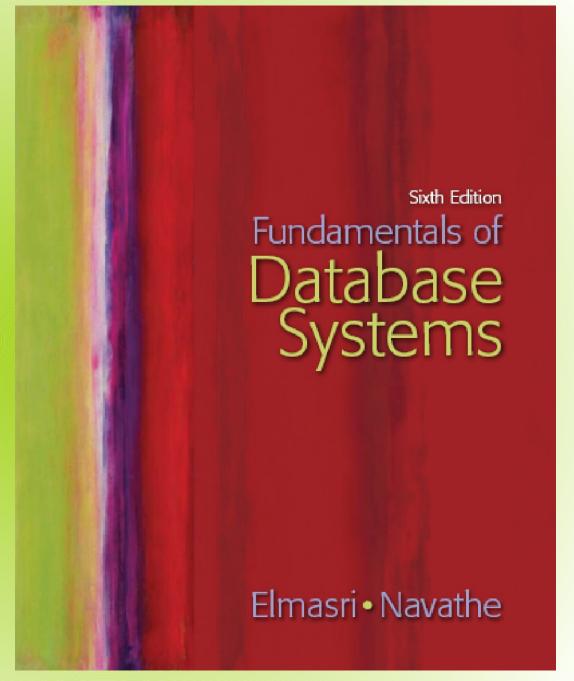
**Chapter 5 More SQL:** Complex Queries, Triggers, Views, and Schema **Modification** 



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### Chapter 5 Outline

More Complex SQL Retrieval Queries

Specifying Constraints as Assertions and Actions as Triggers

Views (Virtual Tables) in SQL

Schema Change Statements in SQL



## More Complex SQL Retrieval<sup>3</sup> Queries

Additional features allow users to specify more complex retrievals from database:

Nested queries, joined tables, outer joins, aggregate functions, and grouping



## Comparisons Involving NULL and Three-Valued Logic

Meanings of NULL

**Unknown value** 

Unavailable or withheld value

Not applicable attribute

Each individual NULL value considered to be different from every other NULL value

SQL uses a three-valued logic:

TRUE, FALSE, and UNKNOWN



## Comparisons Involving NULL and Three-Valued Logic (cont'd.)

Table 5.1	Logical Connectives in	Three-Valued Logic				
(a)	AND	TRUE	FALSE	UNKNOWN		
	TRUE	TRUE	FALSE	UNKNOWN		
	FALSE	FALSE	FALSE	FALSE		
	UNKNOWN	UNKNOWN	FALSE	UNKNOWN		
(b)	OR	TRUE	FALSE	UNKNOWN		
	TRUE	TRUE	TRUE	TRUE		
	FALSE	TRUE	FALSE	UNKNOWN		
	UNKNOWN	TRUE	UNKNOWN	UNKNOWN		
(c)	NOT					
	TRUE	FALSE				
	FALSE	TRUE				
	UNKNOWN	UNKNOWN				

## Comparisons Involving NULL and Three-Valued Logic (cont'd.)

SQL allows queries that check whether an attribute value is NULL

IS or IS NOT NULL

Query 18. Retrieve the names of all employees who do not have supervisors.

Q18: SELECT Fname, Lname

FROM EMPLOYEE

WHERE Super\_ssn IS NULL;

## Nested Queries, Tuples, and Set/Multiset Comparisons

### **Nested queries**

Complete select-from-where blocks within WHERE clause of another query

#### **Outer query**

Comparison operator IN

Compares value v with a set (or multiset) of values V

Evaluates to TRUE if v is one of the elements in V

Q4A: SELECT DISTINCT Pnumber

FROM PROJECT
WHERE Pnumber IN

( SELECT Pnumber

FROM PROJECT, DEPARTMENT, EMPLOYEE

WHERE Dnum=Dnumber AND

Mgr\_ssn=Ssn AND Lname='Smith')

OR

Pnumber IN

( SELECT Pno

FROM WORKS\_ON, EMPLOYEE

WHERE Essn=Ssn AND Lname='Smith');

Use tuples of values in comparisons

Place them within parentheses

FROM WORKS\_ON
WHERE (Pno, Hours) IN ( SELECT

Pno, Hours

FROM WORKS\_ON

WHERE Essn='123456789');

## Use other comparison operators to compare a single value *v*

- = ANY (or = SOME) operator
- Returns TRUE if the value v is equal to some value in the set V and is hence equivalent to IN

Other operators that can be combined with

Avoid potential errors and ambiguities

Create tuple variables (aliases) for all tables
referenced in SQL query

**Query 16.** Retrieve the name of each employee who has a dependent with the same first name and is the same sex as the employee.

Q16: SELECT E.Fname, E.Lname

FROM EMPLOYEE AS E

WHERE E.Ssn IN ( SELECT Essn

FROM DEPENDENT AS D

WHERE E.Fname=D.Dependent\_name

AND E.Sex=D.Sex );

### Correlated Nested Queries

Correlated nested query

Evaluated once for each tuple in the outer query

## The EXISTS and UNIQUE Functions in SQL

EXISTS function

Check whether the result of a correlated nested query is empty or not

EXISTS and NOT EXISTS

Typically used in conjunction with a correlated nested query

SQL function UNIQUE (Q)

Returns TRUE if there are no duplicate tuples in the result of query Q



## The EXISTS and UNIQUE Functions in SQL

Q16B: SELECT E.Fname, E.Lname

FROM EMPLOYEE AS E

WHERE EXISTS ( SELECT

FROM DEPENDENT AS D

WHERE E.Ssn=D.Essn AND E.Sex=D.Sex

AND E.Fname=D.Dependent\_name);

Query 6. Retrieve the names of employees who have no dependents.

Q6: SELECT Fname, Lname

FROM EMPLOYEE

WHERE NOT EXISTS ( SELECT

FROM DEPENDENT WHERE Ssn=Essn ):

Query 7. List the names of managers who have at least one dependent.

Q7: SELECT Fname, Lname

FROM EMPLOYEE

WHERE EXISTS ( SELECT

FROM DEPENDENT WHERE Ssn=Essn )

AND

EXISTS ( SELECT \*

FROM DEPARTMENT WHERE Ssn=Mgr ssn );

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## Explicit Sets and Renaming of Attributes in SQL

Can use explicit set of values in WHERE clause (Example)

Query 17. Retrieve the Social Security numbers of all employees who work on project numbers 1, 2, or 3.

Q17: SELECT DISTINCT Essn

FROM WORKS\_ON

WHERE Pno IN (1, 2, 3);

Use qualifier AS followed by desired new name Rename any attribute that appears in the result of a query

Q8A: SELECT E.Lname AS Employee\_name, S.Lname AS Supervisor\_name

FROM EMPLOYEE AS E, EMPLOYEE AS S

WHERE E.Super\_ssn=S.Ssn;



## Joined Tables in SQL and Outer Joins

#### Joined table

Permits users to specify a table resulting from a join operation in the FROM clause of a query

#### The FROM clause in Q1A

Contains a single joined table

Q1: SELECT Fname, Lname, Address

FROM EMPLOYEE, DEPARTMENT

WHERE Dname='Research' AND Dnumber=Dno;

Q1A: SELECT Fname, Lname, Address

FROM (EMPLOYEE JOIN DEPARTMENT ON Dno=Dnumber)

WHERE Dname='Research';



## Joined Tables in SQL and Outer Joins (cont'd.)

Specify different types of join

NATURAL JOIN

Various types of OUTER JOIN

NATURAL JOIN on two relations R and S

No join condition specified

Implicit EQUIJOIN condition for each pair of attributes with same name from R and S

Q1B: SELECT Fname, Lname, Address

FROM (EMPLOYEE NATURAL JOIN

(DEPARTMENT AS DEPT (Dname, Dno, Mssn, Msdate)))

WHERE Dname='Research';



## Joined Tables in SQL and Outer Joins (cont'd.)

#### Inner join

Default type of join in a joined table

Tuple is included in the result only if a matching tuple exists in the other relation

#### LEFT OUTER JOIN

Every tuple in left table must appear in result If no matching tuple

Padded with NULL values for attributes of right table

Q8B: SELECT E.Lname AS Employee\_name,

S.Lname AS Supervisor\_name

FROM (EMPLOYEE AS E LEFT OUTER JOIN EMPLOYEE AS S

ON E.Super\_ssn=S.Ssn);



## Joined Tables in SQL and Outer Joins (cont'd.)

#### RIGHT OUTER JOIN

Every tuple in right table must appear in result If no matching tuple

Padded with NULL values for the attributes of left table

#### **FULL OUTER JOIN**

Can nest join specifications

Q2A: SELECT Pnumber, Dnum, Lname, Address, Bdate

FROM ((PROJECT JOIN DEPARTMENT ON Dnum=Dnumber)

JOIN EMPLOYEE ON Mgr\_ssn=Ssn)

WHERE Plocation='Stafford';



### Aggregate Functions in SQL

Used to summarize information from multiple tuples into a single-tuple summary

### Grouping

Create subgroups of tuples before summarizing

Built-in aggregate functions count, sum, max, min, and avg

Functions can be used in the SELECT clause or in a HAVING clause

## Aggregate Functions in SQL (cont'd.)

NULL values discarded when aggregate functions are applied to a particular column

**Query 20.** Find the sum of the salaries of all employees of the 'Research' department, as well as the maximum salary, the minimum salary, and the average salary in this department.

Q20: SELECT SUM (Salary), MAX (Salary), MIN (Salary), AVG (Salary)

FROM (EMPLOYEE JOIN DEPARTMENT ON Dno=Dnumber)

WHERE Dname='Research';

Queries 21 and 22. Retrieve the total number of employees in the company (Q21) and the number of employees in the 'Research' department (Q22).

Q21: SELECT COUNT (\*)

FROM EMPLOYEE;

Q22: SELECT COUNT (\*)

FROM EMPLOYEE, DEPARTMENT

**WHERE** DNO=DNUMBER **AND** DNAME='Research';

Query 23. Count the number of distinct salary values in the database.

Q23: SELECT COUNT (DISTINCT Salary)

FROM EMPLOYEE;



## Grouping: The GROUP BY and HAVING Clauses

Partition relation into subsets of tuples

Based on grouping attribute(s)

Apply function to each such group independently

GROUP BY clause

Specifies grouping attributes

If NULLs exist in grouping attribute

Separate group created for all tuples with a NULL value in grouping attribute



## Grouping: The GROUP BY and HAVING Clauses

Query 24. For each department, retrieve the department number, the number of employees in the department, and their average salary.

Q24: SELECT Dno, COUNT (\*), AVG (Salary)

FROM EMPLOYEE

GROUPBY Dno;

Fname	Minit	Lname	San		Salary	Super_ssn	Dno			Dno	Count (*)	Avg (Salary)
John	В	Smith	123 456789		30000	333445555	5		_	5	4	33250
Franklin	Т	Wong	333 4455 55		40000	888665555	5		╛	4	3	31000
Ramesh	K	Narayan	66 688 44 44		38000	333445555	5		_	1	1	55000
Joyce	Α	English	453 4534 53		25 000	333445555	5			Result	of Q24	
Alicia	J	Zelaya	999887777		25 000	987654321	4	П				
Jennifer	S	Wallace	987654321		43 000	888665555	4	1  .	_			
Ahmad	٧	Jabbar	987987987		25 000	987654321	4					
James	E	Bong	888665555	1 1	55 000	NULL	1	Ħ.				

## Grouping: The GROUP BY and HAVING Clauses (cont'd.)

#### HAVING clause

Provides a condition on the summary information

**Query 28.** For each department that has more than five employees, retrieve the department number and the number of its employees who are making more than \$40,000.

Q28: SELECT Dnumber, COUNT (\*)

FROM DEPARTMENT, EMPLOYEE

WHERE Dnumber=Dno AND Salary>40000 AND

( SELECT Dno

FROM EMPLOYEE

**GROUP BY Dno** 

**HAVING** COUNT (\*) > 5



## Grouping: The GROUP BY and HAVING Clauses (cont'd.)

Query 25. For each project, retrieve the project number, the project name, and the number of employees who work on that project.

Q25: SELECT Pnumber, Pname, COUNT (\*)

FROM PROJECT, WORKS\_ON

WHERE Pnumber=Pno GROUPBY Pnumber, Pname;

Query 26. For each project on which more than two employees work, retrieve the project number, the project name, and the number of employees who work on the project.

Q26: SELECT Pnumber, Pname, COUNT (\*)

FROM PROJECT, WORKS\_ON

WHERE Pnumber=Pno
GROUPBY Pnumber, Pname
HAVING COUNT (\*) > 2;



Pname	Pnumber		Essn	Pno	Hours		These groups are not sele	
ProductX	1		123456789	1	32.5	П	the HAVING condition of	Q26.
ProductX	1		453453453	1	20.0	∐.	_	
ProductY	2	1	123456789	2	7.5	Ī		
ProductY	2		453453453	2	20.0			
ProductY	2		333445555	2	10.0			
ProductZ	3	1	666884444	3	40.0	Π.	_	
ProductZ	3		333445555	3	10.0	U.	-	
Computerization	10		333445555	10	10.0	П		
Computerization	10		999887777	10	10.0			
Computerization	10		987987987	10	35.0			
Reorganization	20		333445555	20	10.0	П		
Reorganization	20		987654321	20	15.0			
Reorganization	20	1	888665555	20	NULL			
Newbenefits	30	1	987987987	30	5.0			
Newbenefits	30		987654321	30	20.0			
Newbenefits	30		999887777	30	30.0			

After applying the WHERE clause but before applying HAVING

Pname	Pnumber		Essn	Pno	Hours	Pi	name	Count (*)
ProductY	2		123456789	2	7.5	_ ► P	roductY	3
ProductY	2		453453453	2	20.0	-	Computerization	3
ProductY	2		333445555	2	10.0	_ R	eorganization	3
Computerization	10		333445555	10	10.0	_    <b>-</b> N	lewbenefits	3
Computerization	10		999887777	10	10.0		sult of Q26	
Computerization	10		987987987	10	35.0	_   (PT	number not show	/n)
Reorganization	20		333445555	20	10.0	$\neg \parallel$		
Reorganization	20		987654321	20	15.0			
Reorganization	20	1	888665555	20	NULL			
Newbenefits	30		987987987	30	5.0			
Newbenefits	30		987654321	30	20.0			
Newbenefits	30		999887777	30	30.0			

After applying the HAVING clause condition



## Discussion and Summary of SQL Queries

```
SELECT <attribute and function list>
FROM 
[ WHERE <condition> ]
[ GROUP BY <grouping attribute(s)> ]
[ HAVING <group condition> ]
[ ORDER BY <attribute list> ];
```

# Specifying Constraints as Assertions and Actions as Triggers

#### CREATE ASSERTION

Specify additional types of constraints outside scope of built-in relational model constraints

#### CREATE TRIGGER

Specify automatic actions that database system will perform when certain events and conditions occur

## Specifying General Constraints<sup>29</sup> as Assertions in SQL

CREATE ASSERTION

Specify a query that selects any tuples that violate the desired condition

Use only in cases where it is not possible to use CHECK on attributes and domains

```
CREATE ASSERTION SALARY_CONSTRAINT
CHECK ( NOT EXISTS ( SELECT *
```

FROM EMPLOYEE E, EMPLOYEE M,

DEPARTMENT D

WHERE E.Salary>M.Salary

AND E.Dno=D.Dnumber

AND D.Mgr\_ssn=M.Ssn ) );



### Introduction to Triggers in SQL

CREATE TRIGGER statement

Used to monitor the database

Typical trigger has three components:

Event(s)

**Condition** 

**Action** 

```
R5: CREATE TRIGGER SALARY_VIOLATION
BEFORE INSERT OR UPDATE OF SALARY, SUPERVISOR_SSN
ON EMPLOYEE
```

```
FOR EACH ROW
WHEN ( NEW.SALARY > ( SELECT SALARY FROM EMPLOYEE
WHERE SSN = NEW.SUPERVISOR_SSN ) )
INFORM_SUPERVISOR(NEW.Supervisor_ssn,
NEW.Ssn );
```



### Views (Virtual Tables) in SQL

Concept of a view in SQL

Single table derived from other tables

Considered to be a virtual table



### Specification of Views in SQL

#### CREATE VIEW command

Give table name, list of attribute names, and a query to specify the contents of the view

V1: CREATE VIEW WORKS\_ON1

AS SELECT Fname, Lname, Pname, Hours

FROM EMPLOYEE, PROJECT, WORKS\_ON

WHERE Ssn=Essn AND Pno=Pnumber;

V2: CREATE VIEW DEPT\_INFO(Dept\_name, No\_of\_emps, Total\_sal)

AS SELECT Dname, COUNT (\*), SUM (Salary)

FROM DEPARTMENT, EMPLOYEE

WHERE Dnumber=Dno

GROUP BY Dname;



## Specification of Views in SQL (cont'd.)

Specify SQL queries on a view

QV1:

SELECT

Fname, Lname

FROM

WORKS\_ON1

WHERE

Pname='ProductX';

View always up-to-date

Responsibility of the DBMS and not the user

DROP VIEW command

Dispose of a view

V1A:

DROP VIEW

WORKS\_ON1;



## View Implementation, View Update, and Inline Views

Complex problem of efficiently implementing a view for querying

Query modification approach

Modify view query into a query on underlying base tables

Disadvantage: inefficient for views defined via complex queries that are time-consuming to execute

QV1:

SELECT

Fname, Lname WORKS\_ON1

FROM WHERE

Pname='ProductX';

SELECT FROM WHERE Fname, Lname

EMPLOYEE, PROJECT, WORKS\_ON

Ssn=Essn AND Pno=Pnumber

**AND** Pname='ProductX';

### View Implementation

### View materialization approach

Physically create a temporary view table when the view is first queried

Keep that table on the assumption that other queries on the view will follow

Requires efficient strategy for automatically updating the view table when the base tables are updated

### View Implementation (cont'd.)

#### **Incremental update strategies**

DBMS determines what new tuples must be inserted, deleted, or modified in a materialized view table

### View Update and Inline Views

Update on a view defined on a single table without any aggregate functions

Can be mapped to an update on underlying base table

View involving joins

Often not possible for DBMS to determine which of the updates is intended

## View Update and Inline Views<sup>38</sup> (cont'd.)

Clause WITH CHECK OPTION

Must be added at the end of the view definition if a view is to be updated

#### In-line view

Defined in the FROM clause of an SQL query

## Schema Change Statements in SQL

#### Schema evolution commands

Can be done while the database is operational

Does not require recompilation of the database schema



### The DROP Command

DROP command

Used to drop named schema elements, such as tables, domains, or constraint

Drop behavior options:

CASCADE and RESTRICT

Example:

DROP SCHEMA COMPANY CASCADE;

### The ALTER Command

#### Alter table actions include:

Adding or dropping a column (attribute)

Changing a column definition

Adding or dropping table constraints

#### Example:

ALTER TABLE COMPANY.EMPLOYEE ADD COLUMN Job VARCHAR (12);

#### To drop a column

Choose either CASCADE or RESTRICT

ALTER TABLE COMPANY.EMPLOYEE DROP COLUMN Address CASCADE;



### The ALTER Command (cont'd.)

Alter a column definition by dropping an existing default clause by defining a new default clause.

ALTER TABLE COMPANY.DEPARTMENT ALTER COLUMN Mgr\_ssn DROP DEFAULT;

ALTER TABLE COMPANY.DEPARTMENT ALTER COLUMN Mgr\_ssn SET DEFAULT '333445555';

Change constraints specified on a table Add or drop a named constraint

ALTER TABLE COMPANY.EMPLOYEE

DROP CONSTRAINT EMPSUPERFK CASCADE;



### Summary

### Complex SQL:

Nested queries, joined tables, outer joins, aggregate functions, grouping

CREATE ASSERTION and CREATE TRIGGER

**Views** 

Virtual or derived tables

