ASSIGNMENT AIR-4

Roll No: 41205

Problem Statement:

Implement crypt-arithmetic problem or n-queens or graph coloring problem (Branch and Bound and Backtracking)

Objective:

- 1. Understand backtracking and branch and bound algorithms
- 2. Apply the algorithms for popular problems such as n-queens

Outcome: One will be able to apply different techniques to solve the n-queens problem

Pre-requisites:

- 1. 64-bit Linux OS
- 2. Programming Languages: Python

Hardware Specification:

- 1. x86 64 bit
- 2. 2/4 GB DDR RAM
- 3. 80 500 GB SATA HD
- 4. 1GB NIDIA TITAN X Graphics Card

Software Specification:

1. Ubuntu 14.04

Theory:

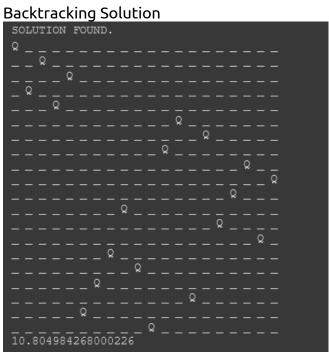
- The N Queen is the problem of placing N chess queens on an N×N chessboard so that no two queens attack each other.
- The expected output is a binary matrix which has 1s for the blocks where queens are placed. For example, following is the output matrix for above 4 queen solution.
 - {0, 1, 0, 0}
 - {0, 0, 0, 1}
 - {1, 0, 0, 0}
 - { 0, 0, 1, 0}
- Backtracking Algorithm
- The idea is to place queens one by one in different columns, starting from the leftmost column. When we place a queen in a column, we check for clashes with already placed queens. In the current column, if we find a row for which there is no clash, we mark this row and column as part of the solution. If we do not find such a row due to clashes, then we backtrack and return false.

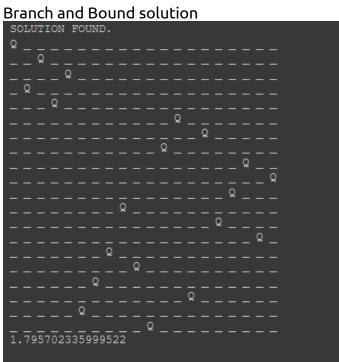
- 1. Start in the leftmost column
- 2. If all queens are placed
 - return true
- 3. Try all rows in the current column.
- 4. Do following for every tried row.
 - a) If the queen can be placed safely in this row then mark this [row, column] as part of the solution and recursively check if placing queen here leads to a solution.
 - b) If placing the queen in [row, column] leads to a solution then return true.
 - c) If placing queen doesn't lead to a solution then unmark this [row, column] (Backtrack) and go to step (a) to try other rows.
- 5. If all rows have been tried and nothing worked,
- 6. Return false to trigger backtracking.
- For branch and bound, we maintain two additional vectors that indicate the diagonals that have been blocked.
- The lookup of any diagonal happens in O(1) time.
- This makes it much quicker for us to check if the position is valid or invalid.
- Thus, the branch and bound solution runs much faster than the backtracking one.

Test Cases:

#	Input	Expected Output	Actual Output	Result
1	N = 8	Solved successfully Branch and bound faster than backtracking	Solved successfully Backtracking: 0.007 secs Branch and Bound: 0.002 secs	Success
2	N = 20	Solved successfully Branch and bound faster than backtracking	Solved successfully Backtracking: 10.8 secs Branch and Bound: 1.8 secs	Success

Output:





Conclusion: Thus, we were able to solve the n-queens problem using both backtracking and branch and bound algorithms.