## Semantics and Rewriting lambda

## Mario Tsatsev s1028415

May 2020

## Exercise 2

a)

$$\begin{array}{l} \mathbf{SKK} & \twoheadrightarrow_{\beta} \mathbf{I} \\ \mathbf{SKK} & \equiv (\lambda xyz.xz(yz))\mathbf{KK} \\ & \rightarrow (\lambda yz.\mathbf{K}z(yz))\mathbf{K} \\ & \rightarrow (\lambda z.\mathbf{K}z(Kz)) \\ & \rightarrow (\lambda z.(\lambda xy.x)z(\mathbf{K}z)) \\ & \rightarrow (\lambda z.(\lambda y.z)(\mathbf{K}z) \\ & \rightarrow (\lambda z.(\lambda.z)) \\ & \rightarrow (\lambda z.z) \\ & \equiv \mathbf{I} \end{array}$$

b)

$$\begin{split} \mathbf{KI} &\twoheadrightarrow_{\beta} K_{*} \\ \mathbf{KI} &\equiv (\lambda x y. x) \mathbf{I} \\ &\rightarrow (\lambda y. \mathbf{I}) \\ &\twoheadrightarrow y \\ &\equiv \mathbf{K}_{*} \end{split}$$

## Exercise 3

To determine that we must reduce until possible.

$\mathbf{SKIKISS} \to {}_{\beta}(\lambda x.x)(\mathbf{KIKISS})$	(1)
$\rightarrow_{\beta} (\lambda x.x)(\lambda xy.x)$ <b>IKISS</b>	(2)
$\rightarrow_{\beta} (\lambda xy.x)$ <b>IKISS</b>	(3)
$\rightarrow_{\beta} (\lambda xy.x)(\lambda x.x)(\lambda xy.x)$ <b>ISS</b>	(4)
$\rightarrow_{\beta} (\lambda y.(\lambda x.x))(\lambda xy.x)$ <b>ISS</b>	(5)
$\rightarrow_{eta} (\Lambda x.x) \mathbf{ISS}$	(6)
$\rightarrow_{\beta} (\Lambda x.x)(\lambda x.x)\mathbf{SS}$	(7)
$\rightarrow_{\beta} (\lambda x.x) \mathbf{SS}$	(8)
$\rightarrow_{\beta} (\lambda x.x)(\lambda xyz.xz(yz))\mathbf{S}$	(9)
$\rightarrow_{\beta} (\lambda xyz.xz(yz))(\lambda xyz.xz(yz))$	(10)
$\rightarrow_{eta} (\lambda yz.(\lambda xyz.xz(yz))z(yz))$	(11)
$\rightarrow_{eta} (\lambda yz.(\lambda yz.zz(yz))(yz))$	(12)
$\rightarrow_{eta} (\lambda yz.(\lambda z.zz(yzz)))$	(13)