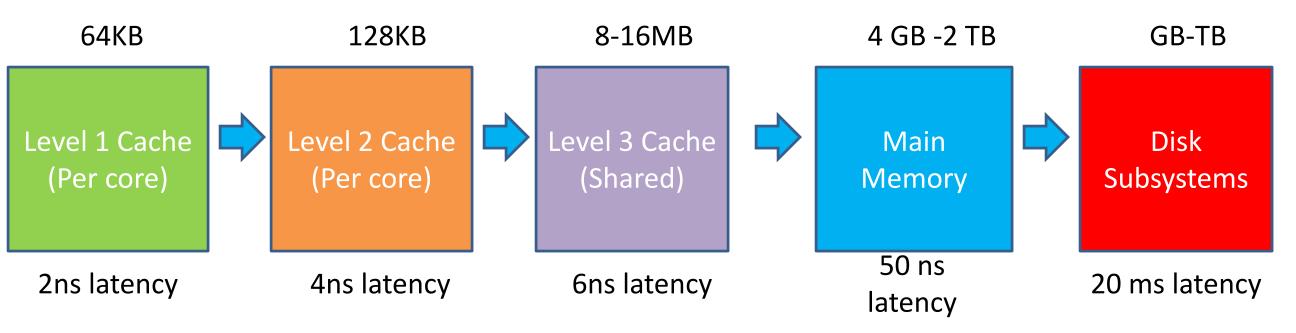








Data Retrieval Time



Instant DB File Initialization



*Hardware: Lenovo ThinkPad i7-2640 Dual Core, 16 GB Memory, 1 TB External HD 5400

Performance Test WITHOUT Instant Database File Initialization

CREATE DATABASE w/ 20 GB Data file = 11:57 min/seconds
ALTER DATABASE by 10 GB = 5:59 min/seconds
BACKUP 30 GB DATABASE(Empty) = 00:17 min/seconds
RESTORE 30 GB DATABASE (Empty Backup) = 17:46 min/seconds
RESTORE 30 GB DATABASE (12 GB Backup) = 25:07 min/seconds
BACKUP 30 GB DATABASE (12 GB of Data) = 17:21 min/seconds

Performance Test WITH Instant Database File Initialization

CREATE DATABASE w/ 20 GB Data file = 00:12 min/seconds

ALTER DATABASE by 10 GB = 00:00 min/seconds

BACKUP 30 GB DATABASE(Empty) = 00:13 min/seconds

RESTORE 30 GB DATABASE (Empty Backup) = 00:07 min/seconds

RESTORE 30 GB DATABASE (12 GB Backup) = 6:50 min/seconds

BACKUP 30 GB DATABASE (12 GB of Data) = 13:43 min/seconds

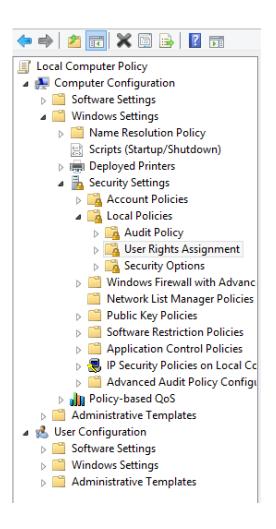
Instant Database File

Initialization

- SQL Service Account
 - Must be in the Local Admin Group
 - Must Have Group Policy Right: Perform Volume Maintenance Tasks
- Allows SQL to call the Windows API SetFileValidData
- Requires a SQL Service Restart to Add
- Requires a Server Restart to Remove







Demo





Agenda



- Wait Stats
- Baseline Disks
- Power Configuration
- Instant Database File Initialization
- Locks and Blocking
- Joins
- Execution Plans



ACID

Atomicity

Consistency

Isolation

Durability

Begin Transaction
Withdraw \$1000 Savings
Deposit \$1000 Checking
Commit Transaction

Transaction Isolation Levels

- Serializable
- Read Committed
- Read Uncommitted
- Repeatable Read
- Snapshot Isolation



Lock Matrix

	NL	SCH-S	SCH-M	s	U	×	IS	IU	IX	SIU	SIX	UIX	BU	RS-S	RS-U	RI-N	RI-S	RI-U	RI-X	RX-S	RX-U	RX-X
NL	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N
SCH-S	N	N	С	N	N	N	N	N	N	N	N	N	N	I	I	I	I	I	I	I	I	I
SCH-M	N	С	C	С	С	С	С	С	С	С	С	С	С	I	I	I	I	I	I	I	I	I
S	N	N	С	N	N	С	N	N	С	N	С	С	С	N	N	N	N	N	С	N	N	С
U	N	N	С	N	С	С	N	С	С	C	С	С	С	N	С	N	N	С	С	N	С	С
X	N	N	С	С	С	С	С	С	С	С	С	С	С	С	С	N	С	С	С	С	С	С
IS	N	N	С	N	N	С	N	N	N	N	N	N	С	I	I	I	I	I	I	I	I	I
IU	N	N	С	N	С	С	N	N	N	N	N	С	С	I	I	I	I	I	I	I	I	I
IX	N	N	C	С	С	С	N	N	N	С	С	С	С	I	I	I	I	I	I	I	I	I
SIU	N	N	С	N	С	С	N	N	С	N	С	С	С	I	I	I	I	I	I	I	I	I
SIX	N	N	C	С	С	С	N	N	С	С	С	С	С	I	I	I	I	I	I	I	I	I
UIX	N	N	С	С	С	С	N	С	С	С	С	С	С	I	I	I	I	I	I	I	I	I
BU	N	N	С	С	С	С	С	С	С	С	С	С	N	I	I	I	I	I	I	I	I	I
RS-S	N	I	I	N	N	С	I	I	I	I	I	I	I	N	N	С	С	С	С	С	С	С
RS-U	N	I	I	N	С	С	I	I	I	I	I	I	I	N	С	С	С	С	С	С	С	С
RI-N	N	I	I	N	N	N	I	I	I	I	I	I	I	С	С	N	N	N	N	С	С	С
RI-S	N	I	I	N	N	С	I	I	I	I	I	I	I	С	С	N	N	N	С	С	С	С
RI-U	N	I	I	N	С	С	I	I	I	I	I	I	I	С	С	N	N	С	С	С	С	С
RI-X	N	I	I	С	С	С	I	I	I	I	I	I	Ĭ	С	С	N	С	С	С	С	С	С
RX-S	N	I	I	N	N	С	I	I	I	I	I	I	I	С	С	С	С	С	С	С	С	С
RX-U	N	I	I	N	С	С	I	I	I	I	I	I	I	С	С	С	С	С	С	С	С	С
RX-X	N	I	I	С	C	С	I	I	I	I	I	I	I	С	С	С	С	С	С	С	С	С

Key

N	No Conflict	SIU	Share with Intent Update
I	Illegal	SIX	Shared with Intent Exclusive
C	Conflict	UIX	Update with Intent Exclusive
		BU	Bulk Update
NL	No Lock	RS-S	Shared Range-Shared
SCH-S	Schema Stability Locks	RS-U	Shared Range-Update
SCH-M	Schema Modification Locks	RI-N	Insert Range-Null
S	Shared	RI-S	Insert Range-Shared
U	Update	RI-U	Insert Range-Update
X	Exclusive	RI-X	Insert Range-Exclusive
IS	Intent Shared	RX-S	Exclusive Range-Shared
IU	Intent Update	RX-U	Exclusive Range-Update
IX	Intent Exclusive	RX-X	Exclusive Range-Exclusive



Locks

Shared - Reader Lock. Universal Giver.

Update - Hybrid lock that starts out with properties of a Shared lock and moves toward an Exclusive.

eXclusive – Data has been modified, within this transaction and cannot be accessed until it is hardened to Disk.

Shared Update eXclusive Shared Y Y N Locks Update Y N N eXclusive N N N N



Locks

Intent Locks- Qualifier to all previous modes. IS, IX, IU.

<u>SCH</u>ema <u>Stability Locks</u> - When a query is being Complied, SS locks is acquired to prevent Schema Modification Locks.

<u>SCH</u>ema <u>Modification Locks</u> – SM Locks are taken as a Table's structure is being modified i.e. DDL Commands.

<u>Bulk Update Locks</u>– When BCP is used or a Bulk Insert Command is Called. Can be set by User (TABLOCK) or Table setting 'table lock on bulk'=TRUE



Locks

Conversion Locks – Requested when SQL Server needs to convert a lock within a Transaction. SIX, SIU, UIX

Shared Intent eXclusive- a resource holding a Shared lock also has a component (page or row) locked with an eXclusive lock.

Shared Intent Update-a resource holding a shared lock also has a component (page or row) locked with an Update lock.

<u>Update Intent eXclusive-resource</u> holding an update lock also has a component (page or row) locked with an eXclusive lock.

Key Range Locks-Locks entire Range of a transaction in Serializable Isolation Level.

Key Locks- Taken on modifications on a Clustered Index instead of Row locks



Lock Matrix

	NL	SCH-S	SCH-M	s	U	×	IS	IU	IX	SIU	SIX	UIX	BU	RS-S	RS-U	RI-N	RI-S	RI-U	RI-X	RX-S	RX-U	RX-X
NL	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N
SCH-S	N	N	С	N	N	N	N	N	N	N	N	N	N	I	I	I	I	I	I	I	I	I
SCH-M	N	С	C	С	С	С	С	С	С	С	С	С	С	I	I	I	I	I	I	I	I	I
S	N	N	С	N	N	С	N	N	С	N	С	С	С	N	N	N	N	N	С	N	N	С
U	N	N	С	N	С	С	N	С	С	C	С	С	С	N	С	N	N	С	С	N	С	С
X	N	N	С	С	С	С	С	С	С	С	С	С	С	С	С	N	С	С	С	С	С	С
IS	N	N	С	N	N	С	N	N	N	N	N	N	С	I	I	I	I	I	I	I	I	I
IU	N	N	С	N	С	С	N	N	N	N	N	С	С	I	I	I	I	I	I	I	I	I
IX	N	N	C	С	С	С	N	N	N	С	С	С	С	I	I	I	I	I	I	I	I	I
SIU	N	N	С	N	С	С	N	N	С	N	С	С	С	I	I	I	I	I	I	I	I	I
SIX	N	N	C	С	С	С	N	N	С	С	С	С	С	I	I	I	I	I	I	I	I	I
UIX	N	N	С	С	С	С	N	С	С	С	С	С	С	I	I	I	I	I	I	I	I	I
BU	N	N	С	С	С	С	С	С	С	С	С	С	N	I	I	I	I	I	I	I	I	I
RS-S	N	I	I	N	N	С	I	I	I	I	I	I	I	N	N	С	С	С	С	С	С	С
RS-U	N	I	I	N	С	С	I	I	I	I	I	I	I	N	С	С	С	С	С	С	С	С
RI-N	N	I	I	N	N	N	I	I	I	I	I	I	I	С	С	N	N	N	N	С	С	С
RI-S	N	I	I	N	N	С	I	I	I	I	I	I	I	С	С	N	N	N	С	С	С	С
RI-U	N	I	I	N	С	С	I	I	I	I	I	I	I	С	С	N	N	С	С	С	С	С
RI-X	N	I	I	С	С	С	I	I	I	I	I	I	Ĭ	С	С	N	С	С	С	С	С	С
RX-S	N	I	I	N	N	С	I	I	I	I	I	I	I	С	С	С	С	С	С	С	С	С
RX-U	N	I	I	N	С	С	I	I	I	I	I	I	I	С	С	С	С	С	С	С	С	С
RX-X	N	I	I	С	C	С	I	I	I	I	I	I	I	С	С	С	С	С	С	С	С	С

Key

N	No Conflict	SIU	Share with Intent Update
I	Illegal	SIX	Shared with Intent Exclusive
C	Conflict	UIX	Update with Intent Exclusive
		BU	Bulk Update
NL	No Lock	RS-S	Shared Range-Shared
SCH-S	Schema Stability Locks	RS-U	Shared Range-Update
SCH-M	Schema Modification Locks	RI-N	Insert Range-Null
S	Shared	RI-S	Insert Range-Shared
U	Update	RI-U	Insert Range-Update
X	Exclusive	RI-X	Insert Range-Exclusive
IS	Intent Shared	RX-S	Exclusive Range-Shared
IU	Intent Update	RX-U	Exclusive Range-Update
IX	Intent Exclusive	RX-X	Exclusive Range-Exclusive

Demo

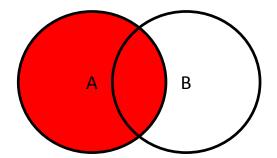




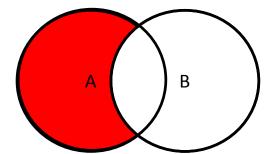
Agenda



- Wait Stats
- Baseline Disks
- Power Configuration
- Instant Database File Initialization
- Locks and Blocking
- Joins
- Execution Plans

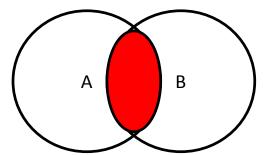


SELECT <select_list>
FROM TableA A
LEFT JOIN TableB B
On A.Key=B.Key

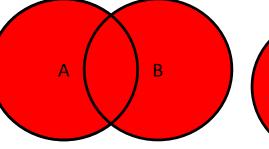


SELECT <select_list>
FROM TableA A
LEFT JOIN TableB B
On A.Key=B.Key
WHERE B.Key IS NULL

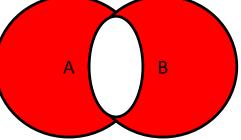
Syntactical Joins



SELECT <select_list>
FROM TableA A
INNER JOIN TableB B
On A.Key=B.Key

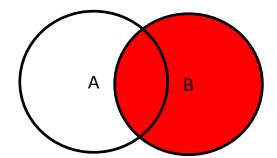


SELECT <select_list>
FROM TableA A
FULL OUTER JOIN TableB B
On A.Key=B.Key

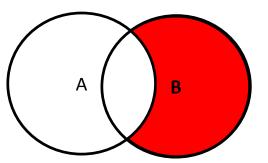


SELECT <select_list>
FROM TableA A
FULL OUTER JOIN TableB B
On A.Key=B.Key
Where A.Key IS NULL OR B.Key IS
NULL





SELECT <select_list>
FROM TableA A
RIGHT JOIN TableB B
On A.Key=B.Key



SELECT <select_list>
FROM TableA A
RIGHT JOIN TableB B
On A.Key=B.Key
WHERE A.KEY IS NULL



Physical Operators

Joins

- Syntactical Joins
 - Inner Join
 - Outer Join
 - Cross Join
 - Cross Apply
 - Outer Apply
 - Semi-Join
 - Anti-Semi Join

- Physical Joins
 - Nested Loop Join
 - Merge Join
 - Hash Join

*If no Syntactical Join Type is specified Inner Join is the default



Nested Loop Join

Compares Each Row from the Outer Table To Each Row in the Inner Table

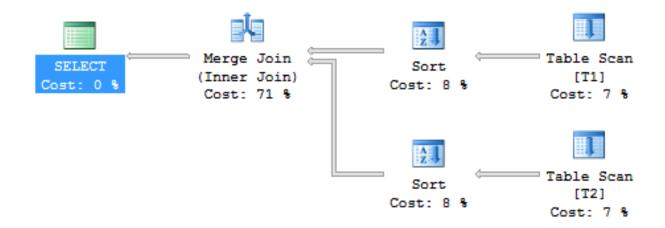
```
for each row R1 in the outer table
  begin
     for each row R2 in the inner table
       if R1 joins with R2
       return (R1, R2) if R1 did not join
                                                              Nested Loops
                                                                                        Table Scan
                                            SELECT
                                                              (Inner Join)
                                                                                     [Customers] [C]
          return (R1, NULL)
                                           Cost: 0 %
                                                               Cost: 1 %
                                                                                        Cost: 48 %
  end
                                                                                        Table Scan
                                                                                        [Sales] [S]
                                                                                        Cost: 51 %
```



Sorted Merge Join

Simultaneously Reads and Compares two Sorted Inputs One Row at a Time

```
get first row R1 from input 1
get first row R2 from input 2
while not at the end of either input
begin
if R1 joins with R2
begin
return (R1, R2)
get next row R2 from input 2
end
else if R1 < R2
get next row R1 from input 1
else
get next row R2 from input 2
end
```

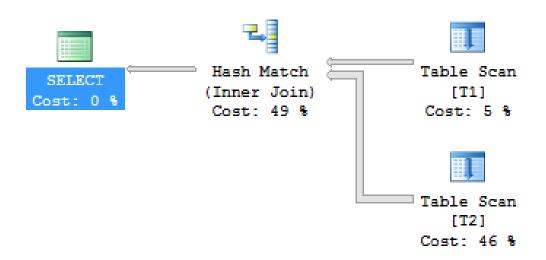




Hash Join

Join Heavy Lifter. Built in Two Phases: Build & Probe.

- Build
- Reads Rows from 1st Input
- Hashes on Equijoin Keys
- Creates In-Memory Hash Table



Probe

```
Hashes on Same Equijoin Keys
Probes for Matching Rows in Hash Table

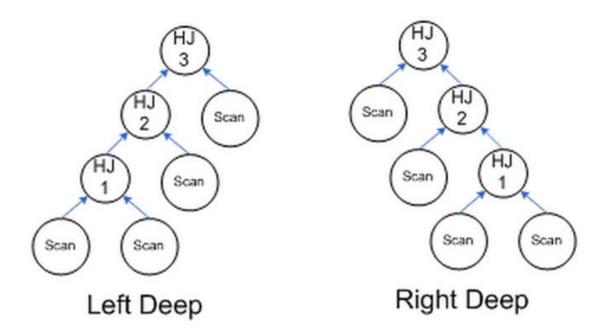
for each row R1 in the build table
begin
calculate hash value on R1 join key(s)
insert R1 into the appropriate hash bucket
end
for each row R2 in the probe table
begin
calculate hash value on R2 join key(s)
for each row R1 in the corresponding hash bucket
if R1 joins with R2
return (R1, R2)
end
```

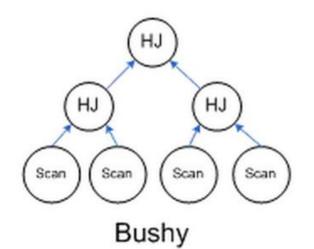
Reads All Rows from 2nd Input



Hash Join

Has 3 Types of Tree Structures

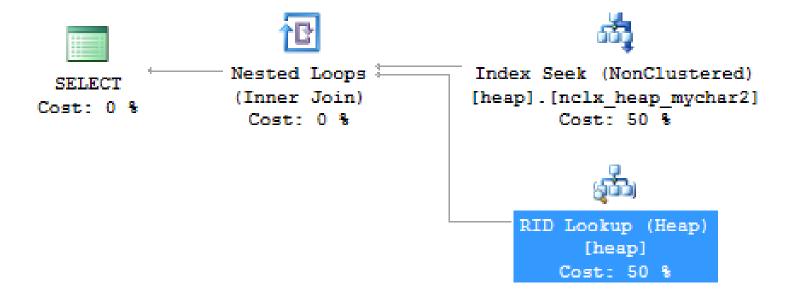






RID Lookup

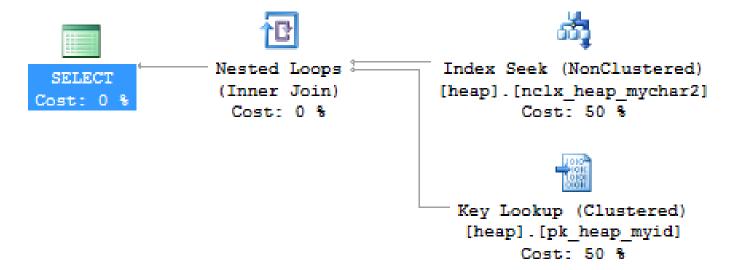
 When a Query Must Point Back to the Clustered Index For Additional Rows





Key Lookup

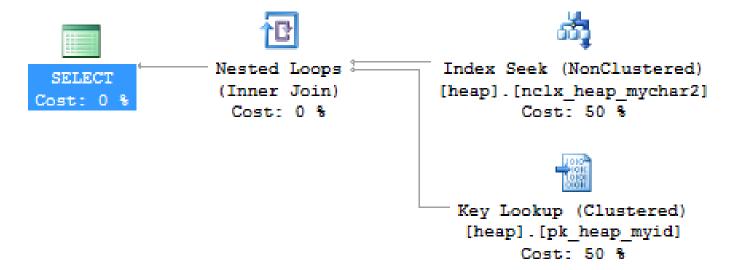
 When a Query Must Point Back to the Clustered Index For Additional Rows





Key Lookup

 When a Query Must Point Back to the Clustered Index For Additional Rows





Warnings

Yield Sign in Query Plan

SELECT	
Cached plan size	24 KB
Degree of Parallelism	4
Estimated Operator Cost	0 (0%)
Estimated Subtree Cost	9.39512
Memory Grant	72
Estimated Number of Rows	1.00006

Statement

select * from students where ssn=@ssn

Warnings

Type conversion in expression (CONVERT_IMPLICIT(nchar(9),[college]. [dbo].[students].[ssn],0)) may affect "CardinalityEstimate" in query plan choice, Type conversion in expression (CONVERT_IMPLICIT(nchar(9),[college]. [dbo].[students].[ssn],0)=[@ssn]) may affect "SeekPlan" in query plan choice



Types

- Implicit Conversion
- Sort
- Unmatched Indexes
- Spill

Demo





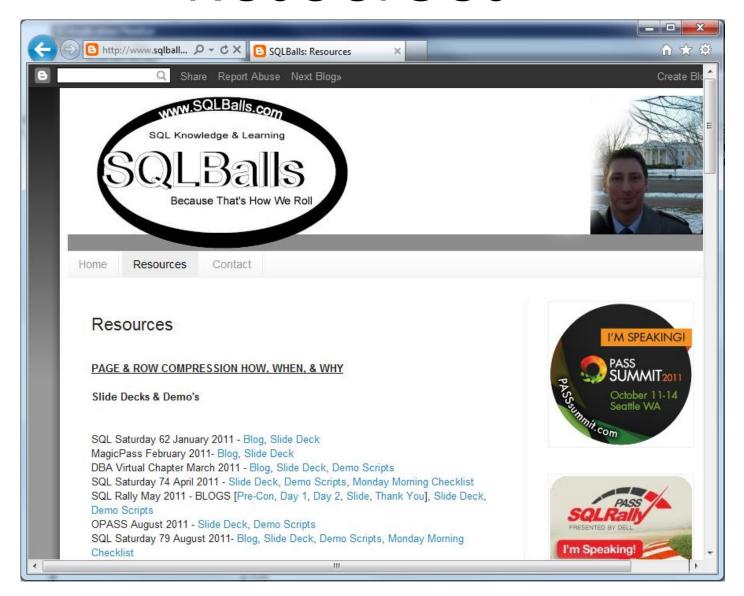
Agenda



- Wait Stats
- Baseline Disks
- Power Configuration
- Instant Database File Initialization
- Locks and Blocking
- Joins
- Execution Plans



Resources





Questions?

