

	$T_1$	$T_2$
Time ↓	read_item(X); $X := X - N;$	
		read_item(X); $X := X + M;$
	write_item(X); read_item(Y);	
		write_item(X);
	$Y := Y + N;$ write_item(Y);	

$S_a: r_1(X); r_2(X); w_1(X); r_1(Y); w_2(X); w_1(Y);$

---

$S_d: r_1(X); r_2(X); w_1(X); r_1(Y); w_2(X); c_2; w_1(Y); c_1;$

$S_c: r_1(X); w_1(X); r_2(X); r_1(Y); w_2(X); c_2; a_1;$

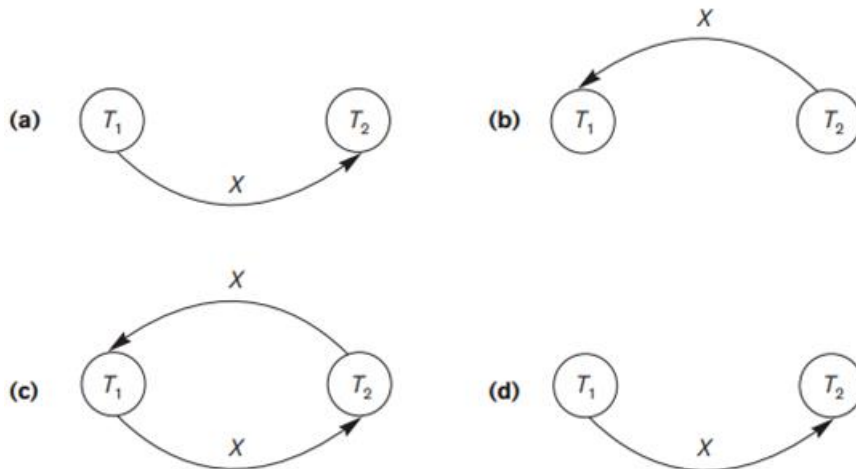
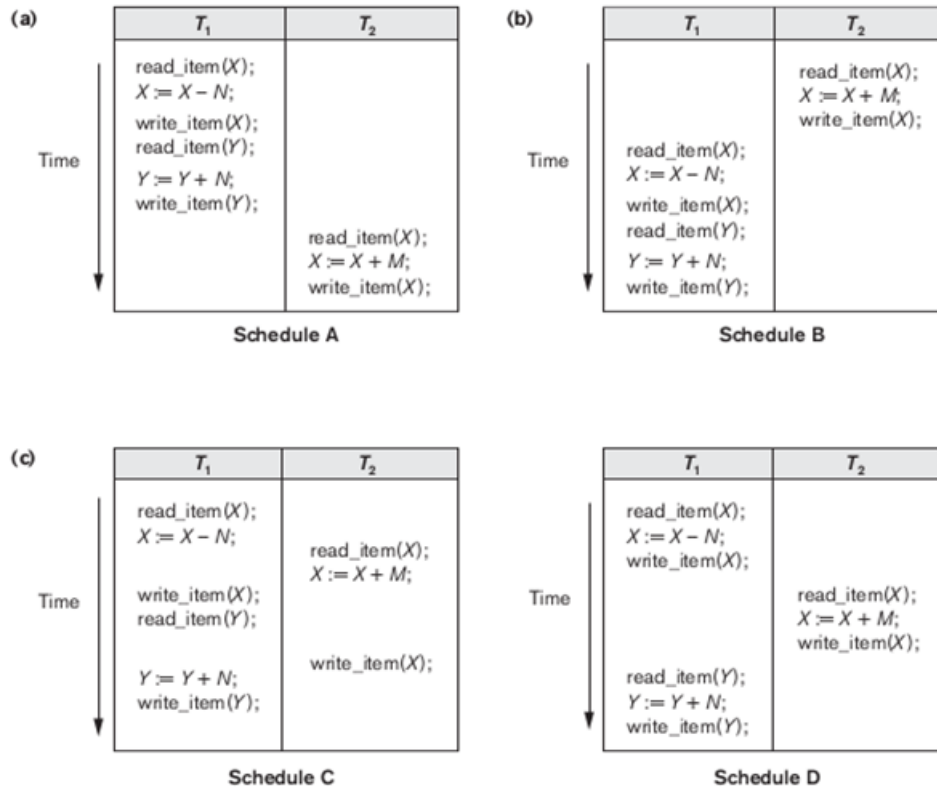
$S_d: r_1(X); w_1(X); r_2(X); r_1(Y); w_2(X); w_1(Y); c_1; c_2;$

$S_e: r_1(X); w_1(X); r_2(X); r_1(Y); w_2(X); w_1(Y); a_1; a_2;$

---

Sf: w1(X, 5); w2(X, 8); a1;

Examples of serial and nonserial schedules involving transactions  $T_1$  and  $T_2$ . (a) Serial schedule A:  $T_1$  followed by  $T_2$ . (b) Serial schedule B:  $T_2$  followed by  $T_1$ . (c) Two nonserial schedules C and D with interleaving of operations.



Constructing the precedence graphs for schedules A to D from to test for conflict serializability. (a) Precedence graph for serial schedule A. (b) Precedence graph for serial schedule B. (c) Precedence graph for schedule C (not serializable). (d) Precedence graph for schedule D (serializable, equivalent to schedule A).

Two schedules that are result equivalent for the initial value of  $X = 100$  but are not result equivalent in general.

$S_1$
<code>read_item(X);</code> <code>X := X + 10;</code> <code>write_item(X);</code>

$S_2$
<code>read_item(X);</code> <code>X := X * 1.1;</code> <code>write_item(X);</code>

Another example of serializability testing. (a) The read and write operations of three transactions  $T_1$ ,  $T_2$ , and  $T_3$ . (b) Schedule E. (c) Schedule F.

(a)

Transaction $T_1$
<code>read_item(X);</code> <code>write_item(X);</code> <code>read_item(Y);</code> <code>write_item(Y);</code>

Transaction $T_2$
<code>read_item(Z);</code> <code>read_item(Y);</code> <code>write_item(Y);</code> <code>read_item(X);</code> <code>write_item(X);</code>

Transaction $T_3$
<code>read_item(Y);</code> <code>read_item(Z);</code> <code>write_item(Y);</code> <code>write_item(Z);</code>

(b)

	Transaction $T_1$	Transaction $T_2$	Transaction $T_3$
Time ↓	<code>read_item(X);</code> <code>write_item(X);</code>	<code>read_item(Z);</code> <code>read_item(Y);</code> <code>write_item(Y);</code>	<code>read_item(Y);</code> <code>read_item(Z);</code>
	<code>read_item(Y);</code> <code>write_item(Y);</code>	<code>read_item(X);</code> <code>write_item(X);</code>	<code>write_item(Y);</code> <code>write_item(Z);</code>

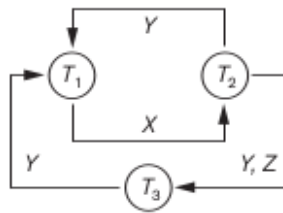
**Schedule E**

(c)

	Transaction $T_1$	Transaction $T_2$	Transaction $T_3$
Time ↓	<code>read_item(X);</code> <code>write_item(X);</code>		<code>read_item(Y);</code> <code>read_item(Z);</code>
	<code>read_item(Y);</code> <code>write_item(Y);</code>	<code>read_item(Z);</code>  <code>read_item(Y);</code> <code>write_item(Y);</code> <code>read_item(X);</code> <code>write_item(X);</code>	<code>write_item(Y);</code> <code>write_item(Z);</code>

**Schedule F**

(d)



Equivalent serial schedules

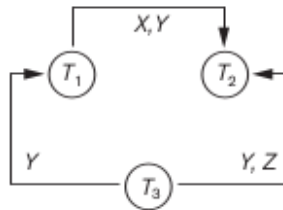
None

Reason

Cycle  $X(T_1 \rightarrow T_2), Y(T_2 \rightarrow T_1)$

Cycle  $X(T_1 \rightarrow T_2), YZ(T_2 \rightarrow T_3), Y(T_3 \rightarrow T_1)$

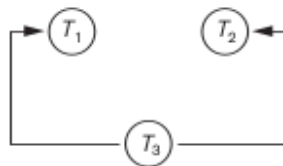
(e)



Equivalent serial schedules

$T_3 \rightarrow T_1 \rightarrow T_2$

(f)



Equivalent serial schedules

$T_3 \rightarrow T_1 \rightarrow T_2$

$T_3 \rightarrow T_2 \rightarrow T_1$

Another example of serializability testing.

(d) Precedence graph for schedule E.

(e) Precedence graph for schedule F.

(f) Precedence graph with two equivalent serial schedules.