

Three Triangles Integer Factorization Algorithm

This is my discovery of an algorithm for integer factorization. To my knowledge it is not based on any existing solutions. I do not claim it to be efficient or useful, I'm only concerned with its correctness and completeness.

Abstract

Factor the composite number $C=(a+1)*(a+b)$ by finding the solution to $C=T(a)+T(a+b)-T(b-1)$ where $T(n)$ generates the nth triangular number.

Definitions

- $a, b :=$ natural numbers
- $T(n) := n*(n+1)/2$
- $S(a, b) := T(a)+T(a+b)-T(b-1)$
- $\text{make_even}(n) :=$ if n is odd return $n-1$ else return n

Proof of $(a+1)*(a+b)=T(a)+T(a+b)-T(b-1)$

1. $(a+1)*(a+b)=T(a)+T(a+b)-T(b-1)$
2. $a*a+a*b+a+b=a*(a+1)/2+(a+b)*(a+b+1)/2-(b-1)*b/2$
3. $2*a*a+2*a*b+2*a+2*b=a*(a+1)+(a+b)*(a+b+1)-(b-1)*b$
4. $2*a*a+2*a*b+2*a+2*b=a*a+a+a*(a+b+1)+b*(a+b+1)-b*b+b$
5. $2*a*a+2*a*b+2*a+2*b=a*a+a+a*a*b+a+b*a+b*b+b-b*b+b$
6. $2*a*a+2*a*b+2*a+2*b=2*a*a+2*a*b+2*a+2*b$
7. $0=0$

Proof that $S(a, b) > S(a-2, b+2)$

This is required to show that the algorithm makes progress on step 3.

1. $S(a, b) > S(a-2, b+2)$
2. $(a+1)*(a+b) > (a+1-2)*(a+b-2+2)$
3. $a*a+a+b > (a-1)*(a+b)$
4. $a*a+a+b > a*a-(a+b)$
5. $a+b > -(a+b)$
6. $1*(a+b) > -1*(a+b)$
7. Since $a, b > 0 \Rightarrow 1 > -1$

Proof that $S(a, b) < S(a, b+\max(1, \text{ceil}((C-S(a, b))/(a+1))))$

This is required to show that the algorithm makes progress on step 4.

1. $S(a, b) < S(a, b+\max(1, \text{ceil}((C-S(a, b))/(a+1))))$
2. $(a+1)*(a+b) < (a+1)*(a+b+\max(1, \text{ceil}((C-S(a, b))/(a+1))))$
3. $a+b < a+b+\max(1, \text{ceil}((C-S(a, b))/(a+1))))$
4. $0 < \max(1, \text{ceil}((C-S(a, b))/(a+1))))$
5. $0 < 1$

Algorithm

input: $C \Rightarrow$ integer greater than 2

output: found factors, or 1 and C if C is a prime number

1. let $a = \text{make_even}(\text{floor}(\text{sqrt}(C)) - 1)$
2. let $b = 1$
3. if $S(a, b) > C$ then $a = a - 2, b = b + 2$
4. if $S(a, b) < C$ then $b = b + \max(1, \text{ceil}((C - S(a, b)) / (a + 1)))$
5. if $S(a, b) = C$ then exit: found factors $(a + 1)$ and $(a + b)$
6. if $a = 0$ then exit: C is a prime number
7. goto step 3.

Examples

Factor 51

1. $a = 6, b = 1$
 $S(6, 1) = 49$
 $\Rightarrow b = b + \max(1, \text{ceil}(2 / 7))$
2. $a = 6, b = 2$
 $S(6, 2) = 56$
 $\Rightarrow a = a - 2, b = b + 2$
3. $a = 4, b = 4$
 $S(4, 4) = 40$
 $\Rightarrow b = b + \max(1, \text{ceil}(11 / 5))$
4. $a = 4, b = 7$
 $S(4, 7) = 55$
 $\Rightarrow a = a - 2, b = b + 2$
5. $a = 2, b = 9$
 $S(2, 9) = 33$
 $\Rightarrow b = b + \max(1, \text{ceil}(18 / 3))$
6. $a = 2, b = 15$
 $S(2, 15) = 51$
 $\Rightarrow 51 = 3 * 17$

Factor 23

1. $a = 2, b = 1$
 $S(2, 1) = 9$
 $\Rightarrow b = b + \max(1, \text{ceil}(14 / 3))$
2. $a = 2, b = 6$
 $S(2, 6) = 24$
 $\Rightarrow a = a - 2, b = b + 2$
3. $a = 0, b = 8$
 $\Rightarrow 23 = 1 * 23$

Conclusion

Instead of focusing on performance of factorization this algorithm tries to explore a new solution. One potential benefit I see is the reduced magnitude of dividends used in divisions. Hope it inspires some ideas.